



Data Structure & Algorithms

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Course Introduction

- Data Structure and Algorithms

- Data Structures: Linked list, Stack, Queue, Binary search tree, Heap, Graph.
- Algorithms: Sorting, Searching, Stack, Queue & Linked list applications, Graph algorithms.

- Course Goals

- Implement each DS & Algorithms from scratch.
- Understand complexity of algorithms.

- Course Schedules

- 5th Sep 2021 to 20th Sep 2020 *
- Mon-Fri: Lecture – 9:00 AM to 1:30 PM
- Mon-Fri: Lab – 2:30 PM to 3:30 PM ✓

- Resource sharing

- <https://gitlab.com/sunbeam-modular/dsa-04> ✓
- Recorded videos will be available for 7 days.
<http://students.sunbeamapps.org> ✓

- Course Format

- Participants are encouraged to code alongside (copy code from code-sharing utility in student portal).
- Post your queries in chat box (on logical end of each topic).
- Practice assignments will be shared. They are optional. If any doubts, share on WA group (possibly with screenshot). Faculty members or peers can help.

- Programming language

- DS & Algorithms are language independent.
- Classroom coding will be in Java (use IDE of your choice).
- Will share C++/Python codes at the end of session.
- Language pre-requisites ?



Course Pre-requisites

Java

- Language Funda
- Methods
- Class & Object
- static members
- Arrays
- Collections

→ ArrayList < >
→ Stack < > , Queue < >

Python

- Language Funda
- Functions
- Class & Object
- Collections

C++

- Language Funda
- Functions
- Class & Object
- Friend class
- Arrays
- Pointers

C

- Language Funda
- Functions
- Structures
- Arrays
- Pointers



Data Structure

- Data Structure
 - Organizing data in memory
 - Processing the data
- Common data structures
 - ✓• Array
 - ✓• Linked List
 - ✓• Stack
 - ✓• Queue
 - ✓• Hash Table
- Advanced data structures
 - ✓• Tree
 - ✓• Heap
 - ✓• Graph



Data Structure

- Data Structure
 - Organizing data in memory
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 - Array
 - Linked List
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 - Queue
- Advanced data structures
 - Tree
 - Heap
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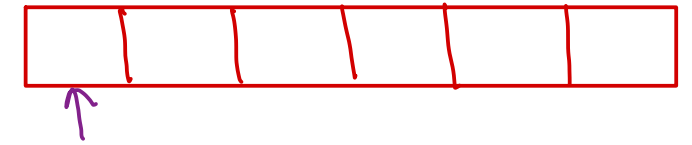
- Asymptotic analysis
 - It is not exact analysis
 - Big' O notation

Exact analysis

- ① How much time?
- ② How many bytes?

* Arch / machine dependent

- Space complexity ✓
 - Unit space to store the data (Input space) and additional space to process the data (Auxiliary space).
 - $O(1)$, $O(n)$, $O(n^2)$
 - $S \propto n$
 - + input space: $O(n)$
 - aux space: $S = k$
 - $O(1)$
- Time complexity
 - Unit time required to complete any algorithm.
 - Approximate measure of time required to complete any algorithm.
 - Depends on loops in the algorithm.
 - $O(n^3)$, $O(n^2)$, $O(n \log n)$, $O(n)$, $O(\log n)$, $O(1)$



$T \propto n$



Time complexity

- Write a program to calculate factorial of given number.

$T \propto n$
 $O(n)$

```
r = 1;
for (i = 1; i <= n; i++)
    r = r * i;
print (r);
```

- Print 2-D matrix of $n \times n$.

$i \times s = n \times n$
 $T \propto n^2$
 $O(n^2)$

```
for (i = 0; i < n; i++) {
    for (j = 0; j < n; j++) {
        print(mat[i][j]);
    }
}
```

- Print given number into binary.

