



# **OBJECT ORIENTED ANALYSIS & DESIGN**

# **DATA STRUCTURES & ALGORITHMS**

## Introduction to Data Structures

# Description of various Data Structures



# Arrays

- An array is defined as a set of finite number of homogeneous elements or same data items.
- It means an array can contain one type of data only, either all integer, all float-point number or all character.

# Arrays

- Simply, declaration of array is as follows:  
`int arr[10]`
- Where int specifies the data type or type of elements arrays stores.
- “arr” is the name of array & the number specified inside the square brackets is the number of elements an array can store, this is also called sized or length of array.

# Arrays

Following are some of the concepts to be remembered about arrays:

- The individual element of an array can be accessed by specifying name of the array, following by index or subscript inside square brackets.
- The first element of the array has index zero[0]. It means the first element and last element will be specified as:arr[0] & arr[9] Respectively.
- The elements of array will always be stored in the consecutive (continues) memory location.
- The number of elements that can be stored in an array, that is the size of array or its length is given by the following equation:  
 $(\text{Upperbound} - \text{lowerbound}) + 1$

## Arrays (Contd.)

- For the above array it would be  $(9-0)+1=10$ , where 0 is the lower bound of array and 9 is the upper bound of array.
- Array can always be read or written through loop. If we read a one-dimensional array it require one loop for reading and other for writing the array.



# Arrays

**For example: Reading an array**

```
For(i=0;i<=9;i++)  
    scanf("%d",&arr[i]);
```

**For example: Writing an array**

```
For(i=0;i<=9;i++)  
    printf("%d",arr[i]);
```



# Arrays

- If we are reading or writing two-dimensional array it would require two loops. And similarly the array of a N dimension would required N loops.
- **Some common operation performed on array are:**
  - ✓ Creation of an array
  - ✓ Traversing an array



# Arrays

- Insertion of new element
- Deletion of required element
- Modification of an element
- Merging of arrays



# Array Usage – A Perspective

Consider the following example:

```
#include<stdio.h>
main( )
{
    int num_array[50],i;
    for( i=0; i < 50; i++ )
    {
        num_array[i]=0;
        scanf( "%d",&num_array[i] );
        fflush(stdin);
    }
}
```



## Array Usage – A Perspective

- When this program is compiled, the compiler estimates the amount of memory required for the variables, and also the instructions defined by you as part of the program.
- The compiler writes this information into the header of the executable file that it creates. When the executable is loaded into memory at runtime, the specified amount of memory is set aside.
- A part of the memory allocated to the program is in an area of memory called a runtime stack or the call stack. Once the size of the stack is fixed, it cannot be changed dynamically.

## Array Usage – A Perspective

- Therefore, arrays present the classic problem of the programmer having allocated too few elements in the array at the time of writing the program, and then finding at run time that more values are required to be stored in the array than what had originally been defined.
- The other extreme is of the programmer allocating too many elements in the array and then finding at run time that not many values need to be stored in the array thereby resulting in wastage of precious memory.

## Array Usage – A Perspective

- Moreover, array manipulation (in terms of insertion and deletion of elements from the array) is more complex and tedious, and a better alternative to all this is to go for dynamic data structures.
- Before venturing into dynamic variables, or dynamic data structures, it would be prudent at this juncture to differentiate between stack and heap variables, and their characteristics.
- Let us begin by understanding the concept of lifetime of variables.

# Lifetime Of A Variable

- is a run-time concept
- period of time during which a variable has memory space associated with it
  - ✓ begins when space is allocated
  - ✓ ends when space is de-allocated
- three categories of "lifetime"
  - ✓ static - start to end of program execution
  - ✓ automatic (stack) - start to end of declaring function's execution
  - ✓ heap (variable declared dynamic at runtime, and also de-allocated dynamically at runtime).





# Java Strings

With Strings in Java you can perform various operations, some of which are:

Search

The quick brown fox jumps over the lazy dog

Create Substring

The quick brown fox jumps over the lazy dog

Create new strings

The quick brown fox jumps over the lazy dog

# Linked Lists



# Self-Referential Structures

- Suppose, you have been given a task to store a list of marks. The size of the list is not known.
- If it were known, then it would have facilitated the creation of an array of the said number of elements and have the marks entered into it.
- Elements of an array are contiguously located, and therefore, array manipulation is easy using an integer variable as a subscript, or using pointer arithmetic.
- However, when runtime variables of a particular type are declared on the heap, let's say a structure type in which we are going to store the marks, each variable of the structure type marks will be located at a different memory location on the heap, and not contiguously located.



# Self-Referential Structures

- Therefore, these variables cannot be processed the way arrays are processed, i.e., using a subscript, or using pointer arithmetic.
- An answer to this is a self-referential structure.



# Self-Referential Structures

A self-referential structure is so defined that one of the elements of the structure variable is able to reference another subsequent structure variable of the same type, wherever it may be located on the heap.

In other words, each variable maintains a link to another variable of the same type, thus forming a non-contiguous, loosely linked data structure.

This self-referential data structure is also called a linked list.

# Declaring a Linked List

Let us define a self-referential structure to store a list of marks the size of which may not be known.

```
struct marks_list
{
    int marks;
    struct marks_list *next;
};
```

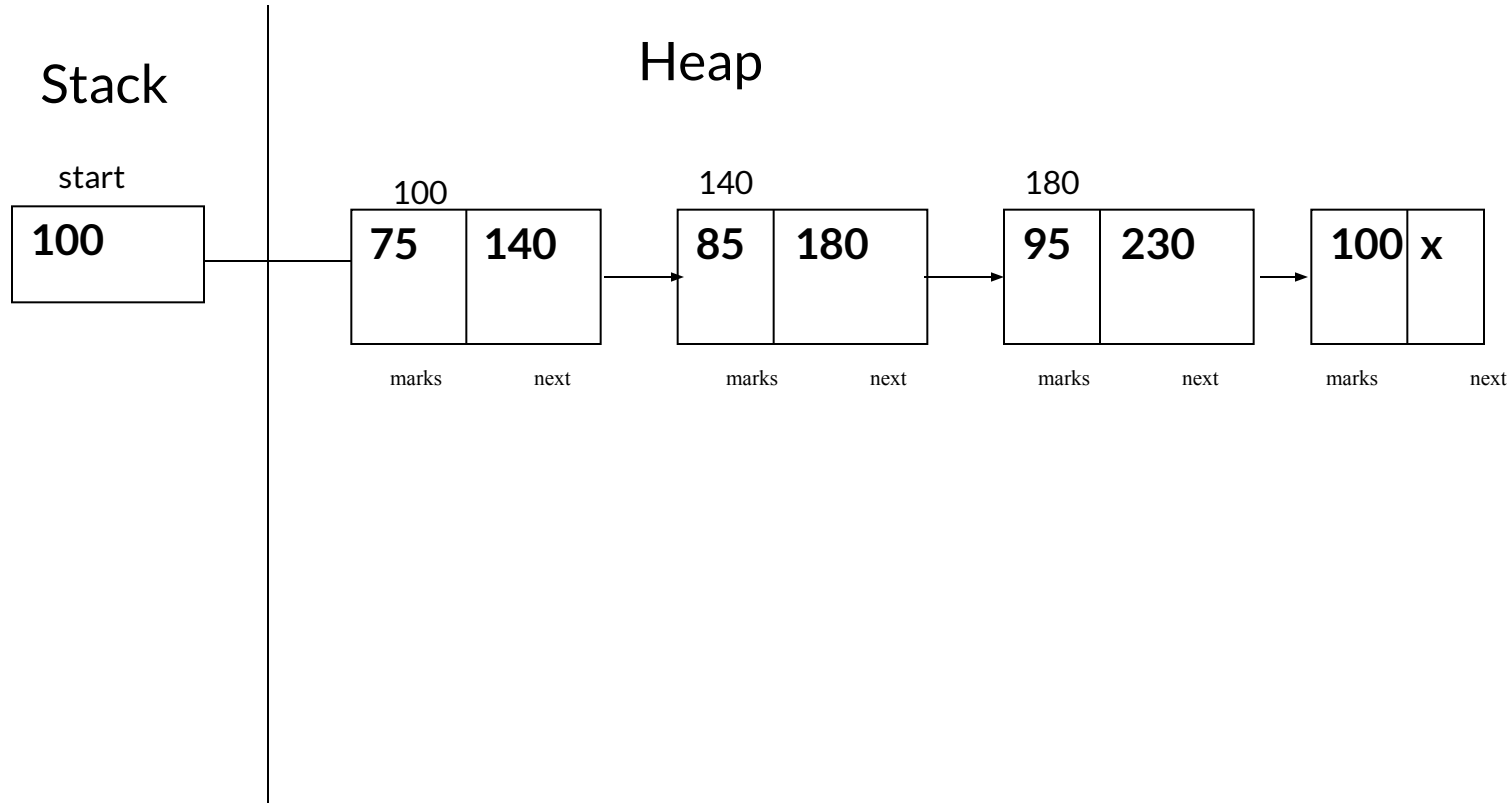
We have defined a structure of type marks\_list. It consists of two elements, one integer element marks and the other element, a pointer next, which is a pointer to a structure of the same type, i.e., of type marks\_list itself.



