

Indexed Search Tree (Trie)



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■ Trie datastructure applications

■ --auto completion

■ --search engines

■ --IP routing

Types of Tries

- **Standard**

- Single character per node

- **Compressed**

- Eliminating chains of nodes

- **Compact**

- Stores indices into original string(s)

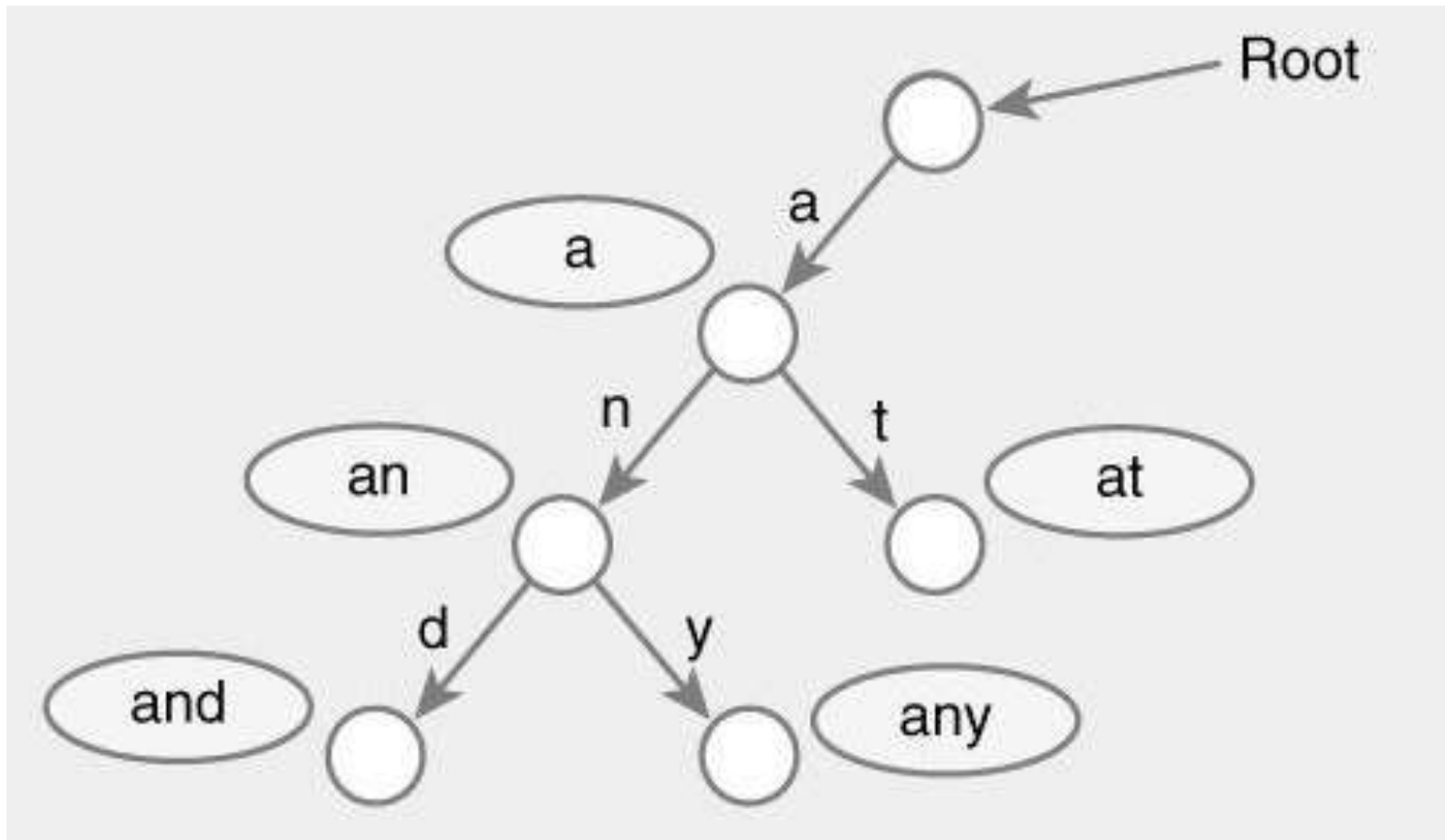
