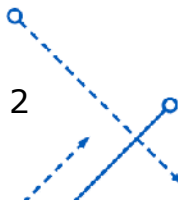


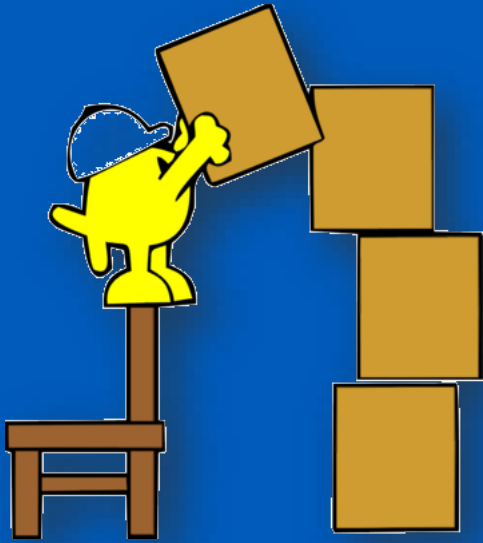
PROGRAMMING TECHNIQUES

Week 5: Dynamic Data Structures Stack & Queue

Today content

- Stacks
- Queues
- Walk through examples

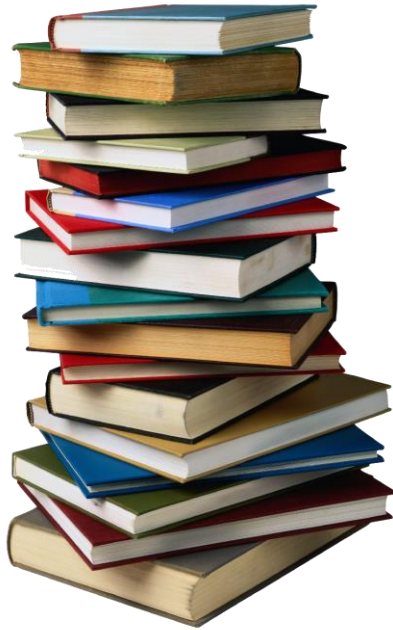




STACKS

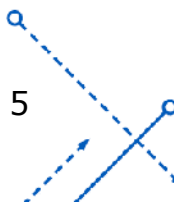
Stacks – Introduction

□ Popular images:



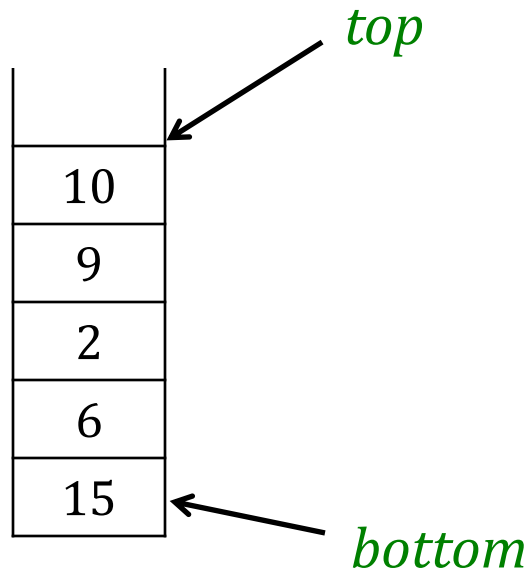
Stacks – Introduction

- We have talked about lists of data much of the term.
- One common use of a list is to represent a stack abstraction.
- Stacks allow us to add data and remove data at only one end (called the **top**).
- We can **push** data onto the top and **pop** data out of the top.



Stacks – Definition

- Stack is a dynamic set in which the element removed from the set is the one **most recently inserted**.
- Last-in, first out (**LIFO**) policy.



