

# Marketplace Operations: From Resource Allocation to Recommendations

Akshit Kumar

Columbia Business School

Joint work with Omar Besbes and Yash Kanoria

# Matching in marketplaces

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**amazon**

**Flipkart**



Match demand location with  
fulfillment center

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**uc** Urban  
Company

**handy**

**Angi**

Match customers with service  
providers

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**NETFLIX**



Recommend movies/songs to  
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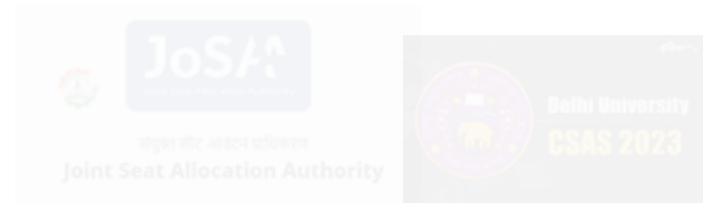
**Angi**

Match customers with service providers

## Part I

### Research Questions

1. How should platforms allocate limited resources dynamically and efficiently?



Match students to college programs

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## Part I

### Research Questions

1. How should platforms allocate limited resources dynamically and efficiently?
2. What are the fundamental driver of algorithmic performances in these marketplaces?

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# Matching in marketplaces

## Part II



Flipkart



Match demand location with fulfillment center



Angi

Match customers with service providers

NETFLIX



Recommend movies/songs to consumers



Match students to college programs

# Matching in marketplaces

## Part II

### Research Questions

1. What is the value of personalized recommendations in marketplaces?

**NETFLIX**



Recommend movies/songs to consumers



Match customers with service providers



Match students to college programs



# Part I in a Nutshell



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We study prototypical models of dynamic resource allocation problem (focus on the multisecretary prob)



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Shape of the reward distribution is a fundamental driver of algorithmic performance



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Workhorse policies can be highly suboptimal and near-optimal algorithms are overly specified



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Workhorse policies can be highly suboptimal and near-optimal algorithms are overly specified



Design a simple and near-optimal policy called Repeatedly Act using Multiple Simulations (**RAMS**)



# Part II in a Nutshell



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We study a prototypical utility model which comprises of public and private utility components



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Identify subtle interplay of personalized recommendations and capacity constraints on social welfare





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Personalized recos may have *little to no* benefit in markets with uncapacitated supply (think movies, songs)



# Part II in a Nutshell



We study a prototypical utility model which comprises of public and private utility components



Identify subtle interplay of personalized recommendations and capacity constraints on social welfare



Personalized recos may have *little to no* benefit in markets with uncapacitated supply (think movies, songs)



Personalized recos unlock *significant* social welfare in capacitated supply settings (think AirBnB, college admits)



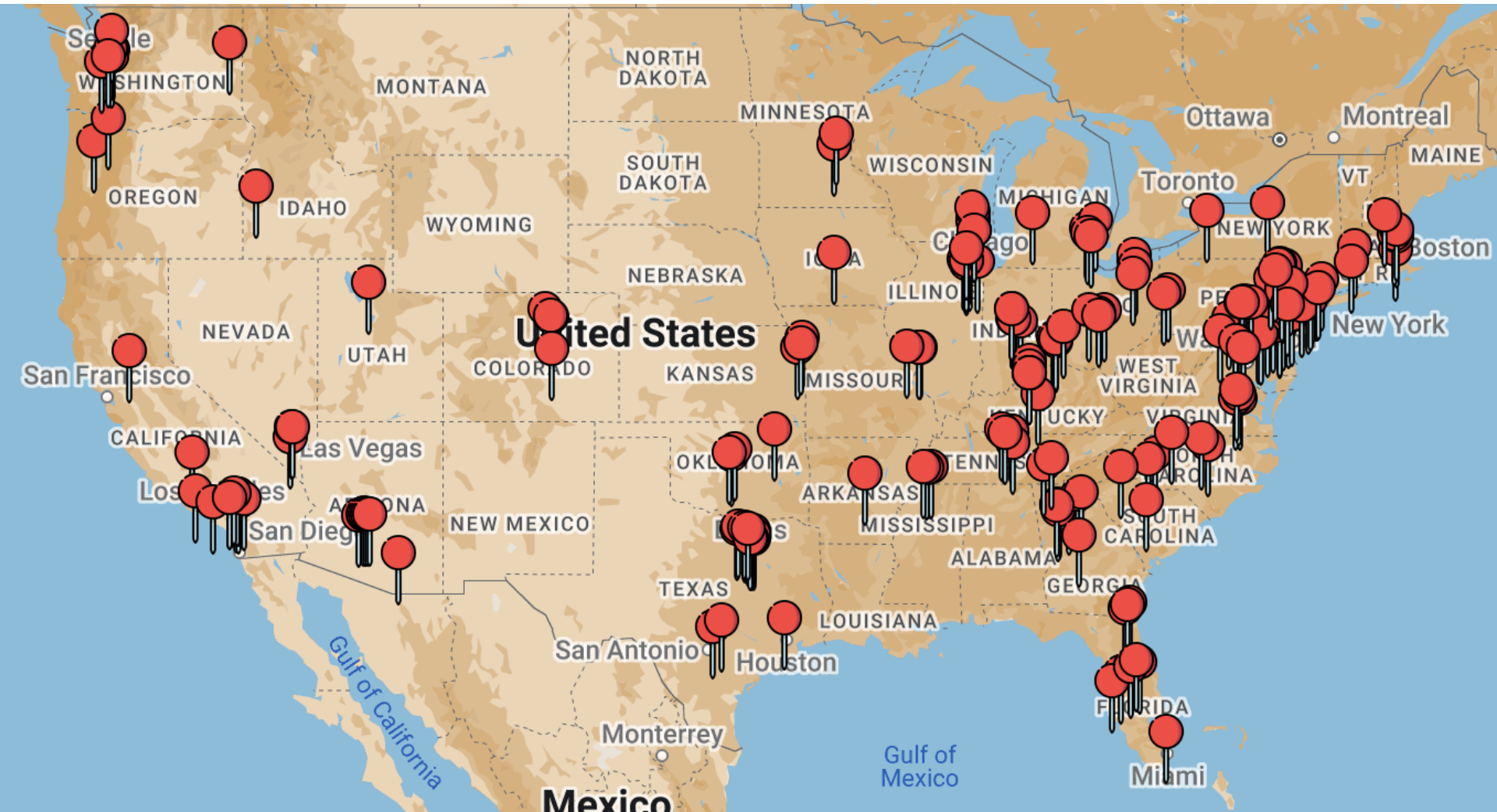
# Part I

## **Dynamic Resource Allocation: Algorithmic Design Principles & Spectrum of Achievable Performances**

*Forthcoming in Operations Research (2024)*

# Dynamic Resource Allocation

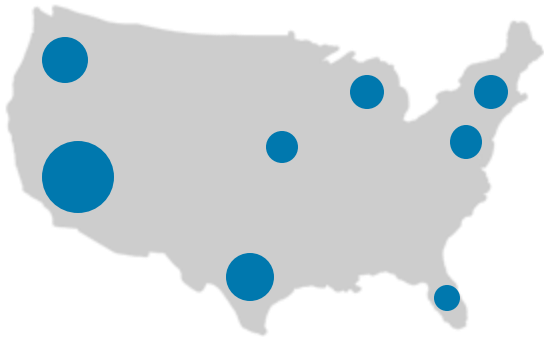
Algorithmic Design Principles & Spectrum of Achievable Performances



# Research Question

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A few (demand)  
types are present



Theory

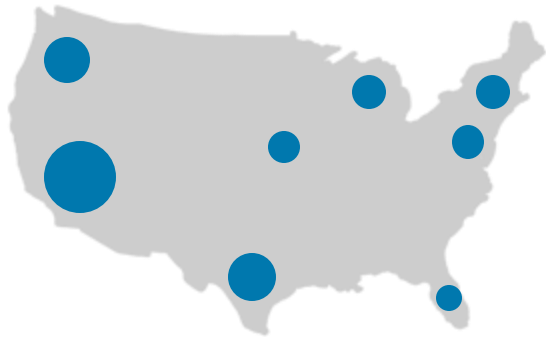
All (demand) types  
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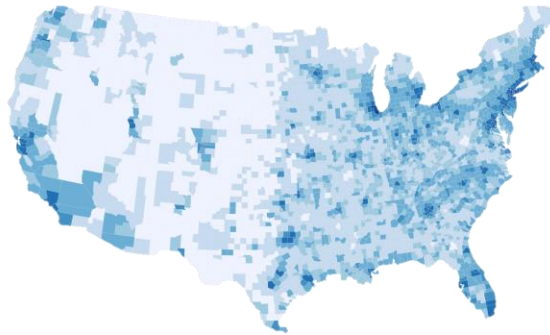
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Practice

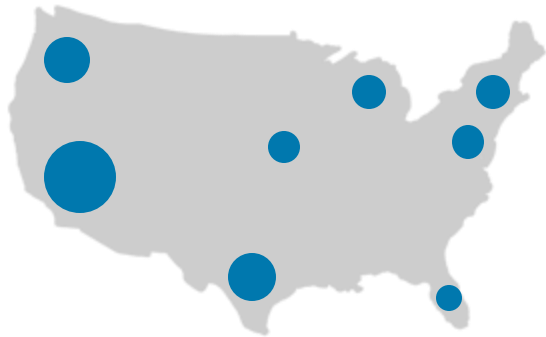
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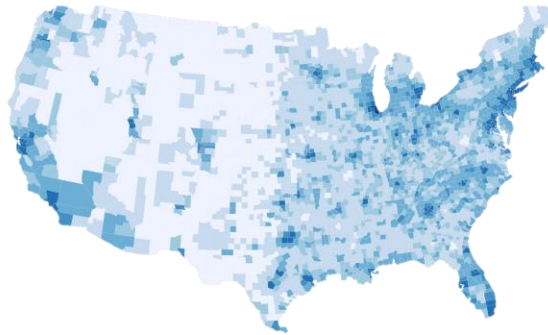
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Theory

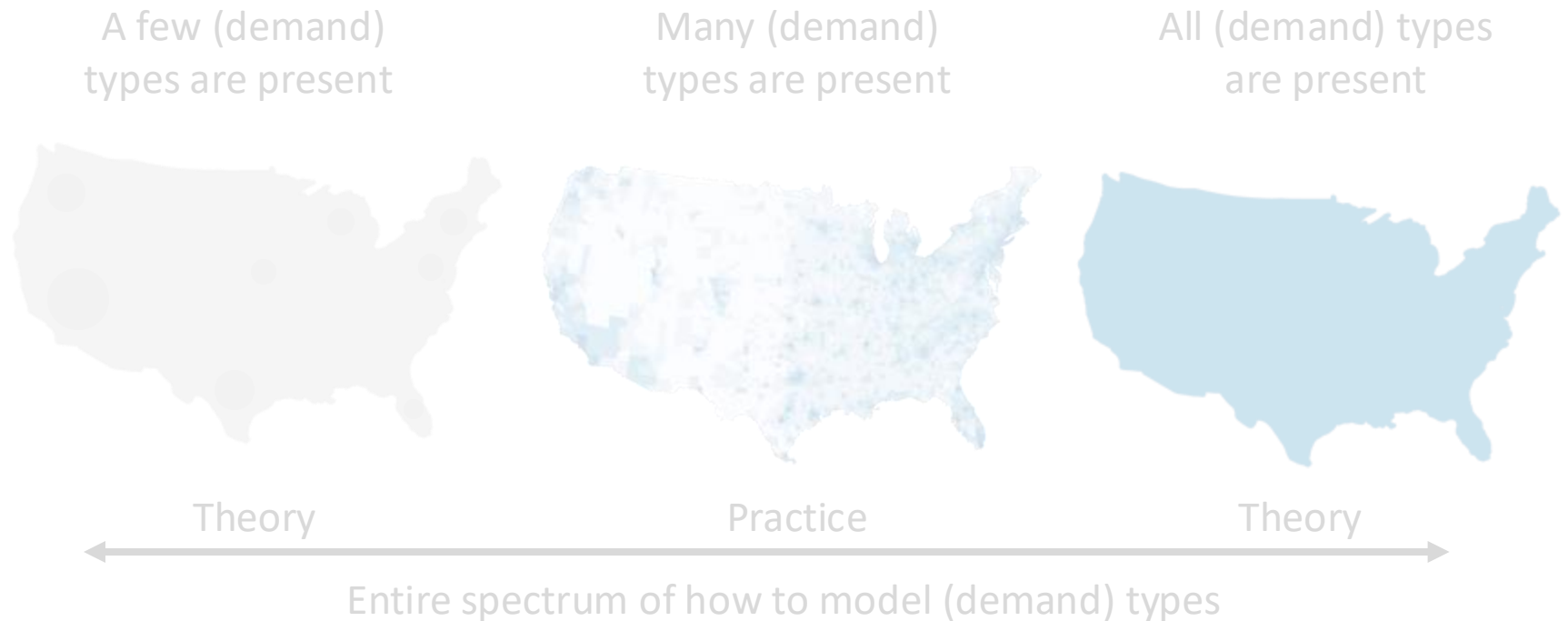
Practice

Theory

Entire spectrum of how to model (demand) types

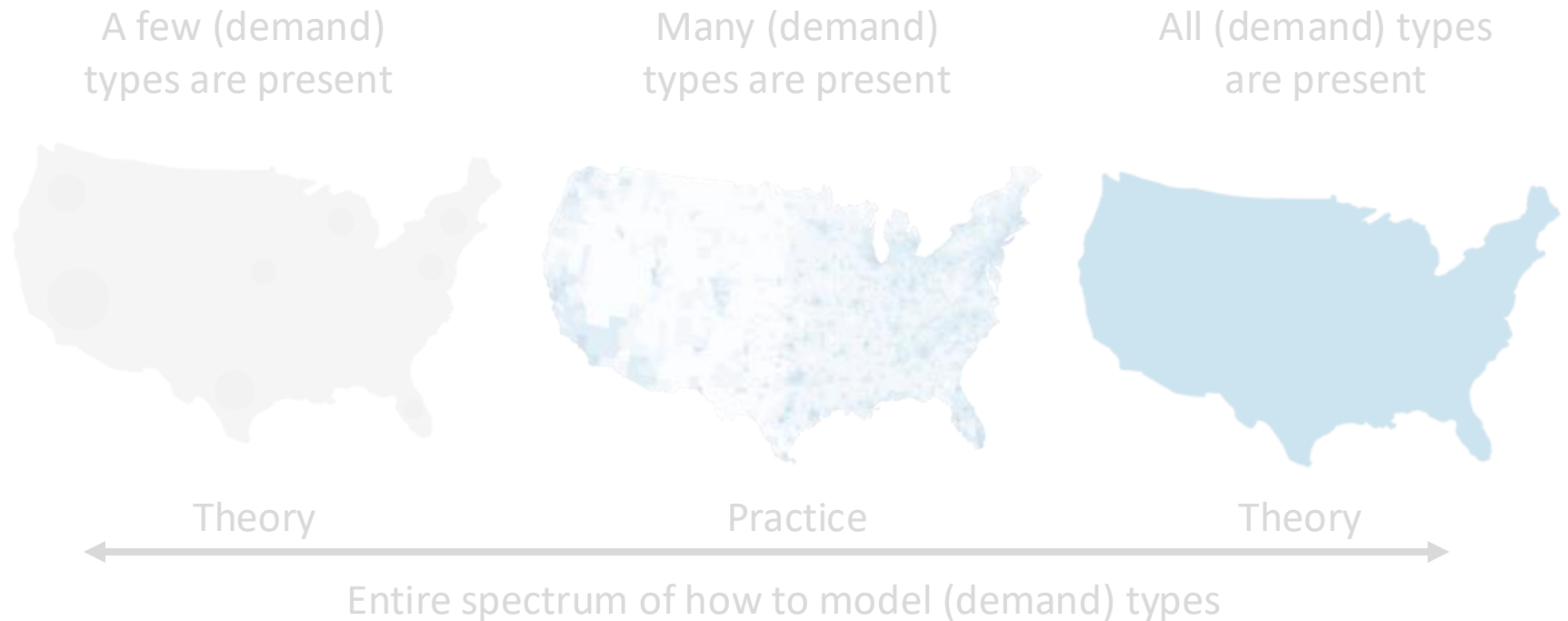


# Research Question



1. What is the interplay between the distribution of request types and achievable algorithmic performance?

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1. What is the interplay between the distribution of request types and achievable algorithmic performance?
2. Can we design a **unified, simple** and **near-optimal** algorithms which works for all type distributions?



