

Q. Given a number N . Find sum of the digits.

Ex $N = 45216$

$$4 + 5 + 2 + 1 + 6 = 18$$

$$\text{sum of digits}(45216) = \text{sum of digits}(4521) + 6$$

$$\text{sumdig}(N) = \text{sumdig}\left(\frac{N}{10}\right) + N \% 10$$

Sum of digits of N

int sumdigits(int N)

if ($N < 10$) return N

OR

if ($N == 0$) return 0

return $\text{sumdig}\left(\frac{N}{10}\right) + N \% 10$

Q. Given a, n . Find a^N . $a > 0$

$$a = 2 \\ n = 5$$

$$a^N = 2^5 = 32$$

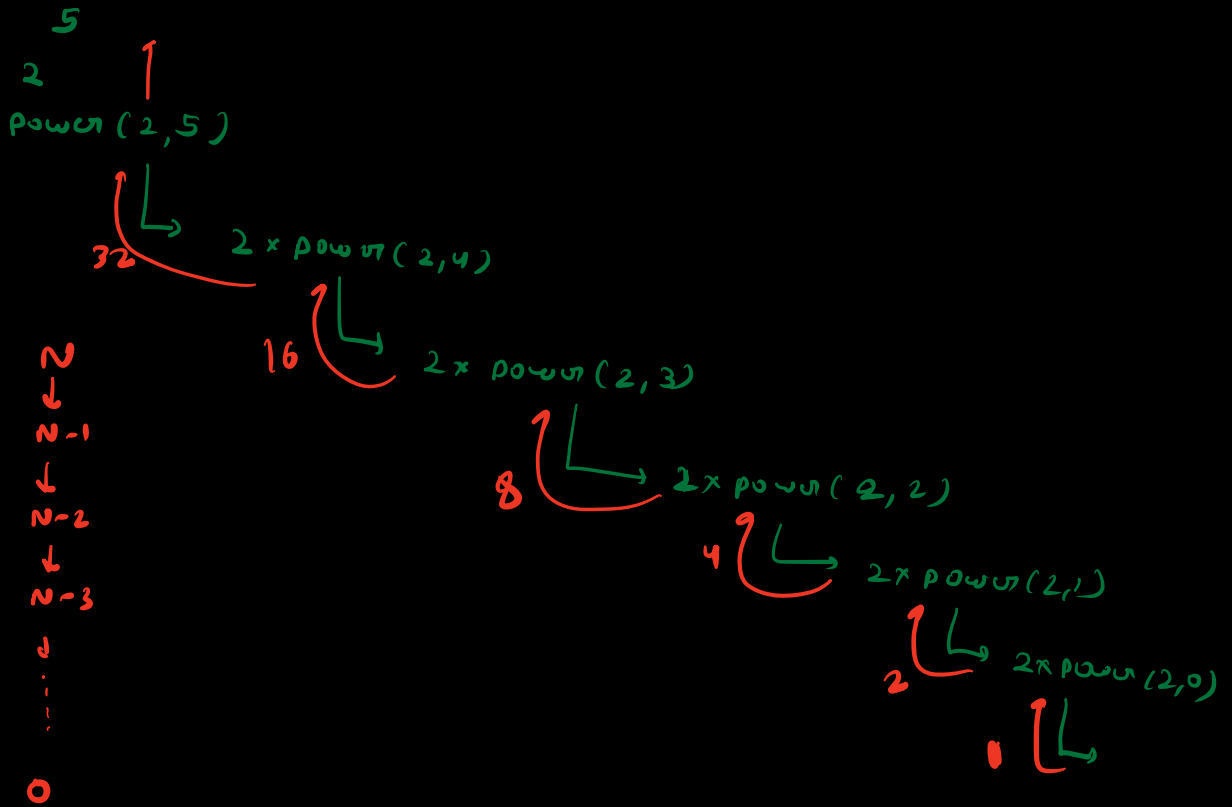
$$a^N = \underbrace{a, a, a, \dots, a}_{N \text{ times}}$$

$$a^N = a \cdot a^{N-1}$$

$$\text{pow}(a, N) = a * \text{pow}(a, N-1)$$

a^N
 int power(a, N)
 {
 if(N==0) return 1
 return a * power(a, N-1)
 }

$$\begin{aligned}
 2^0 &= 1 \\
 3^0 &= 1 \\
 4^0 &= 1 \\
 a^0 &= 1
 \end{aligned}$$



