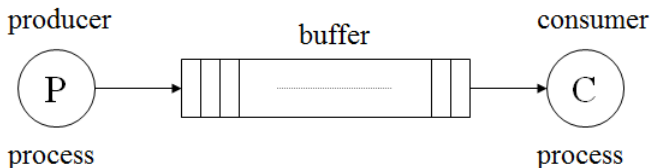


IE411: Operating Systems

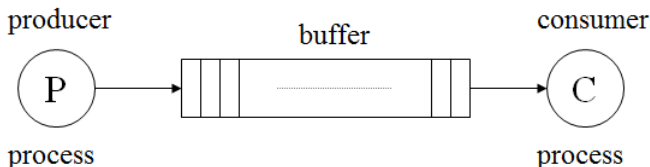
Producer-Consumer Problem

Producer/Consumer (bounded buffer) Problem



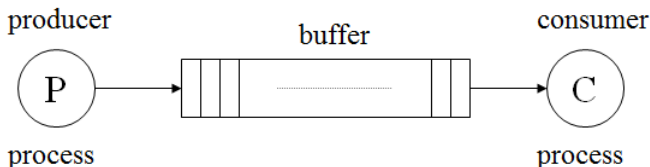
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- from time to time
 - the producer adds items to buffer
 - the consumer removes items from buffer

Producer/Consumer (bounded buffer) Problem



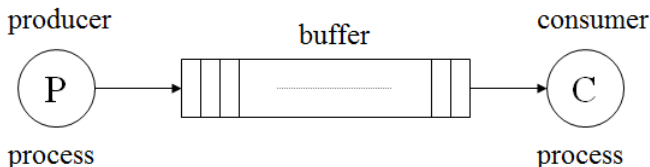
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Producer/Consumer (bounded buffer) Problem



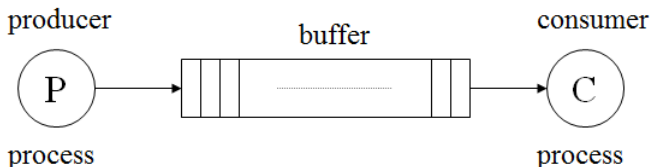
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 - the consumer must wait when the buffer is empty

Producer/Consumer (bounded buffer) Problem



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- from time to time
 - the producer adds items to buffer
 - the consumer removes items from buffer
- buffer is shared resource → synchronized access is required
 - the consumer must wait when the buffer is empty
 - the producer must wait when the buffer is full

Producer/Consumer (bounded buffer) Problem



- Assume shared, finite size buffer
- from time to time
 - the producer adds items to buffer
 - the consumer removes items from buffer
- buffer is shared resource → synchronized access is required
 - the consumer must wait when the buffer is empty
 - the producer must wait when the buffer is full

Bounded buffer example

- A bounded buffer is used when you pipe the output of one program into another
- Example: `grep foo file.txt | wc -l`
 - the grep process is the producer
 - The wc process is the consumer
 - Between them is an in-kernel bounded buffer

The Put and Get Routines (Version 1)

- Assume buffer can hold only one item

```
1      int buffer;
2      int count = 0;    // initially, empty
3
4      void put(int value) {
5          assert(count == 0);
6          count = 1;
7          buffer = value;
8      }
9
10     int get() {
11         assert(count == 1);
12         count = 0;
13         return buffer;
14     }
```

- typical solution would use a shared variable count

The Put and Get Routines (Version 1)

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```

- typical solution would use a shared variable count
- only put data into the buffer when count is zero (buffer is empty)

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8      }
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11         assert(count == 1);
12         count = 0;
13         return buffer;
14     }
```

- typical solution would use a shared variable count
- only put data into the buffer when count is zero (buffer is empty)
- only get data from the buffer when count is one (buffer is full)

Producer/Consumer Threads (Version 1)

- Using a condition variable (and mutex) to synchronise producer and consumer