- **Suffix**

- Stores all suffixes of string

Standard Trie Example

■ For strings

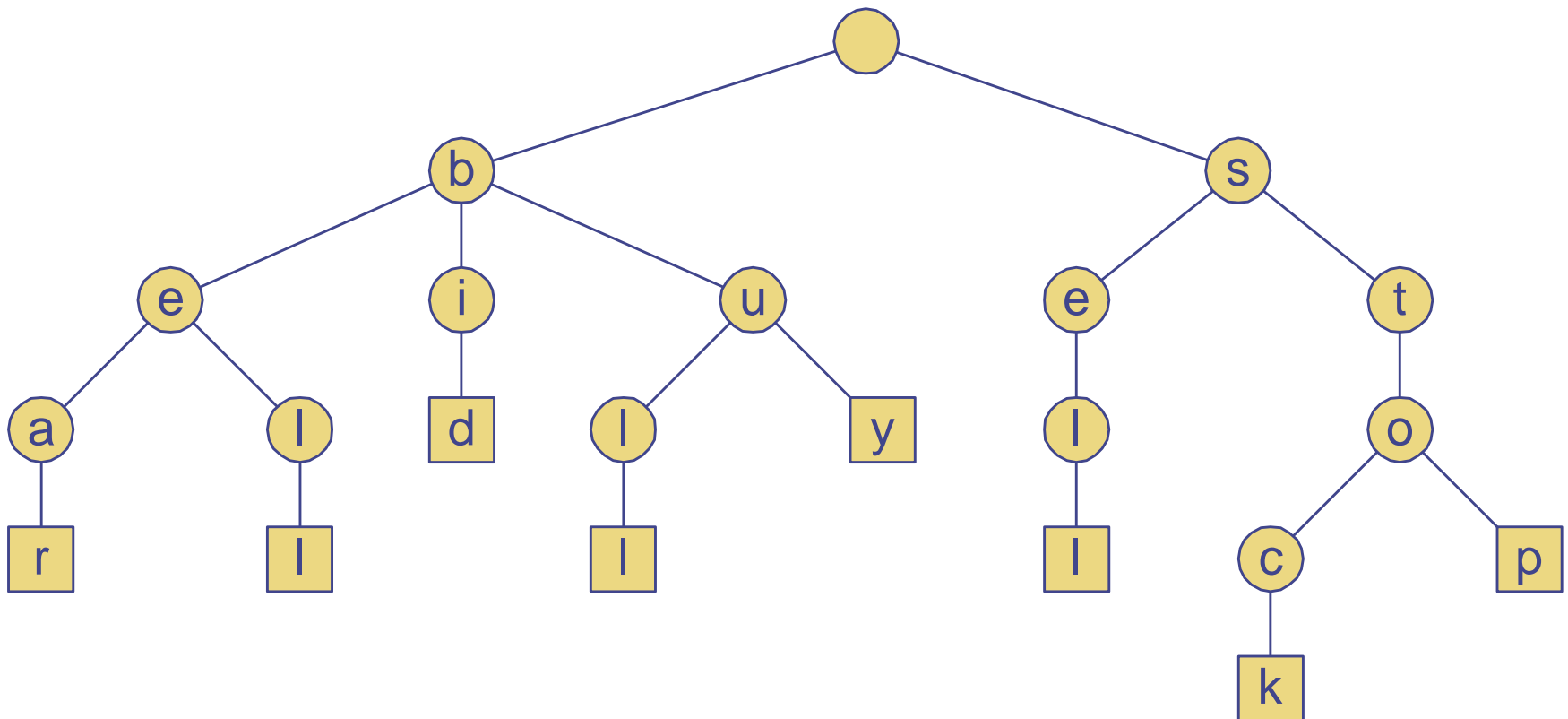
■ { a, an, and, any, at }



Standard Trie Example

■ For strings

■ { bear, bell, bid, bull, buy, sell, stock, stop }



Standard Tries

■ Node structure

- Value between $1 \dots m$
- Reference to m children
 - Array or linked list

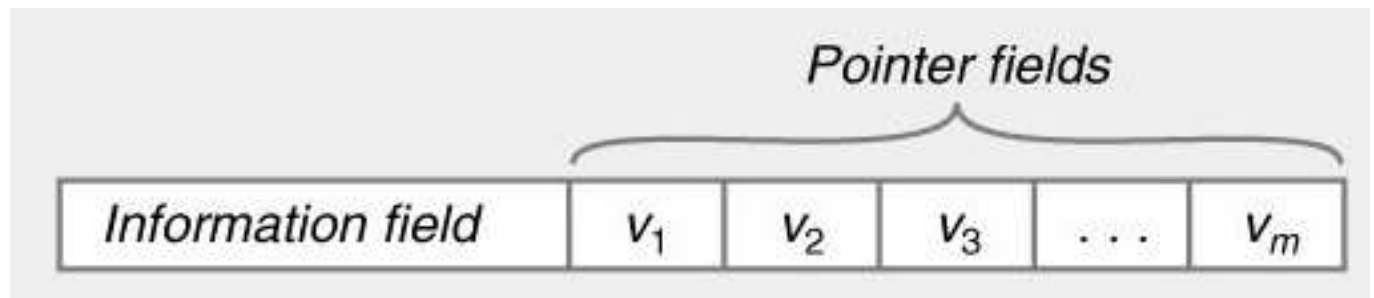
