

# SC402: Introduction to Cryptography

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## Homework 2

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### Details:

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### Q.1

Suppose we consider a random throw of a pair of dice. Let  $X$  be the random variable defined on the set  $X = \{2, \dots, 12\}$ , obtained by considering the sum of two dice. Further, suppose that  $Y$  is a random variable which takes on the value  $D$  if the two dice are the same (i.e., if we throw “doubles”), and the value  $N$ , otherwise. Determine all the joint and conditional probabilities,  $\Pr[x, y]$ ,  $\Pr[x|y]$ , and  $\Pr[y|x]$ , where  $x \in \{2, \dots, 12\}$  and  $y \in \{D, N\}$ .

Ans:

$X = \{1, 2, 3, \dots, 12\}$

$Y = \{D, N\}$

X	Dies Rolls	P(x)	Count of D	Count of N
2	(1,1)	1/36	1	0
3	(1,2)(2,1)	2/36	0	2
4	(1,3)(2,2)(3,1)	3/36	1	2
5	(1,4)(2,3)(3,2)(4,1)	4/36	0	4
6	(1,5)(2,4)(3,3)(4,2)(5,1)	5/36	1	4
7	(1,6)(2,5)(3,4)(4,3)(5,2)(6,1)	6/36	0	6
8	(2,6)(3,5)(4,4)(5,3)(6,2)	5/36	1	4
9	(3,6)(4,5)(5,4)(6,3)	4/36	0	4
10	(4,6)(5,5)(6,4)	3/36	1	2
11	(5,6)(6,5)	2/36	0	2
12	(6,6)	1/36	1	0
			<b>6</b>	<b>30</b>

$P(X, Y) = P(X=2, Y=D) = 1/36$ ,  $P(X=2, Y=N)=0$

$P(X=3, Y=D)=0$	$P(X=3, Y=N)=2/36$
$P(X=4, Y=D)=1/36$	$P(X=4, Y=N)=2/36$
$P(X=5, Y=D)=0$	$P(X=5, Y=N)=4/36$
$P(X=6, Y=D)=1/36$	$P(X=6, Y=N)=4/36$
$P(X=7, Y=D)=0$	$P(X=7, Y=N)=6/36$
$P(X=8, Y=D)=1/36$	$P(X=8, Y=N)=4/36$
$P(X=9, Y=D)=0$	$P(X=9, Y=N)=4/36$
$P(X=10, Y=D)=1/36$	$P(X=10, Y=N)=2/36$
$P(X=11, Y=D)=0$	$P(X=11, Y=N)=2/36$
$P(X=12, Y=D)=1/36$	$P(X=12, Y=N)=0$

$P(Y/X)$  = Probability of Y given X

$P(Y=D/X=2)=1$	$P(Y=N/X=2)=0$
$P(Y=D/X=3)=0$	$P(Y=N/X=3)=1$
$P(Y=D/X=4)=3/3$	$P(Y=N/X=4)=2/3$
$P(Y=D/X=5)=0$	$P(Y=N/X=5)=1$
$P(Y=D/X=6)=3/5$	$P(Y=N/X=6)=4/5$
$P(Y=D/X=7)=0$	$P(Y=N/X=7)=1$
$P(Y=D/X=8)=1/5$	$P(Y=N/X=8)=4/5$
$P(Y=D/X=9)=0$	$P(Y=N/X=9)=1$
$P(Y=D/X=10)=1/3$	$P(Y=N/X=10)=2/3$
$P(Y=D/X=11)=0$	$P(Y=N/X=11)=1$
$P(Y=D/X=12)=1$	$P(Y=N/X=12)=0$

$P(X/Y)$  = Probability of X given Y

Now as  $Y=0$  is only possible for  $X \rightarrow$  even and once in every even value of X.