```
1      cond_t cond;
2      mutex_t mutex;
3
4      void *producer(void *arg) {
5          int i, loops = (int) arg;
6          for (i = 0; i < loops; i++) {
7              Pthread_mutex_lock(&mutex);                // p1
8              if (count == 1)                            // p2
9                  Pthread_cond_wait(&cond, &mutex);    // p3
10             put(i);                                    // p4
11             Pthread_cond_signal(&cond);                // p5
12             Pthread_mutex_unlock(&mutex);              // p6
13         }
14     }
15
16     void *consumer(void *arg) {
17         int i, loops = (int) arg;
18         for (i = 0; i < loops; i++) {
19             Pthread_mutex_lock(&mutex);                // c1
```

Producer/Consumer Threads (Version 1)

```
20         if (count == 0)                                // c2
21             Pthread_cond_wait(&cond, &mutex);           // c3
22         int tmp = get();                                   // c4
23         Pthread_cond_signal(&cond);                      // c5
24         Pthread_mutex_unlock(&mutex);                    // c6
25         printf("%d\n", tmp);
26     }
27 }
```

- p1–p3: producer waits for the buffer to be empty
- c1–c3: consumer waits for the buffer to be full

Producer/Consumer Threads (Version 1)

```
20         if (count == 0)                                // c2
21             Pthread_cond_wait(&cond, &mutex);           // c3
22         int tmp = get();                                  // c4
23         Pthread_cond_signal(&cond);                      // c5
24         Pthread_mutex_unlock(&mutex);                    // c6
25         printf("%d\n", tmp);
26     }
27 }
```

- p1-p3: producer waits for the buffer to be empty
- c1-c3: consumer waits for the buffer to be full
- This code works for 1P and 1C. How about 1P and 2C?

Thread Trace: Broken Solution (Version 1)

```
void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // c1
        if (count == 0)                       // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get();                      // c4
        Pthread_cond_signal(&cond);           // c5
        Pthread_mutex_unlock(&mutex);         // c6
        printf("%d\n", tmp);
    }
}

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // p1
        if (count == 1)                       // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i);                               // p4
        Pthread_cond_signal(&cond);           // p5
        Pthread_mutex_unlock(&mutex);         // p6
    }
}
```

T_{c1}	State	T_{c2}	State	T_p	State	Count	Comment
c1	Running		Ready		Ready	0	Nothing to get
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	

Thread Trace: Broken Solution (Version 1)

```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);          // c1
        if (count == 0)                      // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get();                     // c4
        Pthread_cond_signal(&cond);          // c5
        Pthread_mutex_unlock(&mutex);        // c6
        printf("%d\n", tmp);
    }
}

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);          // p1
        if (count == 1)                      // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i);                             // p4
        Pthread_cond_signal(&cond);          // p5
        Pthread_mutex_unlock(&mutex);        // p6
    }
}

```

T _{c1}	State	T _{c2}	State	T _p	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep		Ready	p1	Running	0	
	Sleep		Ready	p2	Running	0	
	Sleep		Ready	p4	Running	1	Buffer now full
	Ready		Ready	p5	Running	1	T _{c1} awoken
	Ready		Ready	p6	Running	1	
	Ready		Ready	p1	Running	1	
	Ready		Ready	p2	Running	1	

Thread Trace: Broken Solution (Version 1)

```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);          // c1
        if (count == 0)                      // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get();                     // c4
        Pthread_cond_signal(&cond);          // c5
        Pthread_mutex_unlock(&mutex);        // c6
        printf("%d\n", tmp);
    }
}

