

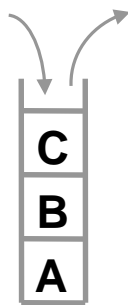
Positional Lists

Outline and Required Reading:

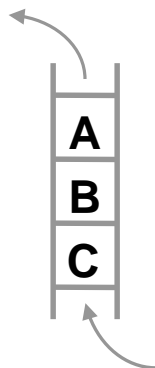
- **Positional Lists (§ 7.3)**

- Lists** – ‘linear container’ which allows direct access to any of its elements
- items can be accessed either through **rank / index** or **position relative to the position of other items in the list**

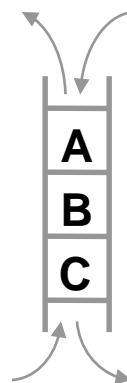
Stacks, Queues, Deques – restricted lists with methods for accessing, inserting, and removing only the first and/or last element



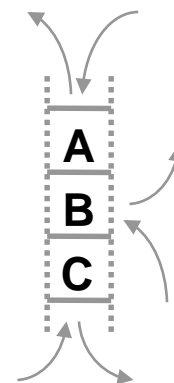
Stack



Queue



Deque



List

ArrayList ADT

```
public interface ArrayList<E> {
```

```
    int size();
```

```
    /* return the # of objects in this list */
```

```
    boolean isEmpty();
```

```
    /* return true if the list is empty */
```

```
    E get(int k) throws IndexOutOfBoundsException;
```

```
    /* return the element at rank k without removing it*/
```

```
    /* error if k<0 or k≥size()=n – current # of elements */
```

```
    E set(int k, E e) throws IndexOutOfBoundsException...;
```

```
    /* replace with e elem. at rank k; return replaced element */
```

```
    /* error if k<0 or k≥size()=n – current # of elements */
```

```
    void add(int k, E e) throws IndexOutOfBoundsException...;
```

```
    /* insert a new element e into list at rank k */
```

```
    /* error if k<0 or k>size()=n – current # of elements */
```

```
    /* rank of all subsequent elements will increase! */
```

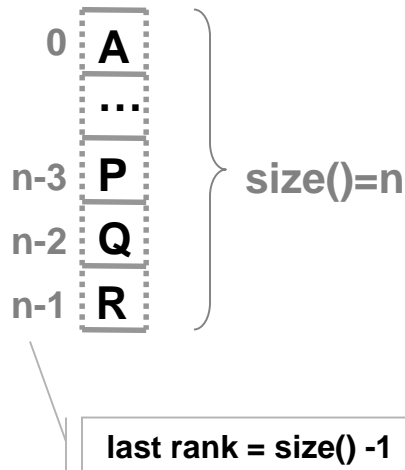
```
    E remove(int k) throws IndexOutOfBoundsException...;
```

```
    /* remove and return the element at rank k */
```

```
    /* error if k<0 or k≥size()=n – current # of elements */
```

```
    /* rank of all subsequent elements will decrease! */
```

```
}
```



Run Times in Array Implementation

Method	Time
size	$O(1)$
isEmpty	$O(1)$
get	$O(1)$
set	$O(1)$
add	$O(n)$
remove	$O(n)$

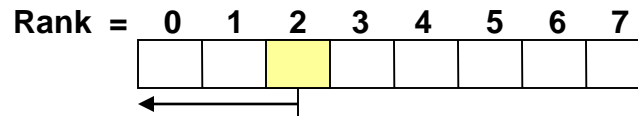
Run Times in Linked List Implementation

Method	Time
size	$O(1)$
isEmpty	$O(1)$
get	$O(n)$
set	$O(n)$
add	$O(n)$
remove	$O(n)$

Array List vs. Positional Lists

Array List – ADT that employs “**sequential allocation**”

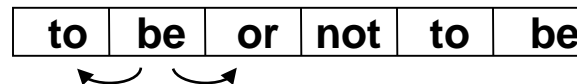
- elements are identified by their rank / index (sequence #)
- no notion of spatial relation among elements, except through rank