Order  
Fulfillment

Network  
Revenue  
Management



**amazon**

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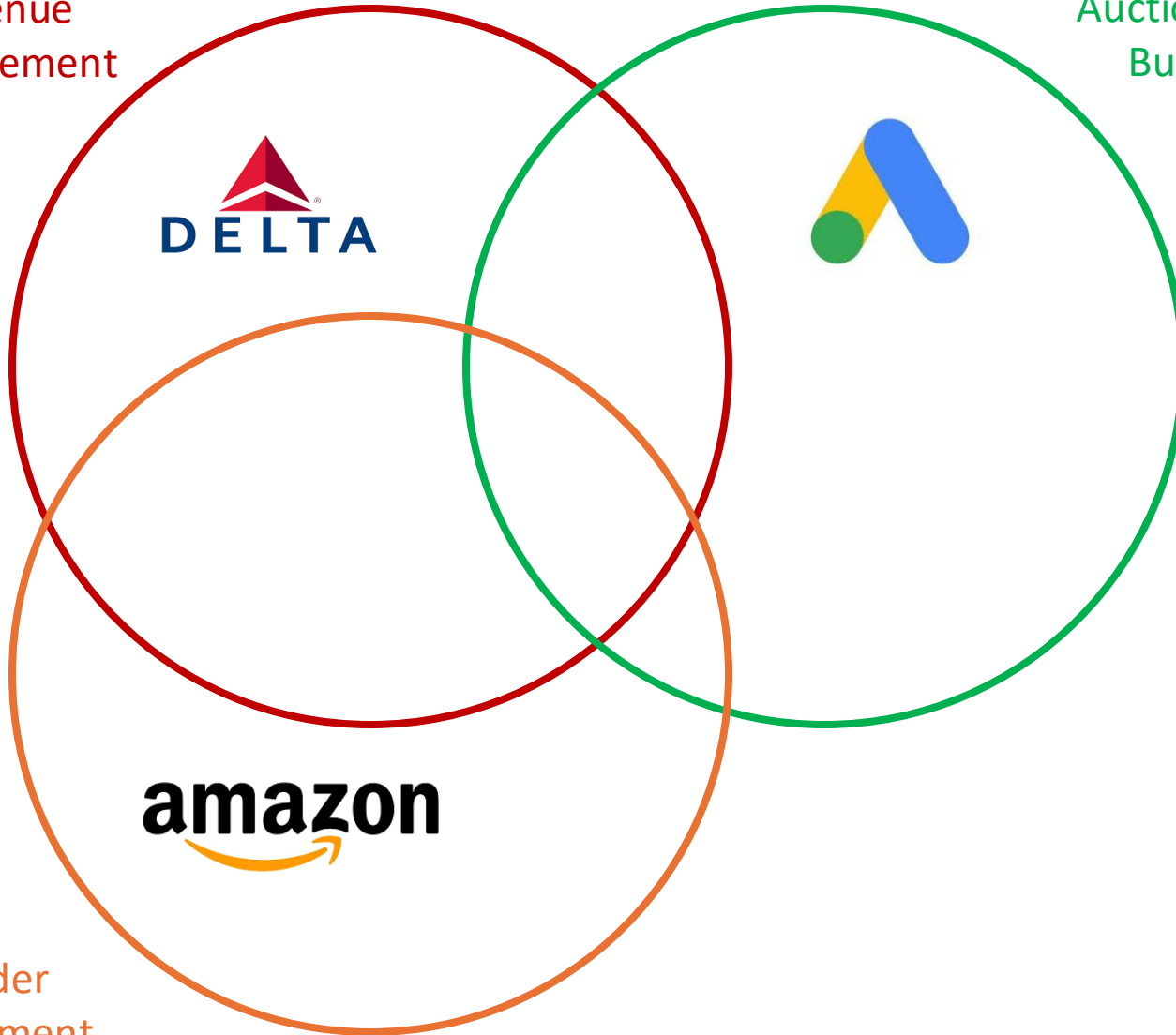


Repeated  
Auctions with  
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Multi-secretary



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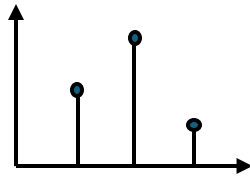
# Multi-secretary Problem

The background of the slide features three thick, curved lines that sweep across the frame. A blue line starts from the left edge and curves upwards towards the top right. A green line starts from the left edge and curves downwards towards the bottom right. A red line starts from the bottom right and curves upwards towards the top right. These lines intersect to form a large, abstract shape in the center of the slide.



# Multi-secretary Problem

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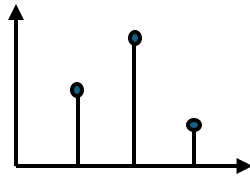
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# Multi-secretary Problem

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Bounded Regret



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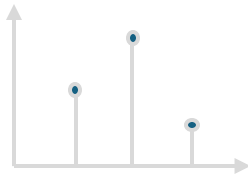
Logarithmic Regret



Regret is the additive gap b/w the value of hindsight opt. and value under some algorithm

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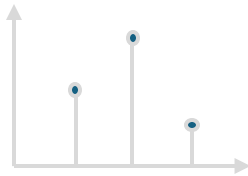


$\beta$ -clustered distributions

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Entire spectrum of regret scalings is possible

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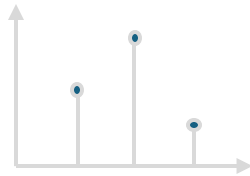
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# Multi-secretary Problem

one algorithm to solve them all

## Repeatedly Act using Multiple Simulations (RAMS)

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**Repeatedly Act using Multiple Simulations (RAMS)**

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# Online selection of top- $B$ values

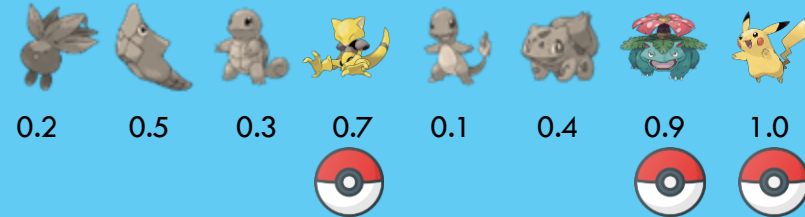
# Online selection of top- $B$ values

- Given a sequence of  $T$  values and budget  $B$



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- Given a sequence of  $T$  values and budget  $B$ , the DM wants to select the top  $B$  values
- The values arrive in an online fashion



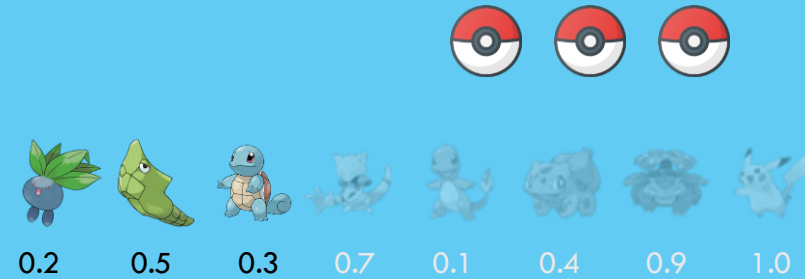
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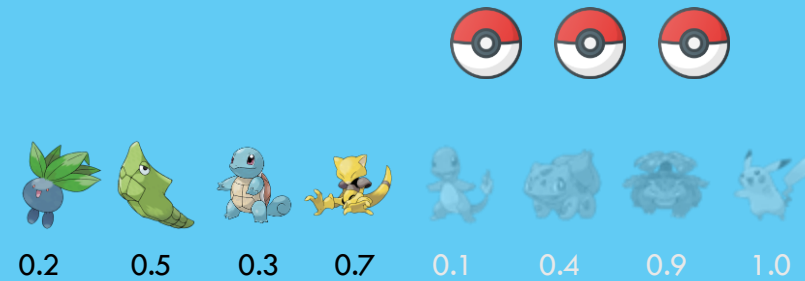
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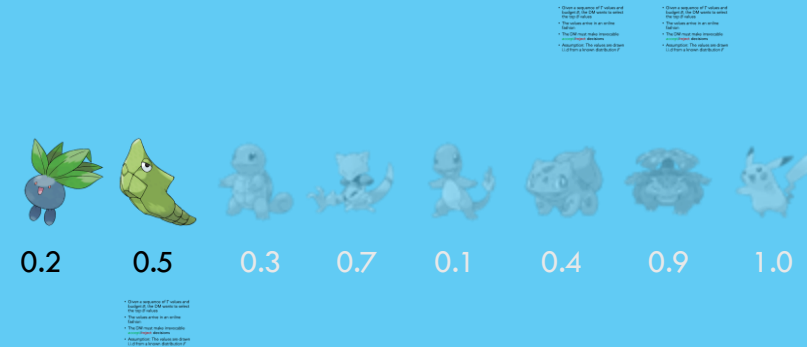
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$$\text{Regret}(\pi) = (\text{Expected Maximum Value in Hindsight}) - (\text{Expected Value under } \pi)$$

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Value in hindsight =  $0.7 + 0.9 + 1.0 = 2.6$



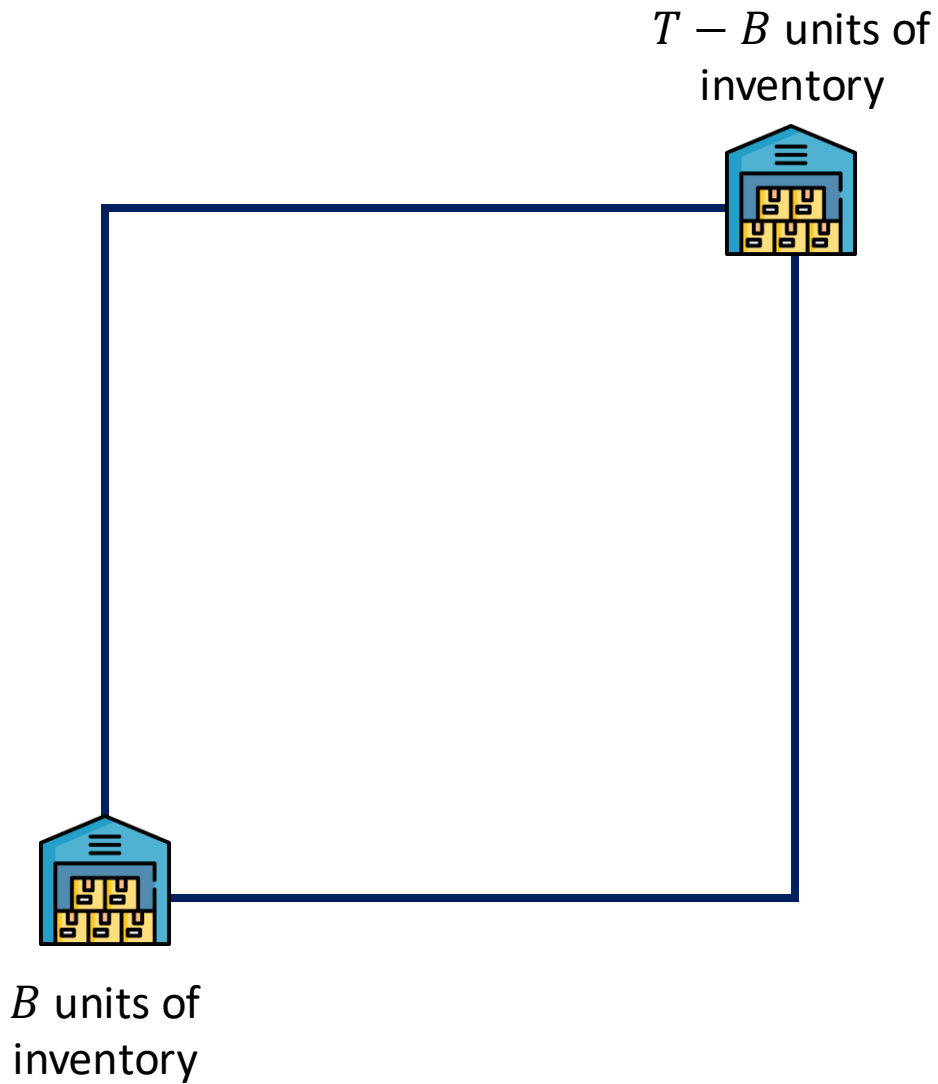
Value for policy  $\pi = 0.5 + 0.4 + 0.9 = 1.8$



