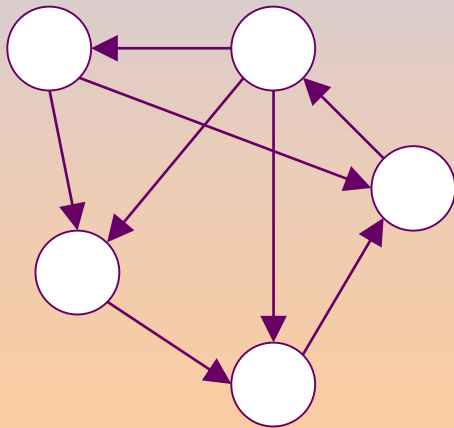


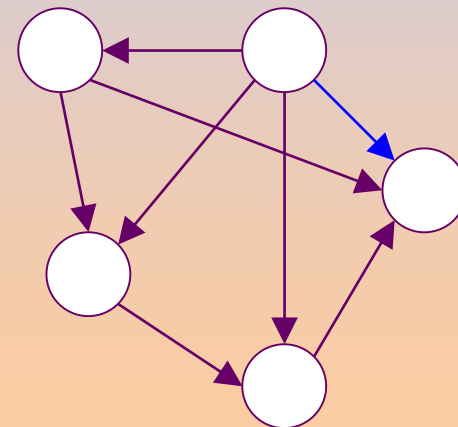
Strongly Connected Components

- Every pair of vertices are reachable from each other
- Graph G is ***strongly connected*** if, for every u and v in V , there is some path from u to v and some path from v to u .

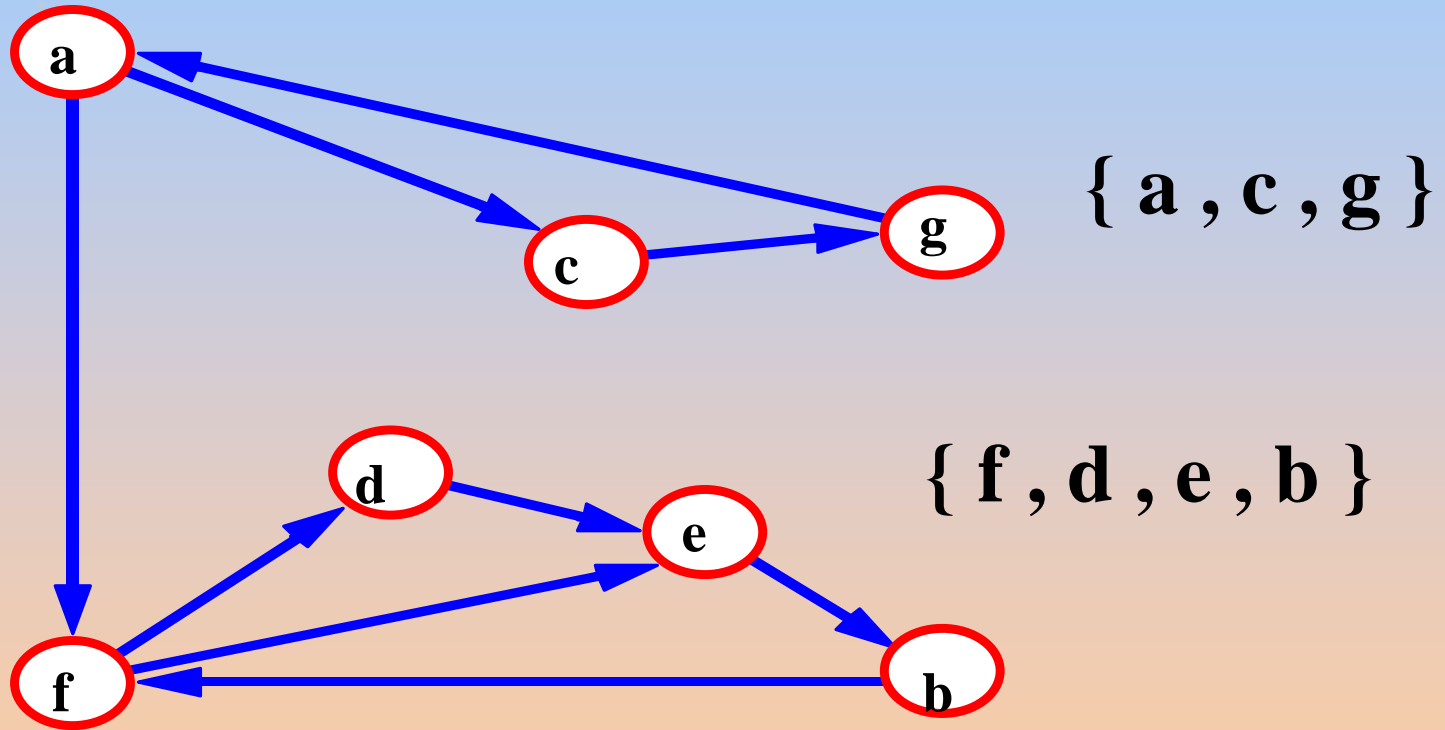
Strongly
Connected



Not Strongly
Connected



Example





Finding Strongly Connected Components

- Input: A directed graph $G = (V, E)$
- Output: a partition of V into disjoint sets so that each set defines a strongly connected component of G



