

## Cyclic Redundancy Check-

- Cyclic Redundancy Check (CRC) is an error detection method.
- It is based on binary division.

## CRC Generator-

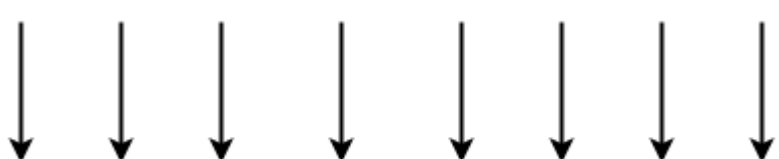
- CRC generator is an algebraic polynomial represented as a bit pattern.
- Bit pattern is obtained from the CRC generator using the following rule-

The power of each term gives the position of the bit and the coefficient gives the value of the bit.

## Example-

Consider the CRC generator is  $x^7 + x^6 + x^4 + x^3 + x + 1$ .

The corresponding binary pattern is obtained as-

$$1x^7 + 1x^6 + 0x^5 + 1x^4 + 1x^3 + 0x^2 + 1x^1 + 1x^0$$


1 1 0 1 1 0 1 1

Thus, for the given CRC generator, the corresponding binary pattern is 11011011.

## **Properties Of CRC Generator-**

The algebraic polynomial chosen as a CRC generator should have at least the following properties-

### **Rule-01:**

- It should not be divisible by  $x$ .
- This condition guarantees that all the burst errors of length equal to the length of polynomial are detected.

### **Rule-02:**

- It should be divisible by  $x+1$ .
- This condition guarantees that all the burst errors affecting an odd number of bits are detected.

## **Important Notes-**

If the CRC generator is chosen according to the above rules, then-

- CRC can detect all single-bit errors
- CRC can detect all double-bit errors provided the divisor contains at least three logic 1's.
- CRC can detect any odd number of errors provided the divisor is a factor of  $x+1$ .
- CRC can detect all burst error of length less than the degree of the polynomial.
- CRC can detect most of the larger burst errors with a high probability.

## **Steps Involved-**

Error detection using CRC technique involves the following steps-

### **Step-01: Calculation Of CRC At Sender Side-**

At sender side,

- A string of  $n$  0's is appended to the data unit to be transmitted.
- Here,  $n$  is one less than the number of bits in CRC generator.

- Binary division is performed of the resultant string with the CRC generator.
- After division, the remainder so obtained is called as **CRC**.
- It may be noted that CRC also consists of n bits.

### **Step-02: Appending CRC To Data Unit-**

At sender side,

- The CRC is obtained after the binary division.
- The string of n 0's appended to the data unit earlier is replaced by the CRC remainder.

### **Step-03: Transmission To Receiver-**

- The newly formed code word (Original data + CRC) is transmitted to the receiver.

### **Step-04: Checking at Receiver Side-**

At receiver side,

- The transmitted code word is received.
- The received code word is divided with the same CRC generator.
- On division, the remainder so obtained is checked.

The following two cases are possible-

#### **Case-01: Remainder = 0**

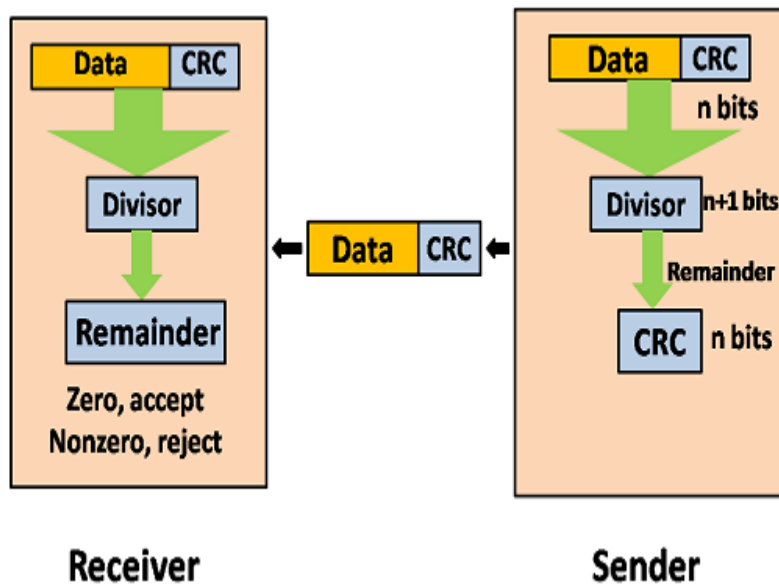
If the remainder is zero,

- Receiver assumes that no error occurred in the data during the transmission.
- Receiver accepts the data.