$$P\left(\frac{X \rightarrow \text{even}}{Y=0}\right) = \frac{1}{6} \quad \{x = 2, 4, 6, 8, 10, 12\}$$

$$P\left(\frac{X \rightarrow \text{odd}}{Y=0}\right) = 0 \quad \{x = 3, 5, 7, 9, 11\}$$

$P'(X/Y=N) \rightarrow$  Prob. of X given  $Y=N$ .

$$P(X=2/Y=N) = 0 \quad P(X=3/Y=N) = \frac{2}{30}$$

$$P(X=4/Y=N) = \frac{2}{30} \quad P(X=5/Y=N) = \frac{4}{30}$$

$$P(X=6/Y=N) = \frac{4}{30} \quad P(X=7/Y=N) = \frac{6}{30}$$

$$P(X=8/Y=N) = \frac{4}{30} \quad P(X=9/Y=N) = \frac{4}{30}$$

$$P(X=10/Y=N) = \frac{2}{30} \quad P(X=11/Y=N) = \frac{2}{30}$$

$$P(X=12/Y=N) = 0$$

## Q.2

Let  $P = \{a, b\}$  and let  $K = \{K_1, K_2, K_3, K_4, K_5\}$ . Let  $C = \{1, 2, 3, 4, 5\}$ , and suppose the encryption functions are represented by the following encryption matrix:

	$a$	$b$
$K_1$	1	2
$K_2$	2	3
$K_3$	3	1
$K_4$	4	5
$K_5$	5	4

Now choose two positive real numbers  $\alpha$  and  $\beta$  such that  $\alpha + \beta = 1$ , and define  $\Pr[K_1] = \Pr[K_2] = \Pr[K_3] = \alpha/3$  and  $\Pr[K_4] = \Pr[K_5] = \beta/2$ .

Prove that this cryptosystem achieves perfect secrecy.

Ans:

$$P = \{a, b\}$$

$$K = \{K_1, K_2, K_3, K_4, K_5\}$$

$$C = \{1, 2, 3, 4, 5\}$$

$$\Rightarrow \alpha + \beta = 1$$

$$\Pr[K_1] = \Pr[K_2] = \Pr[K_3] = \alpha/3$$

$$\Pr[K_4] = \Pr[K_5] = \beta/2$$

$$\text{Let } P(a) = x \text{ and } P(b) = 1-x$$

$$\text{Here we need to prove } P(X) = P(X/4)$$

$$Y=1,2,3$$

$$P(y) = P(X=a, y) + P(X=b, y)$$

$$= x(\alpha/3) + (1-x)(\alpha/3)$$

$$= \alpha/3 \dots\dots\dots(1)$$

$$P(y/x) = P(\text{Key}) \text{ where Key} = K_1, K_2, K_3$$

$$= P(K_1) = \alpha/3 \dots\dots\dots(2)$$

From Eq 1 and 2

$$P(y) = P(y/x)$$

For,

$$P(x/y) = P(x) * P(y/x) / P(y)$$

$$\text{So, } P(x/y) = P(x) \text{ -----} \rightarrow \text{for } Y=1,2,3$$

**Y=4,5**

$$\begin{aligned}P(y) &= P(x=a,y) + P(x=b,y) \\&= x(\beta/2) + (1-x)(\beta/2) \\&= \beta/2 \dots \dots \dots (3)\end{aligned}$$

$$\begin{aligned}P(y/x) &= P(\text{Key}) \text{ where Key}=K4,K5 \\&= P(\text{Key}) = \beta/2 \dots \dots \dots (4)\end{aligned}$$

Similarly,

$$\begin{aligned}P(x/y) &= P(x) * P(y/x) / P(y) \\&= P(x)\end{aligned}$$

$$\mathbf{P(x/y) = P(x)}$$

So now, We can say that, Cryptosystem achieve perfect secrecy.