## Declaring a Linked List

- Therefore, a part of the structure is referencing a structure type of itself, and hence the name **self-referential structure**.
- Such data structures are also popularly known as **linked lists**, since each structure variable contains a link to other structure variables of the same type.
- One can visualize a linked list as shown in the following slide:

# Visualizing a Linked List



# Creating a Sorted Linked List

```
#include<stdio.h>
#include<malloc.h>
struct marks_list *start, *prev;
/* variables declared outside main() are global in nature and can be accessed by other
functions called from main() */
struct marks_list
{
    int marks;
    struct marks_list *next;
};
main()
{
    struct marks_list * makenode();
    /* function prototype declaration */
```

# Creating a Sorted Linked List

```
struct marks_list *new;  
start = NULL;  
char menu = '0';  
while (menu != '4')  
{  
    printf( "Add Nodes  :\n");  
    printf( "Delete Nodes :\n");  
    printf( "Traverse a list :\n");  
    printf( "Exit      :\n");  
    menu = getchar( );
```



# Creating a Sorted Linked List

```
switch (menu)
{
    case '1': addnode( );
                break;
    case '2': deletenode( )
                break;
    case '3': traverse( );
                break;
    case '4': exit( );
                break;
}/* end of switch */
}/* end of main() */
```



# Creating a Sorted Linked List

```
addnode()  
{  
    char ch = 'y';  
    while ( ch == 'y' )  
    {  
        new = makenode();  
        /* creation of a list is treated as a special case of insertion */  
        if ( start == NULL )  
        {  
            start = new;  
        }  
        else  
        {  
            insert();  
            traverse( );  
        }  
    }
```





## Creating a Sorted Linked List

```
printf("%s","Want to add more nodes\n");  
scanf( "%c", &ch );  
fflush( stdin );  
/* end of while */  
/* end of addnode( )
```



## Creating a Sorted Linked List

```
struct marks_list * makenode()
{
    struct marks_list *new;
    new=(struct marks_list *) malloc(sizeof(struct(marks_list)));
    scanf("%d",&new->marks);
    new->next = NULL;
    return(new);
}
```

# Creating a Sorted Linked List

```
insert(struct marks_list *start)
{
    struct marks_list *ptr, *prev;
    for(ptr=start,prev=start;(ptr);prev=ptr,ptr=ptr->next)
    {
        if (new->marks < start->marks)
        {
            /* insertion at the beginning of a list */
            new->next = start;
            start = new;
        }
    }
```

## Creating a Sorted Linked List

```
/* insertion in the middle of a list */  
if(new->marks > ptr->marks)  
{  
    continue;  
}  
else  
{  
    prev->next = new;  
    new->next = ptr;  
}  
}/* end of for loop */
```

## Creating a Sorted Linked List

```
/* insertion at the end of the list */  
if (ptr == null)  
{  
    prev->next = new;  
    new->next = null;  
}  
}/* end of insert */
```

## Searching a Value in a Linked List

```
struct marks_list *search( int val)
{
    for( ptr = start; (ptr); ptr = ptr->next)
    {
        if (val == ptr-> marks)
            return ptr;
    }
}
```



# Deleting a Node From a Linked List

```
delete ()
{
    struct marks_list *ptr, *prev, *temp;
    int score;
    /* search the linked list for the value to be deleted */
    scanf("%d", &score);
    fflush(stdin);
    for (ptr = start, prev = start; (ptr); prev = ptr, ptr = ptr->next)
    {
        /* deletion of the first node in the list */
        if (score == start->marks)
        {
            temp = start;
            start = start->next;
            free(temp);
        }
    }
```



## Deleting a Node From a Linked List

```
/* this code would hold true for deletion in the middle and at the end of a linked list */  
if (score == ptr-> marks)  
{  
    prev-> next = ptr-> next;  
    free(ptr);  
}  
/* end of for loop */  
/* end of delete */
```

# Doubly Linked Lists



## Need For a Doubly Linked List

- The disadvantage with a singly linked list is that traversal is possible in only direction, i.e., from the beginning of the list till the end.
- If the value to be searched in a linked list is toward the end of the list, the search time would be higher in the case of a singly linked list.
- It would have been efficient had it been possible to search for a value in a linked list from the end of the list.

# Properties of a Doubly Linked List

- This would be possible only if we have a doubly linked list.
- In a doubly linked list, each node has two pointers, one say, **next** pointing to the next node in the list, and another say, **prior** pointing to the previous node in the list.
- Therefore, traversing a doubly linked list in either direction is possible, **from the start to the end using next**, and **from the end of the list to the beginning of the list using prior**.

