

Pollard's Rho p Algo

factor n ?

Try ~~2~~ random nos.
Prob. of picking one no. is $= \frac{2}{n}$

$$n = 221 = 13 \times 17$$

① $\{2, 220\} \rightarrow$ 219 nos
search space \rightarrow

only 2 divide 221

Now

$$13 \times 1 = 13$$

$$13 \times 2 = 26$$

$$13 \times 3 = 39$$

$$13 \times 16 = 208$$

16
multiples
of
13

and

$$17 \times 1 = 17$$

$$17 \times 2 = 34$$

$$17 \times 3 = 51$$

$$17 \times 12 = 204$$

12
multiples
of
17

They have one thing common

— They all share a common factor with 221

— ~~∴~~ having found a no. that has a common factor with our n , then that common factor will be a factor of n .

② ∴ Instead of looking 2 nos between 2 & 221 why not look for one of the 28 (16+12) nos which share a common factor.
∴ we have improved a chance by >14 times

③ What is the prob. if we pick 2 nos between 1 & 221 & the difference between them has a common factor with 221?

No. of ways of picking 2 nos $= \binom{220}{2} = 48,841$

— out of 48,841 poss. combinations there are
only 6000 of them which results in the diff.
between ~~the~~ the no's having a factor in
comm. with 221.

— If no. 1 \nmid 221 there are 28 no's that
go with it for which the diff. is \times of either
13 or 17. for eg.

14, 18, 27, 35, 40, 52, 53, 66, 69, 79, 86, 92,
103, 105, 118, 220, 131, 137, 144, 154, 157, 170,
171, 183, 188, 188, 196, 205 & 209.



If you were to pick 2 (different) no's
between 1 & 221, there is a ~~greater~~ greater
~~chance~~ than one in 8 chance that the difference
between those no's would have a common factor
in comm. with 221

p-factoring algo

$$x_0 = \text{random}(0, n-1)$$

$$x_i = f(x_{i-1}) \pmod{n}, \quad i=1, 2, 3, \dots$$

$x_0 \rightarrow$ random starting value

$n \rightarrow$ no. to be factored

$f \in \mathbb{Z}[x]$ poly. with integer coeff.

usually, $f(x) = x^2 \pm a$, $a \neq -2, 0$

— If p is a prime, $\{x_i \pmod{p}\}_{i \geq 0}$ must eventually repeat.