■ Example

Class Node {

Letter value; // Letter $V = \{ V_1, V_2, \dots V_m \}$

Node child[m];

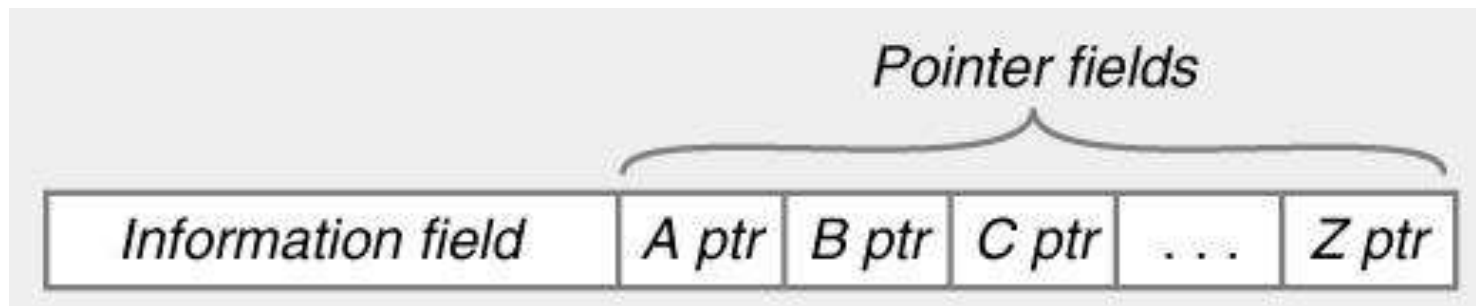
}



Standard Tries

■ Efficiency

- Uses $O(n)$ space
- Supports search / insert / delete in $O(d \times m)$ time
- For
 - n total size of strings indexed by trie
 - d length of the parameter string
 - m size of the alphabet