# Properties of a Doubly Linked List

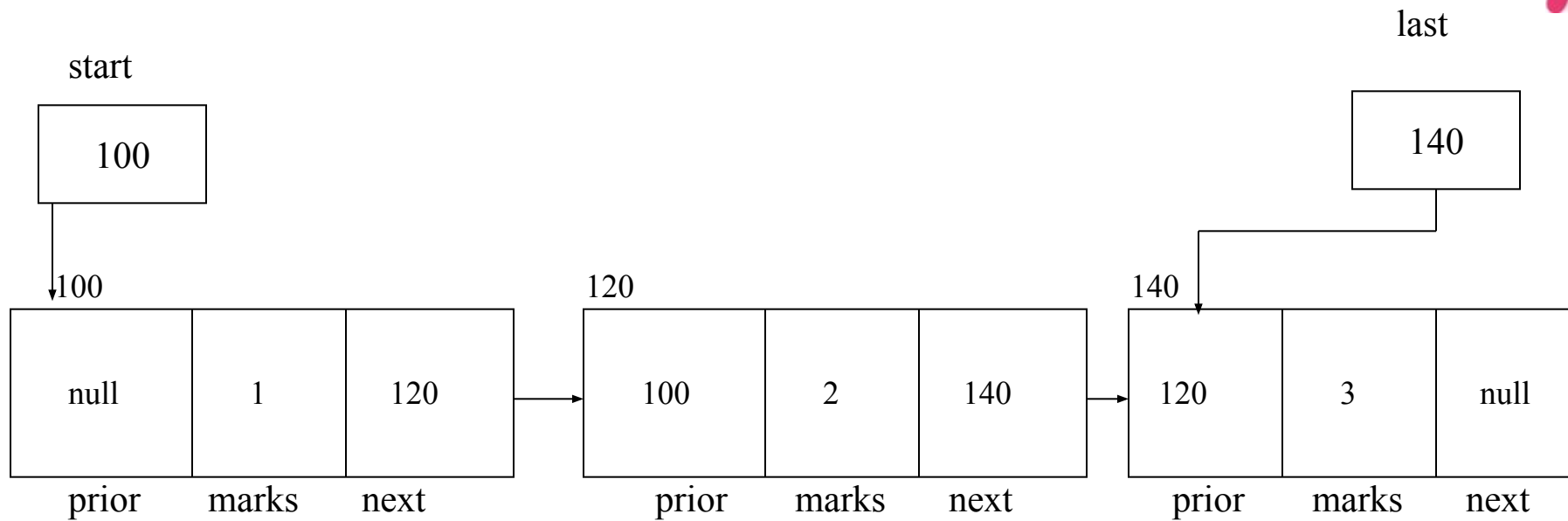
- In a doubly linked list, the **prior pointer of the first node will be null**, as there is no node before the first node.
- In a doubly linked list, the **next pointer of the last node will be null**, as there is no node after this list.
- Bidirectional traversal of a doubly linked list is useful for implementing **page up**, and **page down** functionality when using doubly linked lists to create editors.
- A doubly linked list would have two pointers, **start** and **last** to facilitate traversal from the beginning and end of the list respectively.



# Declaration of a Doubly Linked List

```
struct marks_list
{
    struct double_list *prior;
    int info;
    struct marks_list *next;
}
```

# Visualizing a Doubly Linked List



# Creating a Sorted Doubly Linked List

```
#include<stdio.h>
#include<malloc.h>
struct marks_list *start, *last;
/* variables declared outside main() are global in nature and can be accessed by
other functions called from main() */
struct marks_list
{
    struct marks_list *prior;
    int marks;
    struct marks_list *next;
};
main()
{
    struct marks_list * makenode();
    /* function prototype declaration */
```

# Creating a Sorted Doubly Linked List

```
struct marks_list *new;  
start = NULL;  
char menu = '0';  
while (menu != '5')  
{  
    printf( "Add Nodes   :\n");  
    printf( "Delete Nodes :\n");  
    printf( "Forward Traverse a list :\n");  
    printf( "Reverse Traverse a list :\n");  
    printf( "Exit           :\n");  
    menu = getchar( );
```

# Creating a Sorted Doubly Linked List

```
switch (menu)
{
    case '1': addnode( );
                break;
    case '2': deletenode( )
                break;
    case '3': forward_traverse( );
                break;
    case '4': reverse_traverse( );
                break;
    case '5': exit( );
                break;
}/* end of switch */
}/* end of main( ) */
```

# Creating a Sorted Doubly Linked List

```
addnode()  
{  
    char ch = 'y';  
    while ( ch == 'y' )  
    {  
        new = makenode();  
        /* creation of a list is treated as a special case of insertion */  
        if ( start == NULL )  
        {  
            start = new;  
            start->prior = null;  
        }  
        else  
        {  
            insert();  
            traverse( );  
        }  
    }
```





## Creating a Sorted Doubly Linked List

```
printf("%s","Want to add more nodes\n");  
scanf( "%c", &ch );  
fflush( stdin );  
}/* end of while */  
}/* end of addnode( )
```

# Creating a Sorted Doubly Linked List

```
struct marks_list * makenode()
{
    struct marks_list *new;
    new=(struct marks_list *) malloc(sizeof(struct(marks_list));
    scanf("%d",&new->marks);
    new->prior = NULL;
    new->next = NULL;
    return(new);
}
```

# Creating a Sorted Linked List

```
insert()  
{  
    struct marks_list *ptr, *prev;  
    for(ptr=start,prev=start;(ptr);prev=ptr,ptr=ptr->next)  
    {  
        if (new->marks < start->marks)  
        {  
            /* insertion at the beginning of a list */  
            new->next = start;  
            new->prior = NULL;  
            start->prior = new;  
            last = start;  
            start = new;  
        }  
    }  
}
```

# Creating a Sorted Doubly Linked List

```
/* insertion in the middle of a list */  
if(new->marks > ptr->marks)  
{  
    continue;  
}  
else  
{  
    prev->next = new;  
    new->prior = prev;  
    new->next = ptr;  
    ptr->prior = new;  
}  
}/* end of for loop */
```

## Creating a Sorted Linked List

```
/* insertion at the end of the list */  
if (ptr == null)  
{  
    prev->next = new;  
    new->prior = prev;  
    new->next = null;  
    last = new;  
}  
}/* end of insert */
```

# Searching a Value in a Doubly Linked List



```
struct marks_list *search( int val)
{
    for( ptr = start; (ptr); ptr = ptr->next)
    {
        if (val == ptr-> marks)
            return ptr;
    }
}
```

```
struct marks_list *search( int val)
{
    for( ptr = last; (ptr); ptr = ptr->prior)
    {
        if (val == ptr-> marks)
            return ptr;
    }
}
```



# Deleting a Node From a Doubly Linked List

```
delete ()
{
    struct marks_list *ptr, *prev, *temp;
    int score;
    /* search the linked list for the value to be deleted */
    scanf("%d", &score);
    fflush(stdin);
    for (ptr = start, prev = start; (ptr); prev = ptr, ptr = ptr->next)
    {
        /* deletion of the first node in the list */
        if (score == start-> marks)
        {
            temp = start;
            start = start-> next;
            start->prior = null;
            free(temp);
        }
    }
}
```

# Deleting a Node From a Linked List

```
/* deletion in the middle of a linked list */
if (score == ptr-> marks)
{
    prev-> next = ptr-> next;
    ptr->next->prior = ptr->prior;
    free(ptr);
}
/* deletion at the end of the list */
if (ptr->next == null)
{
    temp = ptr;
    prev->next = null;
    last = ptr->prior;
}
}}
```

# Traversal of a Doubly Linked List

```
/* forward traversal of a doubly linked list */  
for( ptr = start; (ptr); ptr = ptr->next)  
{  
    printf("%d", ptr->marks);  
}
```

```
/* reverse traversal of a doubly linked list */  
for( ptr = last; (ptr); ptr = ptr->prior)  
{  
    printf("%d", ptr->marks);  
}
```

**Stack**



# What is a Stack?

**A stack is a data structure in which insertions and deletions can only be done at the top.**

A common example of a stack, which permits the selection of only its end elements, is a stack of books.

A person desiring a book can pick up only the book at the top, and if one wants to place a plate on the pile of plates, one has to place it on the top.