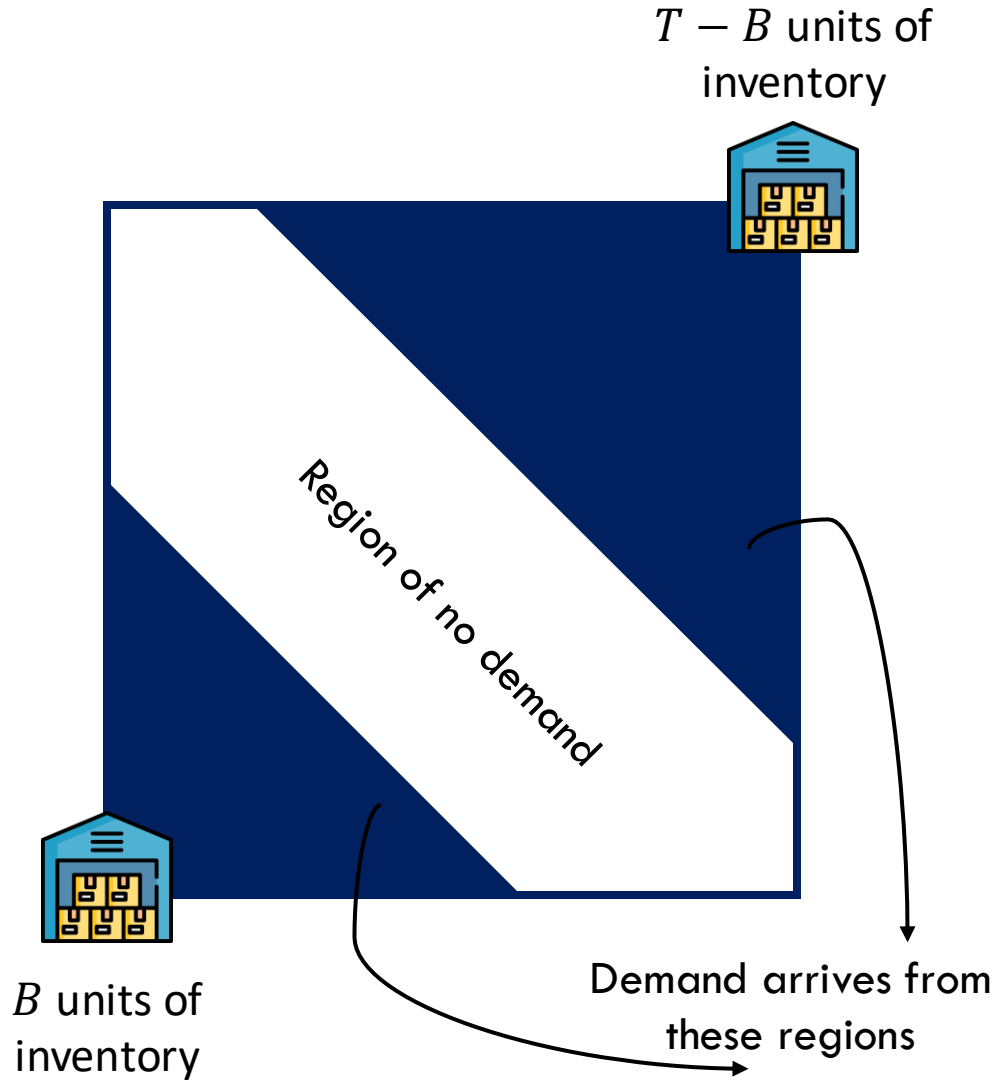


# A tale of two FCs

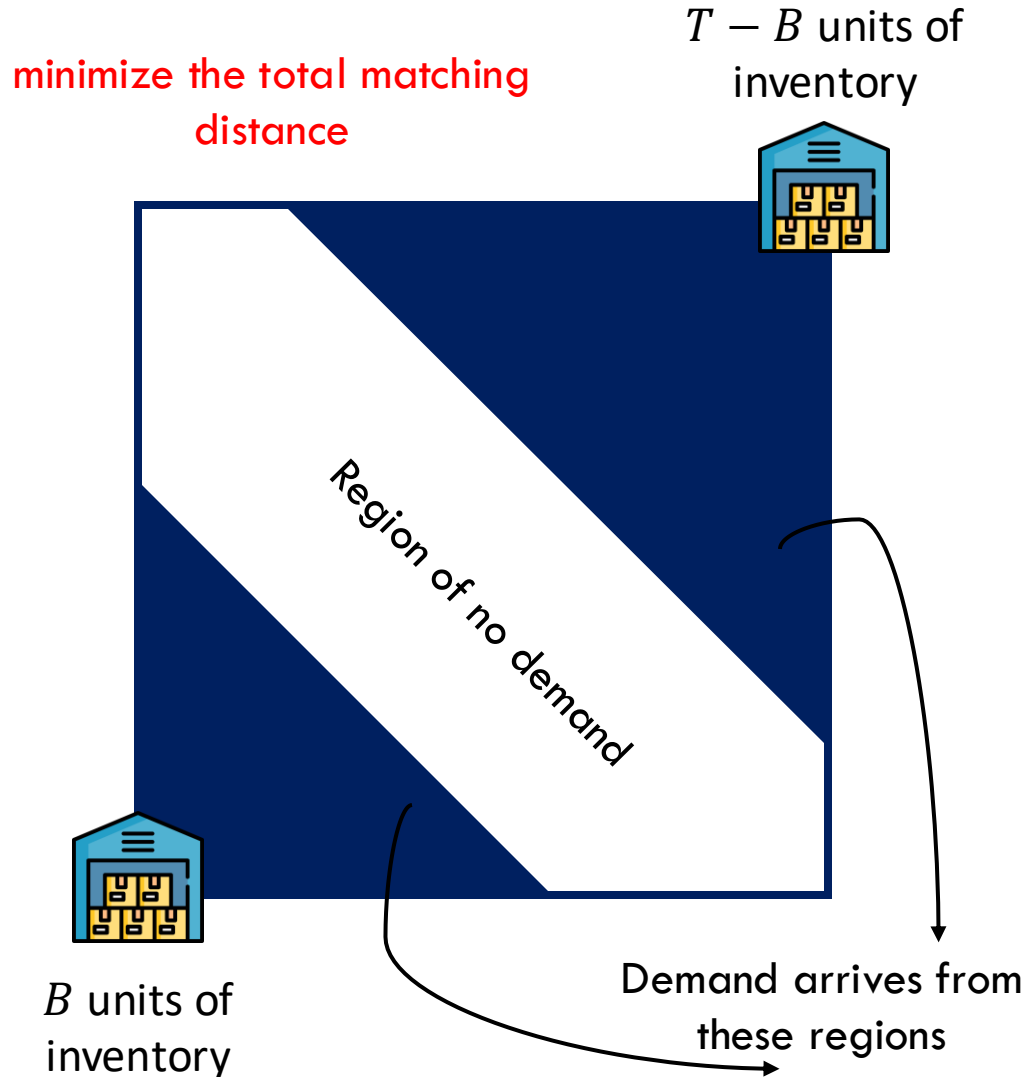
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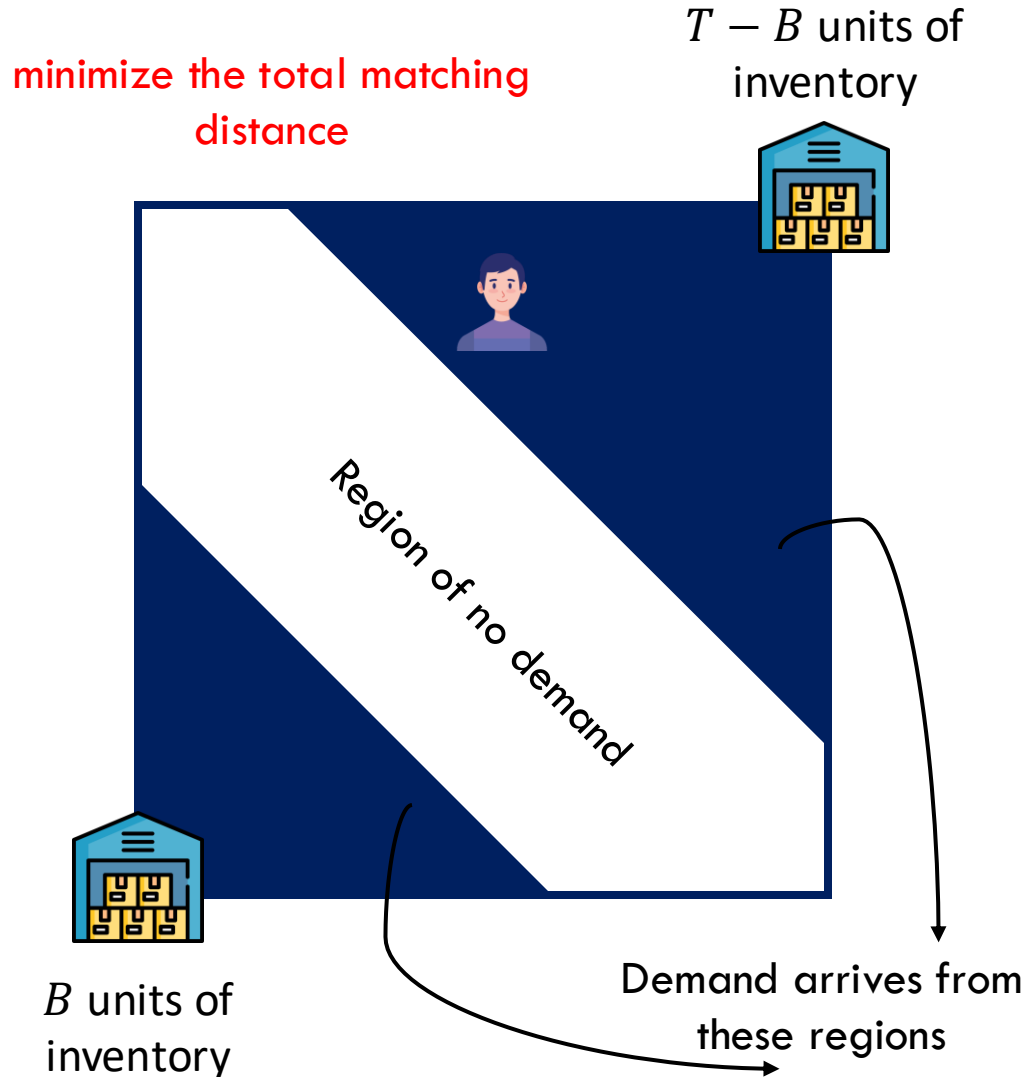
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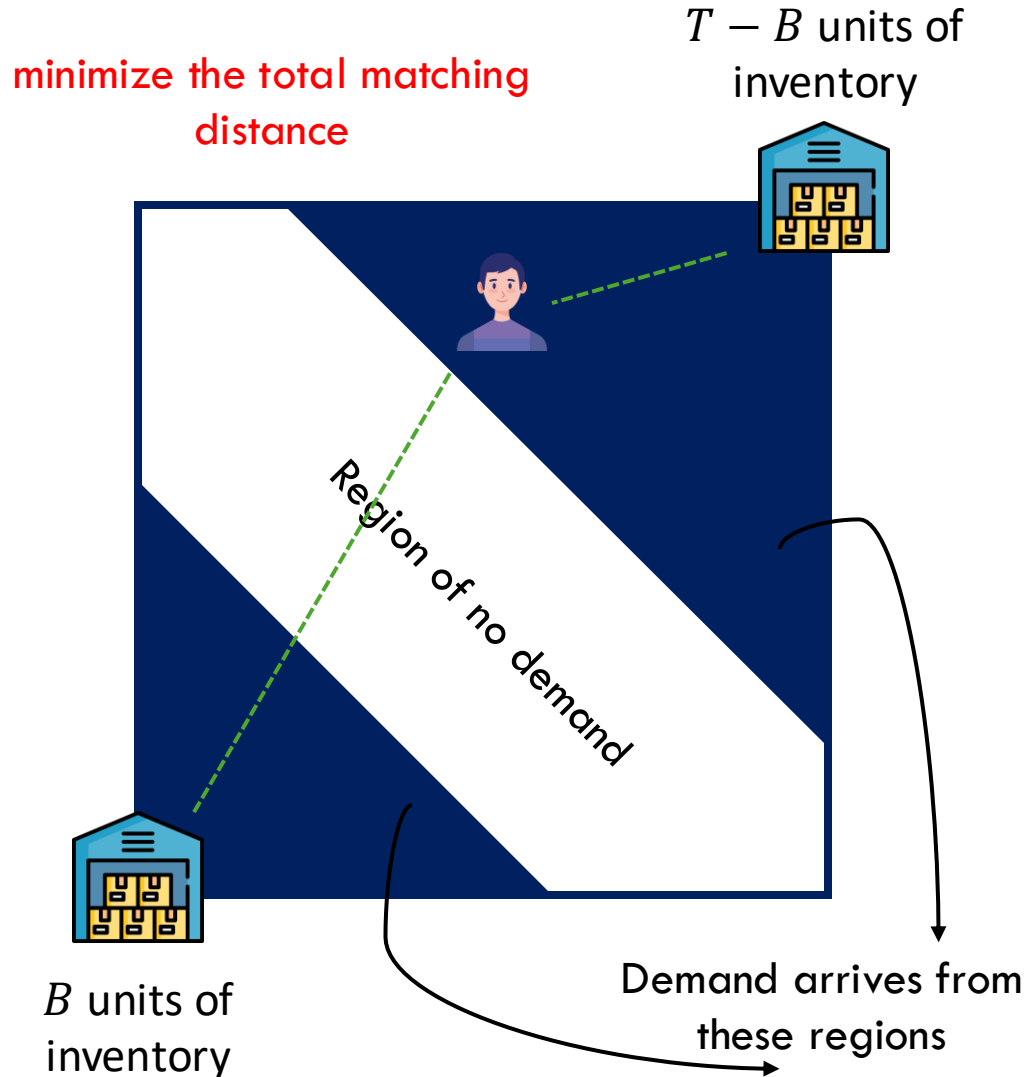
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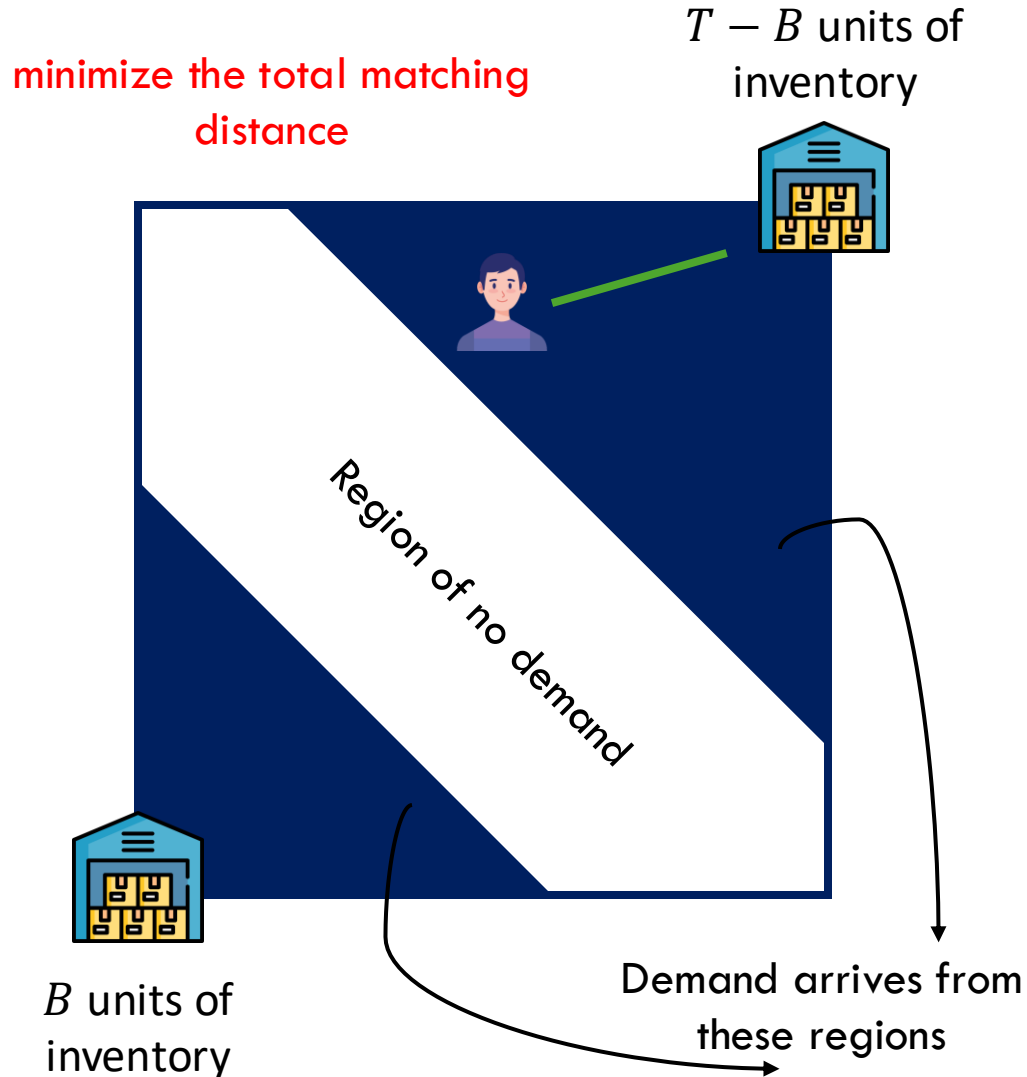
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# A tale of two FCs

