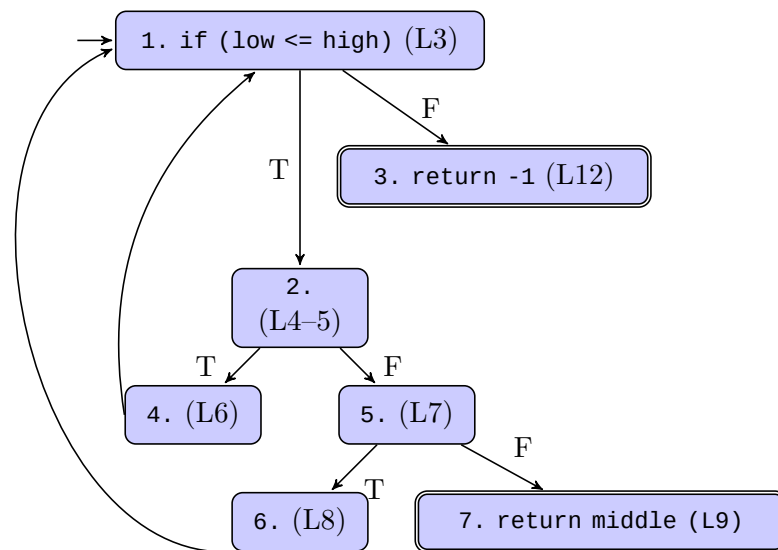


Larger CFG example. You can draw a 7-node CFG for this program:

```

1  /** Binary search for target in sorted subarray a[low..high] */
2  int binary_search(int[] a, int low, int high, int target) {
3      while (low <= high) {
4          int middle = low + (high-low)/2;
5          if (target < a[middle])
6              high = middle - 1;
7          else if (target > a[middle])
8              low = middle + 1;
9          else
10             return middle;
11     }
12     return -1; /* not found in a[low..high] */
13 }

```



Here are more exercise programs that you can draw CFGs for.

```

1  /* effects: if x==null, throw NullPointerException
2             otherwise, return number of elements in x that are odd, positive or both. */
3  int oddOrPos(int[] x) {
4      int count = 0;
5      for (int i = 0; i < x.length; i++) {
6          if (x[i]%2 == 1 || x[i] > 0) {
7              count++;
8          }
9      }
10     return count;
11 }
12
13 // example test case: input: x=[-3, -2, 0, 1, 4]; output: 3

```

Finally, we have a really poorly-designed API (I'd give it a D at most, maybe an F) because it's impossible to succinctly describe what it does. **Do not design functions with interfaces like this.** But we can still draw a CFG, no matter how bad the code is.

```

1  /** Returns the mean of the first maxSize numbers in the array,
2     if they are between min and max. Otherwise, skip the numbers. */
3  double computeMean(int[] value, int maxSize, int min, int max) {
4      int i, ti, tv, sum;
5
6      i = 0; ti = 0; tv = 0; sum = 0;
7      while (ti < maxSize) {
8          ti++;
9          if (value[i] >= min && value[i] <= max) {
10             tv++;
11             sum += value[i];
12         }
13         i++;
14     }
15     if (tv > 0)
16         return (double)sum/tv;
17     else
18         throw new IllegalArgumentException();
19 }

```

Statement and Branch Coverage

We defined Control-Flow Graphs so that we can give principled definitions of statement and branch coverage. We can start with the definition of a test path:

Definition 1 *A test path is a path p (possibly of length 0) that starts at some initial node (i.e. in N_0) and ends at some final node (i.e. in N_f).*

Here's a definition of coverage for graphs:

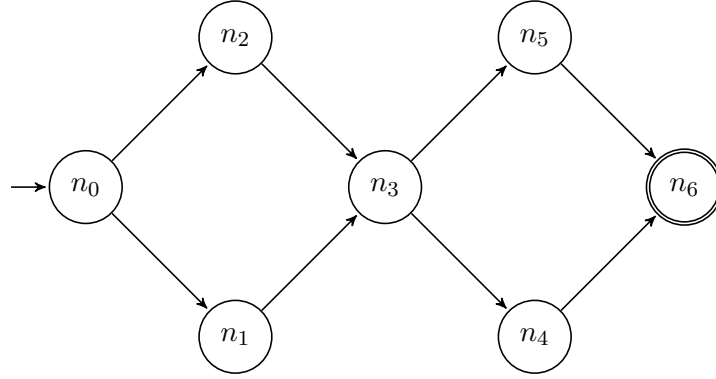
Definition 2 *Given a set of test requirements TR for a graph criterion C , a test set T satisfies C on graph G iff for every test requirement tr in TR , at least one test path p in $path(T)$ exists such that p satisfies tr .*

We'll use this notion to define a number of standard testing coverage criteria. But first, what are test paths?

Test cases and test paths. We connect test cases and test paths with a mapping $path_G$ from test cases to test paths; e.g. $path_G(t)$ is the set of test paths corresponding to test case t .

- usually we just write $path$ since G is obvious from the context.
- we can lift the definition of path to test sets T by defining $path(T) = \{path(t) | t \in T\}$.
- each test case gives at least one test path. If the software is deterministic, then each test case gives exactly one test path; otherwise, multiple test cases may arise from one test path.