Ex. $f(x) = x^2 + 1$, $x_0 = 0$ & $p = 563$

$\{x_i \pmod{p}\}_{i \geq 0}$ given as

$$x_0 = 0$$

$$x_1 = x_0^2 + 1 = 1$$

$$x_2 = x_1^2 + 1 = 2$$

$$x_3 = x_2^2 + 1 = 5$$

$$x_4 = x_3^2 + 1 = \underline{26}$$

$$x_5 = x_4^2 + 1 = 114$$

$$x_6 = x_5^2 + 1 = 48$$

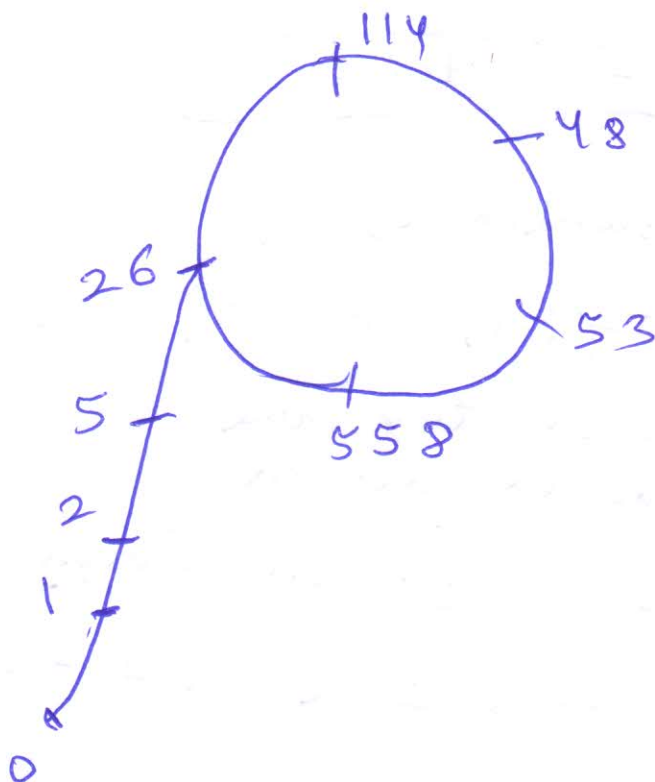
$$x_7 = x_6^2 + 1 = 53$$

$$x_8 = x_7^2 + 1 = 558$$

$$x_9 = x_8^2 + 1 = \underline{26}$$

i.e. $0, 1, 2, 5, \underline{26}, 114, 48, 53, 558$

- Fig 3



Verify! f cycle mod 1951 wrt
 $f(x) = x^2 + 1$
 $x_0 = 0$

To factor $n = 1098413 = 563 \cdot 1951$

we perform (gen. two seqs) $\begin{cases} \{x_i\} \\ \{y_i\} \end{cases} \text{ mod } n$
 $\gcd(x_i - y_i, n)$

$$x_0 = 0$$

$$y_i = x_{2i}$$

$$x_1 = x_0^2 + 1 = 1$$

$$x_2 = x_1^2 + 1 = 2 \quad y_1 = x_2 = 2 \quad \gcd(1-2, n) = 1$$

$$x_3 = x_2^2 + 1 = 5$$

$$x_4 = x_3^2 + 1 = 26$$

$$y_2 = x_4 = 26, \gcd(2-26, n) = 1$$

$$x_5 = x_4^2 + 1 = 677 \equiv 114$$

$$x_6 = x_5^2 + 1 = 458330 \equiv 48$$

$$y_3 = x_6 = 458330$$

$$\gcd(5-458330, n) = 1$$

$$x_7 = x_6^2 + 1 = 394716 \\ \equiv \underline{53}$$

$$x_8 = x_7^2 + 1 = 722324 \\ \equiv \underline{558}$$

$$y_4 = x_8 = 722324 \\ \gcd(26 - 722324, n) = 1$$

$$x_9 = x_8^2 + 1 = 203912 \\ \equiv \underline{26}$$

$$x_{10} = x_9^2 + 1 = 671773 \\ \equiv \underline{114}$$

$$y_5 = x_{10} = 671773 \\ \gcd(677 - 671773, n) \\ = \underline{563} \\ \therefore \text{factor found.}$$

Brent-Pollard's p-Method

Algo: n (composite \mathbb{Z}) > 1

Algo finds a non trivial factor d of n which is small compared with \sqrt{n} . Suppose for $f(x) = x^2 + 1$

① Initialization Choose a seed $x_0 = 2$

— $f(x) = x^2 + 1 \pmod{n}$

— Choose a value k not much bigger than \sqrt{d} perhaps $k < 100\sqrt{d}$

② Iteration & Computation

Compute $\{x_i\}$ & $\{y_i\}$ as follows:

$$x_1 = f(x_0)$$

$$x_2 = f(f(x_0)) = f(x_1)$$

$$x_3 = f(f(f(x_0))) = f(f(x_1)) = f(x_2)$$

$$x_i = f(x_{i-1})$$

⑤

$$y_1 = x_2 = f(x_1) = f(f(x_0)) = f(f(y_0))$$

$$y_2 = x_4 = f(x_3) = f(f(x_2)) = f(f(y_1))$$

$$y_3 = x_6 = f(x_5) = f(f(x_4)) = f(f(y_2))$$

⋮

$$y_i = x_{2i} = f(f(y_{i-1}))$$

and simultaneously compare x_i & y_i by
computing $d = \gcd(x_i - y_i, n)$

③ [factor found?] if $1 < d < n$ then d
is a non-trivial factor of n , print d
and go to ~~step~~ step ⑤

④ [another search?]

if $x_i \equiv y_i \pmod{n}$ for some i or $i \geq \sqrt{n}$

then go to step ① to choose a
new seed & a new gen. & repeat

⑤ [Exit] Terminate the algo.

[Conjecture] p (prime) $| n$ & $p = O(\sqrt{n})$

then f -algo has expected sum. time

$$O(\sqrt{n}) = O(\sqrt{n} (\log n)^2) = O(n^{1/4} (\log n)^2)$$

to find prime factor p of n

⑥