2 elements “before”, but their identity is not known

Positional List – ADT that employs “**position allocation**”

- each position can be used to refer to next / preceding element
- each position can be accessed through one of its neighbors



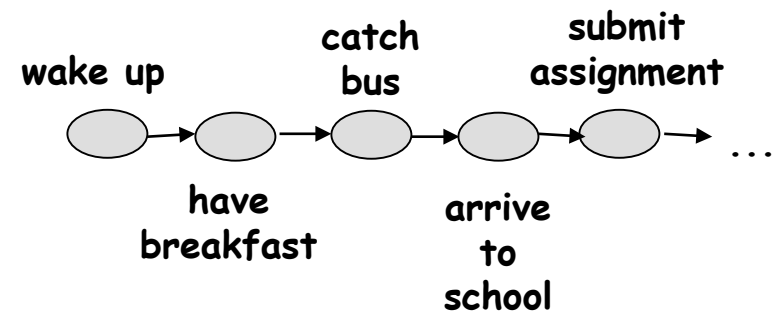
“be” is after “to”, and before “or”

Positional List Application

- whenever it is required that insert / remove run in $O(1)$ time – e.g. Text Editor: insert/remove letters at a cursor

Example [Array vs. Positional List Application – building database]

Group	Name	Login
	Bui, Natalie	NB1511
	Brock, Jayden	JB1511
	Burt, Jasmine	JB1511
	Dibble, Jack	JD1511
	Dixon, Jesse	JD1511
	Dye, Frank	FD1511



Array List example: database of unrelated items - e.g. list of students organized according to their student numbers. The neighbors in the list are not directly related.

Positional List example: database of related items - e.g. list of daily activities. Activities 'next to each other' in the list occur 'next to each other' in real life. New activity B is typically added/referred to as 'activity that happens after A'.

Positional List ADT – abstraction of a linked list with “shielded” internal structure

Positional List ADT Implementation – linked list (SLL or DLL) is the natural choice for implementation of Node List ADT

- however, **direct use of “node-based” operations should be avoided**, because:
 - (1) we should not expose too much information about the implementation of the list
 - (2) **we should not burden the end user with too much implementation details** – e.g what to do with “links” on every insert / removal

SOLUTION: Use **Position ADT** – a helper ADT!

public interface PositionalList<E> {

Generic Methods

public int **size**();
public boolean **isEmpty**();

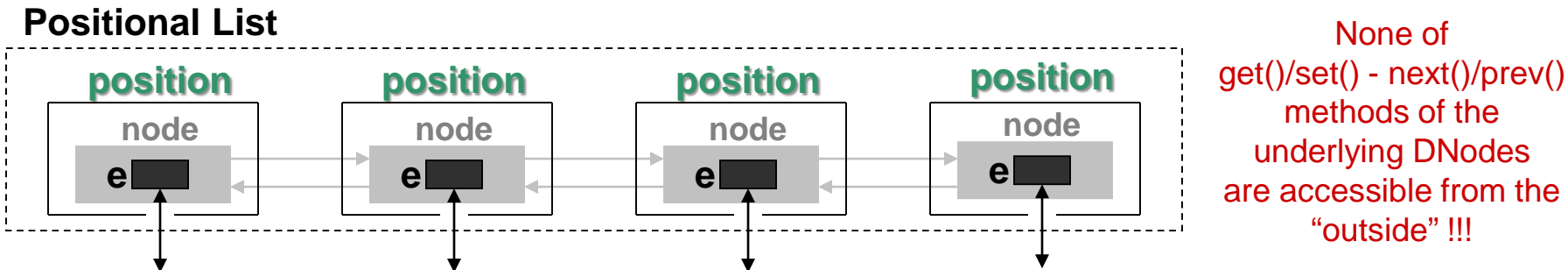
Accessor Methods

public Position<E> **first**();
public Position<E> **last**();
public Position<E> **before**(Position<E> p) throws ...;
public Position<E> **after**(Position<E> p) throws ...;

