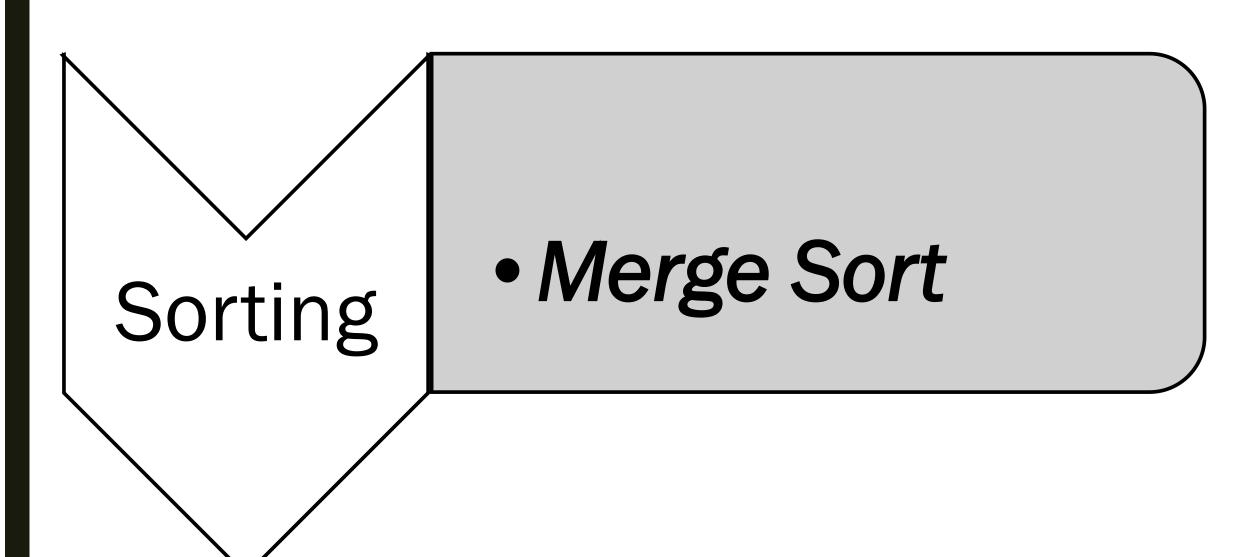
Todays discussion

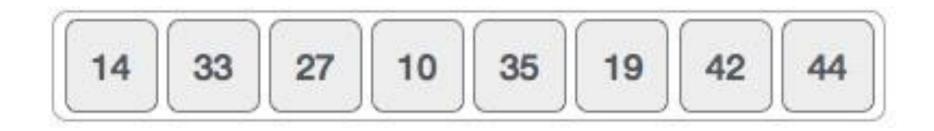


Merge Sort

Merge sort is a sorting technique based on divide and conquer technique. With worst-case time complexity being O(n log n), it is one of the most respected algorithms.

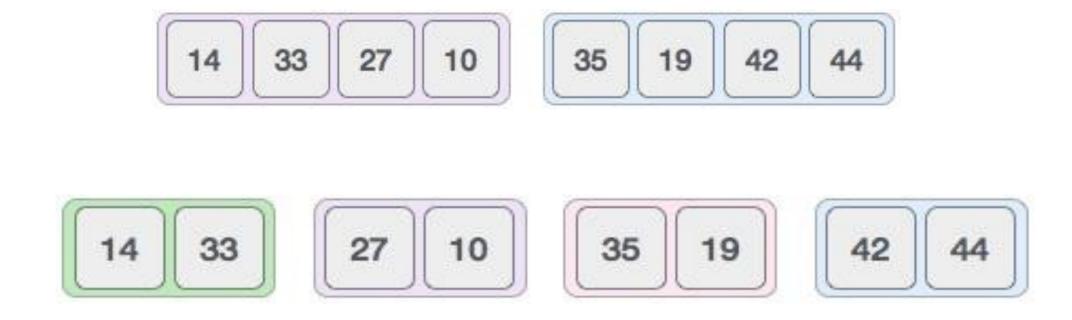
Merge sort first divides the array into equal halves and then combines them in a sorted manner.

■ Take an unsorted array as the following



- Merge sort first divides the whole array iteratively into equal halves unless the atomic values are achieved.
 - here that an array of 8 items is divided into two arrays of size 4

■ This does not change the sequence of appearance of items in the original. Now we divide these two arrays into halves.



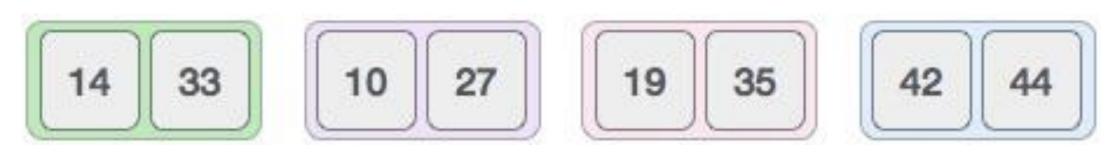
■ Further divide these arrays and we achieve atomic value which can no more be divided.



