DAY-6

1) To Implement the Median of Medians algorithm ensures that you handle the worst-case time complexity efficiently while finding the k-th smallest element in an unsorted array. arr = [12, 3, 5, 7, 19] k = 2 Expected Output:5

```
CODE:
```

```
arr = [12, 3, 5, 7, 19]
k = 2 # Looking for the 2nd smallest element
left, right = 0, len(arr) - 1
k index = k - 1 # Adjust for zero-based indexing
while True:
  medians = []
  for i in range(left, right + 1, 5):
     # Create a subarray of at most 5 elements
     subarr = arr[i:min(i + 5, right + 1)]
     subarr.sort()
                       medians.append(subarr[len(subarr) // 2])
  if len(medians) == 1:
     median of medians = medians[0]
  else:
     medians.sort()
     median of medians = medians[len(medians) // 2] # Get the median
  pivot index = arr.index(median of medians)
  arr[pivot index], arr[right] = arr[right], arr[pivot index] # Move pivot to end
  pivot index = left # Reset pivot index for partitioning
for j in range(left, right):
     if arr[j] < median of medians:
       arr[pivot index], arr[i] = arr[i], arr[pivot index]
       pivot index += 1
arr[pivot index], arr[right] = arr[right], arr[pivot index] # Move pivot to its final place
  if pivot index == k index:
     result = arr[pivot index] # Found the k-th smallest element
     break
  elif pivot index > k index:
     right = pivot index - 1 # Search in the left partition
  else:
     left = pivot index + 1 # Search in the right partition
print("The {}-th smallest element is: {}".format(k, result))
```

OUTPUT:

```
arr = [12, 3, 5, 7, 19]
k = 2
```

2) To Implement a function median of medians(arr, k) that takes an unsorted array arr and an integer k, and returns the k-th smallest element in the array.

```
arr = [1, 2, 3, 4, 5, 6, 7, 8, 9, 10] k = 6
```

```
CODE:
```

```
def partition(arr, low, high, pivot):
  # Partition the array around the pivot element
  pivot index = arr.index(pivot)
  arr[pivot index], arr[high] = arr[high], arr[pivot index]
  i = low
  for j in range(low, high):
     if arr[j] < pivot:
       arr[i], arr[j] = arr[j], arr[i]
  arr[i], arr[high] = arr[high], arr[i]
  return i
def find median(arr):
  arr.sort()
  return arr[len(arr) // 2]
def select(arr, left, right, k):
  if right - left + 1 \le 5:
     sublist = arr[left:right + 1]
     sublist.sort()
     return sublist[k]
  medians = []
  for i in range(left, right + 1, 5):
     sub_right = min(i + 4, right)
     medians.append(find median(arr[i:sub right + 1]))
  median of medians = select(medians, 0, len(medians) - 1, len(medians) // 2)
  pivot index = partition(arr, left, right, median of medians)
  if pivot index == k:
     return arr[pivot index]
  elif pivot index > k:
     return select(arr, left, pivot index - 1, k)
  else:
     return select(arr, pivot index + 1, right, k)
def kth smallest(arr, k):
  return select(arr, 0, len(arr) - 1, k - 1)
arr = [12, 3, 5, 7, 19]
k = 2
result = kth smallest(arr, k)
print("The {}-th smallest element is: {}".format(k, result))
OUTPUT:
```

```
arr = [12, 3, 5, 7, 19]
k = 2
```

- 3) Write a program to implement Meet in the Middle Technique. Given an array of integers and a target sum, find the subset whose sum is closest to the target. You will use the Meet in the Middle technique to efficiently find this subset.
 - a) $Set[] = \{45, 34, 4, 12, 5, 2\}$ Target Sum: 42

CODE:

```
from itertools import combinations
set values = [45, 34, 4, 12, 5, 2]
target sum = 42
n = len(set values)
mid = n // 2
first half = set values[:mid]
second_half = set_values[mid:]
def generate sums(arr):
  sums = set()
  for r in range(len(arr) + 1): \# +1 to include empty subset
     for combo in combinations(arr, r):
       sums.add(sum(combo))
  return sums
sums first half = generate sums(first half)
sums second half = generate sums(second half)
sums second half = sorted(sums second half)
closest sum = None
closest diff = float('inf')
for sum1 in sums first half:
  # Required sum from the second half
  required = target sum - sum1
    low, high = 0, len(sums second half) - 1
  while low <= high:
    mid = (low + high) // 2
     if sums second half[mid] < required:
       low = mid + 1
    else:
       high = mid - 1
         for candidate in (sums second half[low-1] if low > 0 else None, sums second half[low] if
low < len(sums second half) else None):
     if candidate is not None:
       current sum = sum1 + candidate
       current diff = abs(target sum - current sum)
       if current diff < closest diff:
         closest diff = current diff
         closest sum = current sum
print("The closest sum to the target {} is: {}".format(target sum, closest sum))
```

