

Second Labs on Real-Time Scheduling

Report of TP2
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L2

Exercice 1:

Let's assume the following task sharing resources R1, R2 and R3:

	First release	WCET	D	P	Priority
T_1	6	$3: \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \$	6	20	4
T_2	4	$5: \begin{array}{ c c c c c c c c c c c c c c c c c c c$	11	20	3
T_3	2	$5: \boxed{ \mid R_2 \mid R_2 \mid R_2R_3 \mid }$	15	20	2
$oxed{T_4}$	0	$5: \begin{array}{ c c c c c c c c c c c c c c c c c c c$	18	20	1

Fig. 1: Task Configuration - Exercice 1

1. The simulation without a specific protocol for resource allocation is **not schedulable**, as shown in this capture, where task T_1 missed its deadline twice.

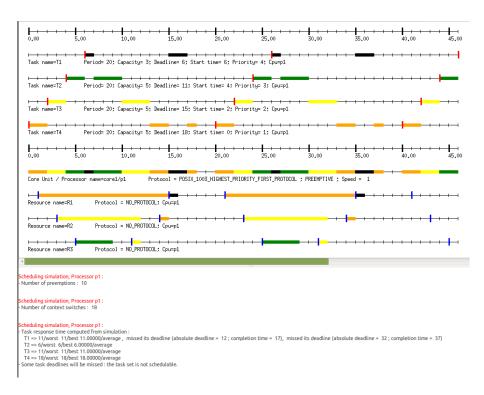


Fig. 2: Simulation of the task configuration without a specific protocol

2. The simulation with the *Priority Inheritance* protocol is **not schedulable**, as both T_1 and T_2 missed their deadlines twice in this case.

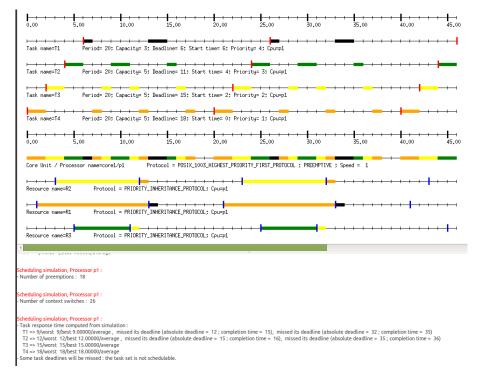


Fig. 3: Simulation of the task configuration with the Priority Inheritance protocol

3. Simulation with the Stack-based Protocol (Immediate Ceiling Inheritance Protocol) is

schedulable.

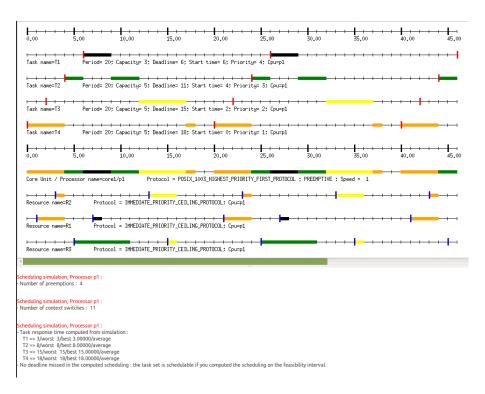


Fig. 4: Simulation of the task configuration with the Stack-based protocol

Exercice 2:

Let's assume the following task sharing resources R1, R2, R3 and R4:

	First release	WCET	D	P	Priority
T_1	6	$4: egin{bmatrix} R_4 & R_4R_3 \end{bmatrix}$	6	20	4
T_2	4	$4: \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \$	9	20	3
T_3	2	$4: \ \ \ \ R_2 \mid R_2R_1 \mid$	13	20	2
T_4	0	$4: egin{array}{ c c c c c c c c c c c c c c c c c c c$	16	20	1

Fig. 5: Task Configuration - Exercice 2

=> The simulation is **not schedulable** due to the capacities of the tasks, as shown in the figure below:



Fig. 6: Simulation of the task configuration

Exercice 3:

1. Simulate the following task configuration on one processor with two cores using a fully global Rate Monotonic scheduler:

	First release	WCET	D	P
T_1	0	2	3	3
T_2	0	2	4	4
T_3	0	7	12	12

Fig. 7: Task Configuration - Exercice 3-1

=> The simulation is **schedulable** because there are no deadline missed if we computed the scheduling on the feasibility interval.