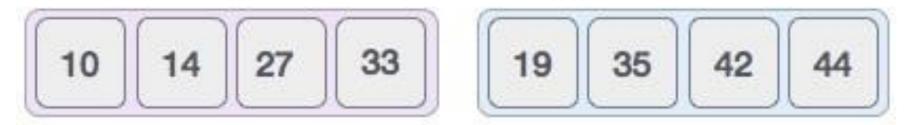
■ Now, we combine them in exactly the same manner as they were broken down. Please note the color codes given to these lists.



- Compare the element for each list and then combine them into another list in a sorted manner.
- 14 and 33 are in sorted positions.
- 27 and 10 and in the target list of 2 values we put 10 first, followed by 27. We change the order of 19 and 35 whereas 42 and 44 are placed sequentially.



■ In the next iteration of the combining phase, we compare lists of two data values, and merge them into a list of found data values placing all in a sorted order.



After the final merging, the list should look like this



14 23 45 98

6 33 42 67

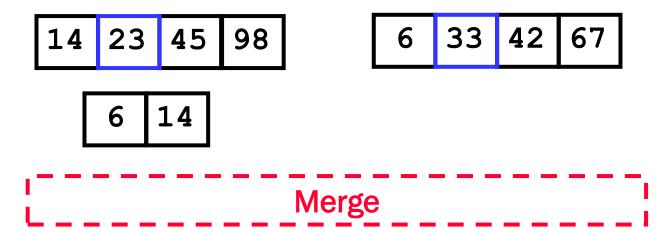
14 23 45 98

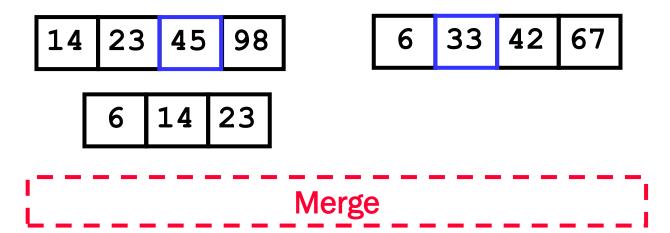
6 33 42 67

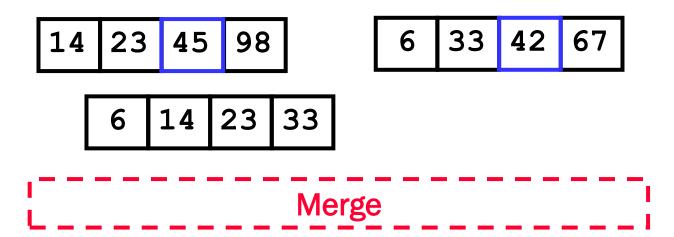
Merge

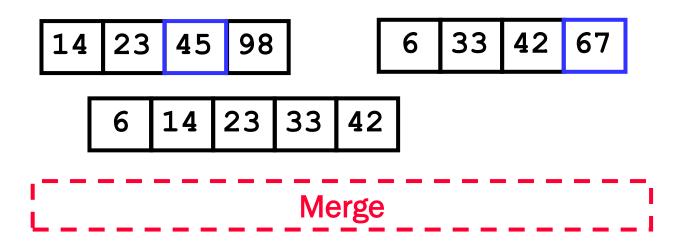
 14
 23
 45
 98
 6
 33
 42
 67

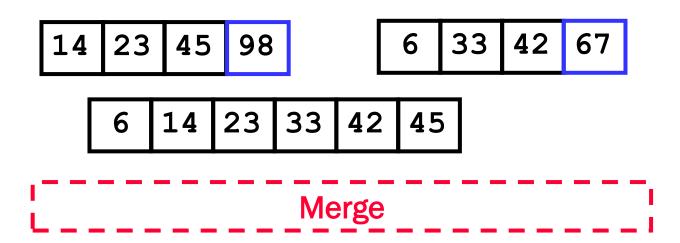
 6
 Merge

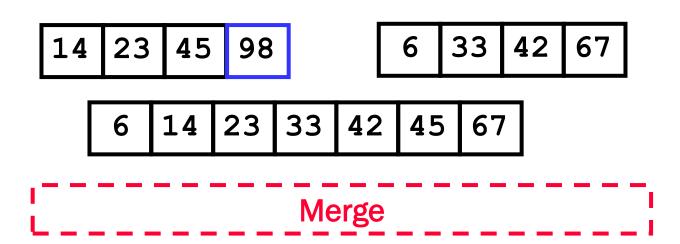


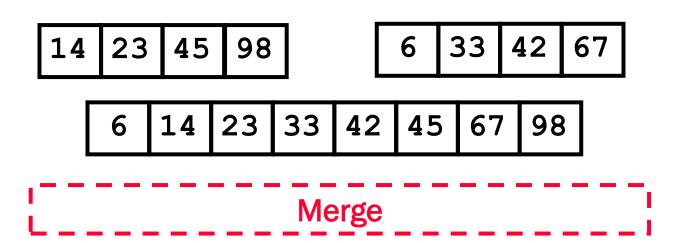












Merge Sort

- Merge sort keeps on dividing the list into equal halves until it can no more be divided.
- By definition, if it is only one element in the list, it is sorted. Then, merge sort combines the smaller sorted lists keeping the new list sorted too.