minimize the total matching  
distance

$T - B - 1$  units of  
inventory



Region of no demand



$B$  units of  
inventory

Demand arrives from  
these regions

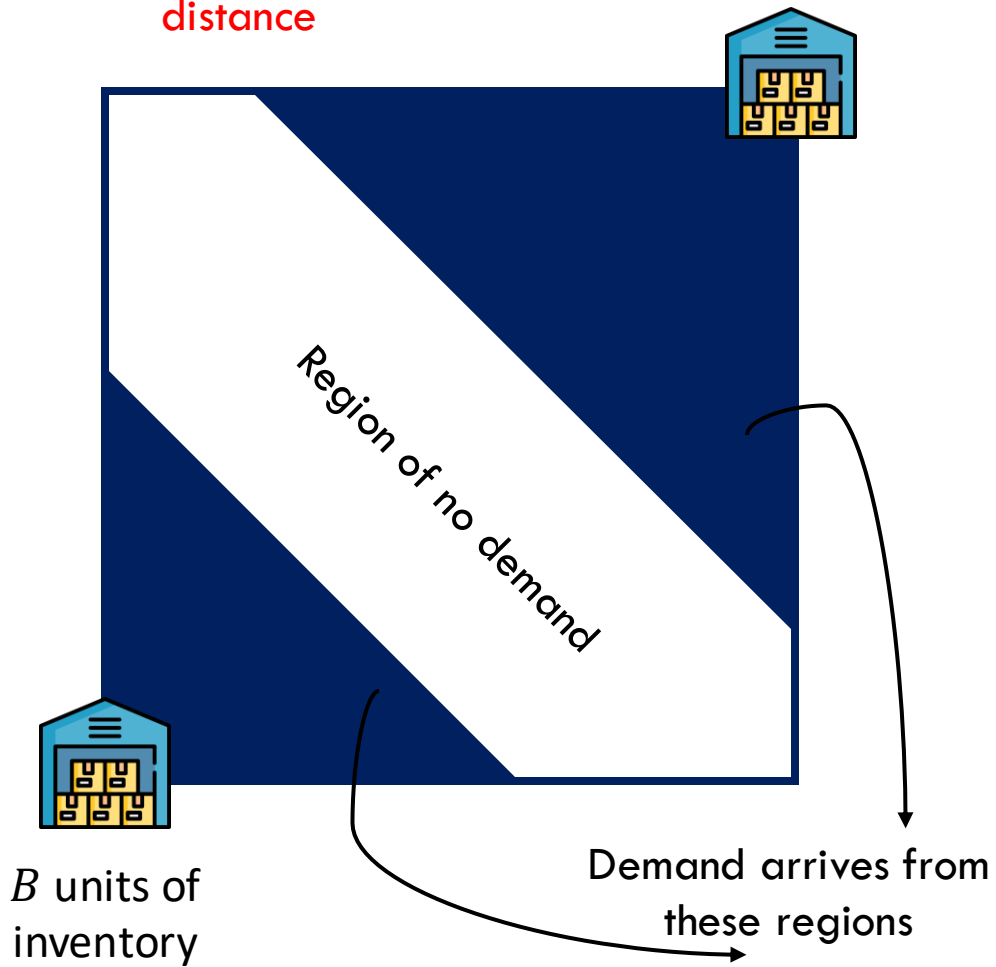




# A tale of two FCs

minimize the total matching  
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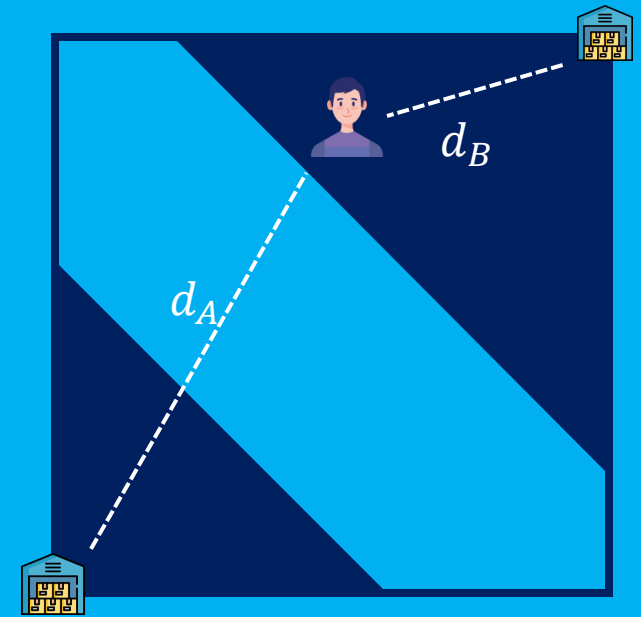
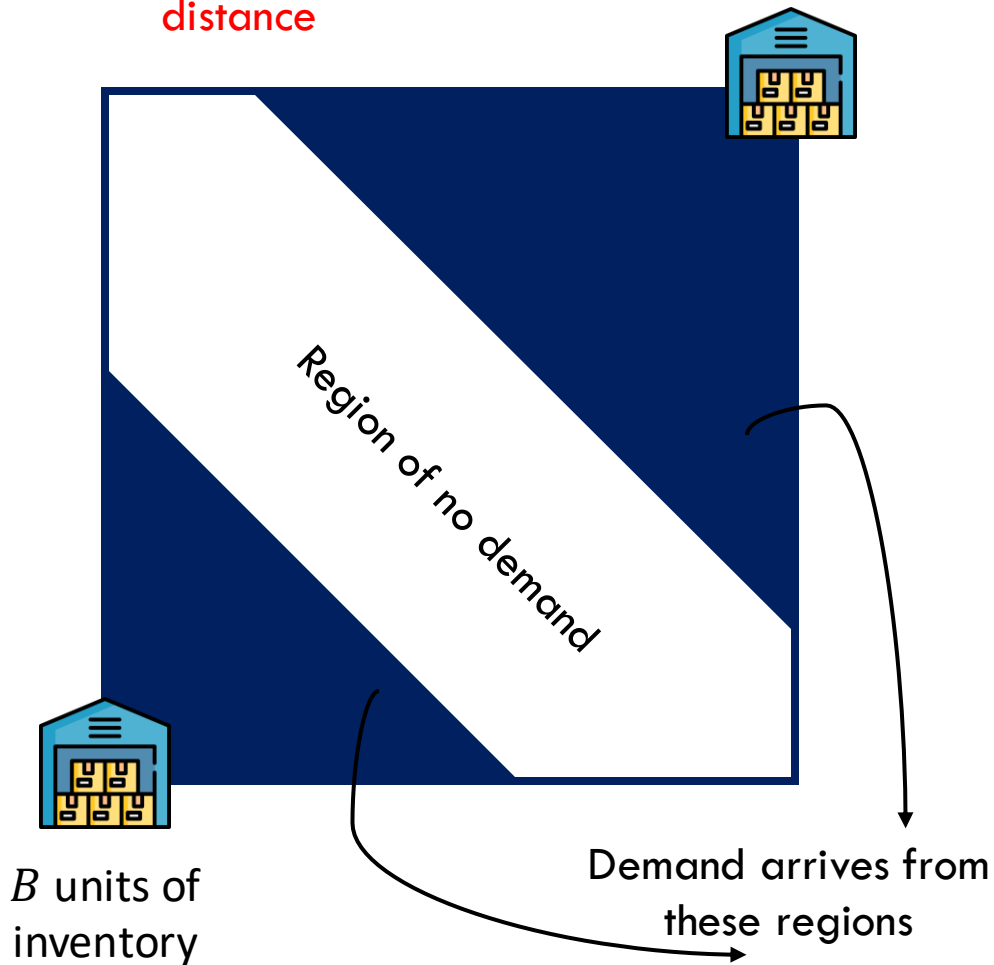
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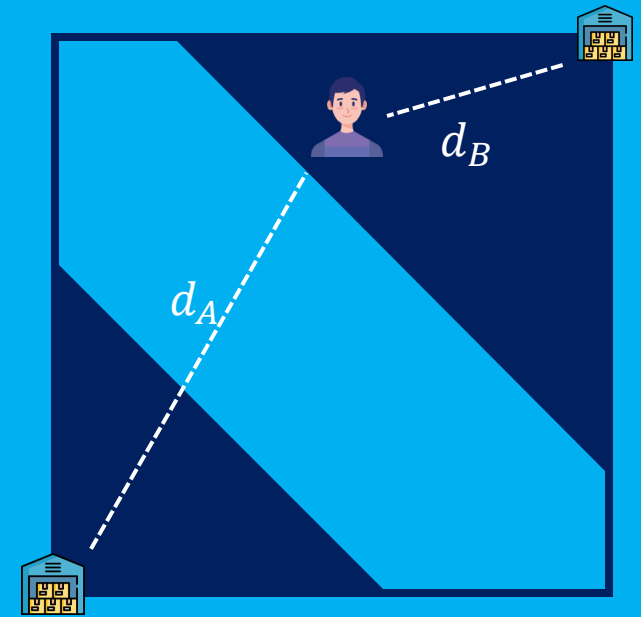
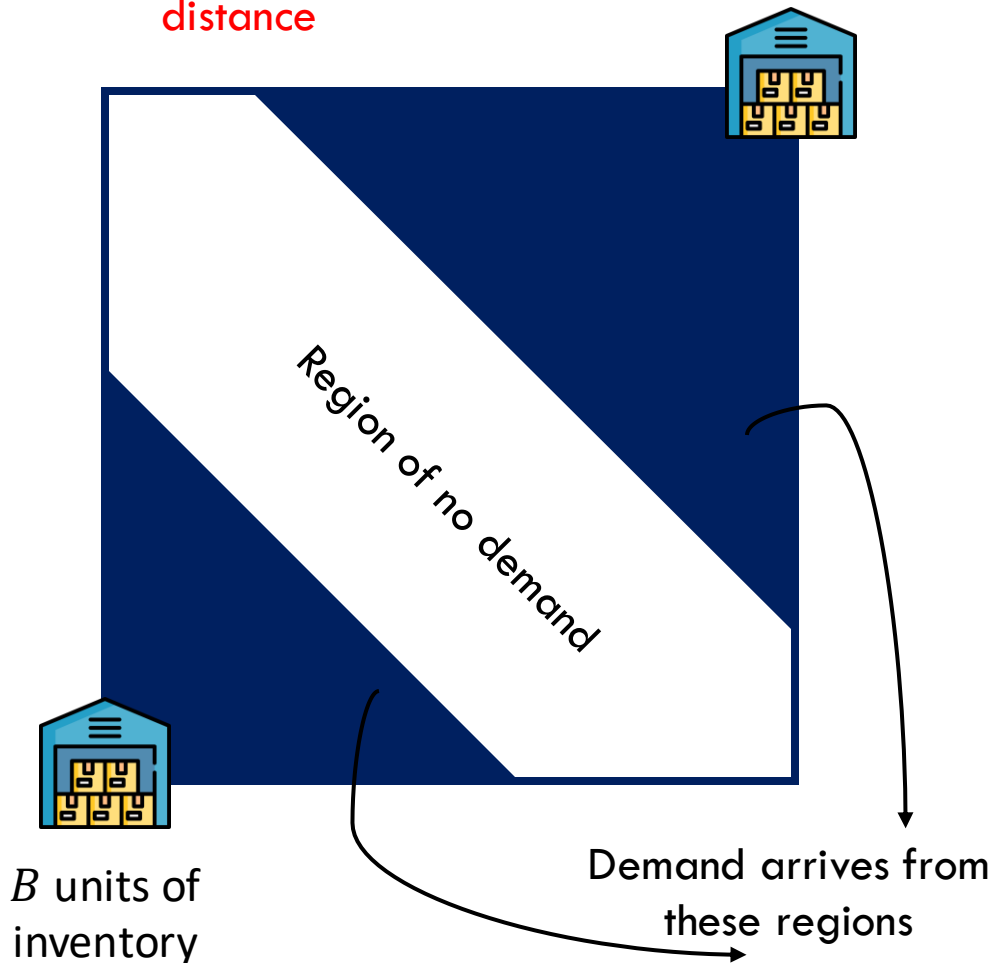
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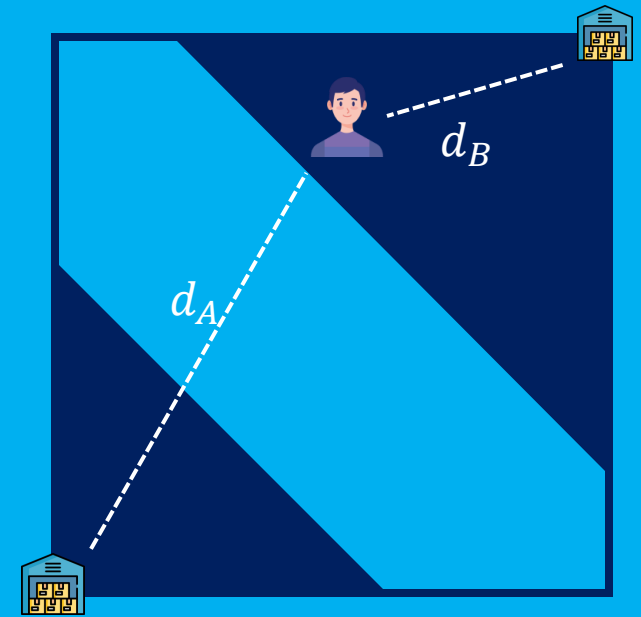
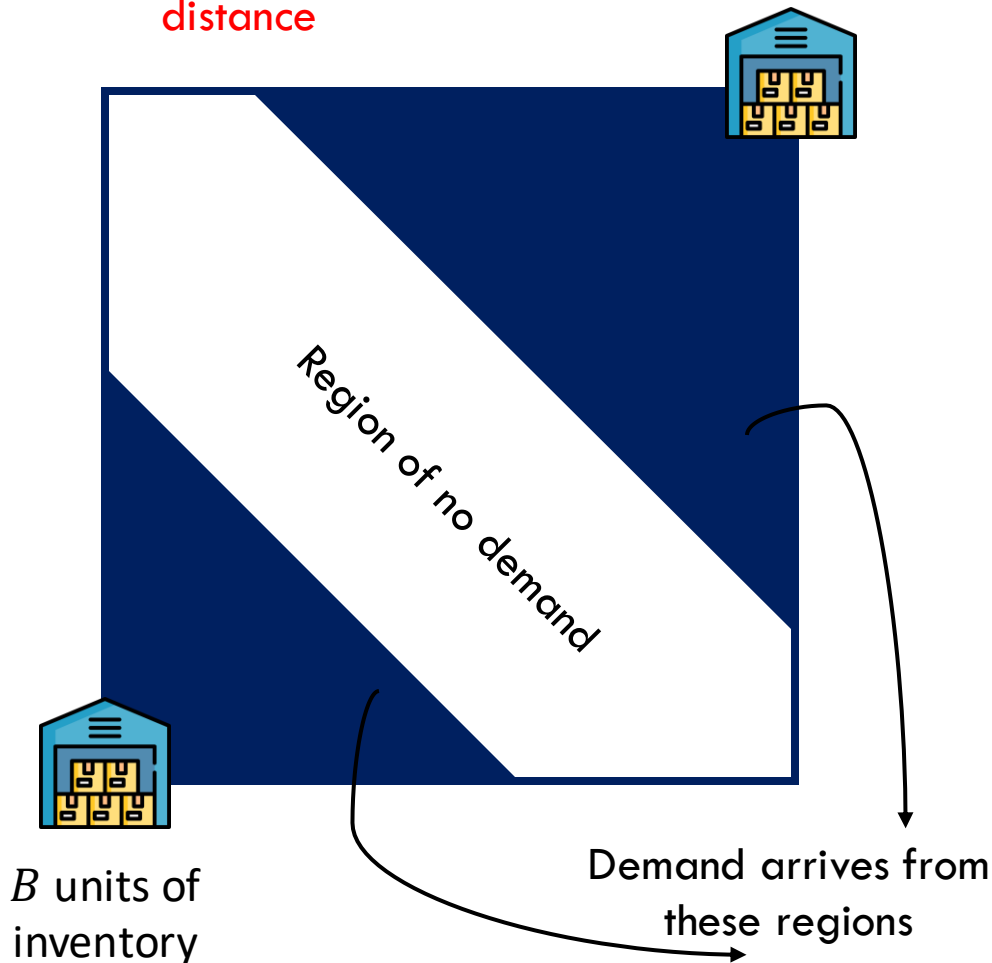
$$r = d_B - d_A + L$$