TC: no of recursive calls × time taken in each call

$$N \times O(1)$$

$$TC: O(N)$$

$$a^N = \begin{cases} a^{N/2} \cdot a^{N/2} & N \text{ is even} \\ a \cdot a^{N/2} \cdot a^{N/2} & N \text{ is odd} \end{cases}$$

```

int power(a, N)
{
    if (N == 0) return 1;

    if (N is even)
        return power(a, N/2) * power(a, N/2);
    else
        return a * power(a, N/2) * power(a, N/2);
}

```

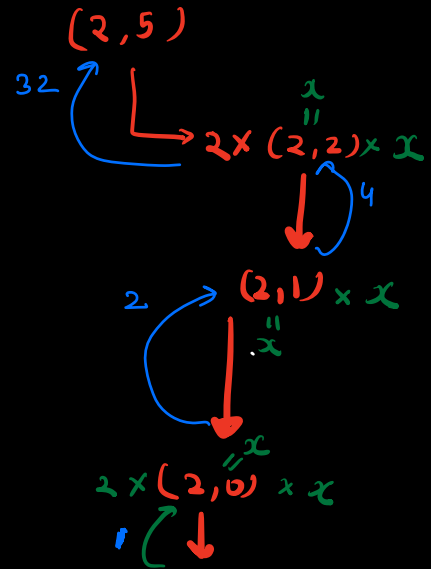
```

int power(a, N)
{
    if (N == 0) return 1;

    int x = power(a, N/2);

    if (N is even)
        return x * x;
    else
        return a * x * x;
}

```



$$T(N) = T(N-1) + 1$$

$$T(N-1) = T(N-2) + 1$$

$$T(N) = [T(N-2) + 1] + 1$$

$$T(N) = T(N-2) + 2$$

$$T(N-2) = T(N-3) + 1$$

$$T(N) = T(N-3) + 3$$

N
 \downarrow
 $N-1$
 \downarrow
 $N-2$
 \downarrow
 $N-3$
 \vdots
 \vdots
 \vdots

After k steps

$$T(N) = T(N-k) + k$$

$$k = N-1$$

$$N \rightarrow k = 1$$

$$k = N-1$$

$$N \times O(1)$$

$$= T(O(N))$$

$$T(N) = T(N - (N-1)) + N-1$$

$$T(N) = T(1) + N-1$$

$$T(N) = N$$

$$TC: O(N)$$

$$T(N) = 2T(N-1) + 1$$

$$T(N-1) = 2T(N-2) + 1$$

$$T(N) = 2 [2T(N-2) + 1] + 1$$

$$T(N) = 4T(N-2) + 2 + 1$$

$$T(N) = 4T(N-2) + 3$$

$$T(N-2) = 2T(N-3) + 1$$

$$T(N) = 8T(N-3) + 7$$

$$T(0) = 1$$

$$T(N) = 2^K T(N-K) + (2^K - 1)$$

$$\begin{aligned}
 T(N) &= 2^N T(0) + 2^{N-1} \\
 &= 2^N + 2^{N-1}
 \end{aligned}$$