Update Methods

public Position<E> **addBefore**(Position<E> p, E e) ...;
public Position<E> **addAfter**(Position<E> p, E e) ...;
public Position<E> **addFirst**(E e);
public Position<E> **addLast**(E e);
public E **remove**(Position<E> p) throws ...;
public E **set**(Position<E> p, E e) throws ...; }

- Position ADT** – encapsulates the idea of “node” in a linked list
- but, has only one public method *public E **getElement()*** which returns element stored at the given position



For outside user, Positional List is viewed as a container of elements, which stores each element at/inside a position, and keeps positions arranged in a linear order.

NOTE: Positions are defined relatively to their neighbours.
The ‘relative’ neighbourhood of *p* does not change even if we replace or swap the element *e* stored at *p* with another element.
A position associated with element *e* does NOT change even if the rank of *e* changes in *S*. (**Positions are NOT tied to rank!**)

```
interface Position<E> {
    E getElement() throws IllegalStateException...; }
```

Wrapper class interface that restricts access to the underlying DLLNode.

```
public static class Node<E> implements Position<E> {
```

```
    private E element;
    private Node<E> prev;
    private Node<E> next;
```

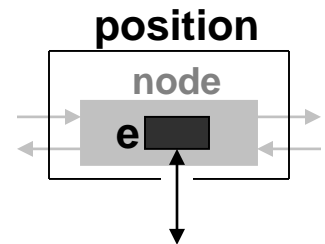
```
    public Node(E e, Node<E> p, Node<E> n) {
        element = e;
        prev = p;
        next = n; }
```

```
    public E getElement() throws IllegalStateException... {
        if (next == null) // convention for defunct node
            throw new IllegalStateException("Position no longer valid");
        return element; }
```

```
    ...
```

```
    // getPrev(), getNext(), setPrev(..), setNext(..), setElement(..)
```

```
}
```