As an example, consider the double-diamond graph D .



Here are the four test paths in D :

$$\begin{aligned} &[n_0, n_1, n_3, n_4, n_6] \\ &[n_0, n_1, n_3, n_5, n_6] \\ &[n_0, n_2, n_3, n_4, n_6] \\ &[n_0, n_2, n_3, n_5, n_6] \end{aligned}$$

For the *statement coverage* criterion, we get the following test requirements:

$$\{n_0, n_1, n_2, n_3, n_4, n_5, n_6\}$$

That is, any test set T which satisfies statement coverage on D must include test cases t ; the cases t give rise to test paths $\text{path}(t)$, and some path must include each node from n_0 to n_6 . (No single path must include all of these nodes; the requirement applies to the set of test paths.)

Let's formally define statement coverage.

Definition 3 *Statement coverage:* For each node $n \in \text{reach}_G(N_0)$, TR contains a requirement to visit node n .

For our example,

$$TR = \{n_0, n_1, n_2, n_3, n_4, n_5, n_6\}.$$

Let's consider an example of a test set which satisfies statement coverage on D .

Start with a test case t_1 ; assume that executing t_1 gives the test path

$$\text{path}(t_1) = p_1 = [n_0, n_1, n_3, n_4, n_6].$$

Then test set $\{t_1\}$ does not give statement coverage on D , because no test case covers node n_2 or n_5 . If we can find a test case t_2 with test path

$$\text{path}(t_2) = p_2 = [n_0, n_2, n_3, n_5, n_6],$$

then the test set $T = \{t_1, t_2\}$ satisfies statement coverage on D .

What is another test set which satisfies statement coverage on D ?

Here is a more verbose definition of statement coverage.

Definition 4 *Test set T satisfies statement coverage on graph G if and only if for every syntactically reachable node $n \in N$, there is some path p in $\text{path}(T)$ such that p visits n .*

A second standard criterion is that of branch coverage.

Criterion 1 Branch Coverage. *TR contains each reachable path of length up to 1, inclusive, in G .*

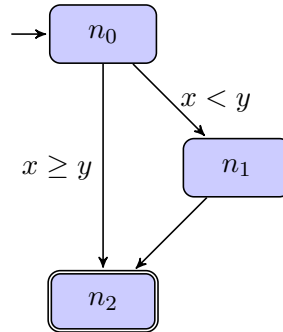
Here are some examples of paths of length ≤ 1 :

Note that since we're not talking about *test paths*, these reachable paths need not start in N_0 .

In general, paths of length ≤ 1 consist of nodes and edges. (Why not just say edges?)

Saying “edges” on the above graph would not be the same as saying “paths of length ≤ 1 ”.

Here is a more involved example:



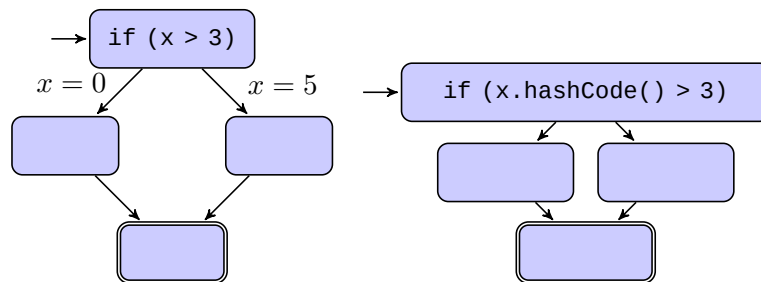
Let's define

$$\begin{aligned} \text{path}(t_1) &= [n_0, n_1, n_2] \\ \text{path}(t_2) &= [n_0, n_2] \end{aligned}$$

Then

$$\begin{aligned} T_1 &= \text{satisfies statement coverage} \\ T_2 &= \text{satisfies branch coverage} \end{aligned}$$

Nondeterminism. I mentioned the mapping between test cases and test paths above. The mapping is not one-to-one for nondeterministic code. Here's an example of deterministic and non-deterministic control-flow graphs:

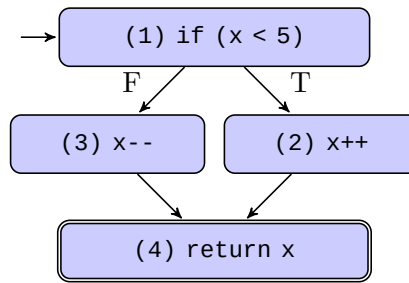


Causes of nondeterminism include dependence on inputs; on the thread scheduler; and on memory addresses, for instance as seen in calls to the default Java `hashCode()` implementation.

Nondeterminism makes it hard to check test case output, since more than one output might be a valid result of a single test input.

Example. Here is a short method, the associated control-flow graph, and some test cases and test paths.

```
1 int foo(int x) {  
2   if (x < 5) {  
3     x ++;  
4   } else {  
5     x --;  
6   }  
7   return x;  
8 }
```



- Test case: $x = 5$; test path: $[(1), (3), (4)]$.
- Test case: $x = 2$; test path: $[(1), (2), (4)]$.

Note that (1) we can deduce properties of the test case from the test path; and (2) in this example, since our method is deterministic, the test case determines the test path.