

Lecture 19

AVL Tree Algorithms

Balanced Trees (Review)

BST Operations take $O(h)$ steps

- in the worst case, h can grow linearly in N

worst case $h = N - 1$
 $\rightarrow O(N) = h$

Strategy: Keep the tree balanced

- similar number of keys in left / right subtrees
- want $h = O(\log N)$

AVL Trees

- balanced by height
- height invariant: heights of left & right subtrees differ by at most 1
- base case: The empty tree is balanced

Minimum Tree Density — AVL Trees

Let $N(h)$ represent the number of nodes

- Q. What's the minimum value of $N(h)$?

$$N(h) \geq N(h-1) + N(h-2) + 1$$

$$N(-1) = 0, N(0) = 1$$

Pattern?

- Fibonacci sequence

$$\text{Claim: } N(h) \geq F_{h+3} - 1$$



h	min $N(h)$
-1	0
0	1
1	2
2	4
3	7
4	12

1 (F_2)
2 (F_3)
3 (F_4)
5 (F_5)
8 (F_6)
13 (F_7)

Minimum Tree Density — AVL Trees Cont'd

An AVL Tree holds ^{at least} $N(h) \geq F_{h+3} - 1$ nodes.

Proof (by induction on h):

Strategy:

- Base case? Both $h=-1$ & $h=0$ satisfy the claim
- Inductive case? Consider an AVL tree of height $k \geq 1$. Then:

$$N(k) \geq N(k-1) + N(k-2) + 1$$

$$\geq (F_{k+2} - 1) + (F_{k+1} - 1) + 1$$

$$= F_{k+2} + F_{k+1} - 1$$

$$= F_{k+3} - 1$$



F_n grows exponentially:

$$F_n \approx \frac{\varphi^n}{\sqrt{5}} \rightarrow N \geq \frac{\varphi^{h+3}}{\sqrt{5}-1}$$

$$\varphi = \frac{\sqrt{5}+1}{2}$$

\hookrightarrow golden ratio

$$h \leq \log_{\varphi} N$$

$$h = O(\log N)$$

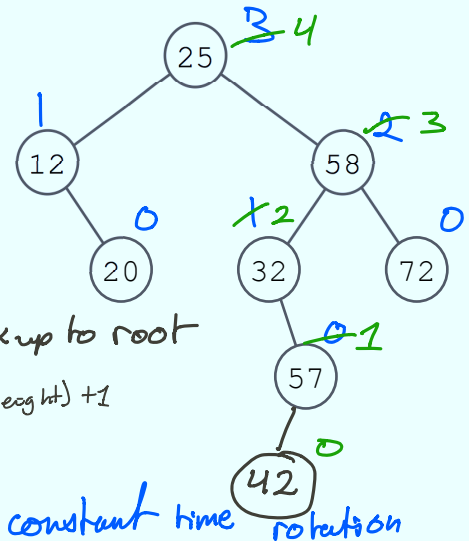
Tree Maintenance

Add `.height` property to each node

For `.insert(x)`, augment tree in the usual manner, and:

- update height fields bottom up (from x up to root)
 - $t \rightarrow \text{height} = \max(b \rightarrow \text{left} \rightarrow \text{height}, b \rightarrow \text{right} \rightarrow \text{height}) + 1$
- E.g., `.insert(42)`
- if node unbalanced, then perform a constant time rotation
 - rotate unbalanced ancestor closest to x

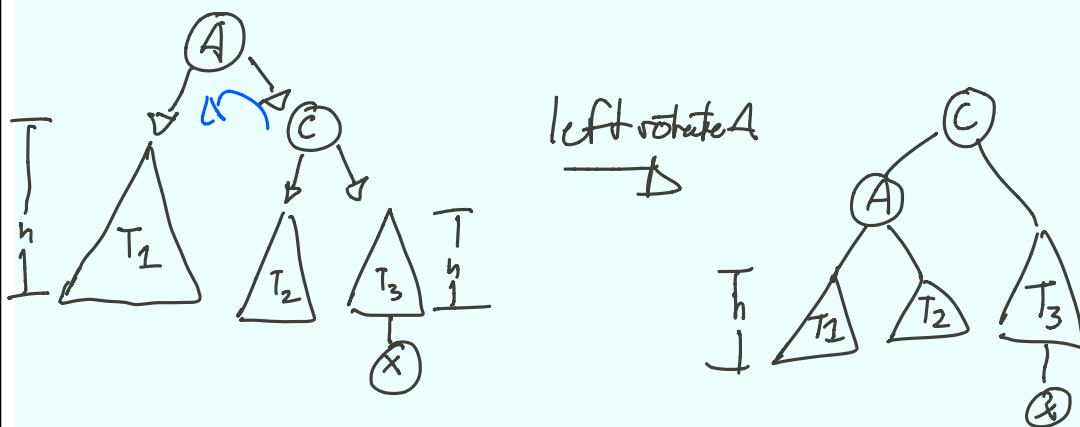
Same idea for `.delete(x)`



Rotations

Without loss of generality, assume left child has height h and x is inserted on the right subtree

Case 1: x is on the right / right grandchild



Case 2: x is on the right / left grandchild