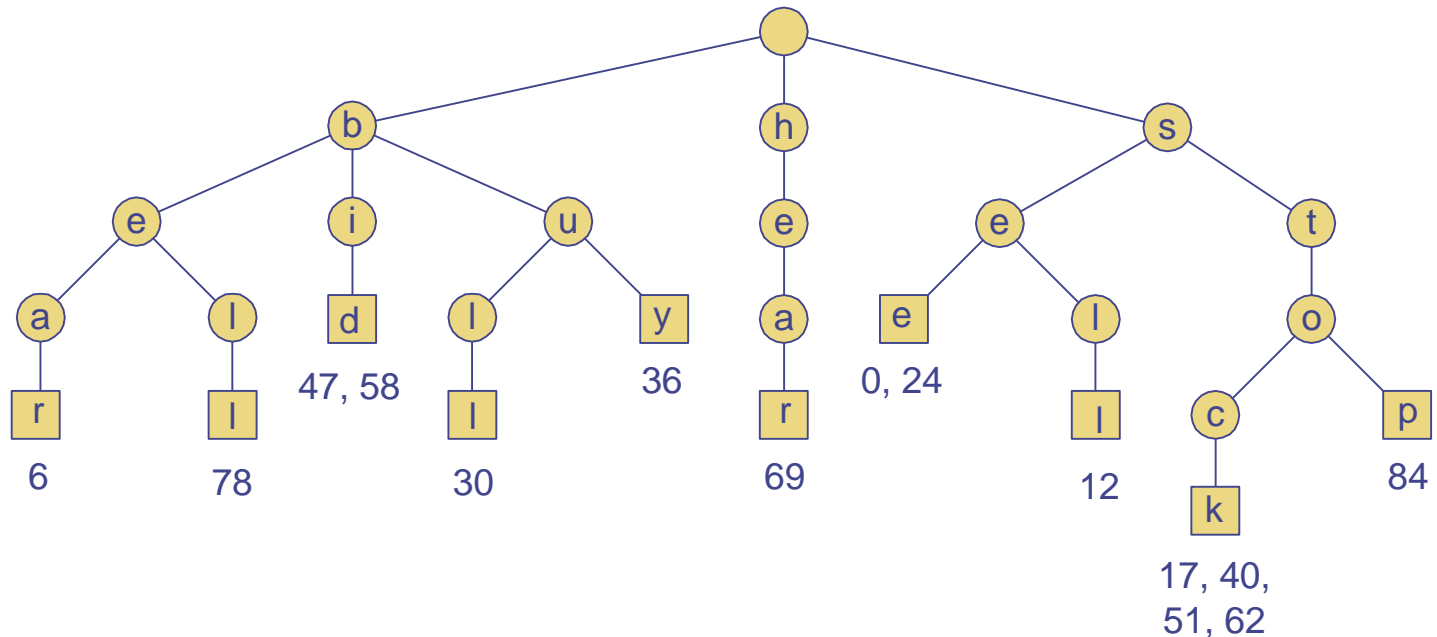


Word Matching Trie

- Insert words into trie
- Each leaf stores occurrences of word in the text

s	e	e		a		b	e	a	r	?		s	e	l	l		s	t	o	c	k	!	
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23
s	e	e		a		b	u	l	l	?		b	u	y		s	t	o	c	k	!		
24	25	26	27	28	29	30	31	32	33	34	35	36	37	38	39	40	41	42	43	44	45	46	
b	i	d		s	t	o	c	k	!		b	i	d		s	t	o	c	k	!			
47	48	49	50	51	52	53	54	55	56	57	58	59	60	61	62	63	64	65	66	67	68		
h	e	a	r		t	h	e		b	e	l	l	?		s	t	o	p	!				
69	70	71	72	73	74	75	76	77	78	79	80	81	82	83	84	85	86	87	88				

Deletion?



Compressed Trie

■ Observation

- Internal node v of T is redundant if v has one child and is not the root

■ Approach

- A chain of redundant nodes can be compressed
 - Replace chain with single node
 - Include concatenation of labels from chain