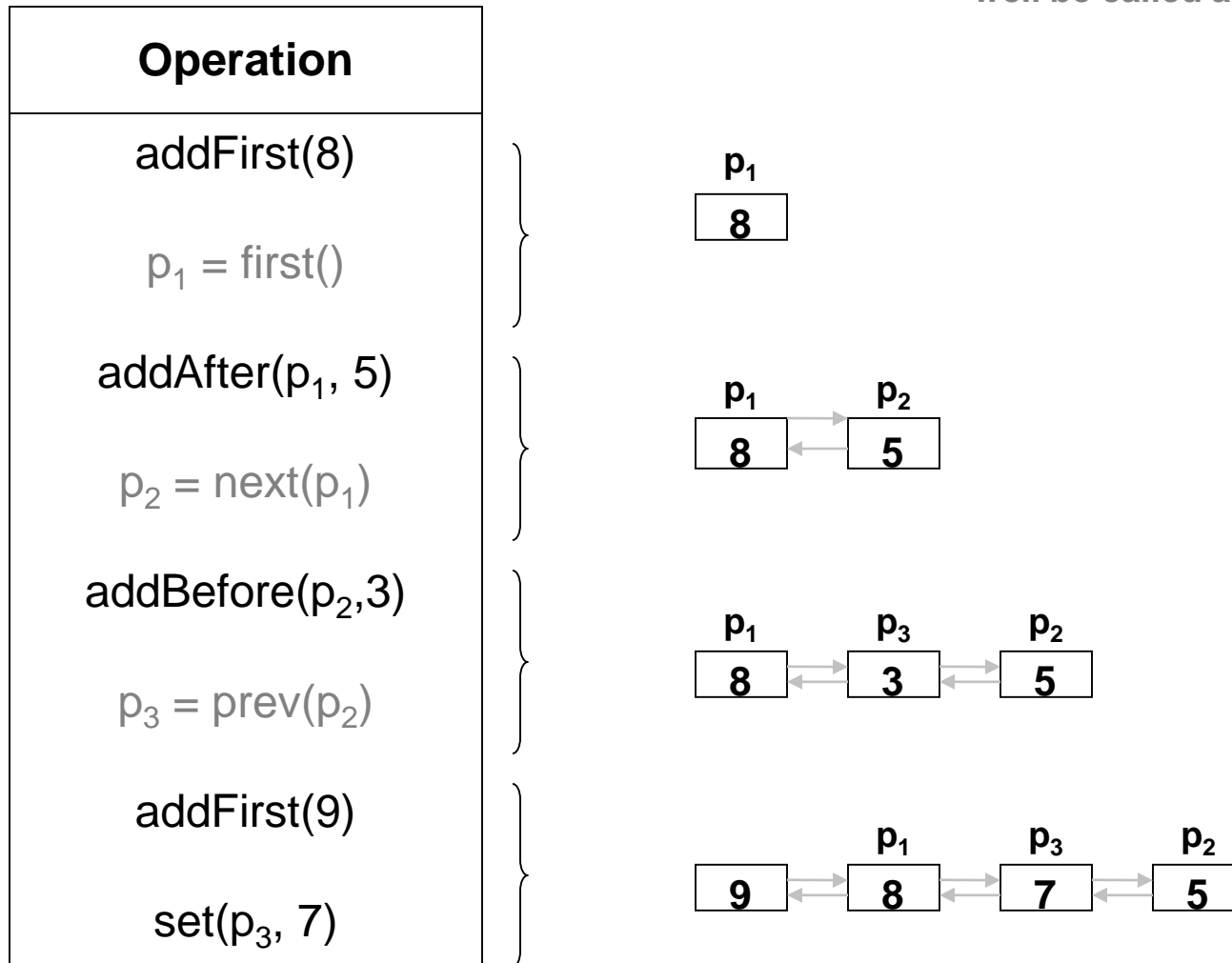


Example [operations on a List]

Method	Return Value	List Contents
addLast(8)	p	(8_p)
first()	p	(8_p)
addAfter(p , 5)	q	$(8_p, 5_q)$
before(q)	p	$(8_p, 5_q)$
addBefore(q , 3)	r	$(8_p, 3_r, 5_q)$
r .getElement()	3	$(8_p, 3_r, 5_q)$
after(p)	r	$(8_p, 3_r, 5_q)$
before(p)	null	$(8_p, 3_r, 5_q)$
addFirst(9)	s	$(9_s, 8_p, 3_r, 5_q)$
remove(last())	5	$(9_s, 8_p, 3_r)$
set(p , 7)	8	$(9_s, 7_p, 3_r)$
remove(q)	“error”	$(9_s, 7_p, 3_r)$

Example [operations on a List]

p_1, p_2, p_3, \dots could as well be called a, b, c, d



Positional List ADT: Doubly Linked List Implementation¹³

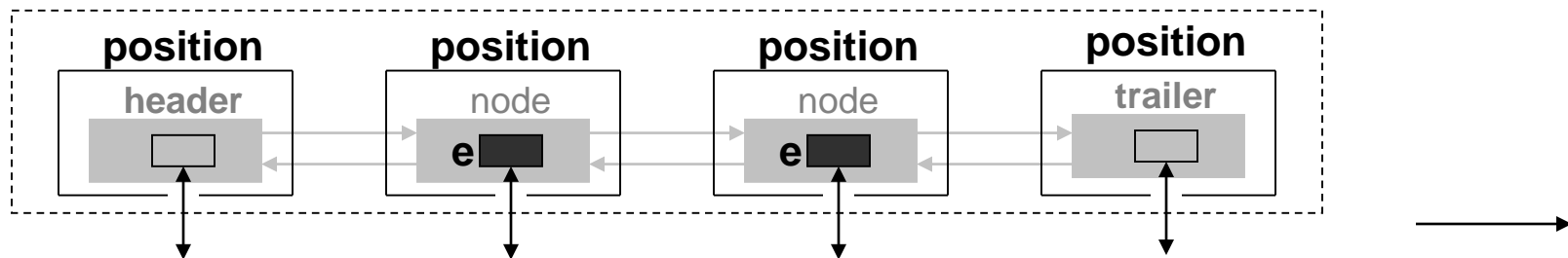
DLL Implementation of Positional List ADT

