

# CSCA67 - Exercises #6

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For each of the following arguments, either prove the argument is valid by using Inference Rules or prove the argument is invalid by providing a counter-example world.

## 1.1

All birds eat at least one species of insect. All species of insects can fly. Therefore, all birds eat at least one flying species.

Let  $I(x)$  be “ $x$  is a species of insect”,  $B(x)$  be “ $x$  is a bird”,  $F(x)$  be “ $x$  flies”, and  $E(x,y)$  be “ $x$  eats  $y$ ”. Universe of discourse is live beings.

$\forall x, B(x) \rightarrow (\exists y, I(y) \wedge E(x,y))$	(1)
$\forall x, I(x) \rightarrow F(x)$	(2)
$\forall x, B(x) \rightarrow (\exists y, F(y) \wedge E(x,y))$	(Conclusion)
Take an arbitrary being $c$	(3)
$B(c) \rightarrow (\exists y, I(y) \wedge E(c,y))$	(4) (1,3, U.I)
Suppose $B(c)$	(5)
$\exists y, I(y) \wedge E(c,y)$	(5) (3,4, Implication)
Choose $d$ such that $I(d) \wedge E(c,d)$	(6)
$I(d) \wedge E(c,d)$	(7) (5, E.I)
$I(d)$	(8) (7, Simp.)
$F(d)$	(9) (2, 8, U.M.P)
$E(c,d)$	(10) (7, Simp.)
$F(d) \wedge E(c,d)$	(11) (9, 10, Conj.)
$\exists y, F(y) \wedge E(c,y)$	(12) (6, 11, E.G)
$B(c) \rightarrow (\exists y, F(y) \wedge E(c,y))$	(13) (5, 12, Implication)
$\forall x, B(x) \rightarrow (\exists y, F(y) \wedge E(x,y))$	(14) (3, 13, U.G)

Therefore the argument is valid. ■

## 1.2

Some smart people make lots of money. Some people who make lots of money buy very big houses. Therefore, some smart people buy very big houses.

Let  $S(x)$  stand for “ $x$  is smart”,  $M(x)$  be “ $x$  makes lots of money”, and  $H(x)$  be “ $x$  buys a big houses”, universe of discourse is people.

The argument given is:

$\exists x S(x) \wedge M(x)$	(1)
$\exists x, M(x) \wedge H(x)$	(2)
$\exists x, S(x) \wedge H(x)$	(Conclusion)

Let  $W \subseteq U$  s.t.  $W = \{a, b, c\}$

$\underline{a}$	$\underline{b}$	$\underline{c}$
$S(a) = \text{True}$	$S(b) = \text{False}$	$S(c) = \text{True}$
$M(a) = \text{True}$	$M(b) = \text{True}$	$M(c) = \text{False}$
$H(a) = \text{False}$	$H(b) = \text{True}$	$H(c) = \text{False}$

**Premise (1)**

$$\exists x, S(x) \wedge M(x)$$

Let  $x = a$

$$S(a) \wedge M(a)$$

$$T \wedge T$$

**T**

**Premise (2)**

$$\exists x, M(x) \wedge H(x)$$

Let  $x = b$ :

$$M(b) \wedge H(b)$$

$$T \wedge T$$

**T**

**Conclusion**

$$\exists x, S(x) \wedge H(x)$$

Let  $x = c$ :

$$S(c) \wedge H(c)$$

$$T \wedge F$$

**F**

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Thus, we have proved that there is a case such that all the premises are true, while the conclusion is false, making this argument invalid. ■

1.3

Everybody knows at least one song. Every song is sung by someone. People only sing songs they like. Therefore, everyone knows at least one song someone likes.  
Let  $S(x)$  be “ $x$  is a song”,  $P(x)$  be “ $x$  is a person”,  $K(x, y)$  be “ $x$  knows  $y$ ”,  $L(x, y)$  be “ $x$  likes  $y$ ”,  $S(x, y)$  be “ $x$  sings  $y$ ”. Universe of discourse is people and songs.

$\forall x, P(x) \rightarrow (\exists y, S(y) \wedge K(x, y))$	(1)
$\forall x, S(x) \rightarrow (\exists y, P(y) \wedge S(y, x))$	(2)
$\forall x \forall y, (P(x) \wedge S(y) \wedge S(x, y)) \rightarrow L(x, y)$	(3)
$\forall x, P(x) \rightarrow (\exists y, S(y) \wedge K(x, y) \wedge \exists z, (P(z) \wedge L(z, y)))$	(Conclusion)
Let a be arbitrary.	(4)
Suppose $P(a)$	(5)
$\exists y, S(y) \wedge K(a, y)$	(6) (1,5, Implication)
Choose b such that $S(b) \wedge K(a, b)$	(7)
$S(b) \wedge K(a, b)$	(8) (7, E.I)
$S(b)$	(9) (8, Simp.)
$K(a, b)$	(10) (8, Simp.)
$\exists y, P(y) \wedge S(y, b)$	(11) (2, 9, U.M.P)
Choose c such that $P(c) \wedge S(c, b)$	(12) (2, 9, U.M.P)
$P(c) \wedge S(c, b)$	(13) (12, E.I)
$S(b) \wedge P(c) \wedge S(c, b)$	(14) (9, 13, Conj.)
$L(c, b)$	(15) (3, 14, U.M.P)
$P(c)$	(16) (13, Simp.)
$P(c) \wedge L(c, b)$	(17) (15, 16, Conj.)
$\exists z, P(z) \wedge L(z, b)$	(18) (17, 12, E.G)
$S(b) \wedge K(a, b) \wedge (\exists z, P(z) \wedge L(z, b))$	(19) (18, 8, conj.)
$\exists y, S(y) \wedge K(a, y) \wedge (\exists z, P(z) \wedge L(z, y))$	(20) (19, 7, E.G)
$P(a) \rightarrow (\exists y, S(y) \wedge K(a, y) \wedge (\exists z, P(z) \wedge L(z, y)))$	(21) (19, 7, Implication)
$\forall x, P(x) \rightarrow (\exists y, S(y) \wedge K(x, y) \wedge \exists z, (P(z) \wedge L(z, y)))$	(22) (4, 21, U.G)

