

Q1. Checking validity

```
>>> tt_true(Expr('|', P, expr('~P')))  
True  
>>>  
>>> tt_true(Expr('>>', P, P))  
True  
>>>  
>>> tt_true(Expr('>>', P, expr('P | Q')))  
True  
>>>  
>>> tt_true(Expr('>>', expr('P | Q'), P))  
False  
>>>  
>>> tt_true(Expr('<=>', Expr('>>', expr('A & B'), expr('C')), Expr('>>', A, expr('B >> C'))))  
True  
>>>  
>>> tt_true(Expr('>>', Expr('>>', expr('A >> B'), A), A))  
True
```

Q2. Satisfiability

```
>>> dpll_satisfiable(Expr('>>', 'Alive', NaryExpr('&', expr('~Dead'), expr('~Alive'), expr('~Dead'))))  
{Alive: False}  
>>>  
>>> dpll_satisfiable(Expr('>>', P, Expr('|', expr('~P'), P)))  
{P: True}  
>>>  
>>> dpll_satisfiable(Expr('~', Expr('|', P, expr('~P'))))  
False
```

Q3. Propositional Consequence

2. $p \models p \wedge q$: False
3. $p \models p \vee q$: True
4. $p \models \neg \neg p$: True
5. $p \rightarrow q \models \neg p \rightarrow \neg q$: False

6. $\neg p \models p \rightarrow q : \text{True}$
7. $\neg q \models p \rightarrow q : \text{False}$
8. $p, p \rightarrow q \models q : \text{True}$
9. $\neg p, q \rightarrow p \models \neg q : \text{True}$

Q4. English to FOL

2. Everything is either dead or alive.
 $\mathbf{A} x \text{ things}(x) \wedge (\text{dead}(x) \vee \text{alive}(x))$
3. Dead things are not animate.
 $\mathbf{A} x \text{ things}(x) \wedge \text{dead}(x) \wedge \sim \text{animate}(x)$
4. Zombies are not alive but they are animate.
 $\mathbf{A} x \text{ zombies}(x) \wedge \sim \text{alive}(x) \wedge \text{animate}(x)$
5. Good food is not cheap and cheap food is not good.
 $\mathbf{A} x \text{ food}(x) \wedge (\text{good}(x) \Leftrightarrow \sim \text{cheap}(x))$
6. John has exactly two brothers.
 $\mathbf{E} x, y \text{ brother}(\text{John}, x) \wedge \text{brother}(\text{John}, y) \wedge \sim(x=y) \wedge \sim(\mathbf{E} z \text{ brother}(\text{John}, z) \wedge \sim(z=x) \wedge \sim(z=y))$
7. No person can have two mothers.
 $\mathbf{A} x, y, z \text{ person}(x) \wedge \text{person}(y) \wedge \text{person}(z) \wedge \text{mother}(x, z) \wedge \text{mother}(y, z) \wedge (y=z)$
8. If John has a sister, she is smart.
 $\mathbf{E} x \text{ sister}(\text{John}, x) \Rightarrow \text{smart}(x)$
9. Every person is either male or female and no person can be both male and female.
 $\mathbf{A} x \text{ person}(x) \Rightarrow (\text{male}(x) \vee \text{female}(x)) \wedge \sim(\text{male}(x) \wedge \text{female}(x))$
10. The enemy of your enemy is your friend.
 $\mathbf{A} x, y (\text{enemy}(\text{you}, x) \wedge \text{enemy}(y, x)) \Rightarrow \text{friend}(\text{you}, y)$
11. An ancestor of your ancestor is your ancestor.

$\forall x, y (\text{ancestor}(\text{you}, x) \wedge \text{ancestor}(x, y)) \Rightarrow \text{ancestor}(\text{you}, y)$

Q5. CNF and Horn Clauses

2. $\forall x \forall y \text{ married}(x, y) \rightarrow \text{loves}(x, y) \vee \text{hates}(x, y)$

Everyone either loves or hates those they have been married to.

