- Left Shift
- Right Shift
- Check bit / Count Bit
- Set bit
- Unset bit
- Negative number
- Significance of MSB
- Datatypes Ranges.

→ Next Class

Bit manipulation-3

int 32 bites

8 bit number

a = 10

2^7 2^6 2^5 2^4 2^3 2^2 2^1 2^0

0 0 0 0 1 0 1 0

7th 6th 5th 4th 3rd 2nd 1st 0th

$2^3 + 2^1 = 8 + 2 = 10$

a<<1 0 0 0 1 0 1 0 0 = $2^4 + 2^2 = 16 + 4 = 20$

a<<2 0 0 1 0 1 0 0 0 = $2^5 + 2^3 = 32 + 8 = 40$

a<<3 0 1 0 1 0 0 0 0 = $2^6 + 2^4 = 64 + 16 = 80$

a<<4 1 0 1 0 0 0 0 0 = $2^7 + 2^5 = 128 + 32 = 160$

a<<5 0 1 0 0 0 0 0 0 = $2^6 = 64$

Overflow = Exceeding the size.

General (No overflow)

int x = 10

$a << 1 = a \times 2$

$a << 2 = a \times 2^2$

$a << 3 = a \times 2^3$

$a << 4 = a \times 2^4$

$a << N = a \times 2^N$.

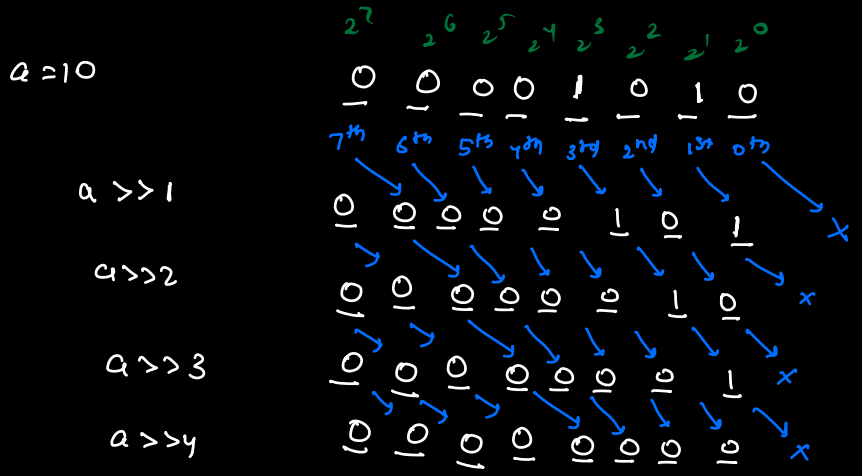
$1 << 1 = 1 \times 2^1 = 2$

$1 << 2 = 1 \times 2^2 = 2^2$

$1 << N = 2^N$.

Right shift

a = 10



$2^3 + 2^1 = 10$ ↓ /2

$2^2 + 2^0 = 5$ ↓ /2

$2^1 = 2$ ↓ /2

$2^0 = 1$ ↓ /2.

0

Generalise

$x >> 1 = \frac{x}{2}$

$x >> 2 = \frac{x}{2^2}$

$x >> 3 = \frac{x}{2^3}$

$x >> N = \frac{x}{2^N}$

int → 4 bytes → 32 bits
[31 ... 2, 1, 0].

long → 8 bytes → 64 bits
[63 ... 0]

Q. Given N, i. Check if i^{th} bit pos in N is set or not

N = 21
i = 2

4 3 2 1 0th
1 0 1 0 1

Bit is set
ans = True

N = 34
i = 3

5 4 3 2 1 0
1 0 0 0 1 0

Bit is unset
ans = False.