minimize total matching distance  $\equiv$  maximize total matching reward

# A tale of two FCs

minimize the total matching distance

$T - B - 1$  units of inventory



$$r = d_B - d_A + L$$

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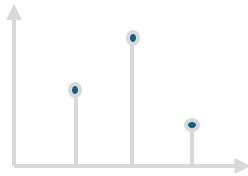


Dynamically choose the highest  $B$  reward values, given the reward distribution

maximize total matching reward  $\equiv$  multi-secretary problem

# Multi-secretary Problem

A few types  
are present  
Bounded Regret



All types are  
present  
Logarithmic Regret



Entire spectrum of regret scalings is possible

$\beta$ -clustered distributions

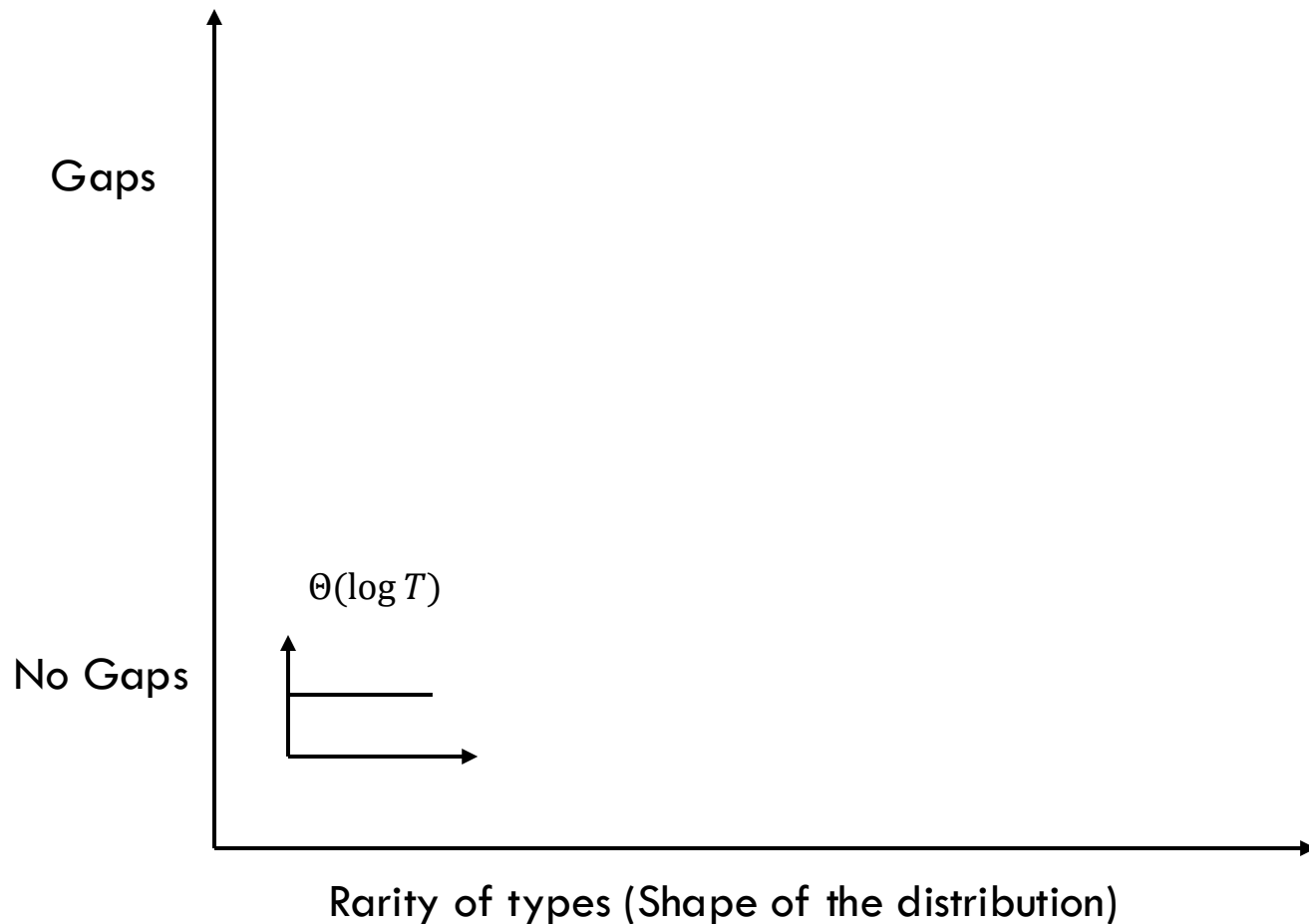
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# Complete Characterization of the Multi-secretary Problem

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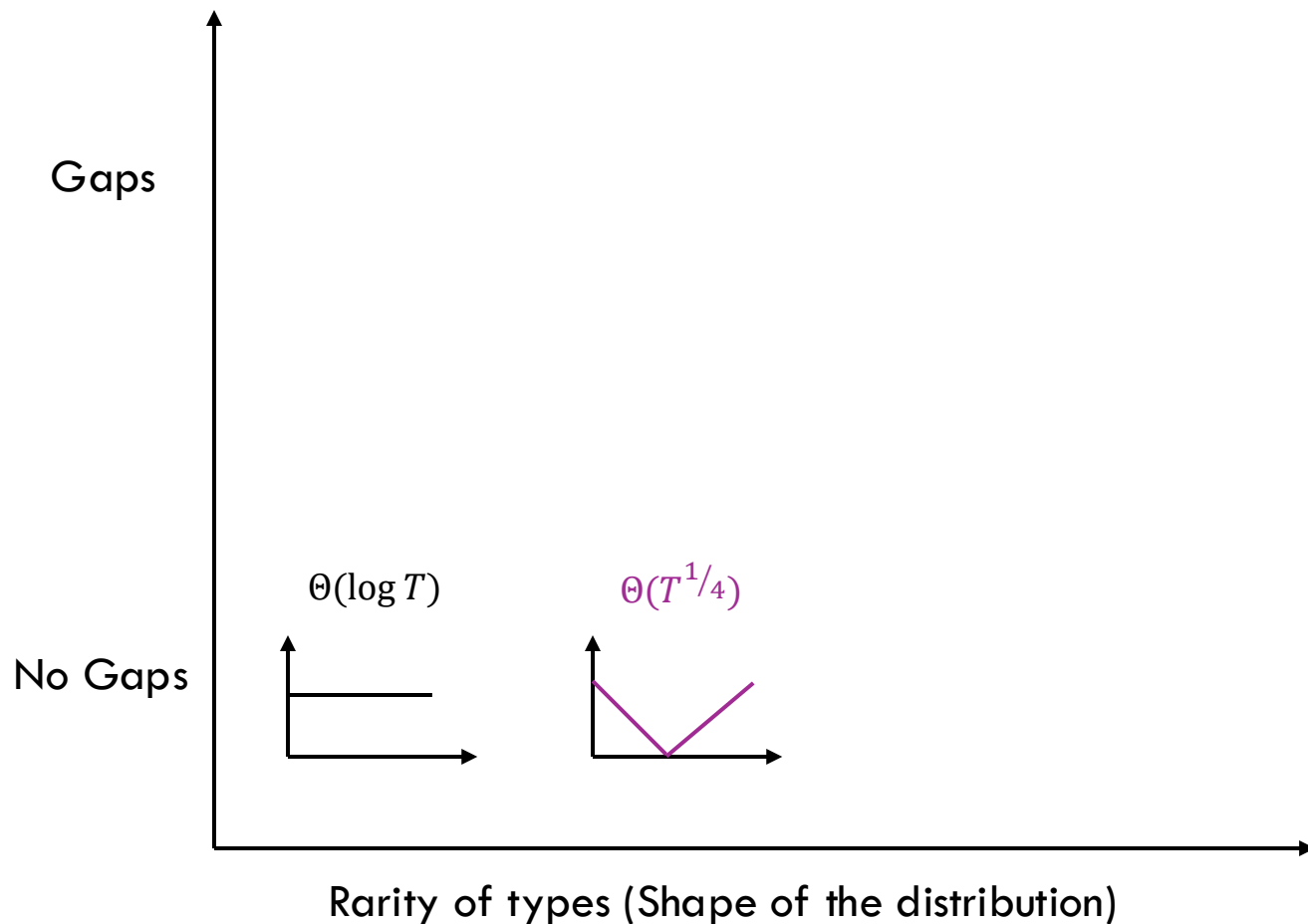


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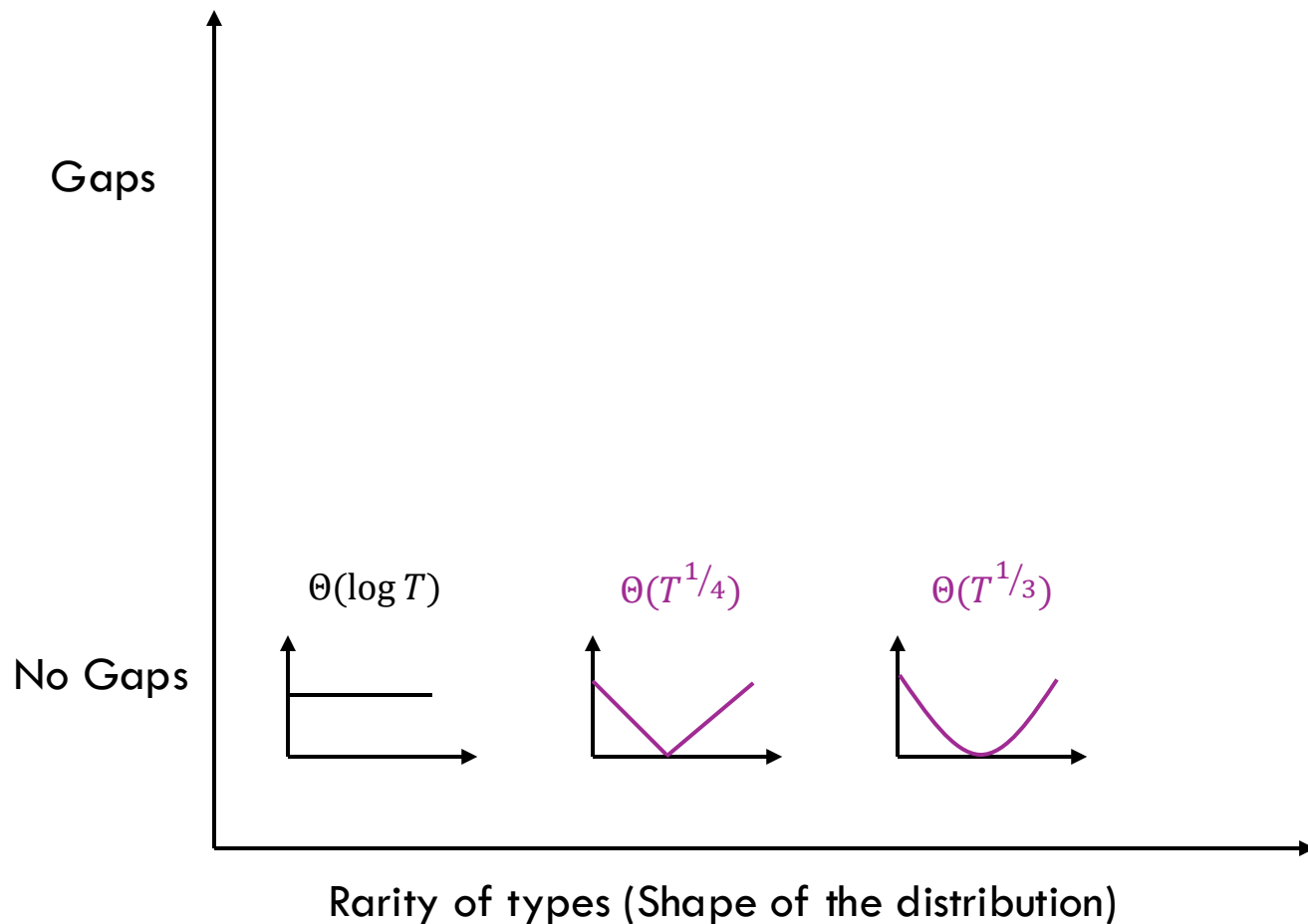