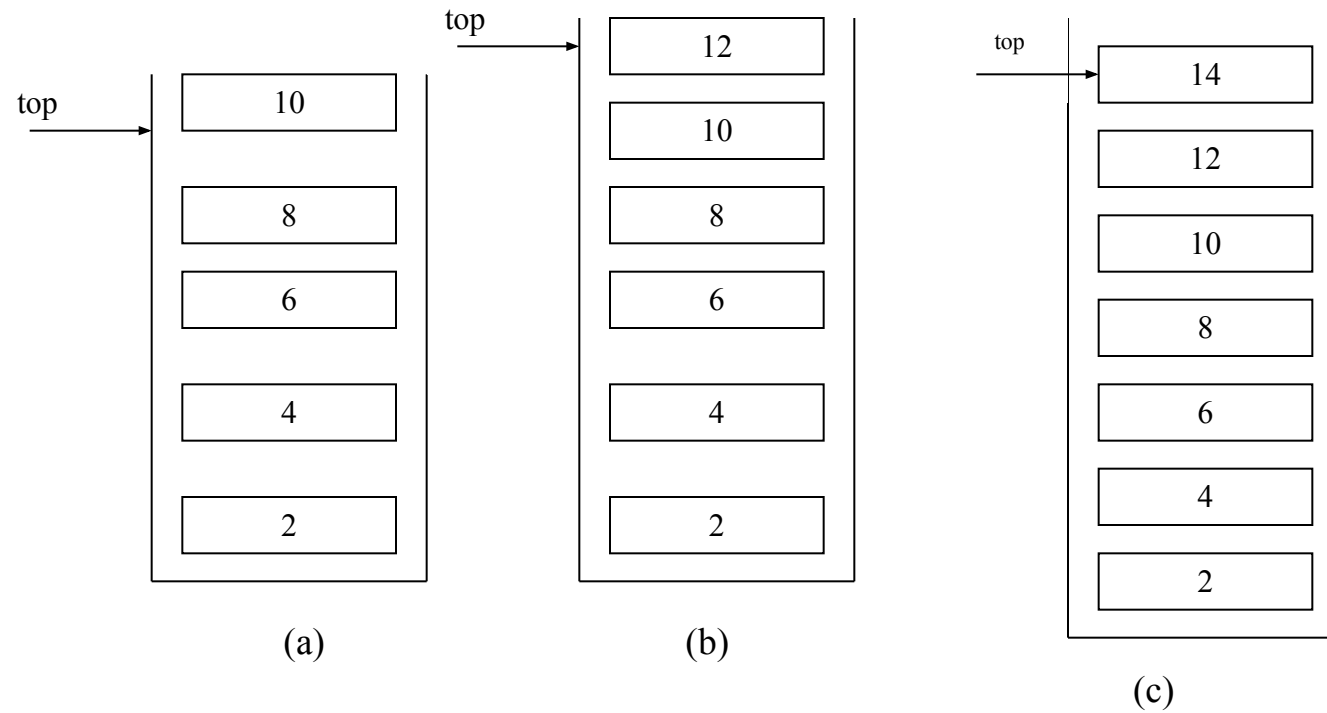
# What is a Stack?

- You would have noticed that whenever a book is placed on top of the stack, the top of the stack moves upward to correspond to the new element (the stack grows).
- And, whenever a book is picked, or removed from the top of the stack, the top of the stack moves downward to correspond to the new highest element (the stack shrinks).





# What is a Stack?



## Characteristics of a Stack

- When you add an element to the stack, you say that you **push** it on the stack, and if you delete an element from a stack, you say that you **pop** it from the stack.
- Since the last item pushed into the stack is always the first item to be popped from the stack, a stack is also called as a **Last In, First Out** or a **LIFO structure**.
- Unlike an array that is static in terms of its size, a **stack is a dynamic data structure**.

## Characteristics of a Stack

- Since the definition of the stack provides for the insertion and deletion of nodes into it, a stack can grow and shrink dynamically at runtime.
- An ideal implementation of a stack would be a special kind of linked list in which insertions and deletions can only be done at the top of the list.

# Operations on Stacks

Some of the typical operations performed on stacks are:

- **create (s)** – to create s as an empty stack
- **push (s, i)** – to insert the element i on top of the stack s
- **pop (s)** – to remove the top element of the stack and to return the removed element as a function value.
- **top (s)** – to return the top element of stack(s)
- **empty(s)** – to check whether the stack is empty or not. It returns true if the stack is empty, and returns false otherwise.

# Implementation of Stacks

- An array can be used to implement a stack.
- But since array size is defined at compile time, it cannot grow dynamically at runtime, and therefore, an attempt to insert an element into a array implementation of a stack that is already full causes a stack overflow.
- A stack, by definition, is a data structure that cannot be full since it can dynamically grow and shrink at runtime.

# Implementation of Stacks

- An ideal implementation for a stack is a linked list that can dynamically grow and shrink at runtime.
- Since you are going to employ a variation of a linked list that functions as a stack, you need to employ an additional pointer (top) that always points to the first node in the stack, or the top of the stack.
- It is using top that a node will either be inserted at the beginning or top of the stack (push a node into the stack), or deleted from the top of the stack (popping a node at the top or beginning of the stack).



# Code Implementing for a Stack

The push( ) and the pop( ) operations on a stack are analogous to insert-first-node and delete-first-node operations on a linked list that functions as a stack.

```
struct stack
{
    int info;
    struct stack *next;
};
/* pointer declaration to point to the top of the stack */
struct stack *top;
main( )
{
```

## Code Implementing for a Stack

```
top = NULL;
char menu = '0';
while (menu != '3')
{
    printf("Add Nodes : \n");
    printf("Delete Nodes : \n");
    printf("Exit : \n");
    menu = getchar();
    switch (menu)
    {
        case '1': push();
                break;
        case '2': pop();
                break;
```

## Code Implementing for a Stack

```
case '3': exit( );  
        break;  
}/* end of switch */  
}/* end of main( ) */
```

# Implementing push( )

```
push()  
{  
    struct stack * new;  
    char ch;  
    ch = 'y';  
    while (ch == 'y')  
    {  
        new = getnode()  
        /* checking for an empty stack */  
        if ( top == null)  
        {  
            top = new;  
        }  
    }  
}
```

