

## How Shor's Algorithm can potentially break RSA generated pins on a quantum system :

RSA is a widely used encryption algorithm that secures data online. It's working goes as follows :

### 1. Key Generation:

- Choose two large prime numbers  $p$  and  $q$ .
- Compute  $N = p \times q \rightarrow$  this is part of the public key.
- Choose  $e$  (public exponent) and compute  $d$  (private key) such that:

$$e \cdot d \equiv 1 \pmod{(p-1)(q-1)}$$

- Public key =  $(N, e)$ , Private key =  $d$

### 2. Encryption:

- Given message  $M$ , the ciphertext is:

$$C = M^e \pmod N$$

### 3. Decryption:

- Given ciphertext  $C$ , the message is recovered by:

$$M = C^d \pmod N$$

RSA relies on the fact that factoring large numbers is hard for classical computers.

Given  $N = p \times q$ , it's easy to multiply  $p$  and  $q$ , but very hard to factor  $N$  back into  $p$  and  $q$ , especially when  $p$  and  $q$  are hundreds of digits long.

Without knowing  $p$  and  $q$ , you can't compute  $(p-1)(q-1)$  and therefore can't compute  $d$ .

## How Shor's Algorithm Breaks RSA :

Shor's Algorithm can factor  $N$  efficiently using a quantum computer.

Here's the connection:

- RSA security : based on factoring being hard
- Shor's Algorithm : factors  $N$  in polynomial time

### Breakdown:

1. Given public key  $N$ , Shor's Algorithm factors it into  $p$  and  $q$ .
2. Compute  $\phi(N) = (p-1)(q-1)$
3. Compute private key  $d$  from public exponent  $e$  using the modular inverse:

$$D \equiv e^{-1} \pmod{\phi(N)}$$

4. Now decrypt any RSA ciphertext intended for that keypair.

Example :

Say :

- $p = 7, q = 13$  , thus  $N = 91$
- $e = 5$

From Shor's Algorithm, you get:

- $\phi(91) = (7-1)(13-1) = 6 \times 12 = 72$
- Compute  $d = 29$  such that  $5 \times 29 \equiv 1 \pmod{72}$

Using Shor's Algorithm, we can recover 7 and 13, and compute d, thus breaking RSA.

Code:

```
import numpy as np
```

```
from math import gcd
```

```
from fractions import Fraction
```

```
from random import randint
```

```
def euclidGCD(n, m):
```

```
    if m == 0:
```

```
        return n
```

```
    return euclidGCD(m, n % m)
```

```
def classical_period_finding(a, N):
```

```
    r = 1
```

```
    while pow(a, r, N) != 1:
```

```
        r += 1
```

```
    if r > N:
```

```
        return None
```

```
    return r
```

```
def shors_classical(N):
```

```

if N % 2 == 0:
    return 2
attempt = 0
while True:
    attempt += 1
    a = randint(2, N - 1)
    print(f"Attempt {attempt}: Trying a = {a}")
    if gcd(a, N) != 1:
        print(f"GCD({a}, {N}) = {gcd(a, N)}. Found factor early!")
        return gcd(a, N)

```

```

r = classical_period_finding(a, N)
if r is None or r % 2 != 0:
    print(f"Invalid period r = {r}")
    continue

```

```

if pow(a, r, N) != 1:
    continue

```

```

plus = gcd(pow(a, r // 2) + 1, N)
minus = gcd(pow(a, r // 2) - 1, N)

```

```

for factor in [plus, minus]:
    if factor != 1 and factor != N and N % factor == 0:
        print(f"Success! Found non-trivial factor: {factor}")
        return factor

```

```

N = 65

```

```

factor = shors_classical(N)
print(f"One factor of {N} is {factor}")

```