Algorithm

- Step 1 if it is only one element in the list it is already sorted, return.
- Step 2 divide the list recursively into two halves until it can no more be divided.
- Step 3 merge the smaller lists into new list in sorted order.

Pseudocode

```
procedure mergesort( var a as array )
 if (n == 1) return a
 var I1 as array = a[0] ... a[n/2]
 var I2 as array = a[n/2+1] ... a[n]
 I1 = mergesort( I1 )
 I2 = mergesort(I2)
 return merge(I1, I2)
end procedure
```

Pseudocode...

```
procedure merge( var a as array, var b as array )
                                                         while (a has elements)
                                                           add a[0] to the end of c
                                                           remove a[0] from a
 var c as array
                                                         end while
 while ( a and b have elements )
                                                         while (b has elements)
   if (a[0] > b[0])
     add b[0] to the end of c
                                                           add b[0] to the end of c
     remove b[0] from b
                                                           remove b[0] from b
                                                         end while
   else
     add a[0] to the end of c
     remove a[0] from a
                                                         return c
   end if
 end while
                                                       end procedure
```

Recap

■ Three intuitive sorting algorithms

Merge Sort

- Divide array in two equal parts
- Separately sort left and right half
- Combine the two sorted halves to get the full array sorted
- Given two sorted lists A and B, combine into a sorted list C
 - Compare first element of A and B
 - Move it into C
 - Repeat until all elements in A and B are over
- Merging A and B

Array Stack

- **■** 32 **■** 21
- **■** 74 **■** 55
- **■** 89 **■** 64

Recap Merge Sort

- Sort A[0] to A[n/2-1]
- \blacksquare Sort A[n/2] to A[n-1]
- Merge sorted halves into B[0..n-1]
- How do we sort the halves?
 - Recursively, using the same strategy!

Divide and conquer

- Break up problem into disjoint parts
- Solve each part separately
- Combine the solutions efficiently

Merge Sort

- To sort A[0..n-1] into B[0..n-1]
- If n is 1, nothing to be done
- Otherwise
 - Sort A[0..n/2-1] into L (left)
 - Sort A[n/2..n-1] into R (right)
 - Merge L and R into B

Merging sorted lists

- Combine two sorted lists L and R into B
 - compare first element of L
 and R and
 - move the smaller of the two into B
 - Repeat until all elements in L and R have been moved
- Otherwise,
 - If L is empty, copy R into C
 - If R is empty, copy L into C

```
function Merge(A,m,B,n,C)
functionMergeSort(A,left,right,B)
                                          // Merge A[0..m-1], B[0..n-1] into C[0..m+n-1]
// Sort the segment A[left..right-1] into B
                                          i = 0; j = 0; k = 0;
if (right - left == 1) // Base case
                                          // Current positions in A,B,C respectively
                                          while (k < m+n)
B[0] = A[left]
                                          // Case 1: Move head of A into C
if (right - left > 1) // Recursive
                                          if (A[i] \leq B[j])
call
                                          \{ C[k] = A[i]; i++; k++; \}
                                          // Case 2: Move head of B into C
mid = (left+right)/2
                                          if (A[i] > B[j])
MergeSort(A,left,mid-1,L)
                                          {C[k] = B[j]; j++; k++;}
                                          // Case 0: One of the two lists is empty
MergeSort(A,mid,right-1,R)
                                          if (i==m) \{j++; k++;\}
Merge(L,mid-left,R,right-
                                          if (j==n) \{i++; k++;\}
mid,B)
```

Variations on merge

- Union of two sorted lists (discard duplicates)
- **1** [1,2,3] [2,3,6]
- If A[i] == B[j], copy A[i] to C[k] and increment i,j,k
- Intersection of two sorted lists
- If A[i] < B[j], increment i
- If B[j] < A[i], increment j
- If A[i] == B[j], copy A[i] to C[k] and increment i,j,k
- **Exercise:** List difference: elements in A but not in B
- \blacksquare [1,2,3] [3,4,6] = [1,2]

Merge Sort: Shortcomings

- Merging A and B creates a new array C
 - No obvious way to efficiently merge in place
- Extra storage can be costly
- Inherently recursive
 - Recursive call and return are expensive

Alternative approach

- Extra space is required to merge
- Merging happens because elements in left half must move right and vice versa
- **[**2,4,6] [1,3,5]
- Can we divide so that everything to the left is smaller than everything to the right?
 - No need to merge!
 - [2,1,3][6,4,5]