#### **Case-02: Remainder $\neq$ 0**

If the remainder is non-zero,

- Receiver assumes that some error occurred in the data during the transmission.
- Receiver rejects the data and asks the sender for retransmission.

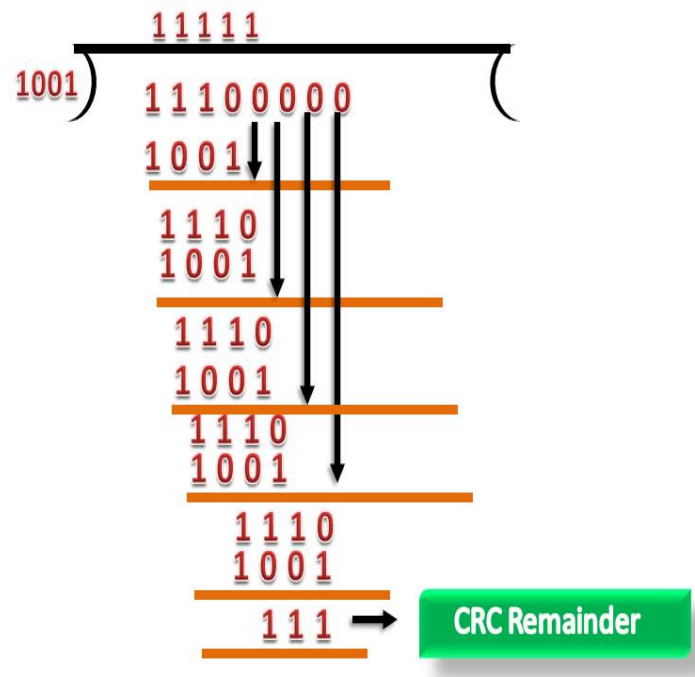


Let's understand this concept through an example:

**Suppose the original data is 11100 and divisor is 1001.**

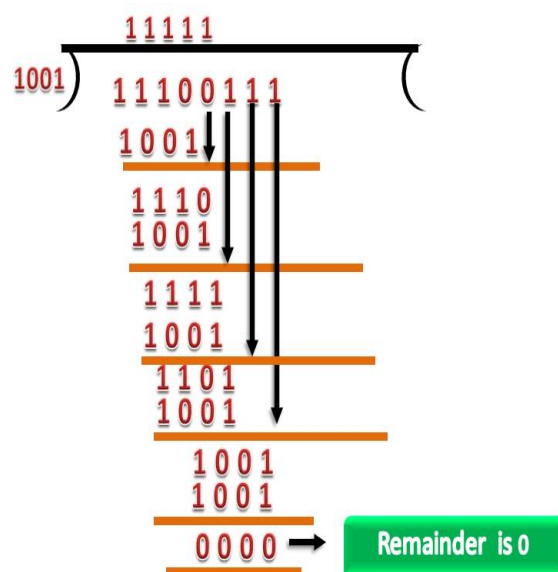
## CRC Generator

- A CRC generator uses a modulo-2 division. Firstly, three zeroes are appended at the end of the data as the length of the divisor is 4 and we know that the length of the string 0s to be appended is always one less than the length of the divisor.
- Now, the string becomes 11100000, and the resultant string is divided by the divisor 1001.
- The remainder generated from the binary division is known as CRC remainder. The generated value of the CRC remainder is 111.
- CRC remainder replaces the appended string of 0s at the end of the data unit, and the final string would be 11100111 which is sent across the network.



## CRC Checker

- The functionality of the CRC checker is similar to the CRC generator.
- When the string 11100111 is received at the receiving end, then CRC checker performs the modulo-2 division.
- A string is divided by the same divisor, i.e., 1001.
- In this case, CRC checker generates the remainder of zero. Therefore, the data is accepted.