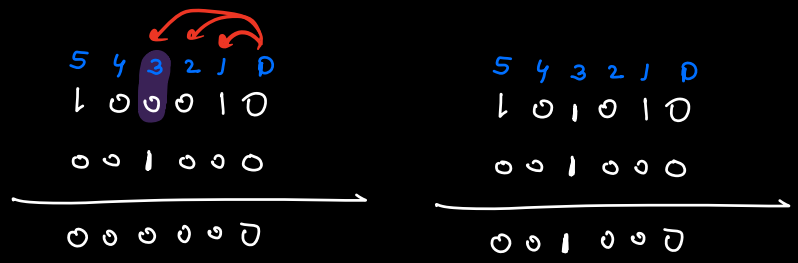
Given N, 0^{th} bit set or not

If (N & 1 == 1) return True
else return False.

```
boolean checkbit ( int N, int i )
{
    If ( (N>>i) & 1 == 1 )  return True
    else  return False
}
```

```
boolean checkbit ( int N, int i )
{
    return (N>>i) & 1 == 1
}
```

N = 39
i = 3




```
bool checkbit ( int N, int i )
{
    If ( (N & (1<<i)) == 0 )  return False
    else  return True.
}
```

```
bool checkbit ( int N, int i )
{
    if ( (N & (1<<i)) != 0 )
        return True
    else return False.
}
```

```
bool checkbit ( int N, int i )
{
    return !( (N & (1<<i)) == 0 )
    return ( (N & (1<<i)) == (1<<i) )
    return ( (N & (1<<i)) != 0 )
}
```

O(1)

Q. Given a number N (integer), cal how many bits are set.

N = 10 → 1 0 1 0 → ans = 2.

N = 27 → 1 1 0 1 1 → ans = 4

Idea 1

32 iterations

```
cnt = 0
for( i=0 ; i<=31 ; i++ )
{
    if ( checkbit( N, i ) ) cnt++
}
return cnt
```

TC: O(1)

SC: O(1)

Idea 2.

N = 10 | 0 | 0

N >> 1 | 0 | 0 | 1

N >> 2 | 0 | 0 | 0 | 1

N >> 3 | 0 | 0 | 0 | 0

N >> 4 | 0 | 0 | 0 | 0

N & 1 == 0

N & 1 == 1 cnt = 1

N & 1 == 0

N & 1 == 1 cnt = 2

```
int cnt = 0
while ( N > 0 )
{
    if( N & 1 == 1 ) cnt++
    N = N >> 1
}
return cnt
```

TC: O(log N)

SC: O(1).

Break

10:24 — 10:32.

N is capped to 2^{32}

O(log N) approach better than O(1) approach.

32 iterations.

Q. Given N, set the i^{th} bit.

N = 10 (3 2 1 0) 1 0 1 0 → 1 1 1 0 ans = 14

i = 2

N = 10 1 0 1 0 → 1 0 1 1 → ans = 11.

i = 0

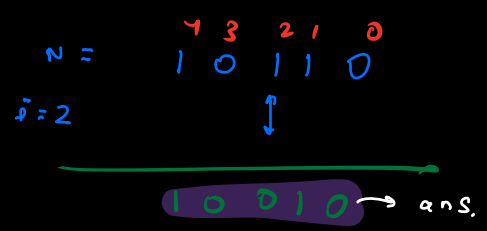
```
int setbit ( int N, int i)
{
    if ( checkbit(N, i) ) return N.
    else return N + 2^i
}
```

$(2^i = 1<<i)$

```
   9 8 7 6 5 4 3 2 1 0
   1 0 1 x 1 1 1 0 0 1
OR 0 0 0 1 0 0 0 0 0 0      <- 2^i (1<<i)
   ---------------------
   1 0 1 1 1 1 1 0 0 1
```

```
boolean set(int N, int i)
{
    return N | (1<<i)
}
```

Q. Given N, i. Unset ith bit.




```
boolean unset(N, i)
{
    if(check(N,i)) return N - 2^i
    else return N
}
```

$2^i = 1<<i$

```
  1 0 1 0/1 0 1 0        & 0
& 1 1 1 0 1 1 1 1      ~(0 0 0 1 0 0 0)
-----------------
  1 0 1 0 0 1 0
```

```
boolean unset(N, i)
{
    return N & (~(1<<i))
}
```