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);          // p1
        if (count == 1)                     // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i);                             // p4
        Pthread_cond_signal(&cond);          // p5
        Pthread_mutex_unlock(&mutex);        // p6
    }
}

```

T _{c1}	State	T _{c2}	State	T _p	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep		Ready	p1	Running	0	
	Sleep		Ready	p2	Running	0	
	Sleep		Ready	p4	Running	1	Buffer now full
	Ready		Ready	p5	Running	1	T _{c1} awoken
	Ready		Ready	p6	Running	1	
	Ready		Ready	p1	Running	1	
	Ready		Ready	p2	Running	1	
	Ready		Ready	p3	Sleep	1	Buffer full; sleep
	Ready	c1	Running		Sleep	1	T _{c2} sneaks in ...
	Ready	c2	Running		Sleep	1	
	Ready	c4	Running		Sleep	0	... and grabs data
	Ready	c5	Running		Ready	0	T _p awoken
	Ready	c6	Running		Ready	0	

Thread Trace: Broken Solution (Version 1)

```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // c1
        if (count == 0)                       // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get();                      // c4
        Pthread_cond_signal(&cond);           // c5
        Pthread_mutex_unlock(&mutex);         // c6
        printf("%d\n", tmp);
    }
}

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // p1
        if (count == 1)                       // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i);                              // p4
        Pthread_cond_signal(&cond);           // p5
        Pthread_mutex_unlock(&mutex);         // p6
    }
}

```

T_{c1}	State	T_{c2}	State	T_p	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep		Ready	p1	Running	0	
	Sleep		Ready	p2	Running	0	
	Sleep		Ready	p4	Running	1	Buffer now full
	Ready		Ready	p5	Running	1	T_{c1} awoken
	Ready		Ready	p6	Running	1	
	Ready		Ready	p1	Running	1	
	Ready		Ready	p2	Running	1	
	Ready		Ready	p3	Sleep	1	Buffer full; sleep
	Ready	c1	Running		Sleep	1	T_{c2} sneaks in ...
	Ready	c2	Running		Sleep	1	
	Ready	c4	Running		Sleep	0	... and grabs data
	Ready	c5	Running		Ready	0	T_p awoken
	Ready	c6	Running		Ready	0	
c4	Running		Ready		Ready	0	Oh oh! No data

Thread Trace: Broken Solution (Version 1)

- After the producer woke T_{c1} , but before T_{c1} ever ran, the state of the bounded buffer changed by T_{c2}

Thread Trace: Broken Solution (Version 1)

- After the producer woke T_{c1} , but before T_{c1} ever ran, the state of the bounded buffer changed by T_{c2}
- There is no guarantee that when the woken thread runs, the state will still be as desired \rightarrow Mesa semantics
 - Virtually every system ever built employs Mesa semantics

Thread Trace: Broken Solution (Version 1)

- After the producer woke T_{c1} , but before T_{c1} ever ran, the state of the bounded buffer changed by T_{c2}
- There is no guarantee that when the woken thread runs, the state will still be as desired \rightarrow Mesa semantics
 - Virtually every system ever built employs Mesa semantics
- Hoare semantics provides a stronger guarantee that the woken thread will run immediately upon being woken

Producer/Consumer (Version 2)

- Consumer T_{c1} wakes up and re-checks the state of the shared variable
- If the buffer is empty, the consumer simply goes back to sleep

```
1      cond_t cond;
2      mutex_t mutex;
3
4      void *producer(void *arg) {
5          int i, loops = (int) arg;
6          for (i = 0; i < loops; i++) {
7              Pthread_mutex_lock(&mutex);           // p1
8              while (count == 1)                    // p2
9                  Pthread_cond_wait(&cond, &mutex); // p3
10             put(i);                                // p4
11             Pthread_cond_signal(&cond);            // p5
12             Pthread_mutex_unlock(&mutex);          // p6
13         }
14     }
15
```

Producer/Consumer (Version 2)

```
(Cont.)
16  void *consumer(void *arg) {
17      int i, loops = (int) arg;
18      for (i = 0; i < loops; i++) {
19          Pthread_mutex_lock(&mutex);           // c1
20          while (count == 0)                    // c2
21              Pthread_cond_wait(&cond, &mutex); // c3
22          int tmp = get();                       // c4
23          Pthread_cond_signal(&cond);           // c5
24          Pthread_mutex_unlock(&mutex);         // c6
25          printf("%d\n", tmp);
26      }
27  }
```

- A simple rule to remember with condition variables is to always use while loops
- However, this code still has a bug!

Thread Trace: Broken Solution (Version 2)