## **PRACTICE PROBLEMS BASED ON CYCLIC REDUNDANCY CHECK (CRC)-**

### **Problem-01:**

A bit stream 1101011011 is transmitted using the standard CRC method. The generator polynomial is  $x^4+x+1$ . What is the actual bit string transmitted?

### **Solution-**

- The generator polynomial  $G(x) = x^4 + x + 1$  is encoded as 10011.
- Clearly, the generator polynomial consists of 5 bits.
- So, a string of 4 zeroes is appended to the bit stream to be transmitted.
- The resulting bit stream is 1101011011**0000**.

Now, the binary division is performed as-



A bit stream 10011101 is transmitted using the standard CRC method. The generator polynomial is  $x^3+1$ .

1. What is the actual bit string transmitted?
2. Suppose the third bit from the left is inverted during transmission. How will receiver detect this error?

## **Solution-**

### **Part-01:**

- The generator polynomial  $G(x) = x^3 + 1$  is encoded as 1001.
- Clearly, the generator polynomial consists of 4 bits.
- So, a string of 3 zeroes is appended to the bit stream to be transmitted.
- The resulting bit stream is 10011101**000**.

Now, the binary division is performed as-



$$\begin{array}{r}
 \phantom{1001} \overline{) 10001100} \\
 1001 \overline{) 10011101000} \\
 \underline{1001} \phantom{0000} \\
 00001 \phantom{000} \\
 \underline{0000} \phantom{000} \\
 00011 \phantom{000} \\
 \underline{0000} \phantom{000} \\
 00110 \phantom{000} \\
 \underline{0000} \phantom{000} \\
 01101 \phantom{000} \\
 \underline{1001} \phantom{000} \\
 01000 \phantom{000} \\
 \underline{1001} \phantom{000} \\
 00010 \phantom{000} \\
 \underline{0000} \phantom{000} \\
 00100 \phantom{000} \\
 \underline{0000} \phantom{000} \\
 0100 \leftarrow \text{CRC}
 \end{array}$$

From here, CRC = 100.

Now,

- The code word to be transmitted is obtained by replacing the last 3 zeroes of 10011101**000** with the CRC.
- Thus, the code word transmitted to the receiver = 10011101**100**.

## **Part-02:**

According to the question,

- Third bit from the left gets inverted during transmission.
- So, the bit stream received by the receiver = 10111101100.

Now,

- Receiver receives the bit stream = 10111101100.
- Receiver performs the binary division with the same generator polynomial as-

$$\begin{array}{r}
 \phantom{1001} \overline{) 10101000} \\
 1001 \overline{) 10111101100} \\
 \underline{1001} \phantom{000000000} \\
 00101 \phantom{000000000} \\
 \underline{0000} \phantom{000000000} \\
 01011 \phantom{000000000} \\
 \underline{1001} \phantom{000000000} \\
 00100 \phantom{000000000} \\
 \underline{0000} \phantom{000000000} \\
 01001 \phantom{000000000} \\
 \underline{1001} \phantom{000000000} \\
 00001 \phantom{000000000} \\
 \underline{0000} \phantom{000000000} \\
 00010 \phantom{000000000} \\
 \underline{0000} \phantom{000000000} \\
 00100 \phantom{000000000} \\
 \underline{0000} \phantom{000000000} \\
 \underline{0100} \phantom{000000000} \leftarrow \text{Remainder}
 \end{array}$$

From here,

- The remainder obtained on division is a non-zero value.
- This indicates to the receiver that an error occurred in the data during the transmission.
- Therefore, receiver rejects the data and asks the sender for retransmission.

There are several different standard polynomials used by popular protocols for CRC generation. These are:

Name	Polynomial	Application
CRC-8	$x^8 + x^2 + x + 1$	ATM header
CRC-10	$x^{10} + x^9 + x^5 + x^4 + x^2 + 1$	ATM AAL
CRC-16	$x^{16} + x^{12} + x^5 + 1$	HDLC
CRC-32	$x^{32} + x^{26} + x^{23} + x^{22} + x^{16} + x^{12} + x^{11} + x^{10} + x^8 + x^7 + x^5 + x^4 + x^2 + x + 1$	LANs