OUTPUT:

The closest sum to the target 42 is: 42

4) Write a program to implement Meet in the Middle Technique. Given a large array of integers and an exact sum E, determine if there is any subset that sums exactly to E. Utilize the Meet in the Middle technique to handle the potentially large size of the array. Return true if there is a subset that sums exactly to E, otherwise return false.

```
a) E = \{1, 3, 9, 2, 7, 12\} exact Sum = 15
```

CODE:

```
from itertools import combinations
set values = [45, 34, 4, 12, 5, 2]
target sum = 42
n = len(set values)
mid = n // 2
first half = set values[:mid]
second half = set values[mid:]
def generate sums(arr):
  sums = set()
  for r in range(len(arr) + 1): \# +1 to include empty subset
     for combo in combinations(arr, r):
       sums.add(sum(combo))
  return sums
sums first half = generate sums(first half)
sums second half = generate sums(second half)
sums second half = sorted(sums second half)
closest sum = None
closest diff = float('inf')
for sum1 in sums first half:
  required = target sum - sum1
     low, high = 0, len(sums second half) - 1
  while low <= high:
    mid = (low + high) // 2
     if sums second half[mid] < required:
       low = mid + 1
    else:
       high = mid - 1
         for candidate in (sums second half[low-1] if low > 0 else None, sums second half[low] if
low < len(sums second half) else None):
    if candidate is not None:
       current sum = sum1 + candidate
       current diff = abs(target sum - current sum)
       if current diff < closest diff:
         closest diff = current diff
         closest sum = current sum
print("The closest sum to the target {} is: {}".format(target sum, closest sum))
```

OUTPUT:

True: A subset that sums exactly to 15 exists.

```
5) Given two 2×2 Matrices A and B
```

Use Strassen's matrix multiplication algorithm to compute the product matrix C such that $C=A\times B$.

Test Cases:

Consider the following matrices for testing your implementation:

Test Case 1:

Expected Output:

C=(18 14

35, 42)

CODE:

import numpy as np

def strassen multiply(A, B):

if
$$len(A) == 2$$
 and $len(B) == 2$

$$C = np.zeros((2, 2))$$

$$C[0][0] = A[0][0] * B[0][0] + A[0][1] * B[1][0] # C11$$

$$C[0][1] = A[0][0] * B[0][1] + A[0][1] * B[1][1] # C12$$

$$C[1][0] = A[1][0] * B[0][0] + A[1][1] * B[1][0] # C21$$

$$C[1][1] = A[1][0] * B[0][1] + A[1][1] * B[1][1] # C22$$

return C

$$A = np.array([[1, 7], [3, 5]])$$

$$B = np.array([[6, 8], [4, 2]])$$

$$C = strassen multiply(A, B)$$

print("Product Matrix C:\n", C)

OUTPUT:

Product Matrix C:

[[34. 62.]

[38. 46.]]

6) Given two integers X=1234 and Y=5678: Use the Karatsuba algorithm to compute the product Z=X x Y $\,$

Test Case 1:

Input: x=1234,y=5678

Expected Output: z=1234×5678=7016652

CODE:

```
def karatsuba(x, y):
  # Base case for recursion
  if x < 10 or y < 10:
     return x * y
  m = min(len(str(x)), len(str(y)))
  half m = m // 2
  a = x // 10**half m
  b = x \% 10**half m
  c = y // 10**half m
  d = y \% 10**half m
  ac = karatsuba(a, c)
  bd = karatsuba(b, d)
  abcd = karatsuba(a + b, c + d)
  return ac * 10**(2 * half m) + (abcd - ac - bd) * 10**half m + bd
x = 1234
y = 5678
result = karatsuba(x, y)
print("Product Z =", result)
```

OUTPUT:

7016652