Therefore the argument is valid. ■

## 2.1

Prove: If x, y, and z are integers and  $x + y + z$  is odd, then at least one of x, y, z is odd

Universe of discourse is integers.

$$WTS : O(x + y + z) \rightarrow (O(x) \vee O(y) \vee O(z))$$

Lemmas.

$$(**) : E(x) \leftrightarrow \exists k, x = 2k$$

$$(***) : \neg E(x) = O(x)$$

Suppose $E(x) \wedge E(y) \wedge E(z)$	(1)
$(\exists k, x = 2k) \wedge (\exists j, y = 2j) \wedge (\exists i, z = 2i)$	(2) (1, (**))
$x = 2k$	(3) (2, E.I)
$y = 2j$	(4) (2, E.I)
$z = 2i$	(5) (2, E.I)
$x + y + z = 2(k + j + i)$	(6) 3, 4, 5, algebra
$E(x + y + z)$	(7) (6, (**))
$(E(x) \wedge E(y) \wedge E(z)) \rightarrow E(x + y + z)$	(8) (1, 7, Implication)
$\neg E(x + y + z) \rightarrow \neg(E(x) \wedge E(y) \wedge E(z))$	(9) (8, contra.)
$\neg E(x + y + z) \rightarrow (\neg E(x) \vee \neg E(y) \vee \neg E(z))$	(10) (9, DeM.)
$O(x + y + z) \rightarrow (O(x) \vee O(y) \vee O(z))$	(11) (10, (***))

## 2.2

Prove: If  $n$  is a positive integer, then  $n$  is even if and only if  $7n + 4$  is even.

Universe of discourse is  $\mathbb{Z}^+$ .

$$WTS : \forall n, E(n) \leftrightarrow E(7n + 4)$$

Lemmas.

$$(*) : O(x) \leftrightarrow \exists k, x = 2k + 1$$

$$(**) : E(x) \leftrightarrow \exists k, x = 2k$$

$$(***) : \neg O(x) = E(x)$$

Take an arbitrary positive integer  $n$

Suppose  $E(n)$

$$\exists k, n = 2k$$

$$n = 2k$$

$$7n + 4 = 7(2k) + 4$$

$$7n + 4 = 2(7k + 2)$$

$$E(7n + 4)$$

$$E(n) \rightarrow E(7n + 4)$$

Assume  $O(n)$

$$\exists k, n = 2k + 1$$

$$n = 2k + 1$$

$$7n + 4 = 7(2k + 1) + 4$$

$$7n + 4 = 14k + 7 + 4$$

$$7n + 4 = 14k + 11$$

$$7n + 4 = 2(7k + 5) + 1$$

$$O(7n + 4)$$

$$O(n) \rightarrow O(7n + 4)$$

$$\neg O(n) \rightarrow \neg O(7n + 4)$$

$$E(n) \rightarrow E(7n + 4)$$

$$(E(n) \rightarrow E(7n + 4)) \wedge (O(n) \rightarrow O(7n + 4))$$

$$(E(n) \leftrightarrow E(7n + 4))$$

$$\forall n, E(n) \leftrightarrow E(7n + 4)$$

(1)

(2)

(3)

(4) 3, E.I

(5) 4, algebra

(6) 5, algebra

(7) 6, (\*\*)

(8) (2, 7, Implication)

(9)

(10) (9, (\*))

(11) (10, E.I)

(12) (11, algebra)

(13) (12, algebra)

(14) (13, algebra)

(15) (14, algebra)

(16) (15, (\*\*))

(17) (9, 16, Implication)

(18) (18, Contr.)

(19) (18, (\*\*\*))

(20) (8, 20, Conj.)

(21) (20, bicond.)

(22) (1, 20, U.G)

### 2.3

Prove:  $(A \cap B \neq \emptyset \wedge A \subseteq C) \rightarrow (B \cap C \neq \emptyset)$

Suppose  $A \cap B \neq \emptyset \wedge A \subseteq C$

$$\exists x(x \in A \wedge x \in B)$$

Choose  $x$  such that  $x \in A \wedge x \in B$

$$x \in A$$

$$x \in B$$

$$\forall y(y \in A \rightarrow y \in C)$$

$$x \in C$$

$$x \in B \wedge x \in C$$

$$\exists x(x \in B \wedge x \in C)$$

$$B \cap C \neq \emptyset$$

$$(A \cap B \neq \emptyset \wedge A \subseteq C) \rightarrow (B \cap C \neq \emptyset)$$

(1)

(2) (1, Definition of  $\cap$ )

(3)

(4) (3, Simp.)

(5) (3, Simp.)

(6) (definition of  $\subseteq$ )

(7) (4,6, U.M.P)

(8) (5,7, Conj.)

(9) (3,8, E.G)

(10)(9, definition of  $\cap$ )

(11)(1, 10, Implication)

$\therefore$  As required to show. ■