$k=N$
 $N-k=0$
 $k=N$

$$T(N) = O(2^N)$$

$$T(N) = T\left(\frac{N}{2}\right) + 1$$

$$T(1) = 1$$

$$T\left(\frac{N}{2}\right) = T\left(\frac{N}{4}\right) + 1$$

$$T(N) = T\left(\frac{N}{4}\right) + 2$$

$$T\left(\frac{N}{4}\right) = T\left(\frac{N}{8}\right) + 1$$

$$T(N) = T\left(\frac{N}{8}\right) + 3$$

Break
10:16 10:26

$$T(N) = T\left(\frac{N}{2^k}\right) + K$$

$$\frac{N}{2^k} = 1$$

$$N = 2^k$$

$$2^k = N$$

$$k = \log N$$

$$k = \log N$$

$$T(N) = T(1) + \log N$$

$$T(N) = 1 + \log N$$

$$T(N) = O(\log N)$$

$$T(N) = 2T\left(\frac{N}{2}\right) + 1$$

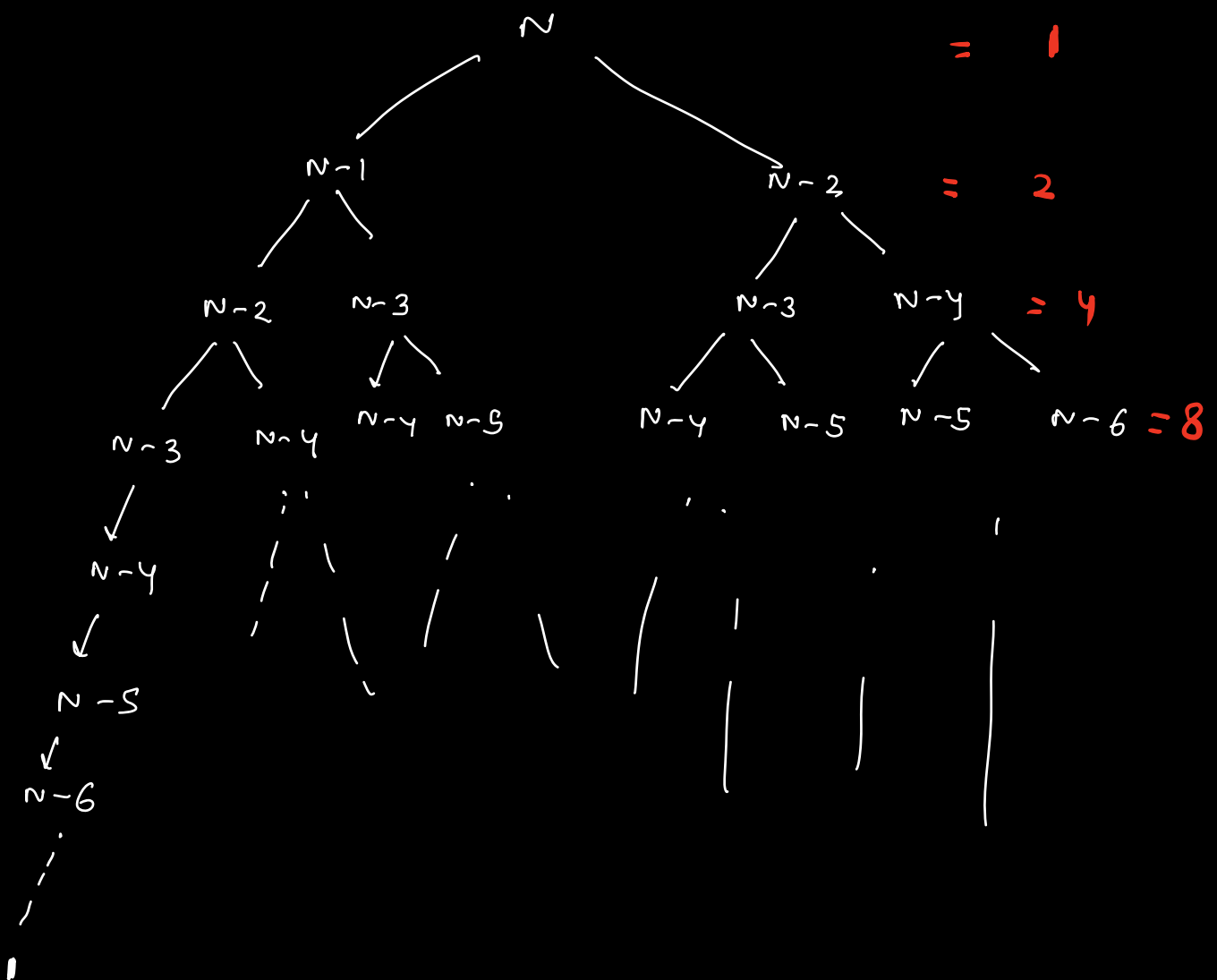
$$T(N) = T(N-1) + T(N-2) + 1$$

$$T.C: O(2^N)$$