- DLL Nodes are viewed internally by List as ‘nodes’, from outside they are viewed only as positions
- through **casting**, given a position p, we can “unwrap” p to reveal the underlying node v

```
public class LinkedPositionalList<E> implements PositionalList<E> {
```

```
    private int size = 0;  
    private Node<E> header, trailer;
```

```
    public LinkedPositionalList() {  
        header = new Node<E>(null, null, null);  
        trailer = new Node<E>(null, header, null);  
        header.setNext(trailer); }  
}
```



Example Methods

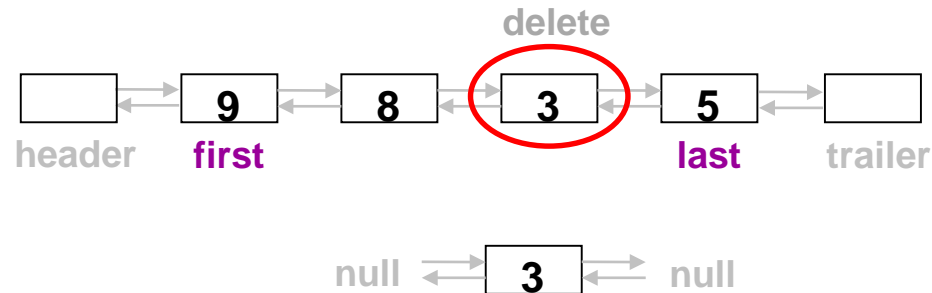
/ checks if position is valid for this list, and if valid converts it to DNode */*

```
private Node<E> validate(Position<E> p) throws IllegalArgumentException. {
    if ( !(p instanceof Node) ) throw new
        IllegalArgumentException("Invalid p.")

    Node<E> node = (Node<E>) p;

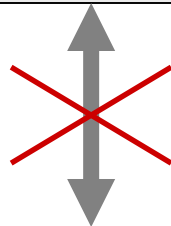
    if ( node.getNext() == null ) throw new
        IllegalArgumentException("p is no longer in the list.")

    return node;
}
```



```
package myList

public class NodePositionalList<E> {
    ...
    private DNode<E> validate(Position<E> p) {
        ...
    }
    ...
}
```



L.validate(p) cannot
be run outside its own
package

```
import myList.NodePositionalList<E>;

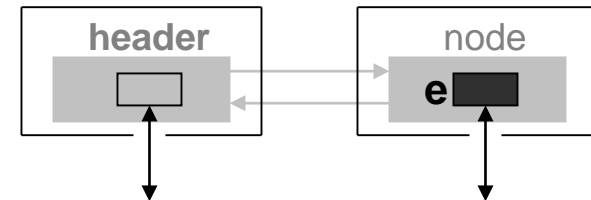
public class ListApplication<E> {
    NodePositionalList<E> L;
    ...
}
```

/ returns the given node as a Position or null, if it is a sentinel */*

```
private Position<E> position(Node<E> node) {  
    if (node == header || node == trailer) return null;  
    return node; }
```