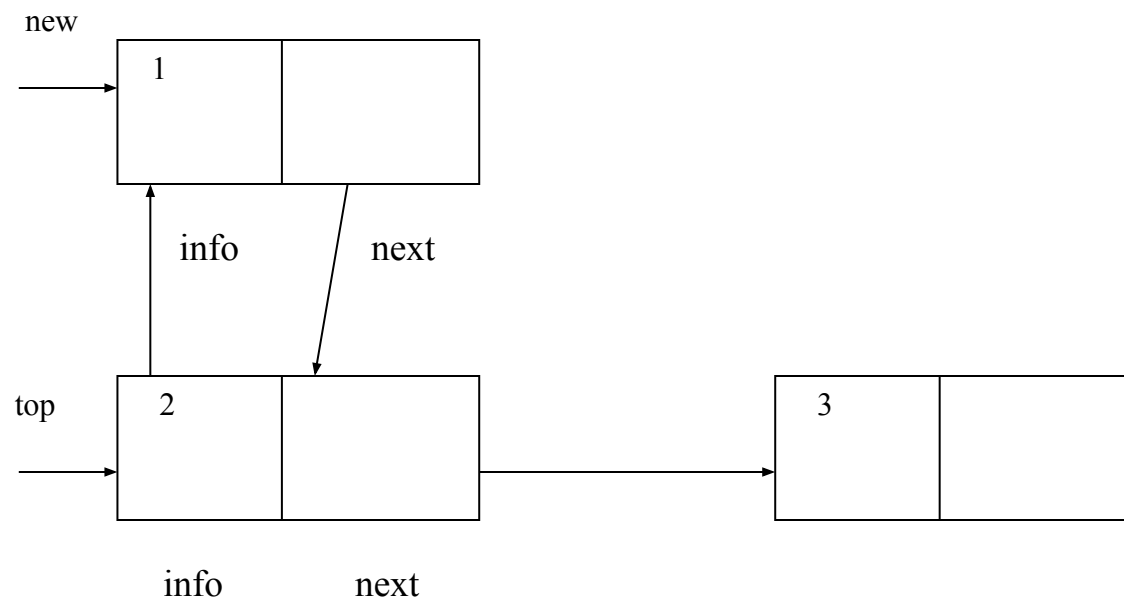


## Implementing push( )

```
else
{
    new->next = top;
    top = new;
}
printf("%s","Want to add more nodes\n");
scanf( "%c", &ch );
fflush( stdin );
}/* end of while */
}/* end of push( )
```



## A View of the Stack After Insertion





## Creating a Node on a Stack

```
struct stack * makenode()
{
    struct stack *new;
    new=(struct stack *) malloc(sizeof(struct(stack)));
    scanf("%d",&new->info);
    new->next = NULL;
    return(new);
}
```

## Implementing pop( )

```
pop()  
{  
    struct stack * temp;  
    int x;  
    /* check for an empty stack */  
    if (top == null)  
    {  
        printf ("Cannot remove nodes from an empty stack */  
        exit( );  
    }  
}
```

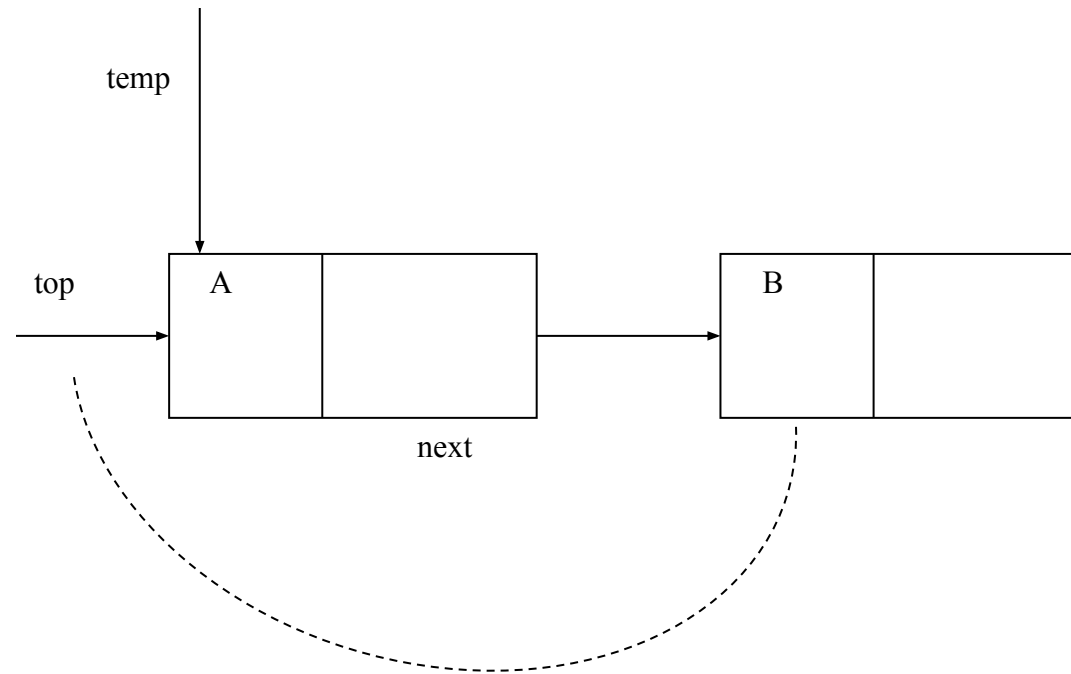


## Implementing pop( )

```
else
{
    temp = top;
    x = top->info;
    top = top->next;
    free( temp);
    return x;
}
```



## A View of the Stack After Deletion



# Applications of Stacks

- As a stack is a LIFO structure, it is an appropriate data structure for applications in which information must be saved and later retrieved in reverse order.
- Consider what happens within a computer when function `main()` calls another function.
- How does a program remember where to resume execution from after returning from a function call?
- From where does it pick up the values of the local variables in the function `main()` after returning from the subprogram?

# Applications of Stacks

- As an example, let **main()** call **a()**. Function **a()**, in turn, calls function **b()**, and function **b()** in turn invokes function **c()**.
- **main()** is the first one to execute, but it is the last one to finish, after **a()** has finished and returned.
- **a()** cannot finish its work until **b()** has finished and returned. **b()** cannot finish its work until **c()** has finished and returned.



# Applications of Stacks

- When **a()** is called, its calling information is pushed on to the stack (calling information consists of the address of the return instruction in **main()** after **a()** was called, and the local variables and parameter declarations in **main()**).
- When **b()** is called from **a()**, **b()**'s calling information is pushed onto the stack (calling information consists of the address of the return instruction in **a()** after **b()** was called, and the local variables and parameter declarations in **a()**).

# Applications of Stacks

- Then, when **b( )** calls **c( )**, **c( )**'s calling information is pushed onto the stack (calling information consists of the address of the return instruction in **b( )** after **c( )** was called, and the local variables and parameter declarations in **b( )**).
- When **c( )** finishes execution, the information needed to return to **b( )** is retrieved by popping the stack.
- Then, when **b( )** finishes execution, its return address is popped from the stack to return to **a( )**

# Applications of Stacks

- Finally, when `a()` completes, the stack is again popped to get back to `main()`.
- When `main()` finishes, the stack becomes empty.
- Thus, a stack plays an important role in function calls.
- The same technique is used in recursion when a function invokes itself.

Queue



## Defining a Queue

- A Queue is a data structure in which elements are added at one end (called the **rear**), and elements are removed from the other end (called the **front**).
- You come across a number of examples of a queue in real life situations.
- For example, consider a line of students at a fee counter. Whenever a student enters the queue, he stands at the end of the queue (analogous to the addition of nodes to the queue)

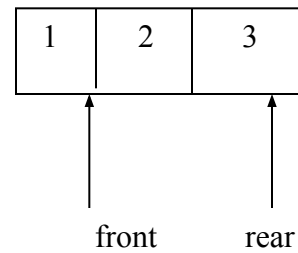
## Defining a Queue

- Every time the student at the front of the queue deposits the fee, he leaves the queue (analogous to deleting nodes from a queue).
- The student who comes first in the queue is the one who leaves the queue first.

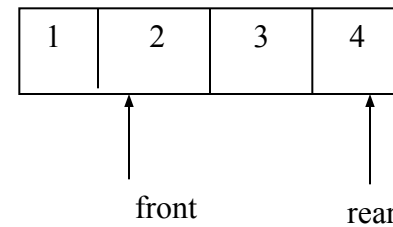
Therefore, a queue is commonly called a first-in-first-out or a FIFO data structure.