Implement a Stack

□ Using an array:

- An array of at most n elements: $S[0 \dots n-1]$
- An attribute top that indexes the most recently inserted element.
- Empty stack: $top = -1$

```
struct stack{  
    video* S; //array that hold the elements  
    int top; //index of the last element  
    int n;   //size of the array  
};
```

Implement a Stack

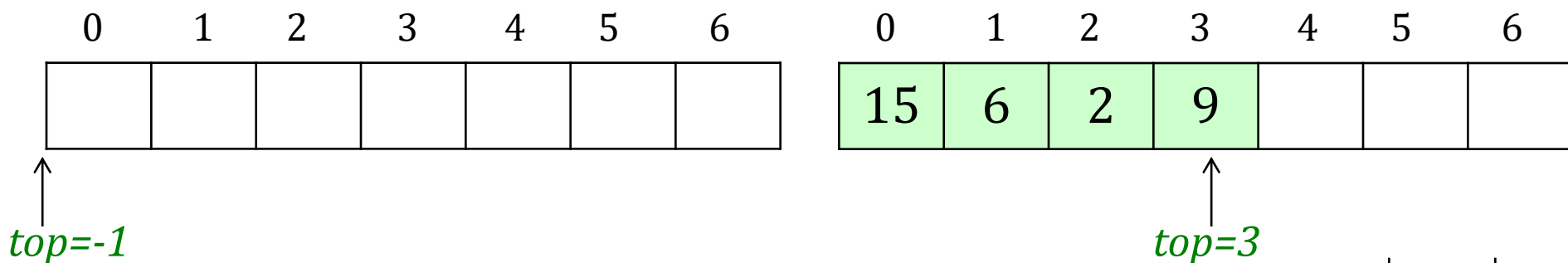
□ Using a linked list

- A pointer *pTop* points to the top of the stack
- Empty stack: *pTop* = NULL

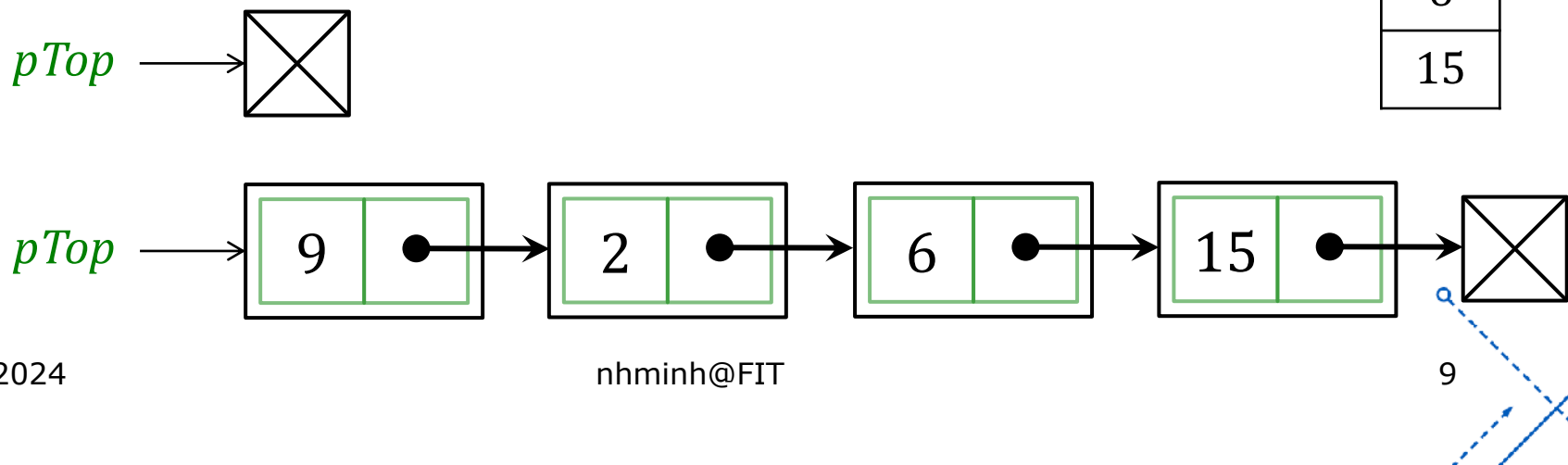
```
struct node {  
    video data;  
    node* pNext;  
};  
  
struct stack{  
    node* pTop; //pointer that points to the top of the stack  
};
```


