# FAT-Pointer based range addresses

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### Abstract

Capability-based addressing in the CHERI architecture is designed to improve hardware-level system security. These security guarantees can degrade runtime performance due to increased L1 TLB misses [1]. Our aim is to achieve capability-based security at reduced performance cost. We will use fat pointers to simplify and implement a recently proposed range TLB scheme [2], with the goal of reducing L1 TLB misses and page-walks for memory-intensive CHERI workloads.

#### Architecture C program Custom Allocator clang-morello march=morello+c64 Clang-morello -O3 -g -W -Wall -Wextra -Wno-unused--mabi=purecap ./<C program>.c parameter -shared -fPIC malloc.c -o malloc.so Malloc Shared Pure Cap binary Object file env LD PRELOAD=malloc.so ./a.out Execution Execution Opportunity Standard Current proposal to store mmap() mmap\_range() Opportunity using ARM 48 bit free Virtual address 48 bit free Virtual address Offset bounds Top [11:3] Permission Virtual Bottom[13:3] Address Virtual Address Physical Offset Physical non-contiguous Address allocation contiguous allocation L1 Cache Page TLB Virtual address offset Bounds Object type Permission Physical address Physical address Virtual Address Hardware Overview

## **Research Questions**

#### Research Questions:

- (RQ1) Can L1 misses and page-walks be reduced for memory intensive CHERI workloads?
- (RQ2) Is it possible to implement a simple Range TLB scheme with 8 bits using region of a fat pointer or ARM TBI?
- (RQ3) Can fat pointer based address translation lower the runtime overheads of hardware-level security in capability hardware like CHERI?

## Our Approach

Range translation is the process of mapping between contigous virtual pages mapped to contigous physical pages [3]. Recent work on FlexPointer[2] extended RMM[3] by encoding RangeIDs to the remaining 16 bits on a 64-bit virtual address. In the absence of fat pointers, their Range TLB was its own data structure in memory. We use fat pointers to simplify the implementation of Range TLB by using the bounds metadata to extract the lower and upper bounds which inturns allows us to encode the offset to the Pointer to map to the Physical page number. We hope to reproduce the FlexPointer results of reducing L1 tlb misses and page walks.

#### References

#### References

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