

Evaluation of Function

Numerical Methods for IT

Polynomials

Given a polynomial $P_n(x) = a_0 + a_1x + a_2x^2 + \dots + a_nx^n$
and $v \in \mathbb{R}$.

Evaluation $P_n(v)$?

We have

$$\begin{aligned} P_n(v) &= a_0 + a_1v + a_2v^2 + \dots + a_{n-1}v^{n-1} + a_nv^n \\ &= a_0 + v(a_1 + a_2v + \dots + a_{n-1}v^{n-2} + a_nv^{n-1}) \\ &= a_0 + v * P_{n-1}(v), \text{ where} \end{aligned}$$

$$P_{n-1}(v) = b_0 + b_1v + \dots + b_{n-2}v^{n-2} + b_{n-1}v^{n-1}, b_i = a_{i+1}, i = 0, \dots, n-2$$

Let $P(n) = a[0 \dots n]$ is a 1-dimensional array that represents the computer representation of the polynomial $P_n(x)$. Given a value v . To evaluate $P_n(v)$, we can represent the recursive form as follows:

$$P(i, v) = \begin{cases} a[0], & i = 0 \\ a[0] + v * P(i, v), & i \geq 1 \end{cases}$$

Rational functions

To evaluate a rational function like:

$$R(x) = \frac{P_\mu(x)}{Q_v(x)} = \frac{p_0 + p_1x + \dots + p_\mu x^\mu}{q_0 + q_1x + \dots + q_v x^v}.$$

Given w , to evaluate $R(w)$, we can use the polynomial evaluation routine to:

- (1) $num \leftarrow P(\mu, w)$
- (2) $dom \leftarrow Q(v, w)$
- (3) Return num/dom

Evaluation of Continued Fractions

Continued fractions are often powerful ways of evaluating functions that occur in scientific applications.

A continued fraction looks like:

$$f(x) = b_0 + \frac{a_1}{b_1 + \frac{a_2}{b_2 + \frac{a_3}{b_3 + \dots}}} \equiv b_0 + \frac{a_1}{b_1 +} \frac{a_2}{b_2 +} \frac{a_3}{b_3 +} \dots$$

For example, the continued fraction representation of tangent function is: $\tan x = \frac{x}{1 -} \frac{x^2}{3 -} \frac{x^2}{5 -} \frac{x^2}{7 -} \dots$

Let f_n denote the result of $f(x)$ with coefficients through a_n and b_n . Then

$$f_n = \frac{A_n}{B_n},$$

where A_n, B_n are given by the following recurrence:

$$A_{-1} \equiv 1 ; B_{-1} \equiv 0$$

$$A_0 \equiv b_0 ; B_0 \equiv 1$$

$$A_j = b_j A_{j-1} + a_j A_{j-2} ; B_j = b_j B_{j-1} + a_j B_{j-2} ;$$

$$j = 1, 2, \dots, n.$$

Steed's method.

- . Set $f_0 = b_0$. If $b_0 = 0$, set $f_0 = \epsilon$.
- . Set $C_0 = f_0, D_0 = 0$.
- . For $j = 1, 2, \dots$:
 - . Set $D_j = b_j + a_j D_{j-1}$
 - . If $D_j = 0$, set $D_j = \epsilon$
 - . Set $C_j = b_j + a_j / C_{j-1}$.
 - . If $C_j = 0$, set $C_j = \epsilon$.
 - . Set $D_j = 1/D_j, \Delta = C_j D_j, f_j = f_{j-1} \Delta_j$.
 - . If $|\Delta_j - 1| < \epsilon_0$, then exit.

Chebyshev Approximation

The Chebyshev polynomial of degree n is denoted $T_n(x) = \cos(n \arccos x)$.

It can be combined with trigonometric identities to yield explicit expression for $T_n(x)$:

$$T_0(x) = 1 ; T_1(x) = x ; T_2(x) = 2x^2 - 1 ; T_3(x) = 4x^3 - 3x^2 ; T_4(x) = 8x^4 - 8x^2 + 1 \dots$$

$$T_{n+1}(x) = 2xT_n(x) - T_{n-1}(x), \forall n \geq 1 \text{ (*)}.$$

Suppose N is so large. We consider the truncated approximation:

$$f(x) \approx \left[\sum_{k=0}^{m-1} c_k T_k(x) \right] - \frac{1}{2} c_0 \text{ (**)} \text{ where } c_j = \frac{2}{N} \sum_{k=0}^{N-1} f \left[\cos \left(\frac{\pi(k+\frac{1}{2})}{N} \right) \right] \cos \left(\frac{\pi j(k+\frac{1}{2})}{N} \right).$$

Now that we have the Chebyshev coefficients, how to evaluate the approximation?

We can use the above recurrence relation (*) to generate values for $T_k(x)$ from $T_0 = 1, T_1 = x$, while also accumulating the sum of (**).

It is better to use Clenshaw's recurrence formula:

$$d_{m+1} \equiv d_m \equiv 0$$

$$d_j = 2xd_{j+1} - d_{j+2} + c_j, j = m-1, m-2, \dots, 1$$

$$f(x) \equiv d_0 = xd_1 - d_2 + \frac{1}{2} c_0.$$

Polynomial Approximation from Cheb. Coeff.

Convert the c_k 's into actual polynomial coefficients in the original variable x and have an approximation of the following form:

$$f(x) \approx \sum_{k=0}^{m-1} g_k x^k, a \leq x \leq b.$$

Given a coefficients $c[0..n-1]$, this routine returns

$d[0..n-1]: \sum_{k=0}^{n-1} d_k y^k = \sum_{k=0}^{n-1} c_k T_k(y) - c_0/2.$

Int k,j;

Doub sv;

VecDoub d(m), dd(m);

For (j=0;j<m;j++): d[0]=c[m-1];

For (j=m-2;j>0;j--):

 for (k=m-2;k>0;k--):

 sv=d[k];

 d[k]=2.0*d[k-1]-dd[k];

 dd[k]=sv;

 sv=d[0];

 d[0]=-dd[0]+c[j]

 dd[0]=sv

For (j=m-2;j>0;j--): d[j]=d[j-1]-dd[j];

d[0]=-dd[0]+0.5*c[0];

Return d

Given a coefficients $c[0..n-1]$, this routine returns

$g[0..n-1]: \sum_{k=0}^{n-1} d_k y^k = \sum_{k=0}^{n-1} g_k x^k - c_0/2.$ The

interval $-1 < y < 1$ is mapped to the interval $a < x < b$.

Int k,j,n=d.size();

Doub cnst=2.0/(b-a), fac=cnst;

VecDoub d(m), dd(m);

For (j=0;j<n;j++):

 d[j] *= fac

 fac *= cnst

cnst = 0.5*(a+b)

For (j=0; j< n-2; j++):

 for (k=n-2; k>=j; k--):

 d[k] -= cnst*d[k+1]