Implement a Stack

□ Using an array:



□ Using a linked list:

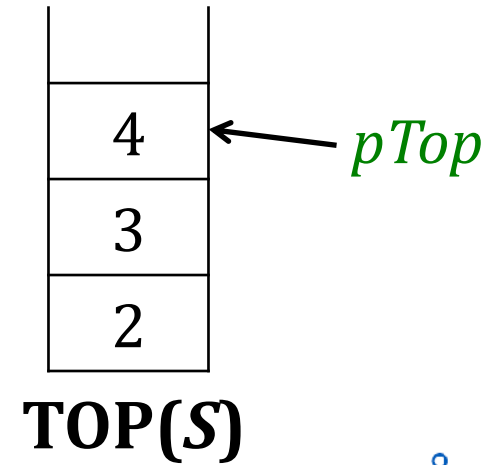
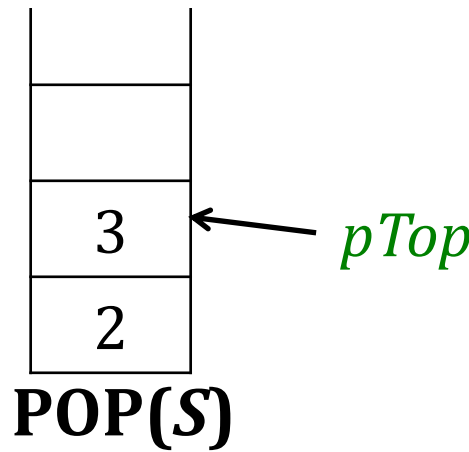
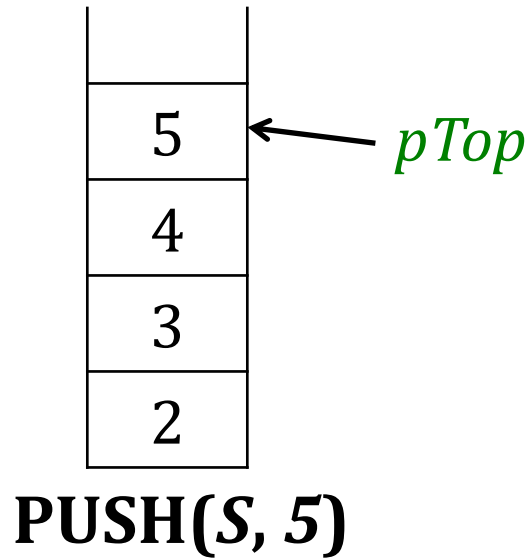
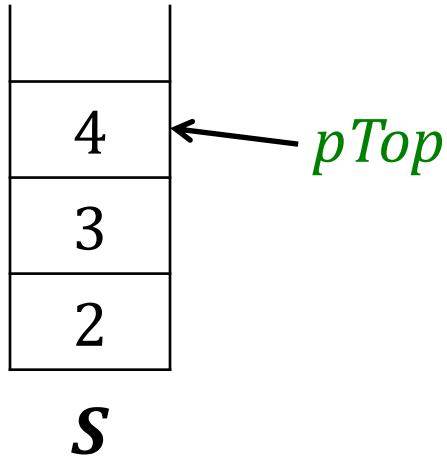


Stack Operations

- **STACK-EMPTY**: Check if a stack is empty
- **PUSH**: Insert a new element into the stack
- **POP**: Delete an element in the stack
- **TOP**: Get the top element without delete it from the stack
- **STACK-FULL**: Check if a stack is full



Stack Operations – Examples



Stack Operations

```
bool IsStackEmpty(stack s);  
bool IsStackFull(stack s);  
int Push(stack& s, const video& data);  
int Pop(stack& s, video& data);  
int Top(stack s, video& data);
```

Implement a Stack – using Linked List

- Implementing **Push** and **Pop** with a linear linked list data structure
 - represent very simple insert and remove algorithms
 - in fact, push is the same as adding at the beginning of a LLL
 - and, pop is the same as removing at the beginning of a LLL
 - why wouldn't we add/remove at the end (or tail) instead?