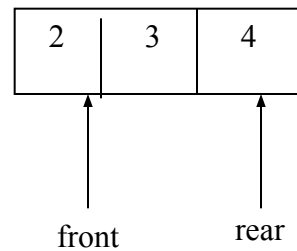
# Queue Insertions and Deletions



(a)



(b)



(c)

# Queue Operations

- To complete this definition of a queue, you must specify all the operations that it permits.
- The first step you must perform in working with any queue is to initialize the queue for further use. Other important operations are to add an element, and to delete an element from a queue.
- Adding an element is popularly known as ENQ and deleting an element is known as DEQ. The following slide lists operations typically performed on a queue.

# Queue Operations

`create(q)` – which creates `q` as an empty queue

`enq(i)` – adds the element `i` to the rear of the queue and returns the new queue

`deq(q)` – removes the element at the front end of the queue (`q`) and returns the resulting queue as well as the removed element

`empty (q)` – it checks the queue (`q`) whether it is empty or not. It returns true if the queue is empty and returns false otherwise

`front(q)` – returns the front element of the queue without changing the queue

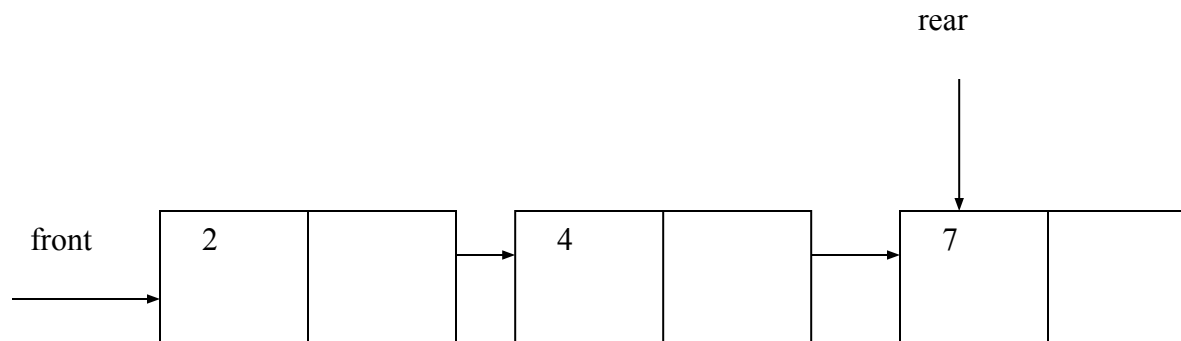
`queuesize (q)` – returns the number of entries in the queue

# Implementing Queues

- Linked lists offer a flexible implementation of a queue since insertion and deletion of elements into a list are simple, and a linked list has the advantage of dynamically growing or shrinking at runtime.
- Having a list that has the functionality of a queue implies that insertions to the list can only be done at the rear of the list, and deletions to the list can only be done at front of the list.

# Implementing Queues

- Queue functionality of a linked list can be achieved by having two pointers **front** and **rear**, pointing to the first element, and the last element of the queue respectively.
- The following figure gives a visual depiction of linked list implementation of a queue.



# Queue Declaration & Operations

```
struct queue
{
    int info;
    struct queue *next;
};
struct queue *front, *rear;
```

An empty queue is represented by  $q \rightarrow \text{front} = q \rightarrow \text{rear} = \text{null}$ . Therefore, `clearq()` can be implemented as follows:

```
void clearq(struct queue * queue_pointer)
{
    queue_pointer->front = queue_pointer->rear = null;
}
```



## Queue Operations

You can determine whether a queue is empty or not by checking its front pointer. The front pointer of a queue can be passed as an argument to `emptyq()` to determine whether it is empty or not.

```
int emptyq(queue_pointer)
{
    if (queue_pointer == null)
        return (1);
    else
        return(0);
}
```

# Insertion into a Queue

```
struct queue
{ int info;
  struct queue *next;
};
/* pointer declarations to point to the front, and rear of the queue */
struct queue *front, *rear;
main()
{
  front = NULL;
  rear = NULL;
```



## Insertion into a Queue

```
char menu = '0';
while (menu != '3')
{
    printf( "Add Nodes   :\n");
    printf( "Delete Nodes :\n");
    printf( "Exit       :\n");
    menu = getchar( );
    switch (menu)
    {
        case '1': enq( );
                    break;
        case '2': deq( )
                    break;
```



## Insertion into a Queue

```
case '3': exit( );  
        break;  
/* end of switch */  
/* end of main( ) */
```



# Insertion into a Queue

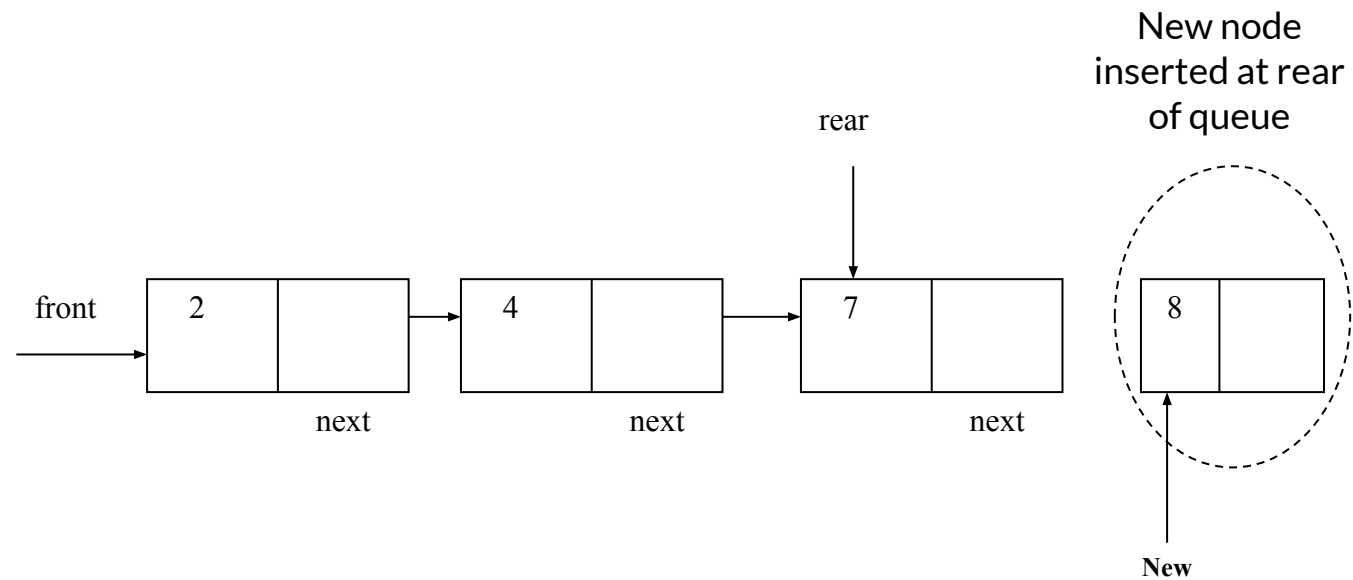
```
void enq()  
{  
    struct queue *new;  
    new = getnode();  
    if(queue_pointer->front == queue_pointer->rear == null)  
    {  
        queue_pointer->front = new;  
        queue_pointer->rear = new;  
    }  
    else  
    {  
        rear->next = new;  
        rear = new;  
    }  
}
```

## Creating a Node on a Queue

```
struct queue * makenode()
{
    struct queue *new;
    new=(struct queue *) malloc(sizeof(struct(queue));
    scanf("%d",&new->info);
    new->next = NULL;
    return(new);
}
```