Algorithm

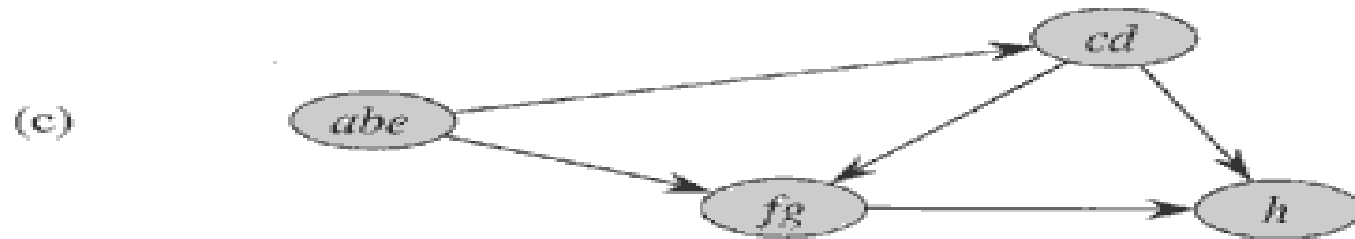
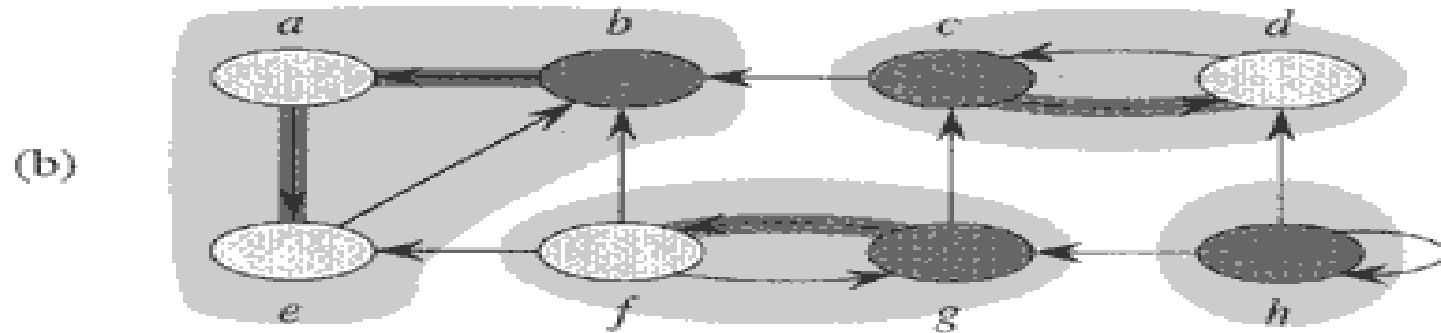
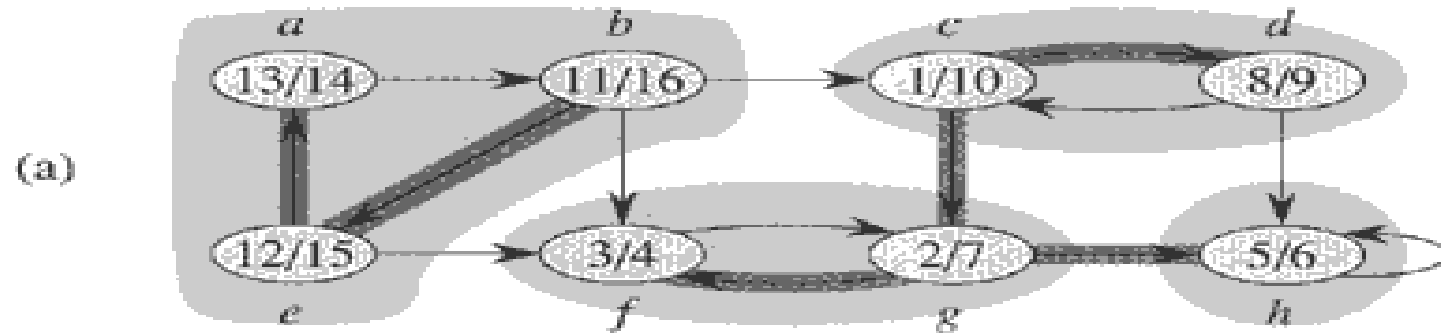
Strongly-Connected-Components(G)

1. call DFS(G) to compute finishing times $f[u]$ for each vertex u . Cost: $O(E+V)$
2. compute G^T Cost: $O(E+V)$
3. call DFS(G^T), but in the main loop of DFS, consider the vertices in order of decreasing $f[u]$ Cost: $O(E+V)$
4. output the vertices of each tree in the depth-first forest of step 3 as a separate strongly connected component.

The graph G^T is the transpose of G , which is visualized by reversing the arrows on the digraph.

■ Cost: $O(E+V)$

Example



Questions

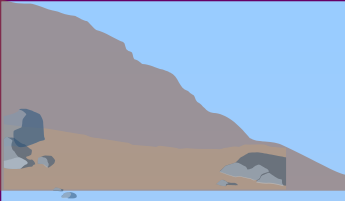
- Let $C=\{C_1, C_2, \dots, C_n\}$ be the set of strongly connected components of $G=(V, E)$.

Let,

$G^{SCC}=(V^{SCC}, E^{SCC})$ where $V^{SCC}=\{v_1, v_2, \dots, v_n\}$ where each vertex v_i in V^{SCC} represents the strongly connected component C_i in C . There is a directed edge (v_i, v_k) in E^{SCC} if there is a directed edge $(x, y) \in E$ such that $x \in C_i$ and $y \in C_k$.

Then, $((G^T)^{SCC})^T = G^{SCC}$ where G^T denotes the transpose graph of G .

True/False ?

- 
- Consider a directed graph $G=(V,E)$ where each node is initially colored white. What should be the minimum number of nodes that we should change to red such that for each white node v in G , there exists at least one red node r such that there is a directed path from r to v and v to r . Assume that for every node there is at least one other node it is reachable from and can reach to.