It cannot be expressed as a set of horn clause.

CNF : $(\text{loves}(x, y) \vee \sim \text{married}(x, y) \vee \text{hates}(x, y))$

3. $\forall x \forall y \text{ loves}(x, y) \leftrightarrow \text{loves}(y, x)$

Everyone loves all those who love them.

It can be expressed as a set of horn clauses.

CNF : $((\text{loves}(x, y) \vee \sim \text{loves}(y, x)) \wedge (\text{loves}(y, x) \vee \sim \text{loves}(x, y)))$

Horn Clause : $\{(\text{loves}(x, y) \vee \sim \text{loves}(y, x)), (\text{loves}(y, x) \vee \sim \text{loves}(x, y))\}$

4. $\forall x \forall y \text{ dating}(x, y) \vee \text{engaged}(x, y) \rightarrow \text{knows}(x, y) \wedge \text{likes}(x, y)$

Everyone knows and likes those they are dating or are engaged to.

It can be expressed as a set of horn clauses.

CNF : $((\sim \text{dating}(x, y) \vee \text{knows}(x, y)) \wedge (\sim \text{engaged}(x, y) \vee \text{knows}(x, y)) \wedge (\sim \text{dating}(x, y) \vee \text{likes}(x, y)) \wedge (\sim \text{engaged}(x, y) \vee \text{likes}(x, y)))$

Horn Clause : $\{(\sim \text{dating}(x, y) \vee \text{knows}(x, y)), (\sim \text{engaged}(x, y) \vee \text{knows}(x, y)), (\sim \text{dating}(x, y) \vee \text{likes}(x, y)), (\sim \text{engaged}(x, y) \vee \text{likes}(x, y))\}$

5. $\forall x \forall y \text{ loves}(x, y) \rightarrow \neg \text{hates}(x, y)$

No one hates those they love.

It can be expressed as a set of horn clauses.

$(\sim \text{hates}(x, y) \vee \sim \text{loves}(x, y))$

Horn Clause : $\{(\sim \text{hates}(x, y) \mid \sim \text{loves}(x, y))\}$

6. $\forall x \forall y \neg \text{knows}(x, y) \rightarrow \neg \text{likes}(x, y)$

No one likes those they don't know.

It can be expressed as a set of horn clauses.

$(\sim \text{likes}(x, y) \mid \text{knows}(x, y))$

Horn Clause : $\{(\sim \text{likes}(x, y) \mid \text{knows}(x, y))\}$

7. $\forall x \exists y \text{knows}(x, y) \wedge \text{hates}(x, y)$

Everyone hates someone they know.

It can be expressed as a set of horn clauses.

CNF : $\text{knows}(x, F(x)) \wedge \text{hates}(x, F(x))$

Horn Clause : $\{\text{knows}(x, F(x)) , \text{hates}(x, F(x))\}$

skolem functions : $\{F(x)\}$

8. $\exists y \forall x \text{knows}(x, y) \wedge \text{hates}(x, y)$

Some are hated by all those who know them.

It can be expressed as a set of horn clauses.

CNF : $\text{knows}(x, A) \wedge \text{hates}(x, A)$

Horn Clause : $\{\text{knows}(x, A) , \text{hates}(x, A)\}$

skolem constants : $\{A\}$

9. $\neg (\forall x \text{loves}(x, x))$

There are some who don't love themselves.

It can be expressed as a set of horn clauses.

CNF : $\sim \text{loves}(A,A)$

Horn Clause : $\{\sim \text{loves}(A,A)\}$

skolem constants : $\{A\}$

10. $\neg (\exists x \forall y \text{ knows}(x, y))$

None knows everyone.

It can be expressed as a set of horn clauses.

CNF : $\sim \text{knows}(x,F(x))$

Horn Clause : $\{\sim \text{knows}(x,F(x))\}$

skolem functions : $\{F(x)\}$