$$T(N) = 2T(N-1) + 1$$

```
int fib(int N)
{
    if (N == 0 || N == 1) return
    return fib(N-1) + fib(N-2)
}
```

$$\text{No of recursion calls} \times O(1) = 2^N$$



N levels.

$$1 + 2 + 4 + 8 + \dots + 2^N$$

$$S_N = \frac{a(r^N - 1)}{r - 1} = \frac{1(2^N - 1)}{2 - 1} = 2^N - 1$$

$$TC: O(2^N)$$

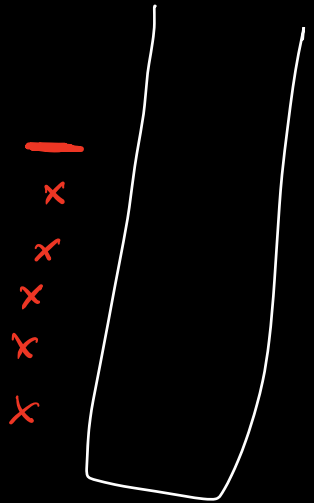
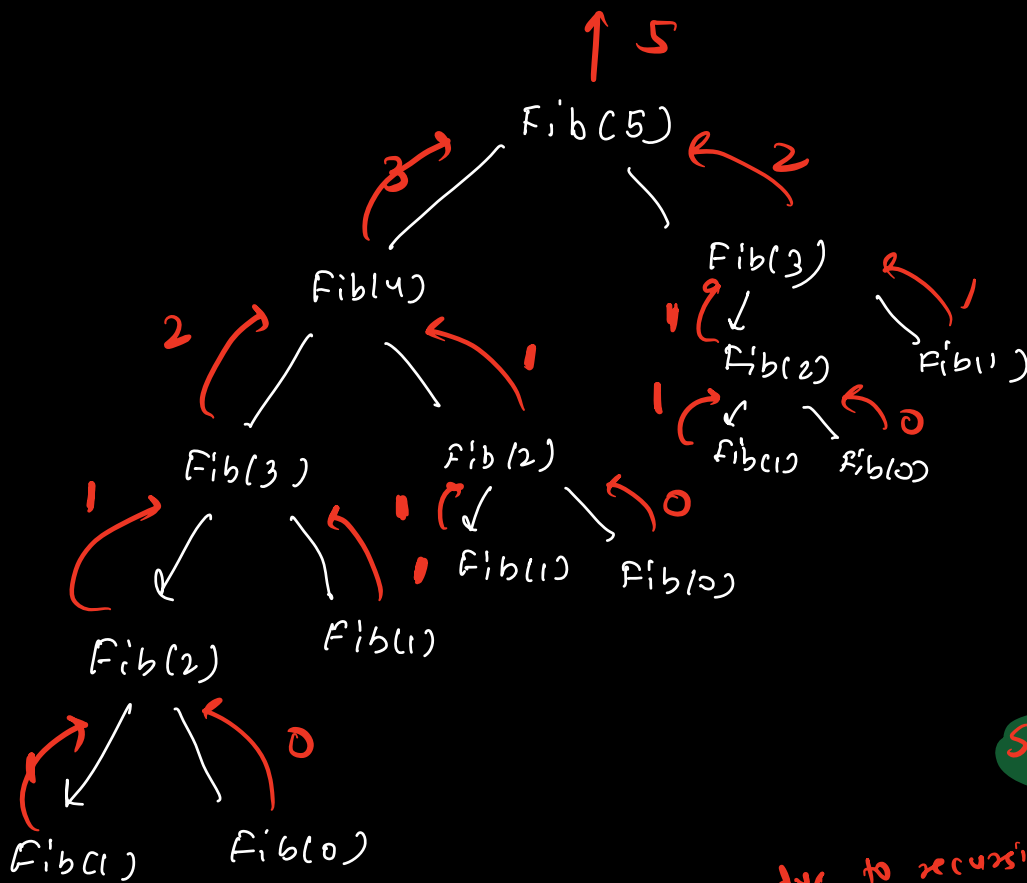
Space complexity. ^{consider both} $\left[\begin{array}{l} \text{function} \\ \text{stack} \end{array} , \begin{array}{l} \text{due to each} \\ \text{call} \end{array} \right]$

```
int sum ( int n )  
{  
    if ( n == 1 ) return 1  
    return sum ( n - 1 ) + n  
}
```

