

Today's Content:

✓ pair sum K

✓ Distinct elements in subarray K.

Q Given an array N. Check if there exists a pair (i, j)

S.T $arr[i] + arr[j] = K$ & $i \neq j$ $\left\{ \begin{array}{l} i, j \text{ are} \\ \text{indexes} \end{array} \right\}$

$arr[] =$

0	1	2	3	4	5	6	7	8	9
8	9	1	-2	4	5	11	-6	7	5

$K = 11$

i	j
4	8

 $arr[i] + arr[j] = 11$ True

$K = 6$

i	j
2	5
0	3
2	9

 $1 + 5 = 6$
 $8 - 2 = 6$
 $1 + 5 = 6$ True

$K = 10$

i	j
5	9
1	2

 $5 + 5 = 10$
 $9 + 1 = 10$ True

$K = 22$

i	j
6	6

 $11 + 11 = 22$ False

Ideal checks all pair

```
for( i=0; i < n; i++)
```

TC: $O(N^2)$
SC: $O(1)$

```
{  
    for( j=0; j < n; j++)  
    {  
        if( i != j && arr[i] + arr[j] == k )  
            return True  
    }  
}
```

return False

```
for( i=0; i < n; i++)
```

```
{  
    for( j=i+1; j < n; j++)  
    {  
        if( arr[i] + arr[j] == k )  
            return True  
    }  
}
```

return False

0 1	12	23 34
0 2	13	24 35
0 3	14	25
0 4	15	
0 5		45

TC: $O(N^2)$
SC: $O(1)$

First element = $arr[0]$

Second element = $K - arr[0]$

i^{th} \rightarrow First element

$K - arr[i]$

$K = 11$

$\{ 8, 9, 1, -2, 4, 7 \}$

1. Create hashset with all elements
2. For every i , search $K - arr[i]$ in the hashset.

$\left\{ \begin{array}{c} 8 \\ 9 \\ 1 \\ -2 \\ 4 \\ 7 \end{array} \right\}$

First element	Second element	Yes/No
8	3	NO
9	2	NO
1	10	NO
-2	13	NO

ans = True

4 | 7 | Yes

$$Ex = \{5, 1\}$$

K = 10

$$\begin{Bmatrix} 5 \\ 1 \end{Bmatrix}$$

False

First elem	Second	Yes/No
5	5	Yes.

$$Ex \{5, 1, 5\}$$

True

1. Create hashmap, by storing frequency.
2. for every i , search $k - arr[i]$

$$arr[i] == k - arr[i] \quad freq \geq 1$$

$$arr[i] != k - arr[i] \quad f > 0$$

$$\begin{Bmatrix} 2, 4, 2, 6, \\ 7, 3 \end{Bmatrix}$$

K = 14

	First	Second	Yes/No
$\langle 2, 1 \rangle$	2	12	NO
$\langle 4, 1 \rangle$	4	10	NO
$\langle 6, 1 \rangle$	2	12	NO
$\langle 7, 1 \rangle$	6	8	NO
$\langle 3, 1 \rangle$	7	7	NO
	3	11	NO

return False

```

bool checkpair ( arr[], int K)
{
    int n = arr.length
    HashMap<int, int> hm
    for (i=0; i<n; i++)
    {
        if ( hm.containsKey( arr[i]) )
        {
            hm.put ( arr[i], hm.get(arr[i]) + 1 )
        }
        else
        {
            hm.put ( arr[i], 1 )
        }
    }
    for ( i=0; i<n; i++)
    {
        int a = arr[i]
        b = K - a
        if ( a != b )
    }
}

```

```

    {
        if ( hm.containsKey(b) )
            return True
        else
            {
                if ( hm.get(b) > 1 )
                    return True
            }
    }
    return False

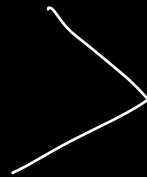
```

TC : $O(N)$

SC : $O(N)$

TC : $O(N)$

SC : $O(N)$



TC : $O(N \log N)$

SC : $O(1)$

Break

10:07

10:15

Q. Given an array N . Cal no of distinct elements in every subarray of size K .

arr[6] = ^{0 1 2 3 4 5}
2 4 3 8 3 4 $K=4$

[0,3] 4
[1,4] 3
[2,5] 3

$s=0, e=K-1$

1, K

2, K+1

3, K+2

⋮

$N-K, N-1$

$s=0, e=K-1$

of subarrays = $N-K+1$

while ($e < N$)

{
 HashSet<int> hs
 for($i=s; i \leq e; i++$) $hs.put(arr[i])$
 Print ($hs.size()$)
 $s++$
 $e++$
}

T.C : $(N-K+1) * K$

$K=1$ $(N-1+1) * 1 = N$

$K=N$ $(N-N+1) * N = N$

$$K = \frac{N}{2} \quad (N - \frac{N}{2} + 1) (\frac{N}{2})$$


$$(\frac{N}{2} + 1) (\frac{N}{2}) = \frac{N^2}{4} + \frac{N}{2} \quad O(N^2)$$

$$TC: O(N^2)$$

$$SC: O(K)$$


arr [6] =

0	1	2	3	4	5
2	4	3	8	3	4



2 sliding window on hashset

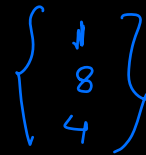
0	1	2	3	4
3	8	4	3	1



$$K = 4$$

0-3 3

1-4 3



Issue: deleting a number in hashset, deletes all the occurrences.

Idea 3

HashMap + Sliding window.

arr[] = 2 4 3 8 3 9 4 9 4 10

(Indices 0-9 are written above the corresponding elements)

(A blue bracket is drawn under the elements 3, 8, 3, 9, 4)

k = 4

0-3	4
1-4	3
2-5	3
3-6	4
4-7	,
5-8	,
6-9	

(A red bracket groups the following pairs)

- <4, 1>
- <9, 1>
- <3, 1>
- <8, 1>

Pseudo Code

HashMap <int, int> hm

for (i = 0; i < n; i++)

{
if (arr[i] in hm) update freq
else hm.put(arr[i], 1)

print (hm.size())

s = 0, e = k

while (e < n)

{
// Remove arr[s-1]
// add arr[e]

