

# Queue – Queue Using Array

## 1. Introduction

A **Queue Using Array** is an implementation of the queue data structure where elements are stored in a **linear array**.

It follows the **FIFO (First In, First Out)** principle, where insertion happens at the **rear** and deletion happens at the **front**.

This is one of the **simplest queue implementations** and is commonly taught in DSA fundamentals.

## 2. What is Queue Using Array?

In this implementation:

- A fixed-size array is used to store queue elements
- Two pointers are maintained:
  - **Front** → points to the first element
  - **Rear** → points to the last element
- Initially:

```
Front = -1  
Rear = -1
```

## 3. Queue Structure (Conceptual)

```
Array: [ 10 | 20 | 30 | | ]  
Index:  0  1  2  
Front → 0  
Rear  → 2
```

Elements are accessed **only from the front or rear**, not randomly.

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## 4. Components of Queue Using Array

- **Array** → Stores queue elements
  - **Front pointer** → Tracks the first element
  - **Rear pointer** → Tracks the last element
  - **Size** → Maximum capacity of the queue
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## 5. Enqueue Operation (Array Queue)

### Logic (Plain English)

1. Check if  $\text{rear} == \text{size} - 1$
2. If true → Queue Overflow
3. If queue is empty:
  - Set  $\text{front} = 0$
4. Increase rear by 1
5. Insert element at  $\text{array}[\text{rear}]$

### Example

Enqueue 40  
Rear moves from 2 → 3

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## 6. Dequeue Operation (Array Queue)

### Logic (Plain English)

1. Check if  $\text{front} == -1$  or  $\text{front} > \text{rear}$
2. If true → Queue Underflow
3. Remove the element at  $\text{array}[\text{front}]$

4. Increase front by 1

## Example

Dequeue → removes 10  
Front moves from 0 → 1

## 7. Peek / Front Operation

### Logic (Plain English)

1. Check if queue is empty
2. If not empty, return array[front]
3. Queue remains unchanged

## 8. isEmpty Operation

The queue is empty when:

`front == -1 OR front > rear`

Used before dequeue and peek operations.

## 9. isFull Operation

The queue is full when:

`rear == size - 1`

Used before enqueue operation.

## 10. Queue Overflow in Array Queue

**Queue Overflow** occurs when:

- Attempting to enqueue an element when  $\text{rear} == \text{size} - 1$

Reason:

- Fixed size of the array
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## 11. Queue Underflow in Array Queue

**Queue Underflow** occurs when:

- Attempting to dequeue an element when the queue is empty
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## 12. Time and Space Complexity

Operation	Time Complexity
Enqueue	$O(1)$
Dequeue	$O(1)$
Peek	$O(1)$

- **Space Complexity:**  $O(n)$   
(where  $n$  is the size of the queue)
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## 13. Advantages of Queue Using Array

- Simple and easy to implement
  - Fast operations
  - Suitable for fixed-size problems
  - Low overhead
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## 14. Limitations of Queue Using Array

- Fixed size
- Space wastage after dequeue operations
- Overflow even when free space exists

- Less flexible than circular queue
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## 15. Comparison with Circular Queue

Feature	Array Queue	Circular Queue
Memory Usage	Inefficient	Efficient
Overflow Issue	Common	Rare
Implementation	Simple	Moderate
Performance	Moderate	Better

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## 16. Real-World Applications

- Basic task scheduling
  - Printer queue simulation
  - Educational demonstrations
  - Small buffering systems
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## 17. Summary

- Queue using array uses fixed-size storage
  - Uses front and rear pointers
  - Enqueue at rear, dequeue at front
  - Simple but has space limitations
  - Time complexity is  $O(1)$
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