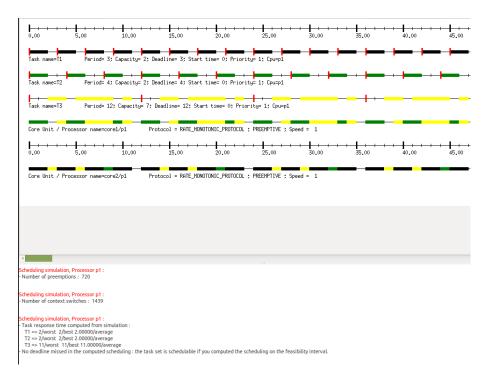


Fig. 8: Simulation of the task configuration with a fully global Rate Monotonic scheduler

2. Same question with the following configuration:

	First release	WCET	D	P
T_1	0	2	4	4
T_2	0	2	4	4
T_3	0	7	12	12

Fig. 9: Task Configuration - Exercice 3-2

^{=&}gt; The simulation is **not schedulable** as shown in this capture, where task T_3 missed its deadline three times.

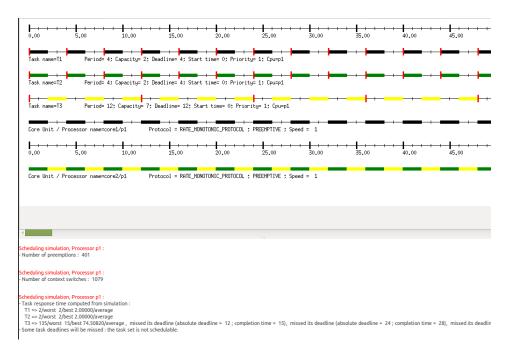


Fig. 10: Simulation of the task configuration with a fully global Rate Monotonic scheduler

3. From the two scheduling simulations, we see that when two tasks have the same periods and are synchronized, both cores of the processor run them at the same time, which affects the performance of the third task, T_3 . Even though T_1 has a lighter load in the second simulation, the tasks need to be different to better alternate and share the work on T_3 .

Exercice 4:

Let's assume the following configuration of tasks:

	First release	WCET	D	P
T_1	0	1	2	2
T_2	0	2	3	3
T_3	0	2	4	4

Fig. 11: Task Configuration - Exercice 4

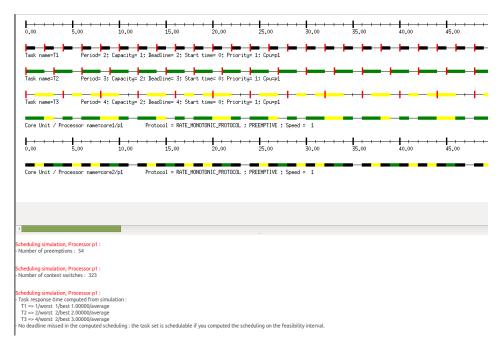


Fig. 12: Simulation of the task configuration

=> The worst response time doesn't occur at time 0, meaning the worst-case scenario isn't when all tasks start simultaneously. This issue arises from synchronized idle times, which lead to wasted processing time. Essentially, when two tasks finish simultaneously, T3 cannot be divided to fully utilize both cores, which contributes to an increase in its response time.

Exercice 5:

Let's assume the following configuration of tasks:

	WCET	D	P
T_1	2	6	6
T_2	4	8	8
T_3	3	10	10
T_4	12	20	20
T_5	1	50	50
T_6	20	50	50
T_7	5	100	100
T_8	1	100	100

Fig. 13: Task Configuration - Exercice 5

1. The configuration in this case is **schedulable** because there are no deadline missed if we computed the scheduling on the feasibility interval.

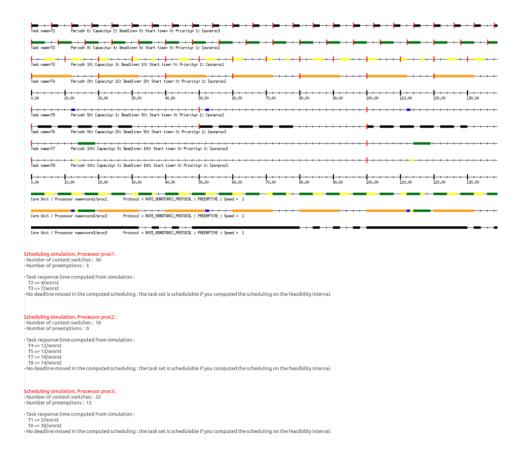


Fig. 14: Simulation of the task configuration with a partitioned First-Fit Rate Monotonic scheduler

2. The configuration in this case is also **schedulable** because there are no deadline missed if we computed the scheduling on the feasibility interval.

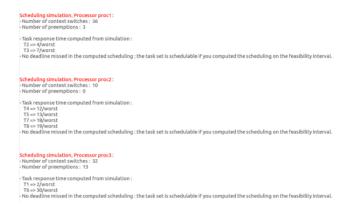


Fig. 15: Simulation of the task configuration with a partitioned First-Fit Earliest Deadline First scheduler