$2^{i \times r} = n$
 $i \times r \log 2 = \log n$
 $i \times r = \frac{\log n}{\log 2}$
 $T \propto \frac{\log n}{\log 2}$
 $T \propto \log n$
 $O(\log n)$

$(10)_{10} = (1010)_2$
 $10 / 2 = 5$
 $5 / 2 = 2$
 $2 / 2 = 1$
 $1 / 2 = 0$

```
while (n > 0) {
    print(n % 2);
    n = n / 2;
}
```

- Print table of given number.

```
for (i = 1; i <= 10; i++)
    print(n * i);
```

$T = k$
 $O(1)$

Linear Search

C → strchr()

array of g des

- Find a number in a list of given numbers (random order).

key = 90

0	1	2	3	4	5	6	7	8
88	33	66	99	11	77	22	55	11 44

- Time complexity

- Worst case $T \propto n$

$O(n)$

- Best case $T = K$

$O(1)$

- Average case

$T \propto \frac{n}{2}$

$O(n)$

$T \propto n$

3

```
int linearSearch(arr, n, key){
```

```
for(i=0; i<n; i++){
```

```
if(arr[i] == key)
```

```
return i;
```

3

```
return -1;
```

Binary Search

key = 30

```
left = 0;
right = n-1;
while (left <= right) {
    mid = (left + right) / 2;
    if (key == arr[mid])
        return mid;
    if (key < arr[mid])
        right = mid - 1;
    else
        left = mid + 1;
}
return -1;
```

0	1	2	3	4	5	6	7	8
11	22	33	44	55	66	77	88	99

R

L

m

best case (middle)
 $O(1)$

$$2^{\text{itr}} \approx n$$

$$\text{itr} \log 2 = \log n$$

$$\text{itr} = \frac{\log n}{\log 2}$$

$$T \propto \frac{\log n}{\log 2}$$

$$T \propto \log n$$

$$O(\log n) \leftarrow$$

avg/worst case



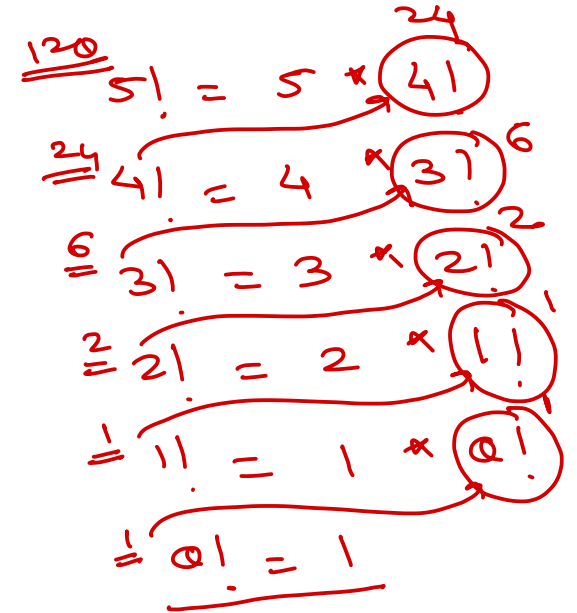
Recursion

- Function calling itself is called as recursive function.
- To write recursive function consider
 - Explain process/formula in terms of itself
 - Decide the end/terminating condition
- Examples: (base)
 - $n! = n * (n-1)!$ $0! = 1$
 - $x^y = x * x^{y-1}$ $x^0 = 1$
 - $T_n = T_{n-1} + T_{n-2}$ $T_1 = T_2 = 1$
- $\text{factors}(n) = 1^{\text{st}} \text{ prime factor of } n * \text{factors}(n)$

- On each function call, function activation record or stack frame will be created on stack.

```
int fact(int n) {  
    int r;  
    if(n==0)  
        return 1;  
    r = n * fact(n-1);  
    return r;  
}
```

res=fact(5);