```
void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);          // c1
        while (count == 0)                   // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get();                     // c4
        Pthread_cond_signal(&cond);          // c5
        Pthread_mutex_unlock(&mutex);        // c6
        printf("%d\n", tmp);
    }
}

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);          // p1
        while (count == 1)                   // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i);                             // p4
        Pthread_cond_signal(&cond);          // p5
        Pthread_mutex_unlock(&mutex);        // p6
    }
}
```

T_{c1}	State	T_{c2}	State	T_p	State	Count	Comment
c1	Running		Ready		Ready	0	Nothing to get
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	

- T_{c1} finds the buffer empty so it is waiting (line c3)

Thread Trace: Broken Solution (Version 2)

```
void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // c1
        while (count == 0)                    // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get();                      // c4
        Pthread_cond_signal(&cond);           // c5
        Pthread_mutex_unlock(&mutex);          // c6
        printf("%d\n", tmp);
    }
}

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // p1
        while (count == 1)                   // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i);                              // p4
        Pthread_cond_signal(&cond);           // p5
        Pthread_mutex_unlock(&mutex);          // p6
    }
}
```

T_{c1}	State	T_{c2}	State	T_p	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep	c1	Running		Ready	0	
	Sleep	c2	Running		Ready	0	
	Sleep	c3	Sleep		Ready	0	Nothing to get

- T_{c1} finds the buffer empty so it is waiting (line c3)
- T_{c2} finds the buffer empty so it is waiting (line c3)

Thread Trace: Broken Solution (Version 2)

```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // c1
        while (count == 0)                    // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get();                      // c4
        Pthread_cond_signal(&cond);           // c5
        Pthread_mutex_unlock(&mutex);         // c6
        printf("%d\n", tmp);
    }
}

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // p1
        while (count == 1)                    // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i);                               // p4
        Pthread_cond_signal(&cond);           // p5
        Pthread_mutex_unlock(&mutex);         // p6
    }
}

```

T_{c1}	State	T_{c2}	State	T_p	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep	c1	Running		Ready	0	
	Sleep	c2	Running		Ready	0	
	Sleep	c3	Sleep		Ready	0	Nothing to get
	Sleep		Sleep	p1	Running	0	
	Sleep		Sleep	p2	Running	0	
	Sleep		Sleep	p4	Running	1	Buffer now full
	Ready		Sleep	p5	Running	1	T_{c1} awoken
	Ready		Sleep	p6	Running	1	
	Ready		Sleep	p1	Running	1	
	Ready		Sleep	p2	Running	1	
	Ready		Sleep	p3	Sleep	1	Must sleep (full)

- Both T_{c1} and T_{c2} initially find buffer empty so they are waiting (line c3)
- T_p adds an item to buffer (line p4), signals cond (line p5), waking up T_{c1} , waits on cond until signaled (line p3)

Thread Trace: Broken Solution (Version 2)

```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // c1
        while (count == 0)                    // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get();                      // c4
        Pthread_cond_signal(&cond);           // c5
        Pthread_mutex_unlock(&mutex);         // c6
        printf("%d\n", tmp);
    }
}