### Q.3

**A) Prove that the Affine Cipher achieves perfect secrecy if every key is used with equal probability 1/312.**

**Ans:**

Probability = 1/312

Since equal probability for each k belong to K

We get  $P_k(K) = 1/312$

We show that there are exactly 12 keys that encrypt x to y for any combination of plain text – cipher text letter (x,y)

For every different value of 'a', the key (a, y - ax) encrypts 'x' to 'y'.

There are 12 keys that map a given plaintext letter to a given cipher text letter since there are 12 possibilities for a,

$$\begin{aligned}P_k(K) * P_p(d_k(y)) &= \frac{12}{312} P_p(a) + \frac{12}{312} (P_p(b)) + \dots + \frac{12}{312} P_p(z) \\&= \frac{12}{312} * 1 \\&= \frac{1}{26}\end{aligned}$$

$$\begin{aligned}\text{Also, } P_c(y/x) &= \sum_{K=x=dk(y)} P_k(K) \\ &= 12/312 \\ &= 1/26\end{aligned}$$

$$\begin{aligned}P_p(x/y) &= \frac{P_p(x) * P_c(y/x)}{P_c(y)} \\ &= \frac{P_p(x) * 1/26}{1/26} \\ &= P_p(x)\end{aligned}$$

So using Baye's theorem, we see that if we use every key with equal probability  $1/312$  then we can achieve perfect secrecy.

**B) More generally, suppose we are given a probability distribution on the set**

$$\{a \in \mathbb{Z}_{26} : \gcd(a, 26) = 1\}.$$

**Suppose that every key  $(a, b)$  for the Affine Cipher is used with probability  $\Pr[a]/26$ .**

**Prove that the Affine Cipher achieves perfect secrecy when this probability distribution is defined on the keyspace.**

**Ans:**

The probability of key  $(a, b)$  should be  $\Pr[a]/26$ .

Proceeding as in part (a), for any  $y \in \{1, \dots, n\}$ ,

$$\begin{aligned}P_x[Y=y] &= \sum_{x \in \{1, \dots, n\}} \sum_{a \in \mathbb{Z}_{26}^*} P_x[K=(a, b(x, y, a))] P_x[X=x] \\ &= \sum_{x \in \{1, \dots, n\}} \sum_{a \in \mathbb{Z}_{26}^*} \left( \Pr[a]/26 \right) \times P_x[X=x] \\ &= \sum_{x \in \{1, \dots, n\}} \frac{1}{26} \times P_x(X=x) \\ &= \frac{1}{26}\end{aligned}$$

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Then for any  $x, y \in \mathbb{Z}_{26}$ , we compute

$$\begin{aligned}
 P_x [y=y / x=x] &= \sum_{a \in \mathbb{Z}_{26}} P_x [K=(a, b(a, y, a))] \\
 &= \sum_{a \in \mathbb{Z}_{26}} P_x [a] / 26 \\
 &= \frac{1}{26}
 \end{aligned}$$

Finally using Bay's theorem, we see that

$$P_x [x=x / y=y] = P_x [x=x] \text{ for all } x, y$$

#### Q.4

Suppose that  $S$  is a random variable representing the sum of a pair of dice. Compute  $H(S)$ .

Ans:

$S$  is sum of pair of dies

So  $S \in \{2, 3, 4, 5, \dots, 12\}$

$$P(S=2) = 1/36 = p(S=12)$$

$$P(S=3) = 2/36 = p(S=11)$$

$$P(S=4) = 3/36 = p(S=10)$$

$$P(S=5) = 4/36 = p(S=9)$$

$$P(S=6) = 5/36 = p(S=8)$$

$$P(S=7) = 6/36$$

$$H(S) = \sum_{i=2}^{12} P(S=i) * \log_2 \frac{1}{P(S=i)}$$

$$H(S) = \frac{2 * \log_2 36}{36} + \frac{2 * 2 * \log_2 36/2}{36} + \frac{2 * 3 * \log_2 36/3}{36} + \frac{2 * 4 * \log_2 36/4}{36} + \frac{2 * 5 * \log_2 36/5}{36} + \frac{6 * \log_2 36/6}{36}$$