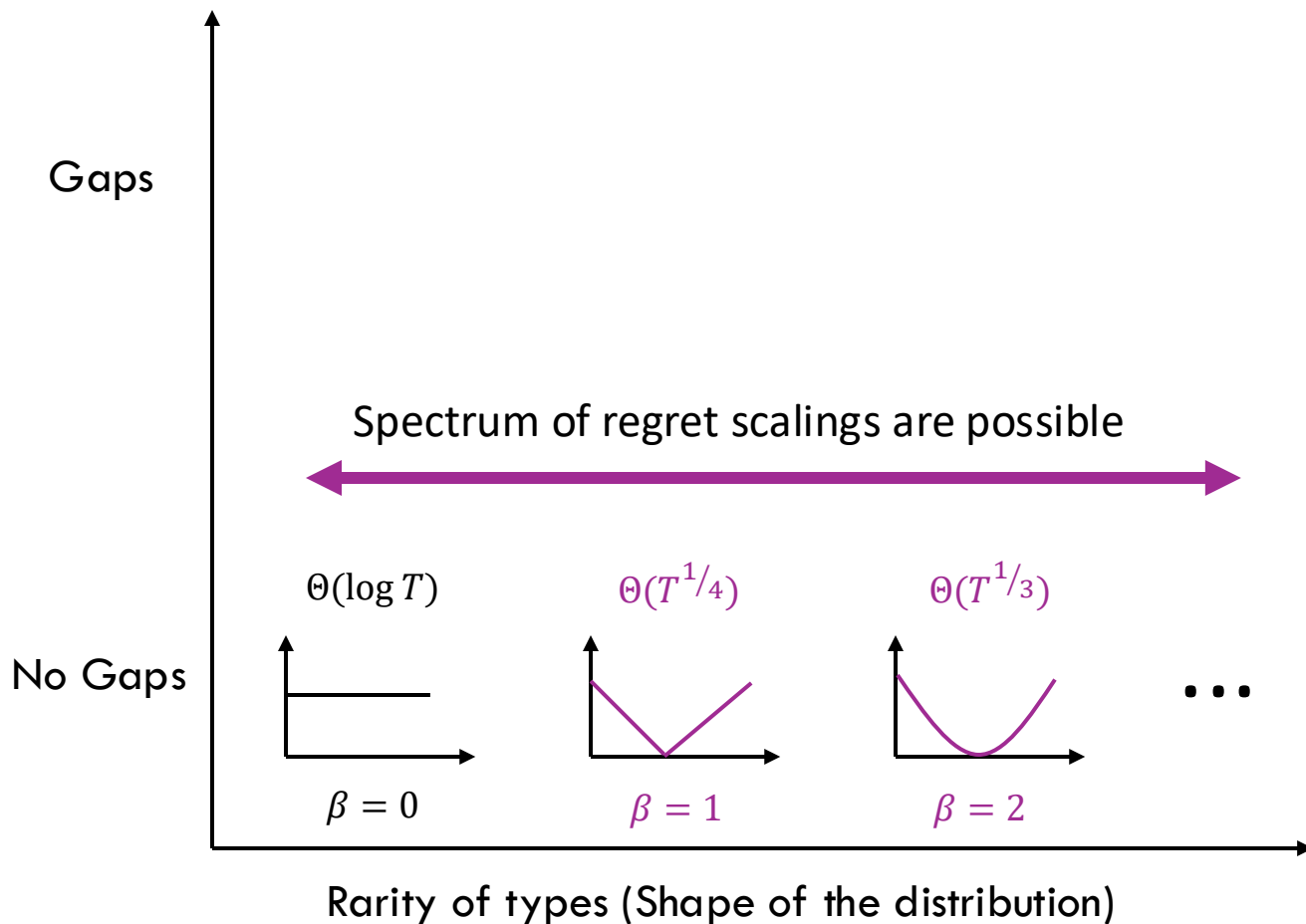
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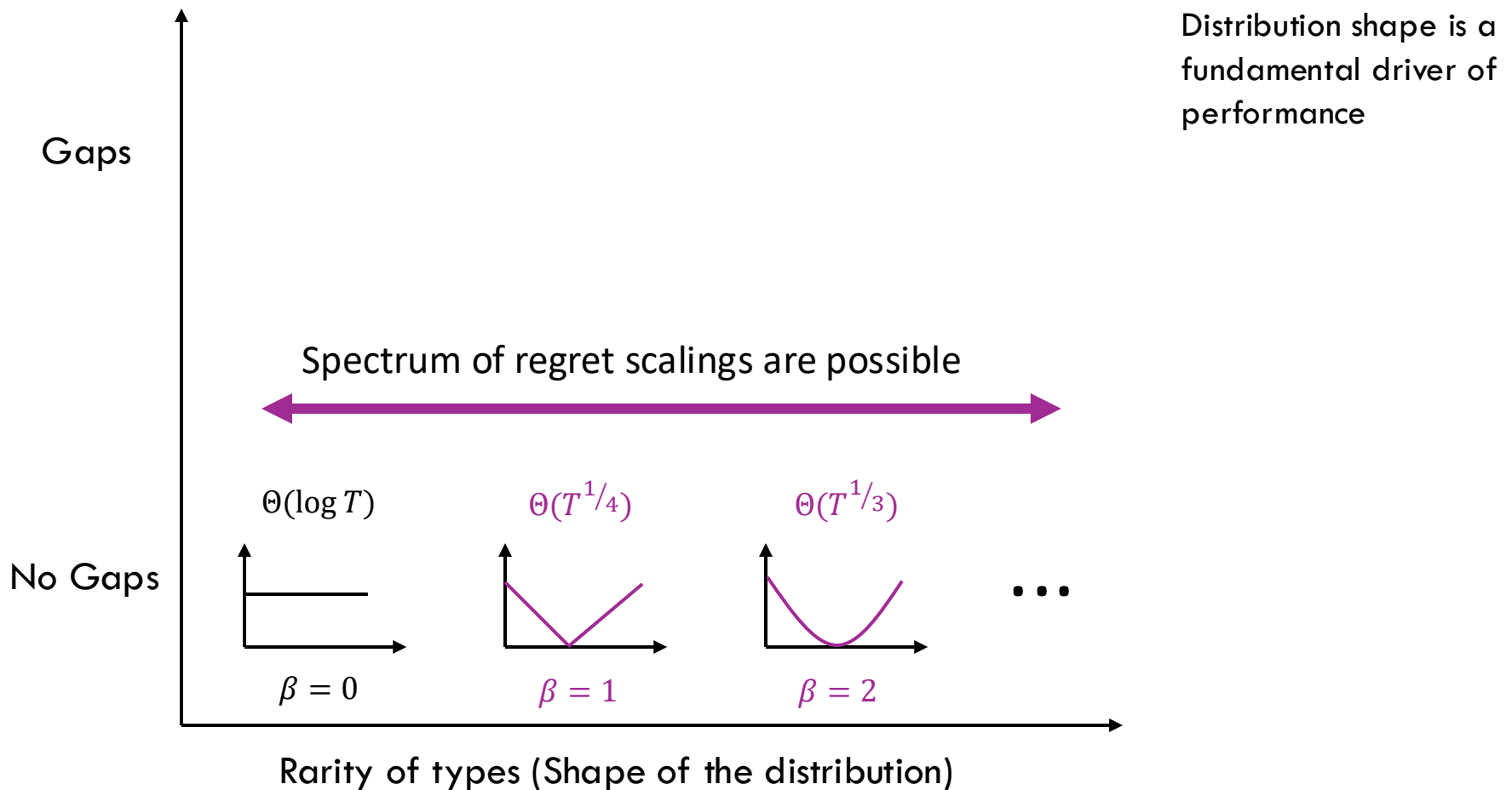
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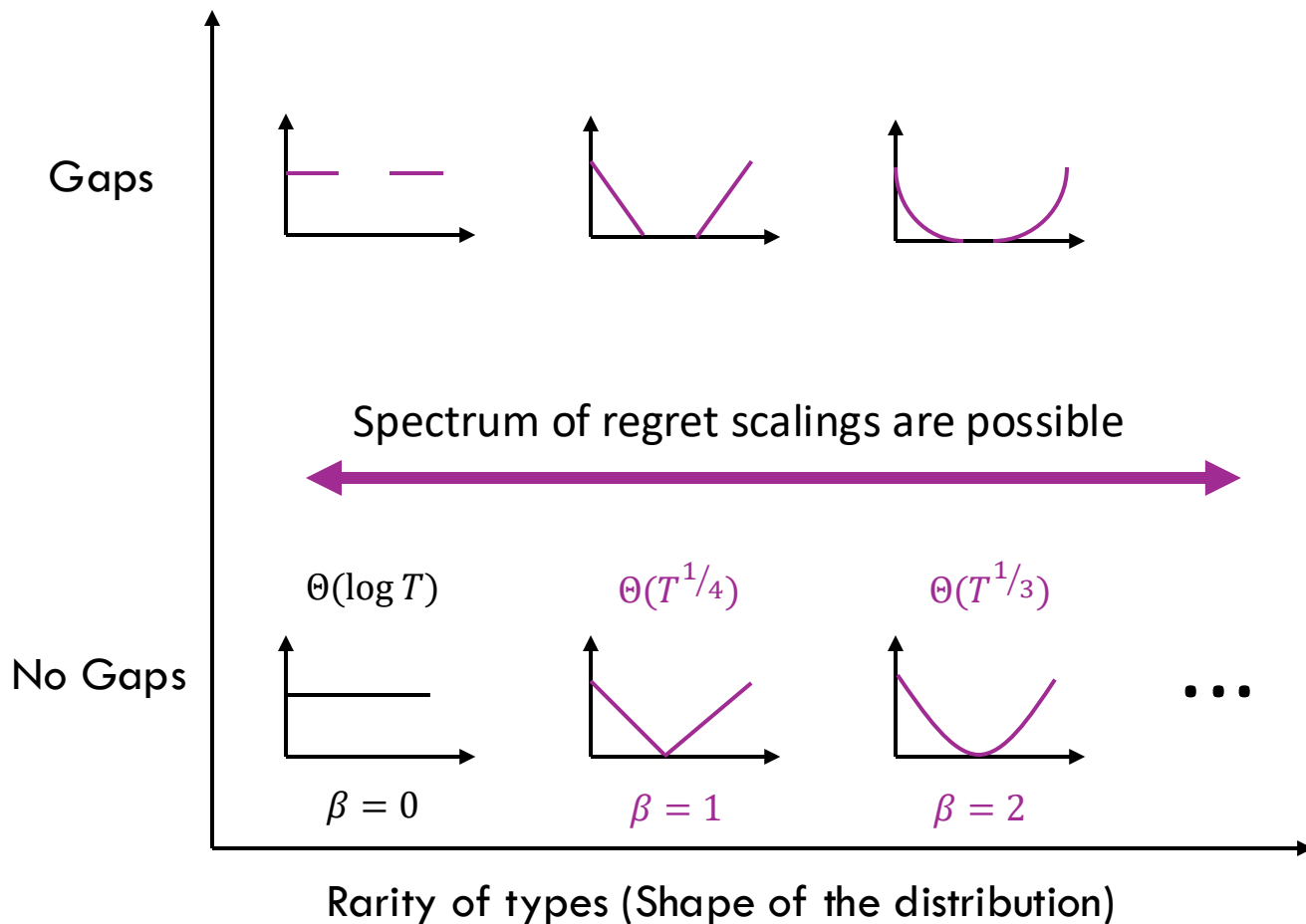
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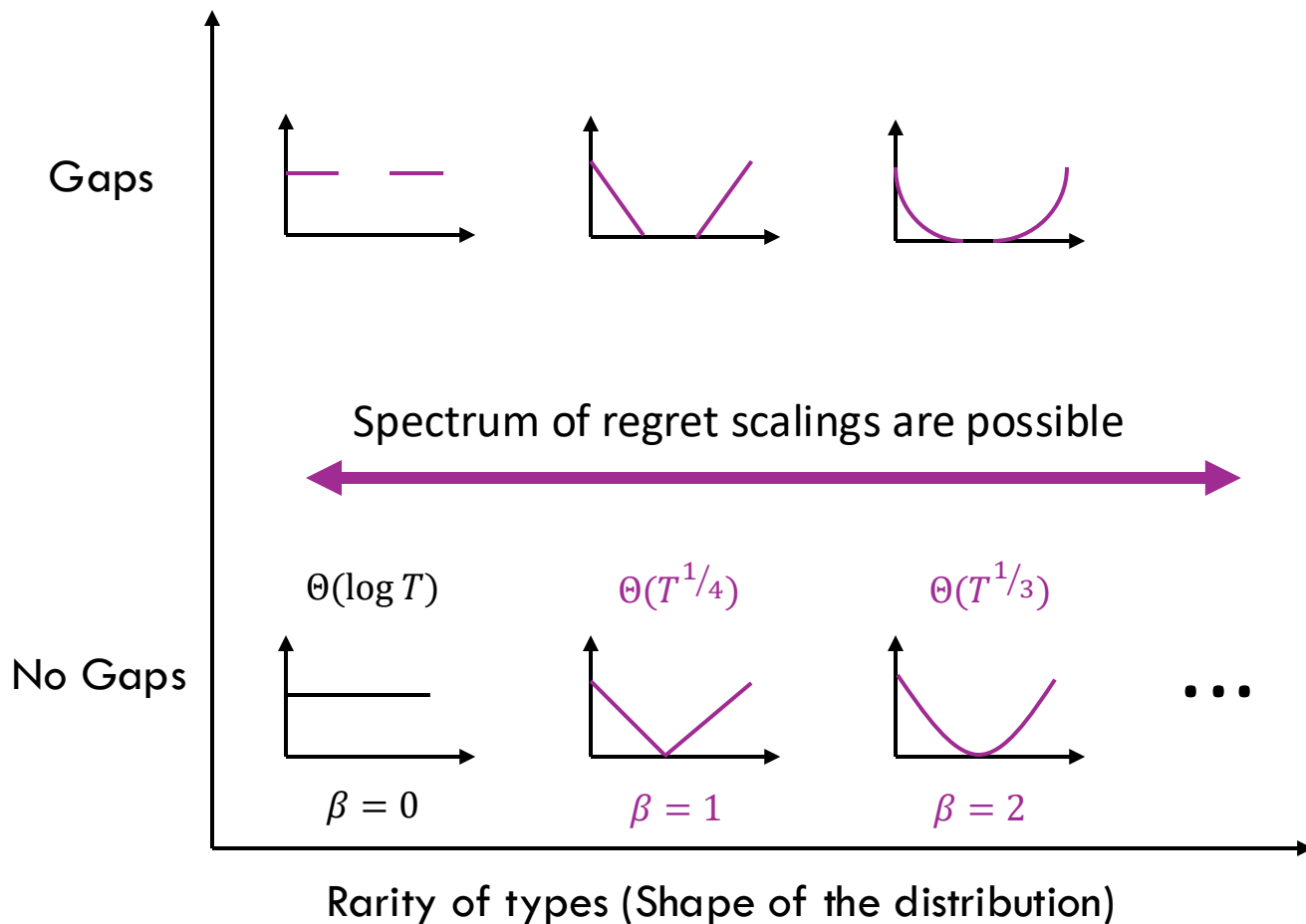
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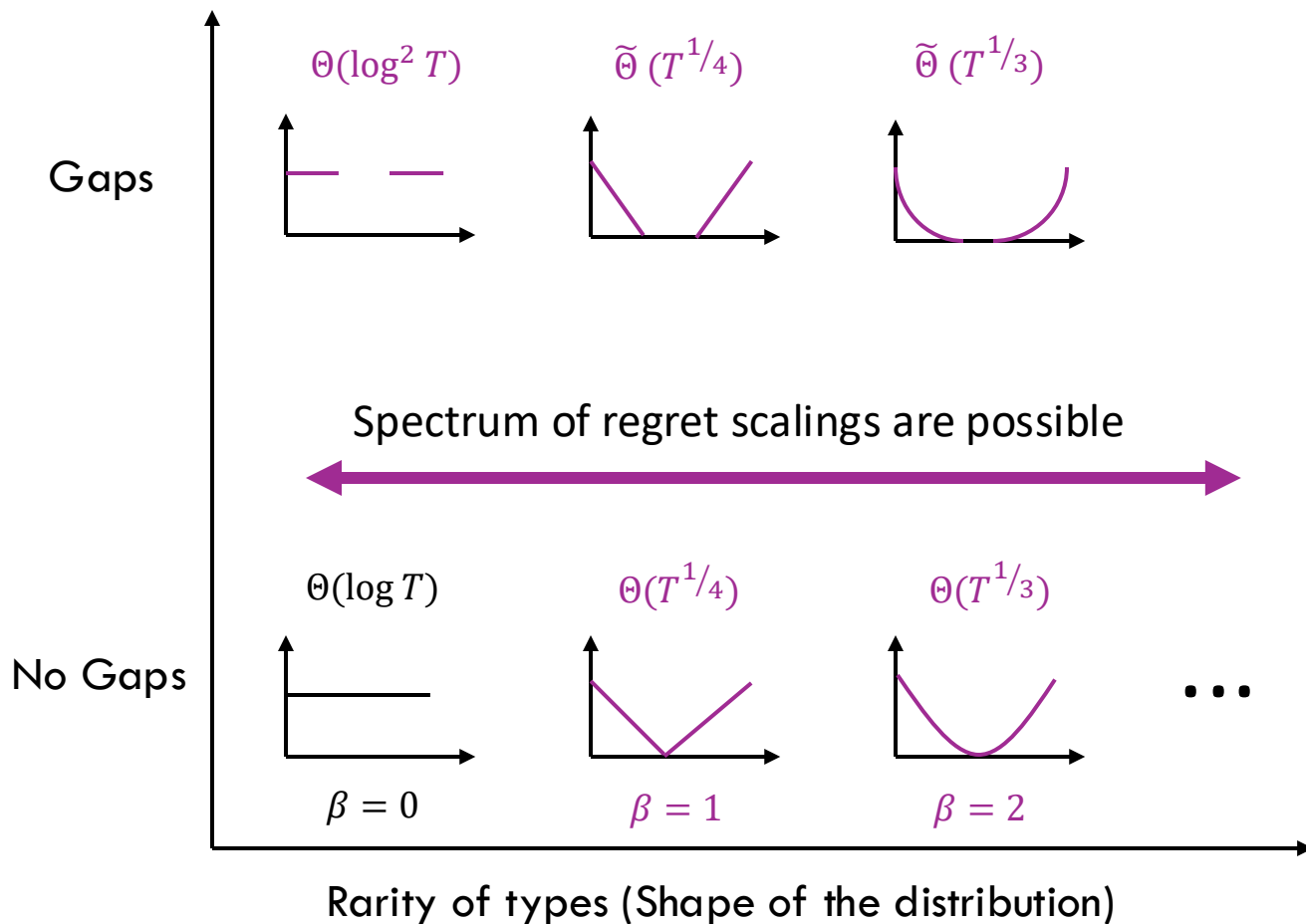
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Distribution shape is a fundamental driver of performance