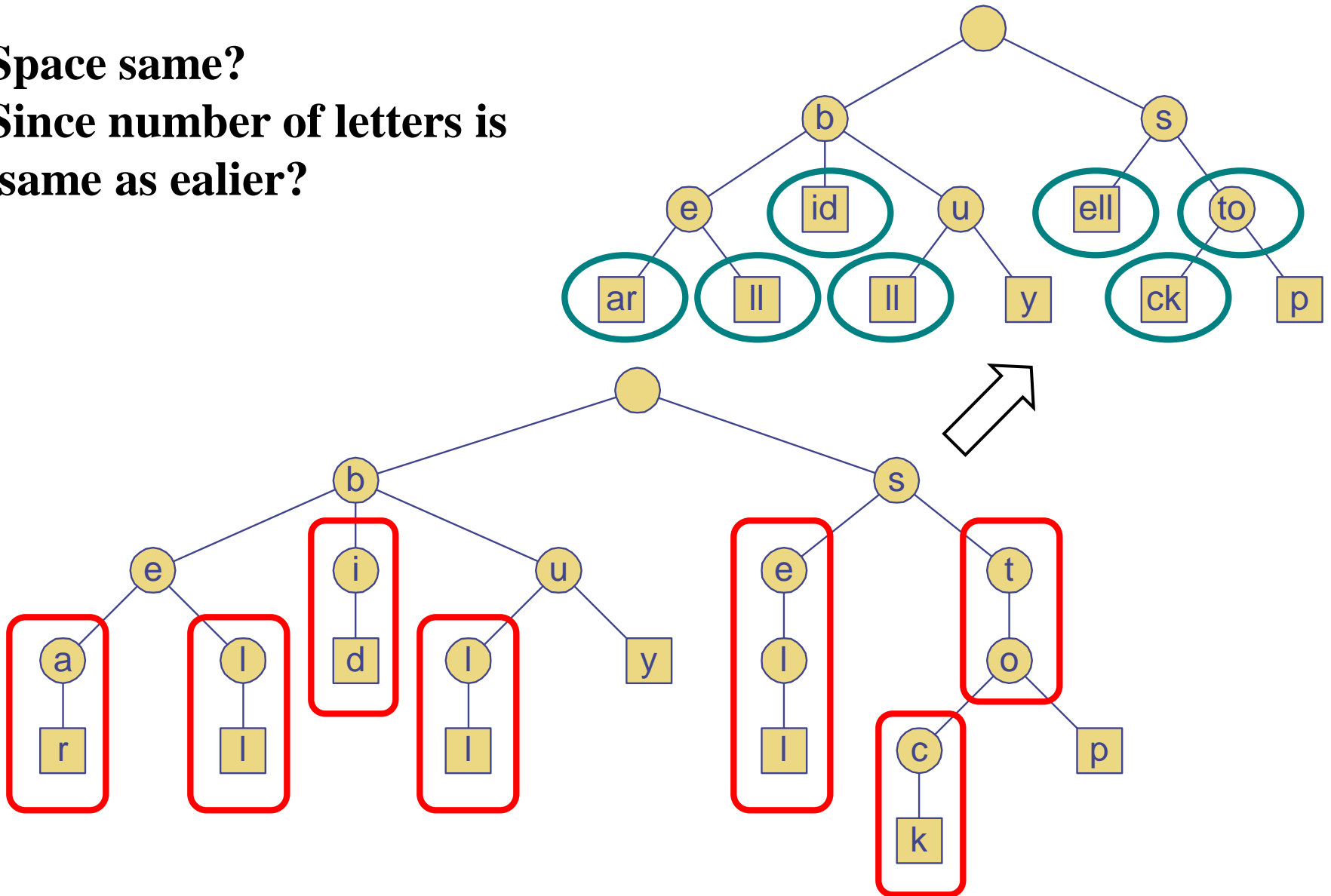
■ Result

- Internal nodes have at least 2 children
- Some nodes have multiple characters

Compressed Trie

Space same?

Since number of letters is
same as ealier?