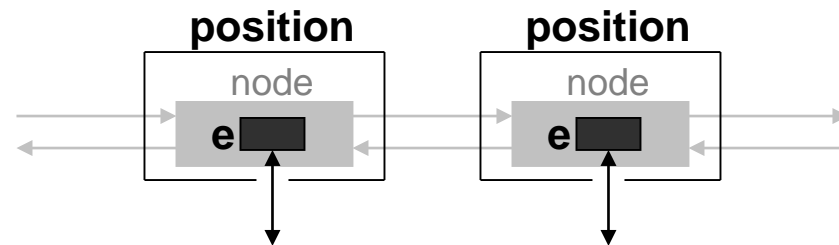
/ returns the first Position in the list or null if empty */*

```
public Position<E> first() {  
    return position(header.getNext()); }
```



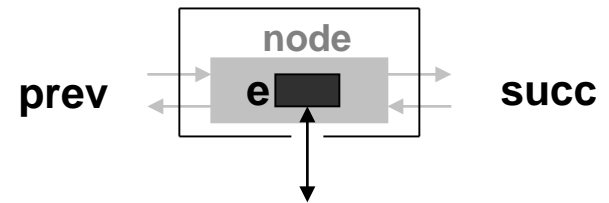
/ returns Position immediately before Position p or null if p is first */*

```
public Position<E> before(Position<E> p) throws IllegalArgumentException {  
    Node<E> node = validate(p);  
    return position(node.getPrev()); }
```



/* add element e to the linked list between the given nodes */

```
private Position<E> addBetween(E e, Node<E> prev, Node<E> succ) {  
    Node<E> newest = new Node<>(e, prev, succ);  
    prev.setNext(newest);  
    succ.setPrev(newest);  
    size++;  
    return newest;  
}
```



/* inserts element e at the front of the linked list and returns ... */

```
public Position<E> addFirst(E e) {  
    return addBetween(e, header, header.getNext());  
}
```

/ inserts element e immediately before Position p and returns ... */*

```
public Position<E> addBefore(Position<E> p, E e) throws IllegalArg.. {  
    Node<E> node = validate(p);  
    return addBetween(e, node.getPrev(), node);  
}
```

/ replaces element stored at Position p and returns replaced element */*

```
public E set(Position<E> p, E e) throws IllegalArg.. {  
    Node<E> node = validate(p);  
    E answer = node.getElement(p);  
    node.setElement(e);  
    return answer;  
}
```

/ removes element e stored at p, returns e, and invalidates p */*

```
public E remove(Position<E> p) throws IllegalArgumentException {
```

```
    Node<E> node = validate(p);
```

```
    Node <E> pred = node.getPrev();
```

```
    Node <E> succ = node.getNext();
```

```
    pred.setNext(succ);
```

```
    succ.setPrev(pred);
```

```
    size --;
```

```
    E answer = node.getElement();
```

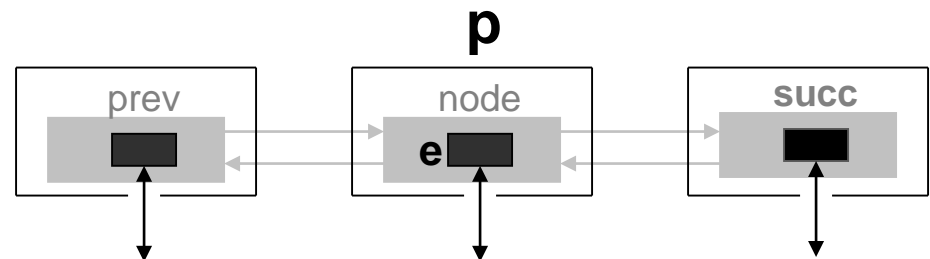
```
    node.setElement(null);
```

```
    node.setNext(null);
```

```
    node.setPrev(null);
```

```
    return answer;
```

```
}
```



Run Time – Good! all methods run in constant $O(1)$ time

List Method	Time
size	$O(1)$
first	$O(1)$
last	$O(1)$
before	$O(1)$
after	$O(1)$
addBefore	$O(1)$
addAfter	$O(1)$
...	...

NOTE:

All methods run in $O(1)$ time since they assume that the reference to the Position in question is given – no searching through DLL is required.