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // p1
        while (count == 1)                   // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i);                              // p4
        Pthread_cond_signal(&cond);           // p5
        Pthread_mutex_unlock(&mutex);         // p6
    }
}

```

T_{c1}	State	T_{c2}	State	T_p	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep	c1	Running		Ready	0	
	Sleep	c2	Running		Ready	0	
	Sleep	c3	Sleep		Ready	0	Nothing to get
	Sleep		Sleep	p1	Running	0	
	Sleep		Sleep	p2	Running	0	
	Sleep		Sleep	p4	Running	1	Buffer now full
	Ready		Sleep	p5	Running	1	T_{c1} awoken
	Ready		Sleep	p6	Running	1	
	Ready		Sleep	p1	Running	1	
	Ready		Sleep	p2	Running	1	
	Ready		Sleep	p3	Sleep	1	Must sleep (full)
c2	Running		Sleep		Sleep	1	Recheck condition
c4	Running		Sleep		Sleep	0	T_{c1} grabs data
c5	Running		Ready		Sleep	0	Oops! Woke T_{c2}
c6	Running		Ready		Sleep	0	
c1	Running		Ready		Sleep	0	
c2	Running		Ready		Sleep	0	
c3	Sleep		Ready		Sleep	0	Nothing to get

- T_{c1} removes item to buffer (line c4), signals cond (line c5), waking up T_{c2} , finds buffer empty, goes to sleep waiting on cond until signaled (line c3)

Thread Trace: Broken Solution (Version 2)

```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // c1
        while (count == 0)                    // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get();                     // c4
        Pthread_cond_signal(&cond);           // c5
        Pthread_mutex_unlock(&mutex);         // c6
        printf("%d\n", tmp);
    }
}

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // p1
        while (count == 1)                   // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i);                              // p4
        Pthread_cond_signal(&cond);           // p5
        Pthread_mutex_unlock(&mutex);         // p6
    }
}

```

T_{c1}	State	T_{c2}	State	T_p	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep	c1	Running		Ready	0	
	Sleep	c2	Running		Ready	0	
	Sleep	c3	Sleep		Ready	0	Nothing to get
	Sleep		Sleep	p1	Running	0	
	Sleep		Sleep	p2	Running	0	
	Sleep		Sleep	p4	Running	1	Buffer now full
	Ready		Sleep	p5	Running	1	T_{c1} awoken
	Ready		Sleep	p6	Running	1	
	Ready		Sleep	p1	Running	1	
	Ready		Sleep	p2	Running	1	
	Ready		Sleep	p3	Sleep	1	Must sleep (full)
c2	Running		Sleep		Sleep	1	Recheck condition
c4	Running		Sleep		Sleep	0	T_{c1} grabs data
c5	Running		Ready		Sleep	0	Oops! Woke T_{c2}
c6	Running		Ready		Sleep	0	
c1	Running		Ready		Sleep	0	
c2	Running		Ready		Sleep	0	
c3	Sleep		Ready		Sleep	0	Nothing to get
	Sleep	c2	Running		Sleep	0	
	Sleep	c3	Sleep		Sleep	0	Everyone asleep... ¹⁰

- T_{c2} , being woken up by T_{c1} , finds buffer empty, goes to sleep waiting on cond (line c3)
- Everyone is sleeping \rightarrow P can't produce \rightarrow no forward progress

Approach 1: Wake up everyone

- When not sure if next waiting thread is the right one to wake up, just wake up all

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- When not sure if next waiting thread is the right one to wake up, just wake up all
- Not the most elegant solution
 - Probably bad for performance: all awoken threads will compete for mutex again
 - But a good fallback mechanism to ensure correctness

Approach 1: Wake up everyone

- When not sure if next waiting thread is the right one to wake up, just wake up all
- Not the most elegant solution
 - Probably bad for performance: all awoken threads will compete for mutex again
 - But a good fallback mechanism to ensure correctness
- Need a new API: `cond_broadcast(cv)`
 - Semantics: wakes up all the threads waiting on `cv`

Approach 2: Use multiple CVs

- Two different conditions in bounded buffer problem
 - waiting for buffer to become empty (parent)
 - waiting for buffer to become full (child)

Approach 2: Use multiple CVs

- Two different conditions in bounded buffer problem
 - waiting for buffer to become empty (parent)
 - waiting for buffer to become full (child)
- Use a separate CV for each condition using `cond_wait()` and `cond_signal()`
- More elegant, better-performing solution than using `cond_broadcast()`

Single buffer solution: 2 CVs

- Producer threads wait on the condition `empty` and signal `fill`
- Consumer threads wait on `fill` and signal `empty`