Compact Tries

■ Compact representation of a compressed trie

■ Approach

- For an array of strings $S = S[0], \dots S[s-1]$
- Store ranges of indices at each node
 - Instead of substring
- Represent as a triplet of integers (i, j, k)
 - Such that $X = s[i][j..k]$
- Example: $S[0] = \text{"abcd"}$, $(0,1,2) = \text{"bc"}$

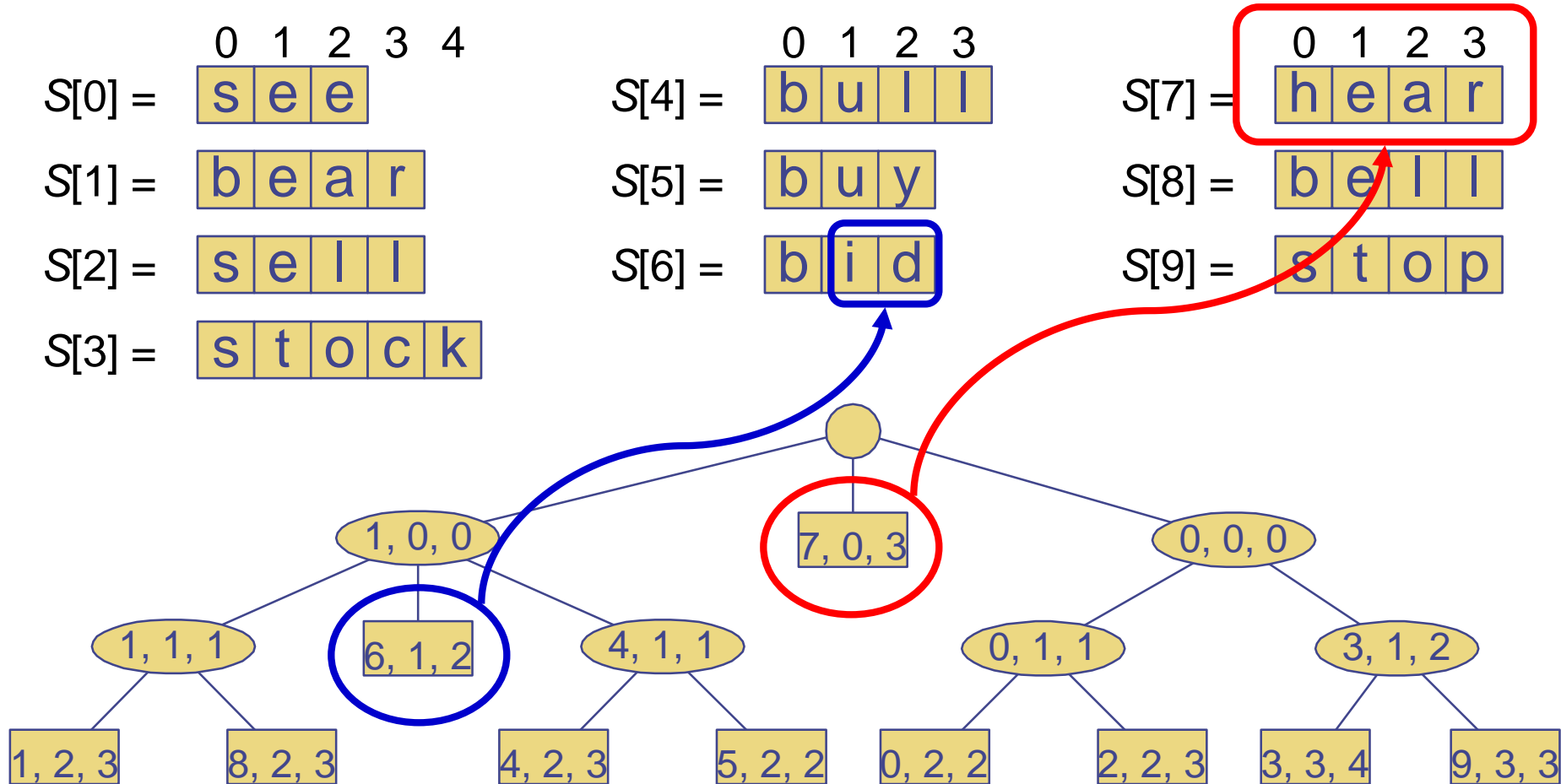
■ Properties

- Uses $O(s)$ space, where $s = \#$ of strings in the array
- Serves as an auxiliary index structure