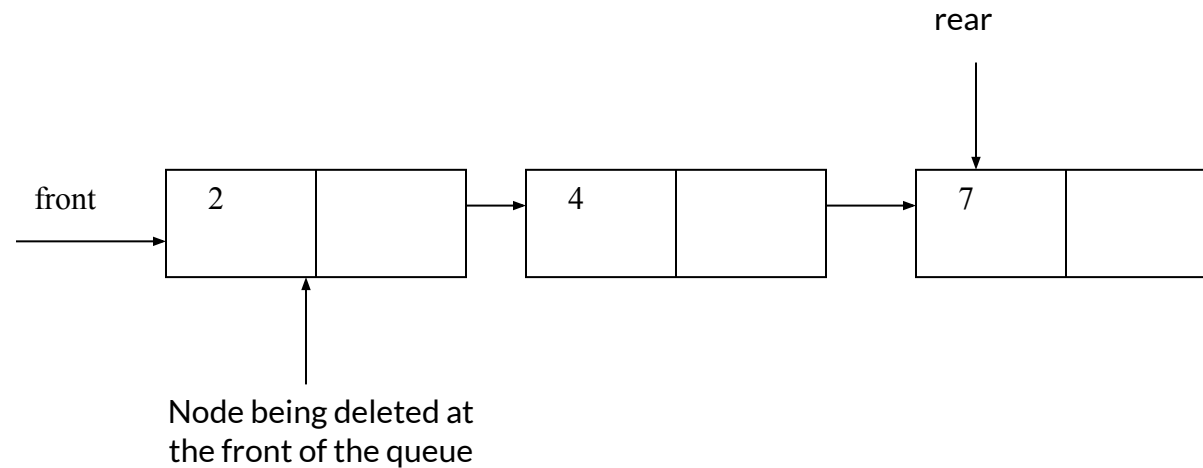
# Insertion into a Queue



# Deletion from a Queue

```
int deq()
{
    struct queue *temp;
    int x;
    if(queue_pointer->front == queue_pointer->rear == null)
    {
        printf("Queue Underflow\n");
        exit(1);
    }
    temp = front;
    x = temp->info;
    front = front->next;
    free(temp);
    if(front == null) /* check for queue becoming empty after node deletion */
        rear = null;
    return(x);
}
```

# Deletion of a Node From a Queue



# Applications of Queues

- Queues are also very useful in a time-sharing multi-user operating system where many users share the CPU simultaneously.
- Whenever a user requests the CPU to run a particular program, the operating system adds the request ( by first of all converting the program into a process that is a running instance of the program, and assigning the process an ID).

**This process ID is then added at the end of the queue of jobs waiting to be executed.**

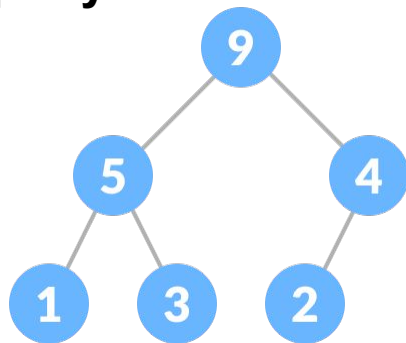
# Applications of Queues

- Whenever the CPU is free, it executes the job that is at the front of the job queue.
- Similarly, there are queues for shared I/O devices. Each device maintains its own queue of requests.
- An example is a print queue on a network printer, which queues up print jobs issued by various users on the network.
- The first print request is the first one to be processed. New print requests are added at the end of the queue.

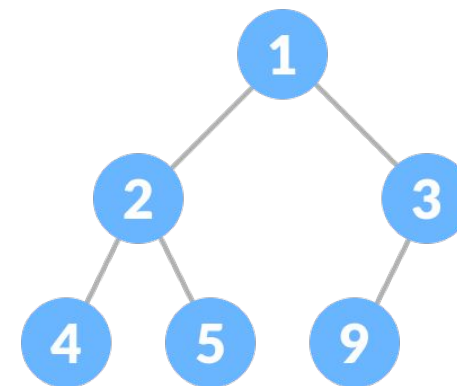
# Heap

Heap data structure is a complete binary tree that satisfies **the heap property**, where any given node is

- always greater than its child node/s and the key of the root node is the largest among all other nodes. This property is also called **max heap property**.
- always smaller than the child node/s and the key of the root node is the smallest among all other nodes. This property is also called **min heap property**.



Max heap



Min heap



# Maps

A *Map* is a type of fast key lookup data structure that offers a flexible means of indexing into its individual elements. Unlike most array data structures that only allow access to the elements by means of integer indices, the indices for a Map can be nearly any scalar numeric value or a character vector.

Indices into the elements of a Map are called *keys*. These keys, along with the data *values* associated with them, are stored within the Map. Each entry of a Map contains exactly one unique key and its corresponding value.

# Dictionaries

Collection of pairs.

(key, value)

Each pair has a unique key.

Operations.

Get(theKey)

Delete(theKey)

Insert(theKey, theValue)



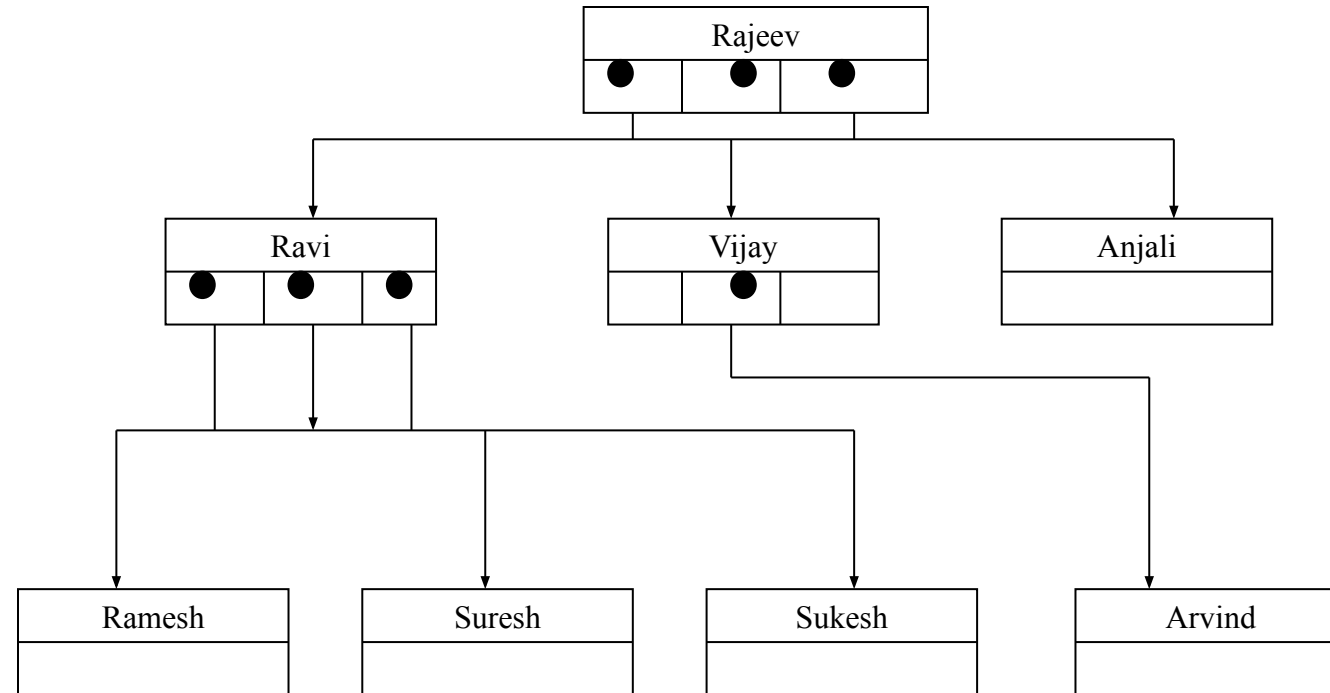
# Introduction of hash table

- Data structure that offers very fast insertion and searching, almost  $O(1)$ .
- Relatively easy to program as compared to trees.
- Based on arrays, hence difficult to expand.
- No convenient way to visit the items in a hash table in any kind of order.