```
1      cond_t empty, fill;
2      mutex_t mutex;
3
4      void *producer(void *arg) {
5          int i;
6          for (i = 0; i < loops; i++) {
7              Pthread_mutex_lock(&mutex);
8              while (count == 1)
9                  Pthread_cond_wait(&empty, &mutex);
10             put(i);
11             Pthread_cond_signal(&fill);
12             Pthread_mutex_unlock(&mutex);
13         }
14     }
15
```

Single buffer solution: 2 CVs

(Cont.)

```
16     void *consumer(void *arg) {
17         int i;
18         for (i = 0; i < loops; i++) {
19             Pthread_mutex_lock(&mutex);
20             while (count == 0)
21                 Pthread_cond_wait(&fill, &mutex);
22             int tmp = get();
23             Pthread_cond_signal(&empty);
24             Pthread_mutex_unlock(&mutex);
25             printf("%d\n", tmp);
26         }
27     }
```

Multiple buffer solution

```
1      int buffer[MAX];
2      int fill = 0;
3      int use = 0;
4      int count = 0;
5
6      void put(int value) {
7          buffer[fill] = value;
8          fill = (fill + 1) % MAX;
9          count++;
10     }
11
12     int get() {
13         int tmp = buffer[use];
14         use = (use + 1) % MAX;
15         count--;
16         return tmp;
17     }
```

The Final Put and Get Routines

Multiple buffer solution

```
1      cond_t empty, fill;
2      mutex_t mutex;
3
4      void *producer(void *arg) {
5          int i;
6          for (i = 0; i < loops; i++) {
7              Pthread_mutex_lock(&mutex);           // p1
8              while (count == MAX)                 // p2
9                  Pthread_cond_wait(&empty, &mutex); // p3
10             put(i);                               // p4
11             Pthread_cond_signal(&fill);           // p5
12             Pthread_mutex_unlock(&mutex);         // p6
13         }
14     }
15
16     void *consumer(void *arg) {
17         int i;
18         for (i = 0; i < loops; i++) {
19             Pthread_mutex_lock(&mutex);           // c1
20             while (count == 0)                    // c2
21                 Pthread_cond_wait(&fill, &mutex); // c3
22             int tmp = get();                       // c4
```

Multiple buffer solution

(Cont.)

```
23         Pthread_cond_signal(&empty);           // c5
24         Pthread_mutex_unlock(&mutex);           // c6
25         printf("%d\n", tmp);
26     }
27 }
```

The Final Working Solution (Cont.)

- p2: a producer sleeps only if all buffers are currently filled
- c2: a consumer sleeps only if all buffers are currently empty

One more example

- Assume there are 0 bytes free
- Thread T_a calls `allocate(100)`
- Thread T_b calls `allocate(10)`
- Both T_a and T_b wait on the condition and go to sleep
- Thread T_c calls `free(50)`
- Q: Which waiting thread should be woken up?

One more example

```
1 // how many bytes of the heap are free?
2 int bytesLeft = MAX_HEAP_SIZE;
3
4 // need lock and condition too
5 cond_t c;
6 mutex_t m;
7
8 void *
9 allocate(int size) {
10     Pthread_mutex_lock(&m);
11     while (bytesLeft < size)
12         Pthread_cond_wait(&c, &m);
13     void *ptr = ...; // get mem from heap
14     bytesLeft -= size;
15     Pthread_mutex_unlock(&m);
16     return ptr;
17 }
18
19 void free(void *ptr, int size) {
20     Pthread_mutex_lock(&m);
21     bytesLeft += size;
22     Pthread_cond_signal(&c); // whom to signal??
23     Pthread_mutex_unlock(&m);
24 }
```

Solution

- Wake up all waiting threads!
- Replace `pthread_cond_signal()` with `pthread_cond_broadcast()`