~~sum(1)~~
~~sum(2)~~
~~sum(3)~~
~~sum(4)~~
~~sum(5)~~

SC: $O(N)$

SC = No of ~~function calls~~



$SC: O(N)$

due to recursion stack

1

SC: max stack size at any point
or
No of level recursive tree.

~~1~~
~~2~~
~~...~~
 ~~$n/8$~~
 ~~$n/4$~~
 ~~$n/2$~~
 ~~n~~

$2 \leftarrow 2 \mid 2 \leftarrow 2$
 $3 \leftarrow 3 \mid 3 \leftarrow 3$
 $4 \leftarrow 4 \mid 4 \leftarrow 4$

SC: $O(\log N)$

Doubts

int factors ()
 {
 int arr[N]
 }
 int fib(n)

{
 int x = factors(x)
 return fib(n-1) + fib(n-2)
 }

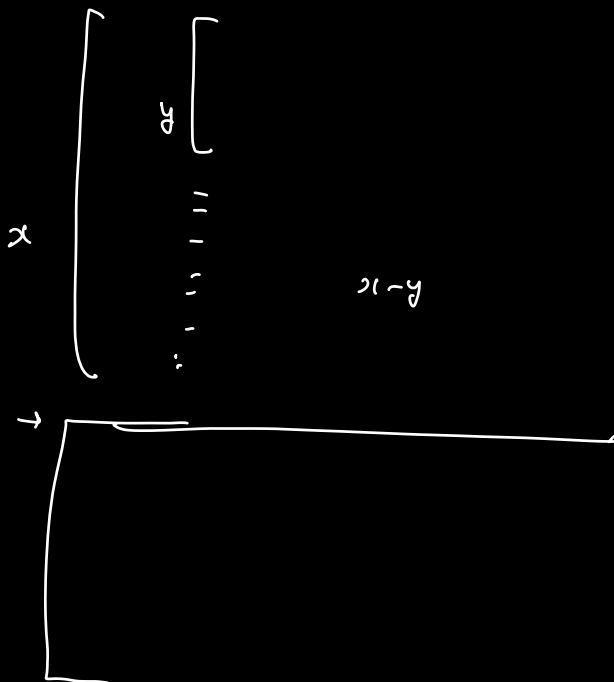
0
 1
 ...
 $n-4$
 $n-3$
 $n-2$
 $n-1$
 n

SC: $N + N = O(N)$

~~2^N~~

$N \times N$

$\tau_c: O(1)$



23rd - 3rd Jan
dec