- **A tree with L leaf nodes in which every node has at least 2 children except the leaf nodes has at most $L-1$ internal nodes.**

Compact Representation

■ Example



Suffix Trie

- Compressed trie of all suffixes of text

- Example: “IPDPS”

- Suffixes

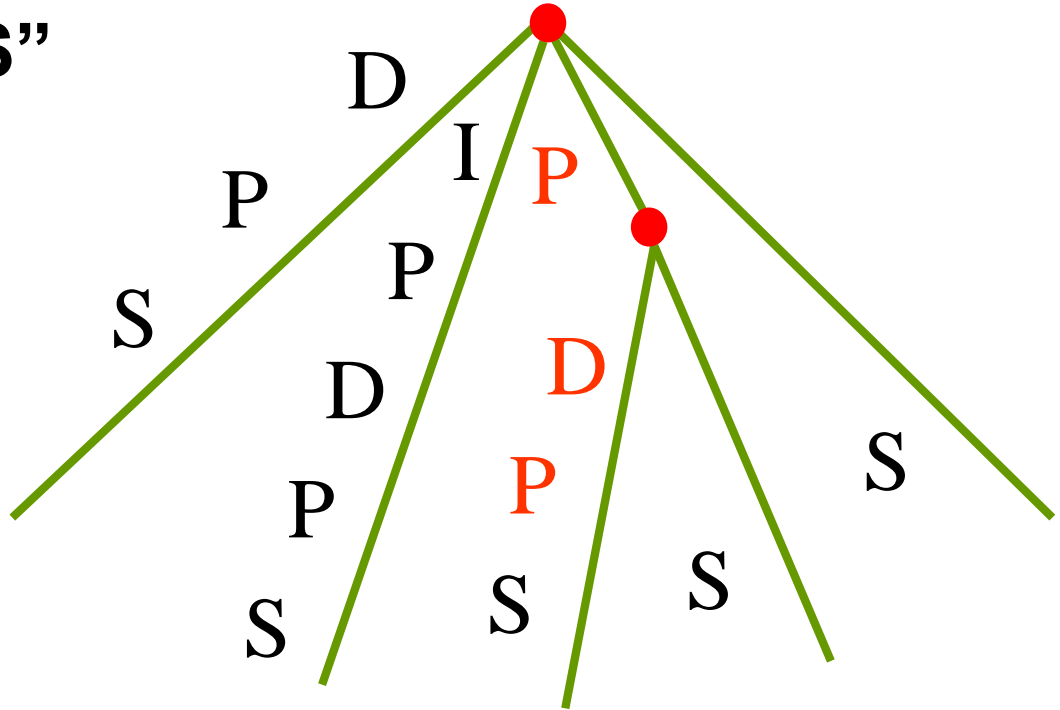
- IPDPS

- PDPS

- DPS

- PS

- S



- Useful for finding pattern in any part of text

- Occurrence \Rightarrow prefix of some suffix

- Example: find **PDP** in **IPDPS**

Suffix Trie

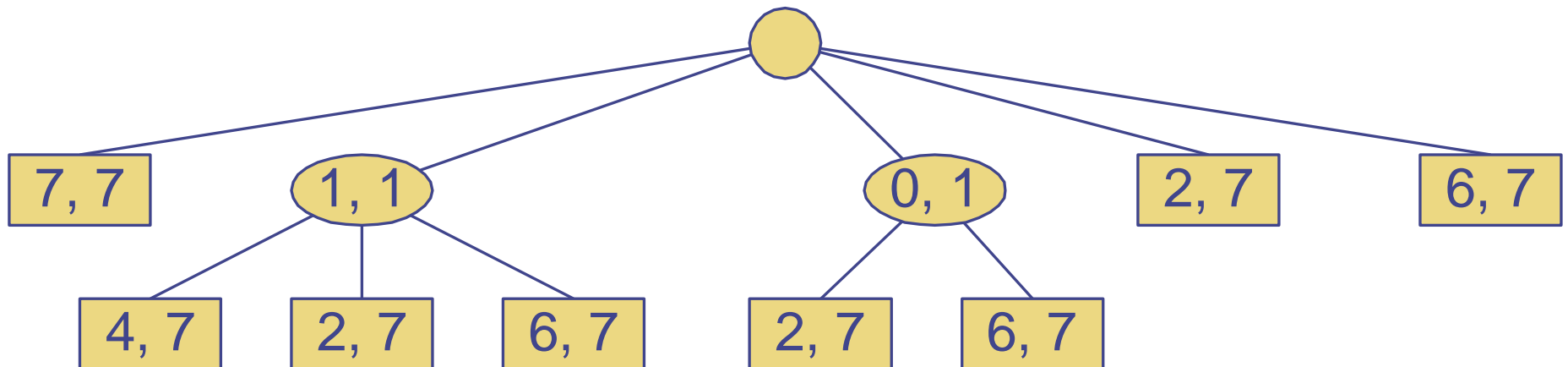
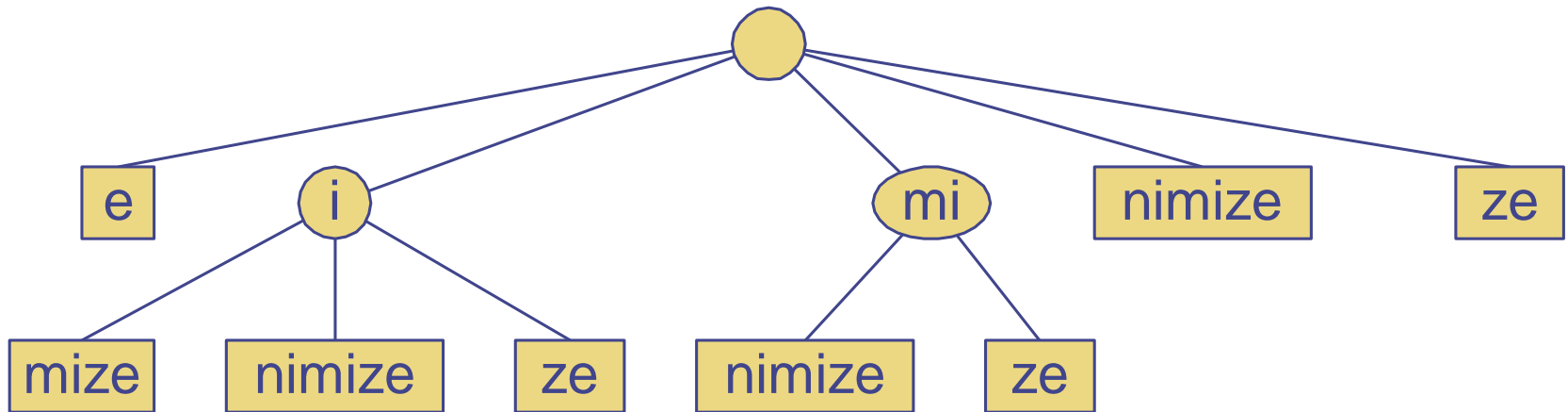
■ Properties

■ For

- String X with length n
 - Alphabet of size m
 - Pattern P with length d
- Uses $O(n)$ space (since we have $O(n)$ leaves)
 - Can be constructed in $O(n)$ time
 - Find pattern P in X in $O(d \times m)$ time
 - Proportional to length of pattern, not text

Suffix Trie Example

m	i	n	i	m	i	z	e
0	1	2	3	4	5	6	7



Tries and Web Search Engines

- **Search engine index**
 - Collection of all searchable words
 - Stored in compressed trie
- **Each leaf of trie**
 - Associated with a word
 - List of pages (URLs) containing that word
 - Called occurrence list
- **Trie is kept in memory (fast)**
- **Occurrence lists kept in external memory**
 - Ranked by relevance