Implement a Stack – using Linked List

- Why wouldn't we add/remove at the end (or tail) instead?
 - pushing at the tail would be just as easy and efficient as pushing at the head iff we kept a tail pointer.
 - but, popping at the tail would require that we traverse to the (tail-1) node...regardless of whether or not there was a tail pointer.
 - a doubly linked list is not the answer



Stack Applications

☐ Delimiter Matching (part of any compiler)

- C++: () { } /* */ []
- Expression: $((a + b) \times ((c - d) + e))$
- Delimiters can be nested

☐ Idea:

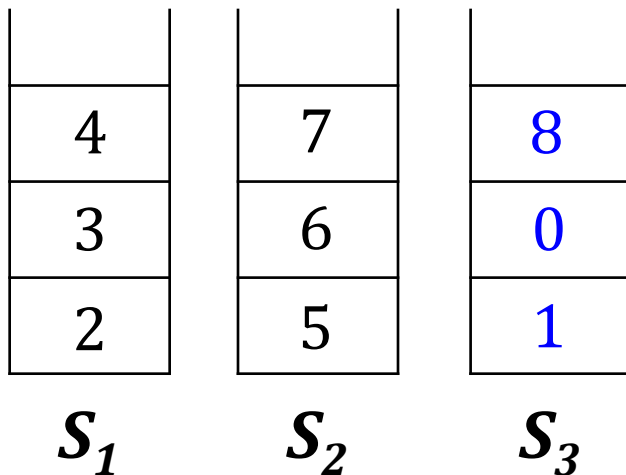
- Read the characters from left to right.
- Whenever you see a left (opening) delimiter, push it to the stack.
- Whenever you see a right (close) delimiter, pop the delimiter from the stack.
 - ☐ If you can't pop (because stack is empty) or they don't match, return false
 - ☐ Otherwise, continue reading characters
- If stack is empty when we finish reading string, return true
- Otherwise, return false.

Stack Applications

□ Adding two large number

- Treat these numbers as strings of numerals, store the numbers corresponding to these numerals on 2 stacks.
- Perform addition by popping numbers from the stacks

□ Example: 234 + 567



- $\text{Pop}(S_1) + \text{Pop}(S_2) = 4 + 7 = 11$
 → Add 1 to S_3
- $\text{Pop}(S_1) + \text{Pop}(S_2) = 3 + 6 + 1 = 10$
 → Add 0 to S_3
- $\text{Pop}(S_1) + \text{Pop}(S_2) = 2 + 5 + 1 = 8$
 → Add 8 to S_3
- S_1 S_2 are empty
 → Pop(S_3) and get the result: 801

Stack Applications

☐ Reverse Polish Notation (RPN):

- Every operator follows all its operands.

- Example:

 - ☐ $3 + 4 \rightarrow 3\ 4\ +$

 - ☐ $(3 - 4) \times 5 \rightarrow 3\ 4\ -\ 5\ \times$

- Implementations:

 - ☐ Mac OS X Calculator

 - ☐ iPhone apps: RPN Calculator

 - ☐ Android apps: RealCalc

Stack Applications

☐ Backtracking:

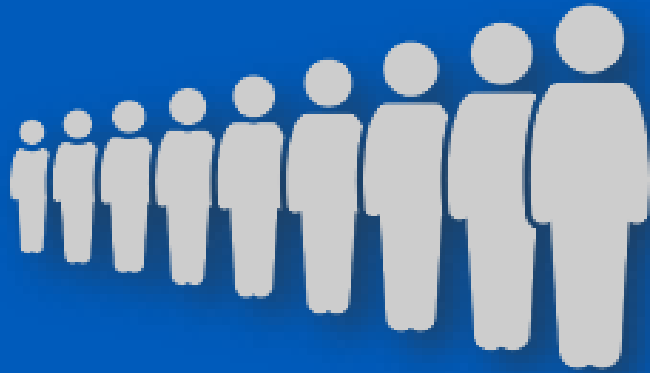
- Used in algorithms in which there are steps along some path from a starting point to some goal.
- Example:
 - ☐ Find your way through a maze
 - ☐ Find a route from one point to another point
 - ☐ Games in which there are moves to be made (chess, tic-tac-toe)



Stack Applications

□ Converting Decimal to Binary:

- Read a number
- Loop while number > 0
 - digit = number % 2
 - push digit to stack
 - number = number/2
- While stack is not empty
 - pop and print the top element of stack



QUEUES

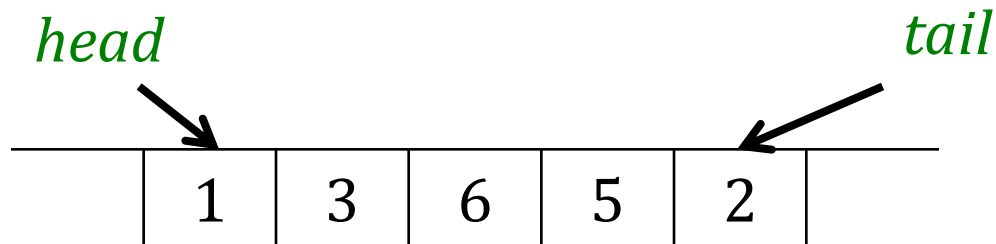
Queues – Introduction

- Another common use of a list is to represent a queue abstraction
- Queues allow us to add data at one end (the rear) and remove data at the other end (at the front)
- We can enqueue data at the rear and dequeue data at the front (FIFO)



Queue – Definition

- Queue is a dynamic set in which the element removed from the set is the one that has been inserted to the set for the longest time.
- First-in, First out (FIFO) policy.



Implement a Queue

□ Using Array

```
struct queue {  
    video* S; //list of elements  
    int n;    //max size of the list  
    int front; //the first element id  
    int tail;  //the last element id  
};
```

Implement a Queue

□ Using Linked List

```
struct node {  
    video data;  
    node* pNext;  
};  
  
struct queue{  
    node* pFront; //the head pointer  
    node* pRear; //the tail pointer  
};
```


Queue Operations

- ❑ **QUEUE-EMPTY**: Check if a queue is empty
- ❑ **ENQUEUE**: Insert a new element into the queue
- ❑ **DEQUEUE**: Delete an element in the queue
- ❑ **QUEUE-FULL**: Check if a queue is full
- ❑ **FIRST**: Get the head element without delete it from the queue



Queue Operations

```
bool IsQueueEmpty(queue s);  
bool IsQueueFull(queue s);  
int Enqueue(queue& s, const video& data);  
int Dequeue(queue& s, video& data);  
int First(queue s, video& data);
```

Queues

- Implementing **enqueue** and **dequeue** with a linear linked list data structure
 - are also simple
 - but, should the “rear” pointer point to the first or the last node? And, should the “front” pointer point to the first or the last node?
 - draw the pointer diagrams for either way and decide which would traverse less...



