# CSCI-570: Homework # 3

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# Contents

HW3																				3
(2)		 						 												3
(3:	Ch#3 Ex#3)	 						 												3
(4)		 						 												3
(5:	Ch#4 Ex#3)	 						 												
(6)		 						 												3
(7:	Ch#4 Ex#4)	 						 												4
(8)		 						 												4

#### HW3

### (2)

A) The intersections can be viewed as the nodes of a directed graph. An intersection  $I_i$  can be reached from  $I_j$  given that there is an edge incident from  $I_j$  to  $I_i$  that is  $(I_j, I_i) \in E$  where E= set of edges of Graph G(V, E)

If such a directed graph allows to reach from any point to any other point, it needs to be strongly connected implying there is a path from  $I_i$  to  $I_j$  and from  $I_j$  to  $I_i$ . Checking if a path from  $I_j$  exists to  $I_i$  will involve reversing the edge directions in the directed graph and checking if  $I_i$  can be reached from  $I_j$ . So if the mayor is right, it should be possible to traverse from  $I_j$  to  $I_i$  with the edges inverted.

Such a strongly connected directed graph can be traversed in linear time using DFS with a run time of O(n+m)

B) Since the mayor's original claim is false  $\implies$  G is not strongly connected. However the focus shifts to the town hall being "strongly connected" with the rest of the nodes. In order for this to happen the node for town hall say T must be a sink of a strongly connected component. The approach would involve determining all such components containing T such that they are strongly connected. Next run a DFS in O(n+m) to determine all nodes reachable from T if all these nodes belong to to same strongly connected component, it should be possible to travel from T to these and back.

#### (3: Ch#3 Ex#3)

(4)

A) The intersections can be viewed as nodes of a directed

#### (5: Ch#4 Ex#3)

A) The intersections can be viwe

#### (6)

Given two sets A and B, each containing n positive integers. Perform reordering maximising  $\Pi_{i=1}^n a_i^{b_i}$ . There are 4 ways to proceed. Consider: A = 1, 2, 3; B = 4, 5, 6 1. Larger values of a raised to larger values of b

- 2. Smaller values of a raised to larger values of b
- 3. Larger values of a raised to smaller values of b
- 4. Smaller values of a raised to smaller values of b

It is easy to rule out Case 4, since it would the smallest possible product. The case maximising the payoff is Case 1.

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Saket Choudhary CSCI-570 : Homework # 3 HW3 (continued)

## (7: Ch#4 Ex#4)

Given two sequence S' of size m and S of size n, to determine if  $S' \subset S$ 

S' can be visualised as a DAG or more specifically as a topological ordering. Also treat S as a DAG. We keep two pointers, one for S' and one for S and perform a DFS on S till we hit a element from S' starting with S'[0], once S'[0] is found in S we start a DFS again after deleting all preceding elements of S now continuing till we find S[1] and deleting the intermediate hits if any.

(8)

A) The intersections can be viwe