$$= \frac{2}{36} (\log_2 36 + 2 * \log_2 \frac{36}{2} + 3 * \log_2 \frac{36}{3} + 4 * \log_2 \frac{36}{4} + 5 * \log_2 \frac{36}{5} + 3 * \log_2 \frac{36}{6})$$

$$= 2.5686$$

**Q.5** Consider a cryptosystem in which  $P = \{a, b, c\}$ ,  $K = \{K_1, K_2, K_3\}$  and  $C = \{1, 2, 3, 4\}$ .

Suppose the encryption matrix is as follows:

	<i>a</i>	<i>b</i>	<i>c</i>
<i>K</i> <sub>1</sub>	1	2	3
<i>K</i> <sub>2</sub>	2	3	4
<i>K</i> <sub>3</sub>	3	4	1

Given that keys are chosen equiprobably, and the plaintext probability distribution is

$\Pr[a] = 1/2$ ,  $\Pr[b] = 1/3$ ,  $\Pr[c] = 1/6$ , compute  $H(P)$ ,  $H(C)$ ,  $H(K)$ ,  $H(K|C)$ , and

$H(P|C)$ .

**Ans.**

$$P = \{a, b, c\}$$

$$\Pr[a] = 1/2, \Pr[b] = 1/3, \Pr[c] = 1/6$$

$$K = \{K_1, K_2, K_3\} = \Pr[K_1] = \Pr[K_2] = \Pr[K_3] = 1/3$$

$$C = \{1, 2, 3, 4\}$$

$$\begin{aligned} H(P) &= \frac{\log_2 2}{2} + \frac{\log_2 3}{3} + \frac{\log_2 6}{6} \\ &= 0.50 + 0.53 + 0.43 \\ &= \mathbf{1.46} \end{aligned}$$

$$\begin{aligned} H(K) &= \frac{\log_2 3}{3} + \frac{\log_2 3}{3} + \frac{\log_2 3}{3} \\ &= \mathbf{1.58} \end{aligned}$$

$$\begin{aligned} \Pr[C=1] &= \Pr[a] * \Pr[K_1] + \Pr[c] * \Pr[K_3] \\ &= 1/6 + 1/18 \Rightarrow 2/9 \end{aligned}$$

$$\Pr[C=2] = 5/18 \quad \Pr[C=3] = 1/3 \quad \Pr[C=4] = 1/6$$

$$\begin{aligned} H(C) &= \frac{2}{9} (\log_2 9 - \log_2 2) + \frac{5}{18} (\log_2 18 - \log_2 5) + \frac{1}{3} (\log_2 3 - \log_2 1) + \frac{1}{6} (\log_2 6 - \log_2 1) \\ &= 0.48 + 0.51 + 0.53 + 0.43 \\ &= \mathbf{1.95} \end{aligned}$$

$$\begin{aligned} H(K|C) &= H(K) + H(P) - H(C) \\ &= \mathbf{1.09} \end{aligned}$$

**For  $H(P|C)$ ,**

$$\begin{aligned} \Pr[p=a, y], y=1,2,3 \\ &= 1/6 \end{aligned}$$

$$\begin{aligned} \Pr[p=b, y], y=2,3,4 \\ &= 1/9 \end{aligned}$$

$$\begin{aligned} \Pr[p=c, y], y=1,3,4 \\ &= 1/18 \end{aligned}$$

$$\begin{aligned} H(P,C) &= 3 \left( \frac{\log_2 6}{6} + \frac{\log_2 9}{9} + \frac{\log_2 18}{18} \right) \\ &\approx 3.044 \end{aligned}$$

$$\begin{aligned} H(P|C) &= H(P,C) - H(C) \\ &= 3.044 - 1.95 \\ &= \mathbf{1.094} \end{aligned}$$