Queues

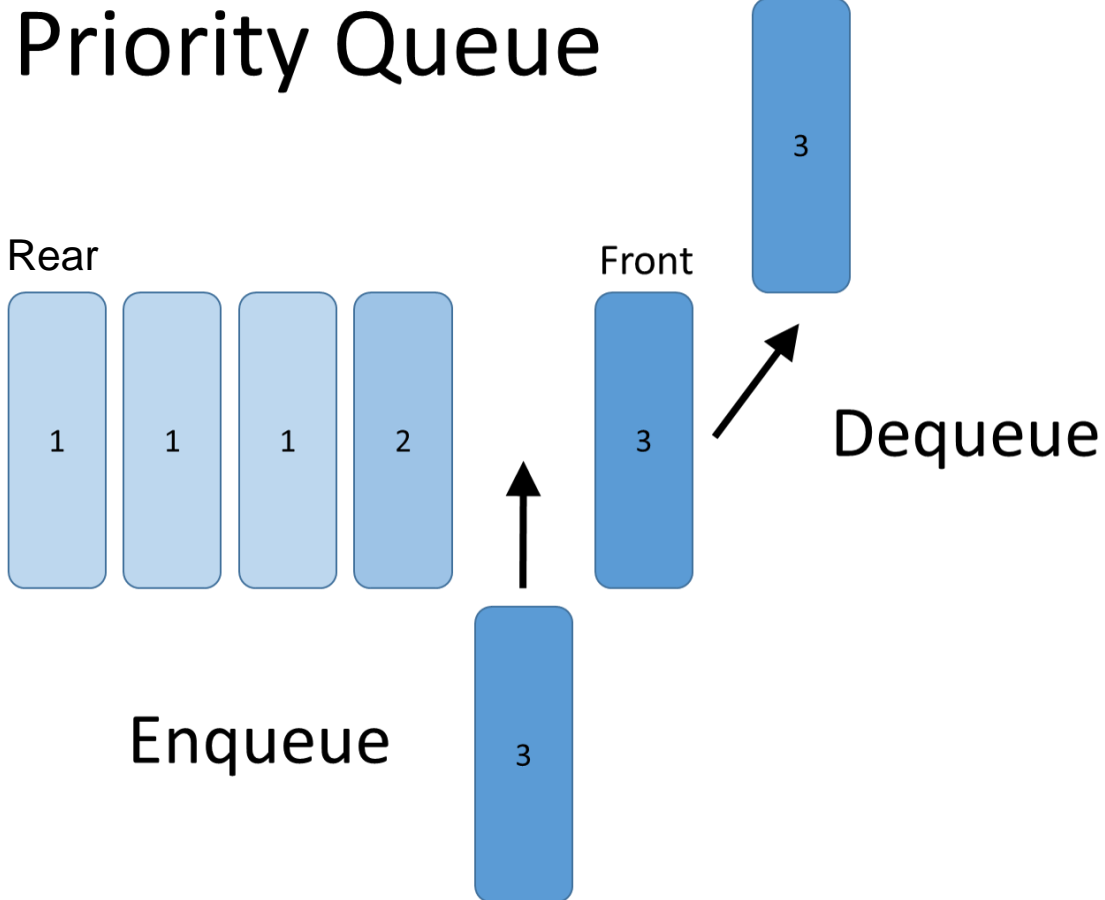
- ❑ enqueue should **add** at *the rear - the tail*
- ❑ dequeue should **remove** at *the front - the head*
- ❑ Why?
 - enqueueing at the front or rear is equally easy and efficient
 - but, dequeuing at the rear requires that we traverse to the “last-1” node (but a doubly linked list would be overkill). Luckily, dequeuing at the front is simple and efficient

Queue Applications

- ☐ Task Scheduling
- ☐ Resource Allocation
- ☐ Message buffering
- ☐ Traffic Management
- ☐ Print jobs, procedures management
- ☐ Download jobs in browsers.
- ☐ ...



Priority Queues



Priority Queues

- Items are ordered by **priority** rather than temporally.
- Queue implementations are modified to acquire a priority queue.
 - Enqueue must scan for the ***correct insertion point***.
- A linked list implementation is preferred.
- A **heap** can be used as the underlying implementation of a priority queue
 - *(to be discussed in heapsort algorithm – Data Structures & Algorithms)*



Next week

□ Midterm Examination