Dealing with gaps in an algorithmic challenge

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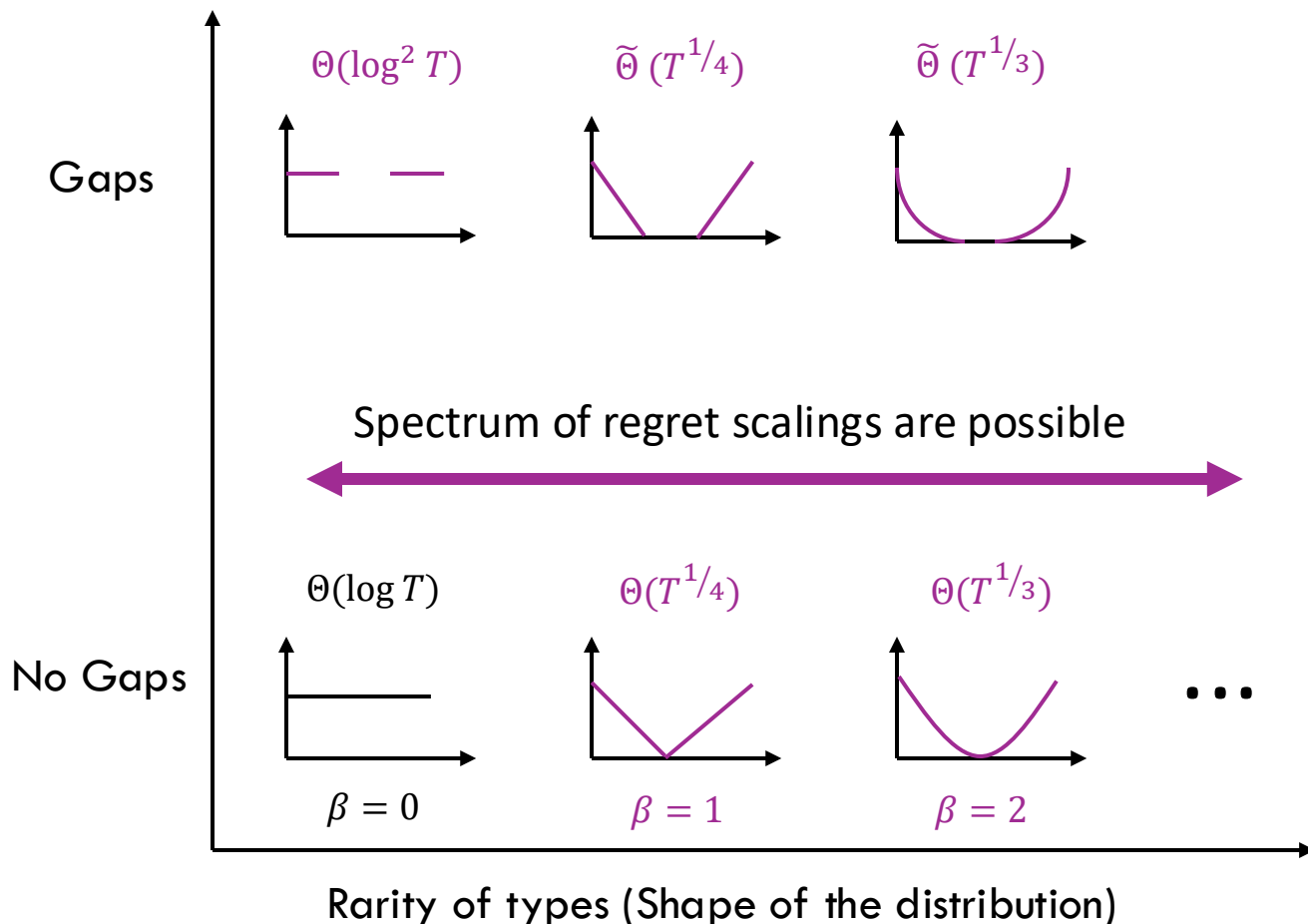


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Conservativeness with respect to gaps (CwG) principle enables near-optimal performance

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Use RAMS to operationalize CwG

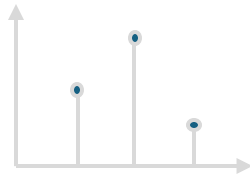


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## Repeatedly Act using Multiple Simulations (RAMS)

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present  
Logarithmic Regret



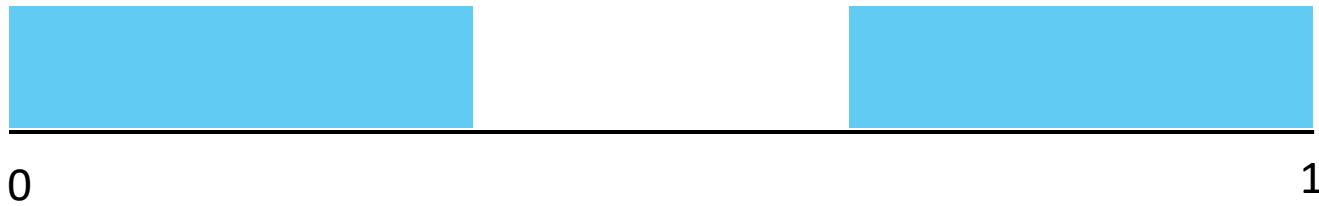
Entire spectrum of regret scalings is possible

$\beta$ -clustered distributions

Regret is the additive gap b/w the value of hindsight opt. and value under some algorithm

Gaps pose an algorithmic challenge

# Gaps pose an algorithmic challenge

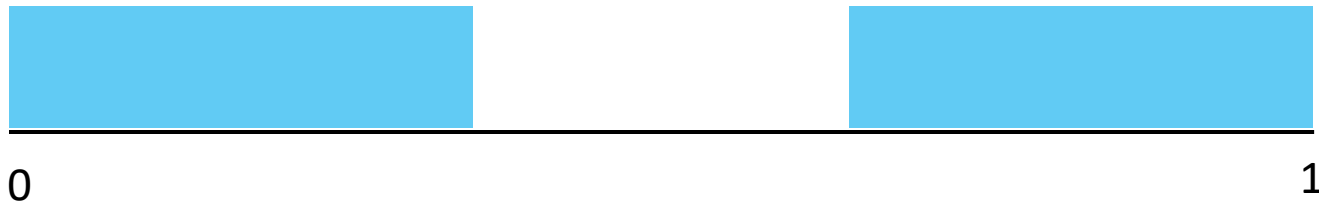


# Gaps pose an algorithmic challenge



Certainty Equivalent Control computes the budget ratio

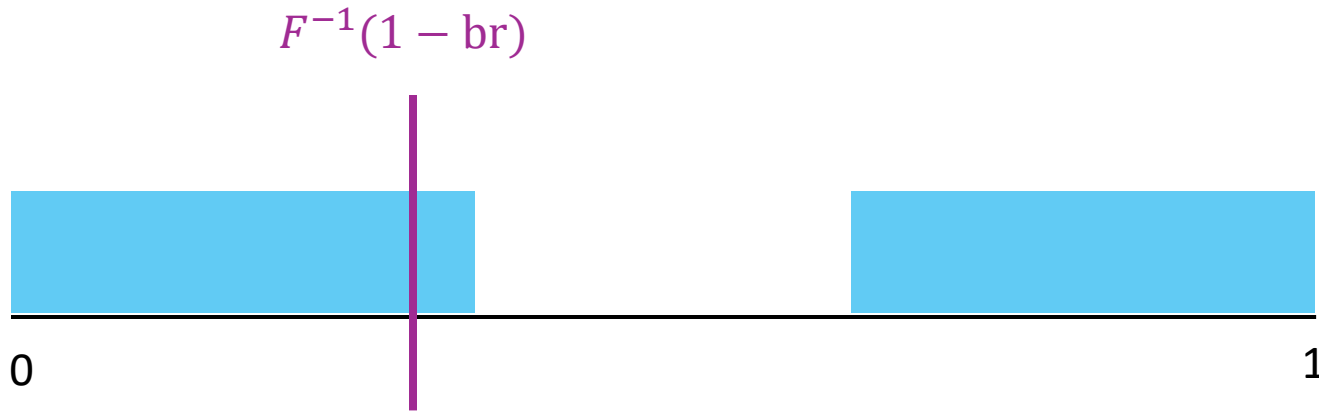
# Gaps pose an algorithmic challenge



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$$br = \text{Budget Ratio} = (\text{Remaining Budget}) / (\text{Remaining Time})$$

# Gaps pose an algorithmic challenge

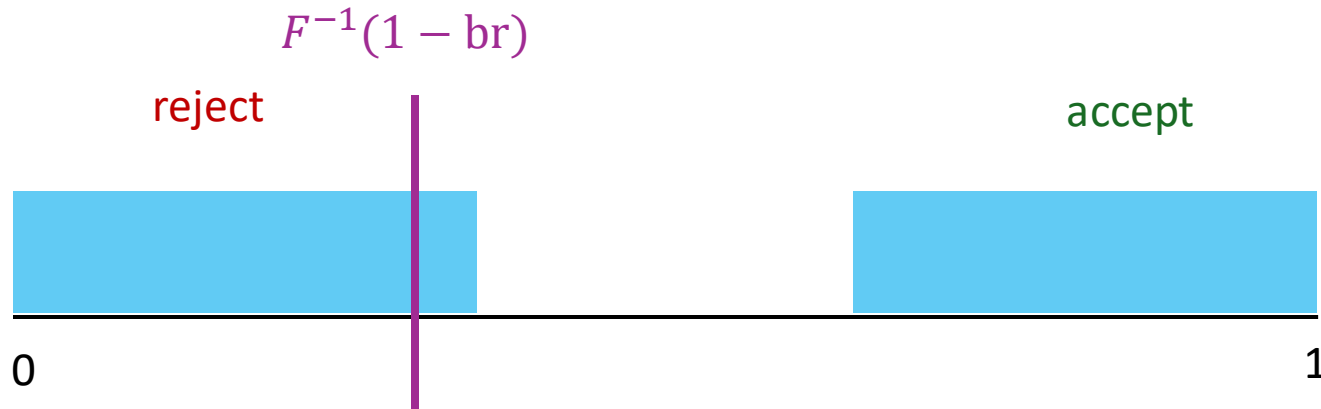


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Accept if the request type value is more than  $F^{-1}(1 - br)$ , else reject the request