Recursion

```
int fact(int n) {  
    int r; ③  
    if(n == 0)  
        return 1;  
    r = n * fact(n-1);  
    return r; 2  
}
```

```
int fact(int n) {  
    int r; ②  
    if(n == 0)  
        return 1;  
    r = n * fact(n-1);  
    return r; 1  
}
```

```
int fact(int n) {  
    int r; ①  
    if(n == 0)  
        return 1;  
    r = n * fact(n-1);  
    return r; 1  
}
```

```
int fact(int n) {  
    int r; ④  
    if(n == 0) ✓  
        return 1;  
    r = n * fact(n-1);  
    return r;  
}
```

```
int fact(int n) {  
    int r; ④  
    if(n == 0)  
        return 1;  
    r = n * fact(n-1);  
    return r; 6  
}
```

```
int fact(int n) {  
    int r; ⑤  
    if(n == 0)  
        return 1;  
    r = n * fact(n-1);  
    return r; 24  
}
```

```
int main() {  
    int res;  
    res = fact(5);  
    printf("%d", res);  
    return 0;  
}
```

Recursion

✗ more space
✗ more time
✓ Simplified coding

stack

fact(6)

fact(5)

fact(4)

fact(3)

fact(2)

fact(1)

fact(0)

main()

Binary Search

Key = 44

0	1	2	3	4	5	6	7	8
11	22	33	44	55	66	77	88	99

L M R

```
int binSearch(arr, left, right, key) {  
    if (left > right)  
        return -1;  
    mid = (left + right) / 2;  
    if (key == arr[mid])  
        return mid;  
    if (key < arr[mid])  
        idx = binSearch(arr, left, mid - 1, key);  
    else  
        idx = binSearch(arr, mid + 1, right, key);  
    return idx;  
}
```

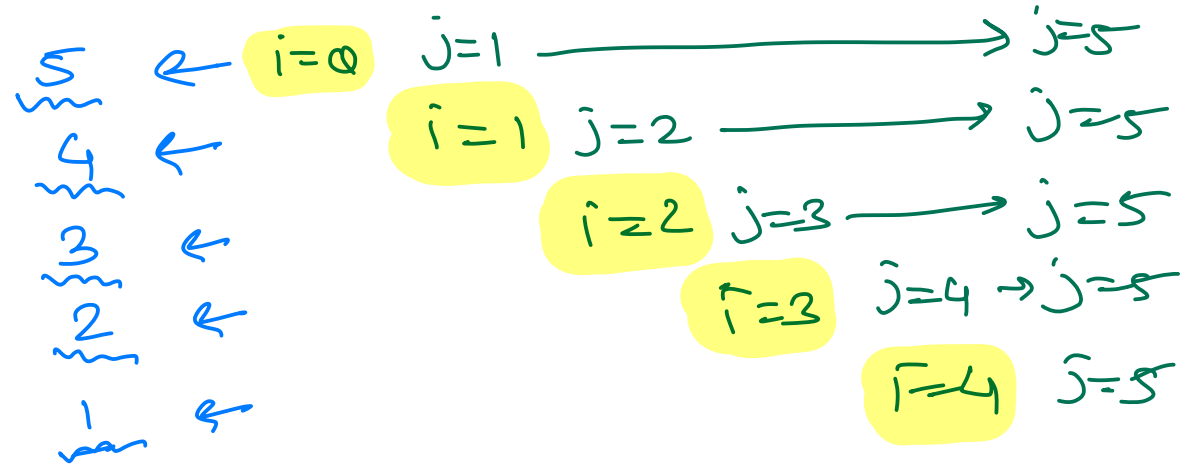
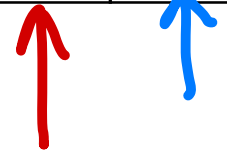


Selection Sort

5 6 3 8 2 4

```
for(i=0; i<n-1; i++){  
    for(j=i+1; j<n; j++){  
        if(a[i] > a[j]){  
            temp = a[i];  
            a[i] = a[j];  
            a[j] = temp;  
        }  
    }  
}
```

0	1	2	3	4	5
2	3	4	5	6	8



3
3





Thank you!

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