# Trees

- Compared to linked lists that are linear data structures, **trees are non-linear data structures.**
- In a linked list, each node has a link which points to another node.
- In a tree structure, however, each node may point to several nodes, which may in turn point to several other nodes.
- Thus, a tree is a very flexible and a powerful data structure that can be used for a wide variety of applications.

# Trees



# Trees

- A tree consists of a collection of nodes that are connected to each other.
- A tree contains a unique first element known as the root, which is shown at the top of the tree structure.
- A node which points to other nodes is said to be the parent of the nodes to which it is pointing, and the nodes that the parent node points to are called the children, or child nodes of the parent node.



# Binary Trees



# Binary Tree

- If you can introduce a restriction that each node can have a maximum of two children or two child nodes, then you can have a binary tree.
- You can give a formal definition of a binary tree as a tree which is either empty or consists of a root node together with two nodes, each of which in turn forms a subtree.
- You therefore have a left subtree and a right subtree under the root node.

## Binary Search Tree

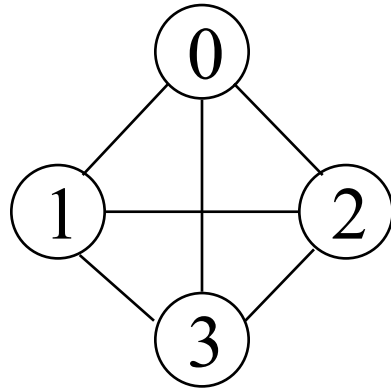
- A complete binary tree can be defined as one whose non-leaf nodes have non empty left and right subtrees and all leaves are at the same level.
- This is also called as a balanced binary tree. If a binary tree has the property that all elements in the left subtree of a node  $n$  are less than the contents of  $n$ , and all elements in the right subtree are greater than the contents of  $n$ , such a tree is called a binary search tree.

# GRAPH

- A graph  $G$  consists of two sets
  - a finite, nonempty set of vertices  $V(G)$
  - a finite, possible empty set of edges  $E(G)$
  - $G(V,E)$  represents a graph
- An undirected graph is one in which the pair of vertices in a edge is unordered,  $(v_0, v_1) = (v_1, v_0)$
- A directed graph is one in which each edge is a directed pair of vertices,  $\langle v_0, v_1 \rangle \neq \langle v_1, v_0 \rangle$

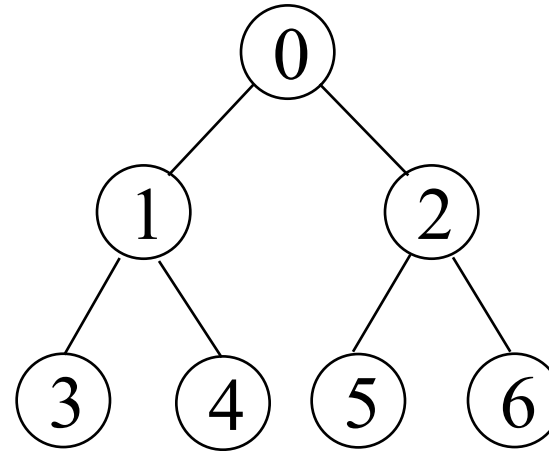
tail  $\xrightarrow{\hspace{1.5cm}}$  head

# Examples for Graph



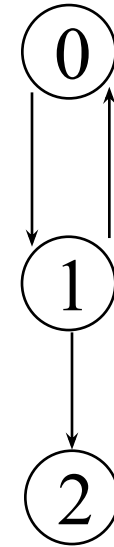
$G_1$

complete graph



$G_2$

incomplete graph



$G_3$

$$V(G_1) = \{0, 1, 2, 3\}$$

$$E(G_1) = \{(0, 1), (0, 2), (0, 3), (1, 2), (1, 3), (2, 3)\}$$

$$V(G_2) = \{0, 1, 2, 3, 4, 5, 6\}$$

$$E(G_2) = \{(0, 1), (0, 2), (1, 3), (1, 4), (2, 5), (2, 6)\}$$

$$V(G_3) = \{0, 1, 2\}$$

$$E(G_3) = \{<0, 1>, <1, 0>, <1, 2>\}$$

complete undirected graph:  $n(n-1)/2$  edges

complete directed graph:  $n(n-1)$  edges

# Some Graph Operations

- Traversal

Given  $G=(V,E)$  and vertex  $v$ , find all  $w \in V$ , such that  $w$  connects  $v$ .

Depth First Search (DFS)

preorder tree traversal

Breadth First Search (BFS)

level order tree traversal

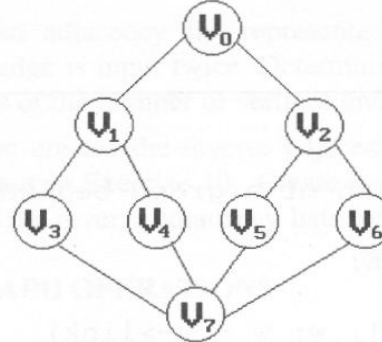
- Connected Components
- Spanning Trees



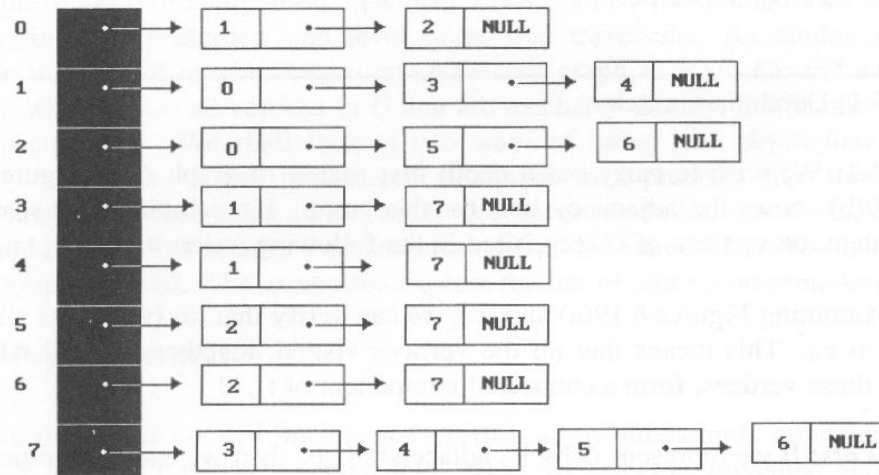


# Graph G and its adjacency lists

depth first search:  $v_0, v_1, v_3, v_7, v_4, v_5, v_2, v_6$



(a)



(b)

breadth first search:  $v_0, v_1, v_2, v_3, v_4, v_5, v_6, v_7$

Thank You!