# Gaps pose an algorithmic challenge

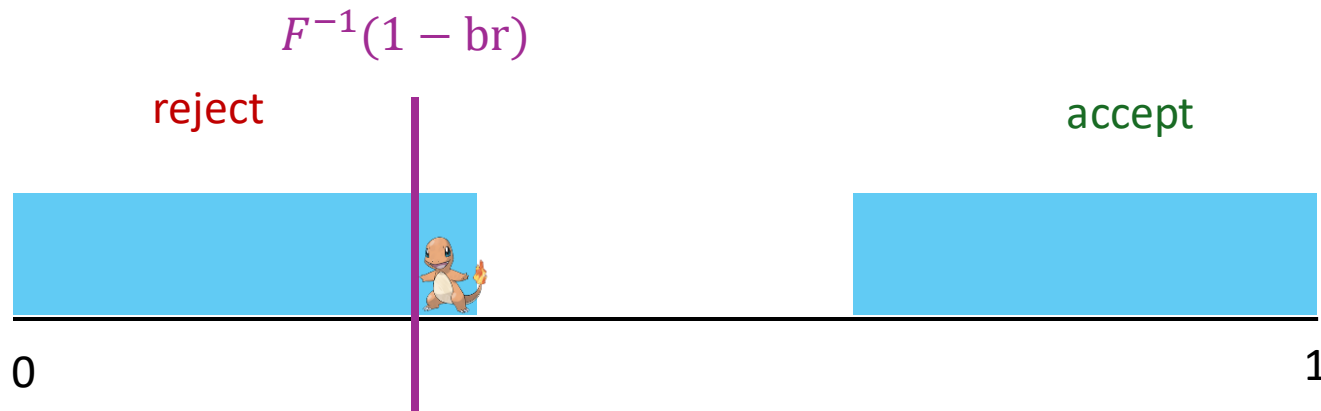


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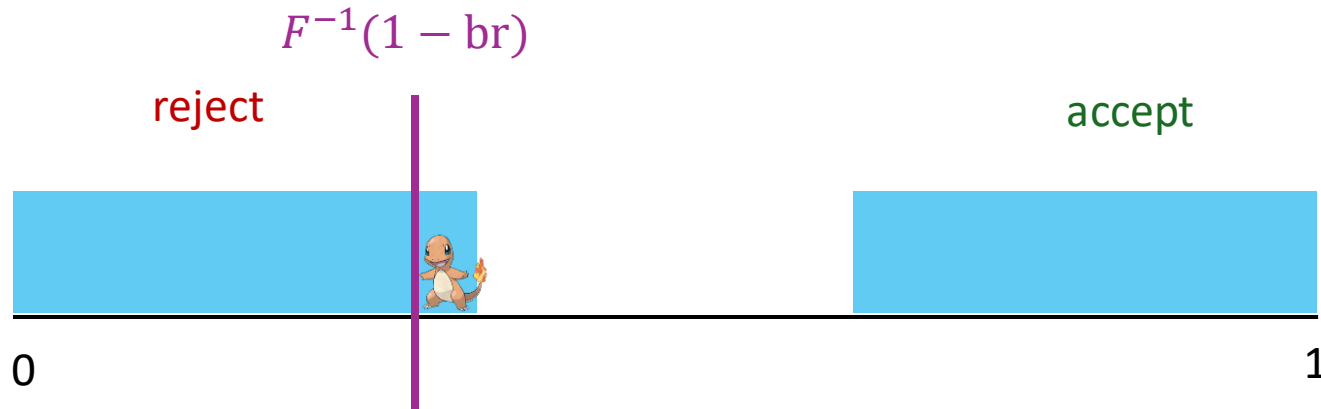
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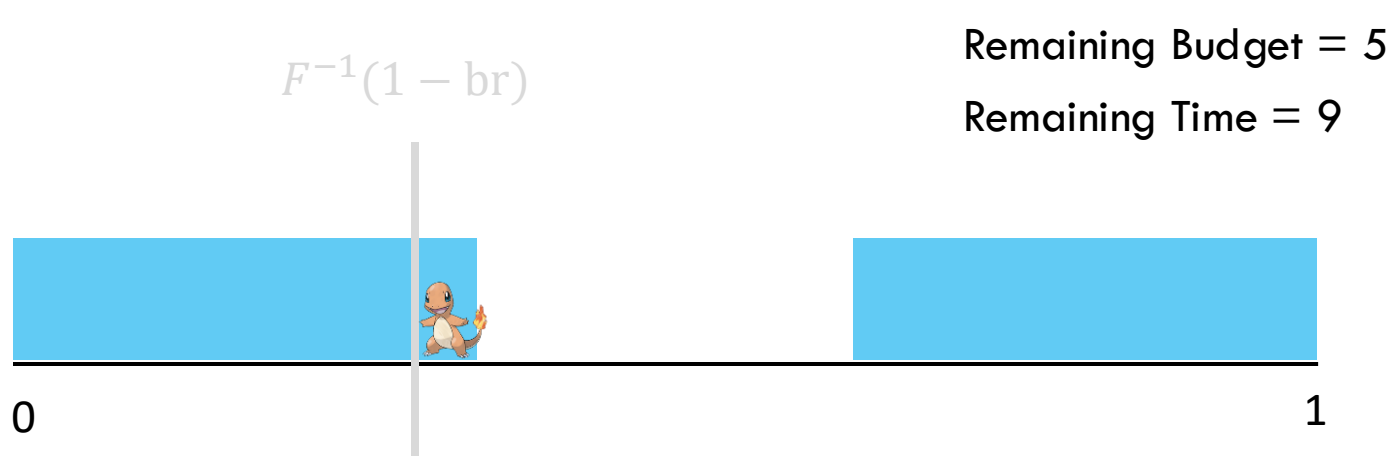
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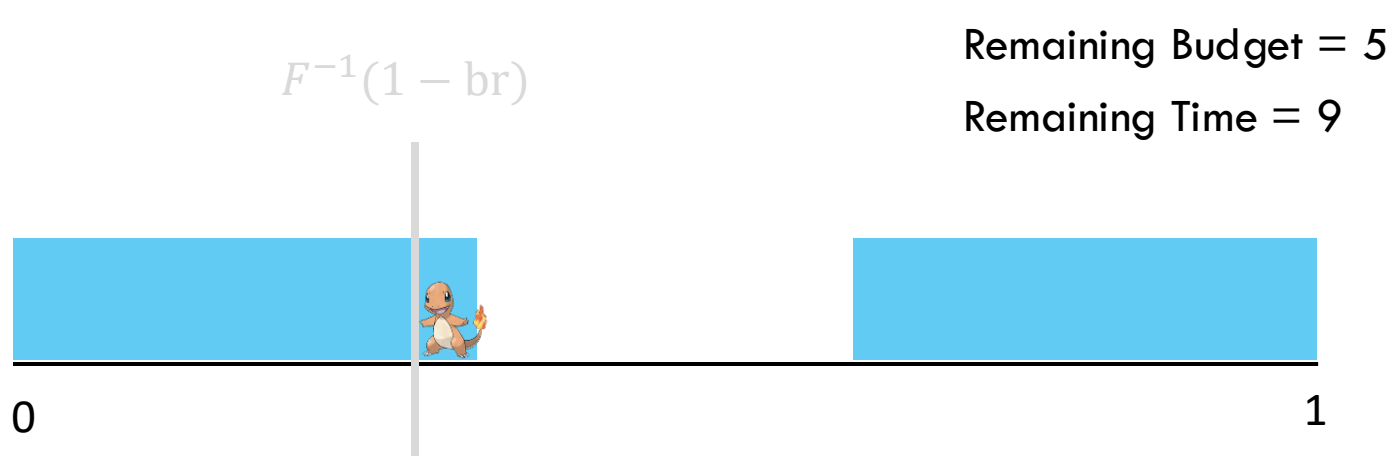
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$\text{Regret(CE)} = \Omega(\sqrt{T})$  (highly sub-optimal regret scaling)

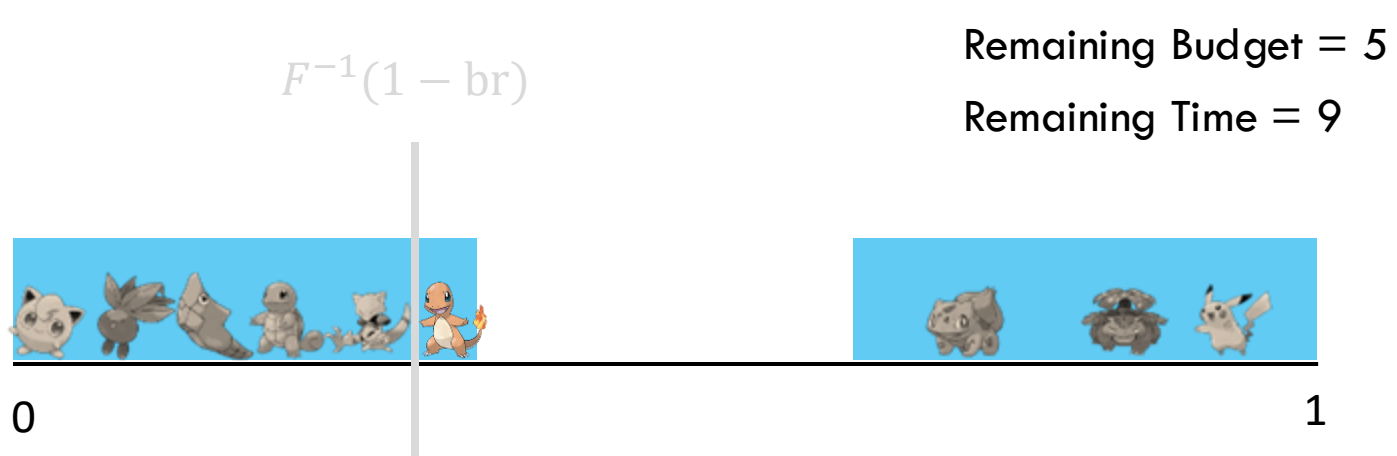
# CwG via multiple simulations



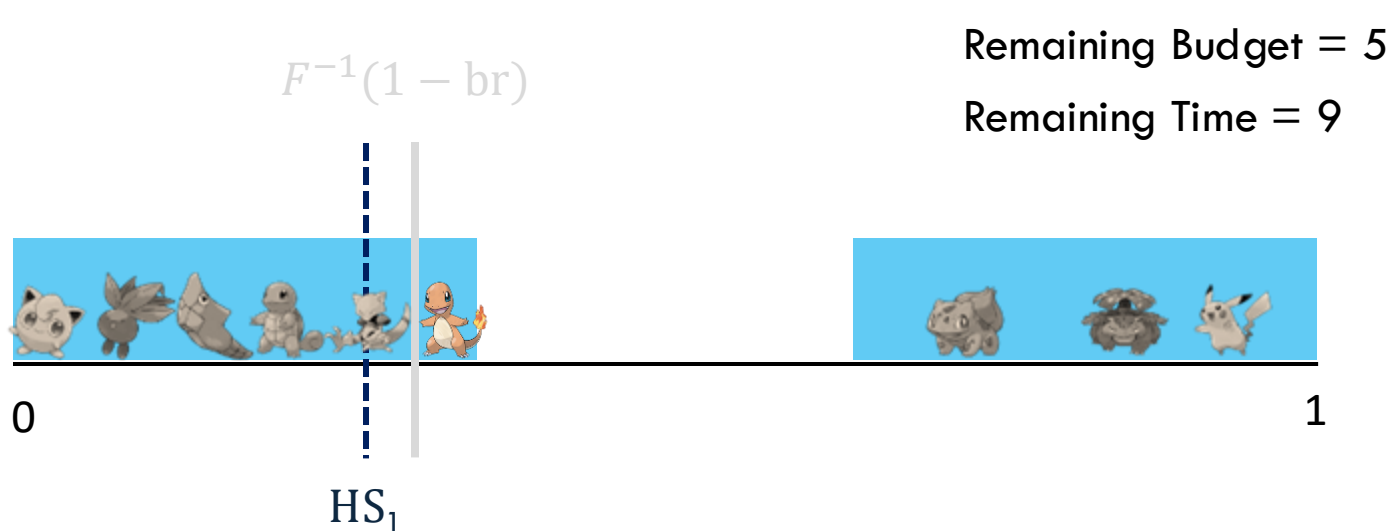
# CwG via multiple simulations



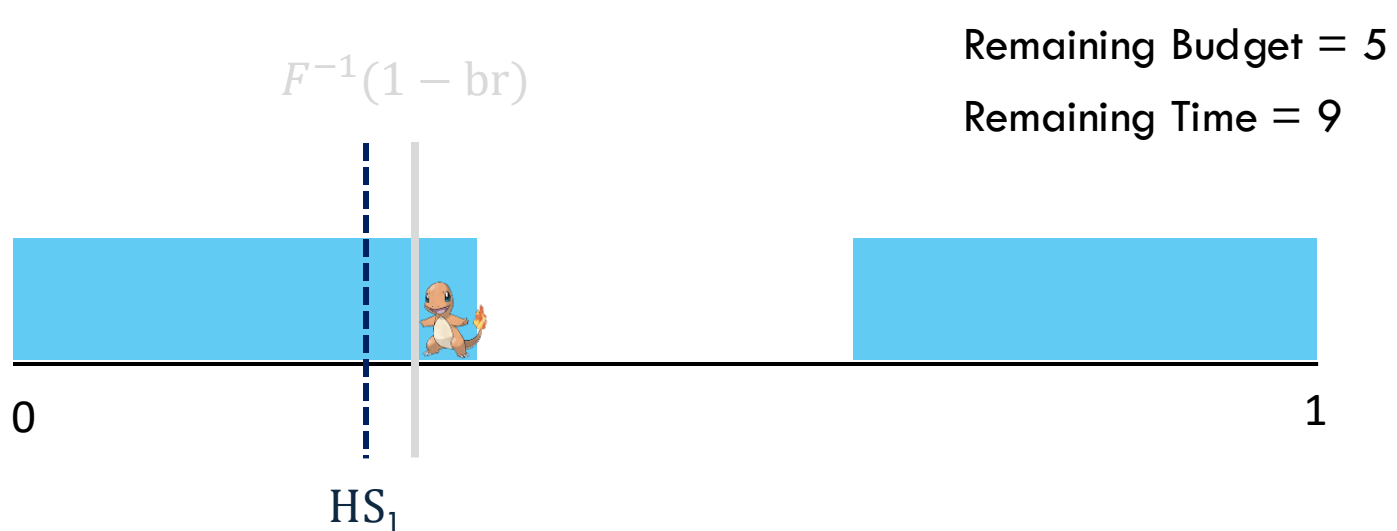
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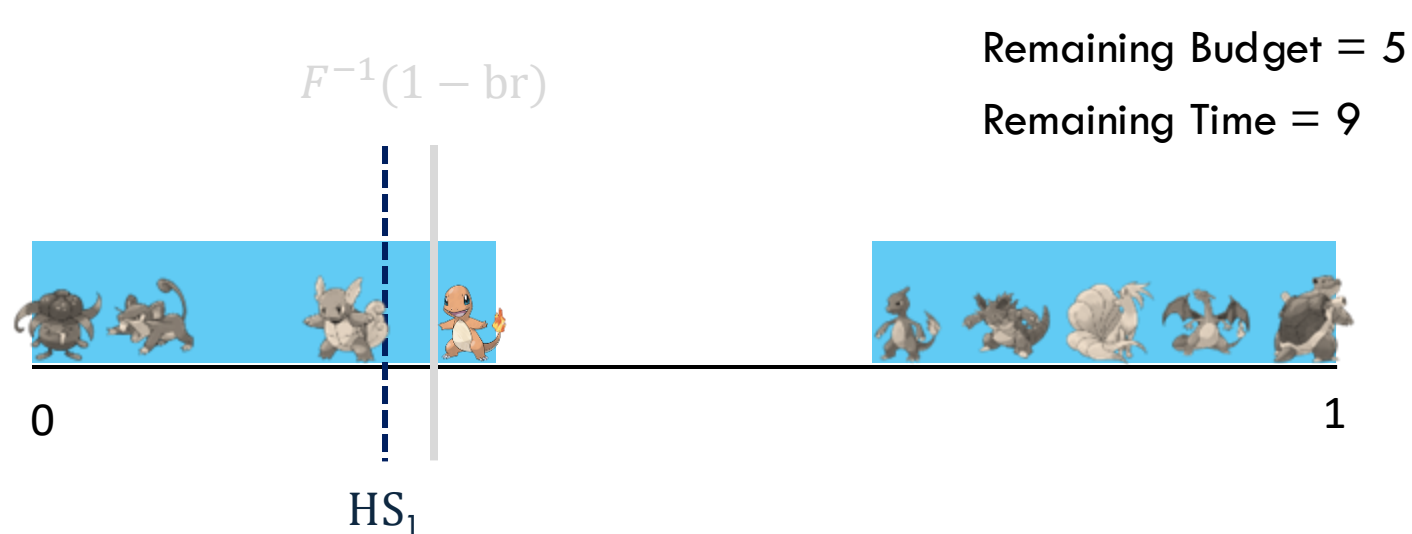
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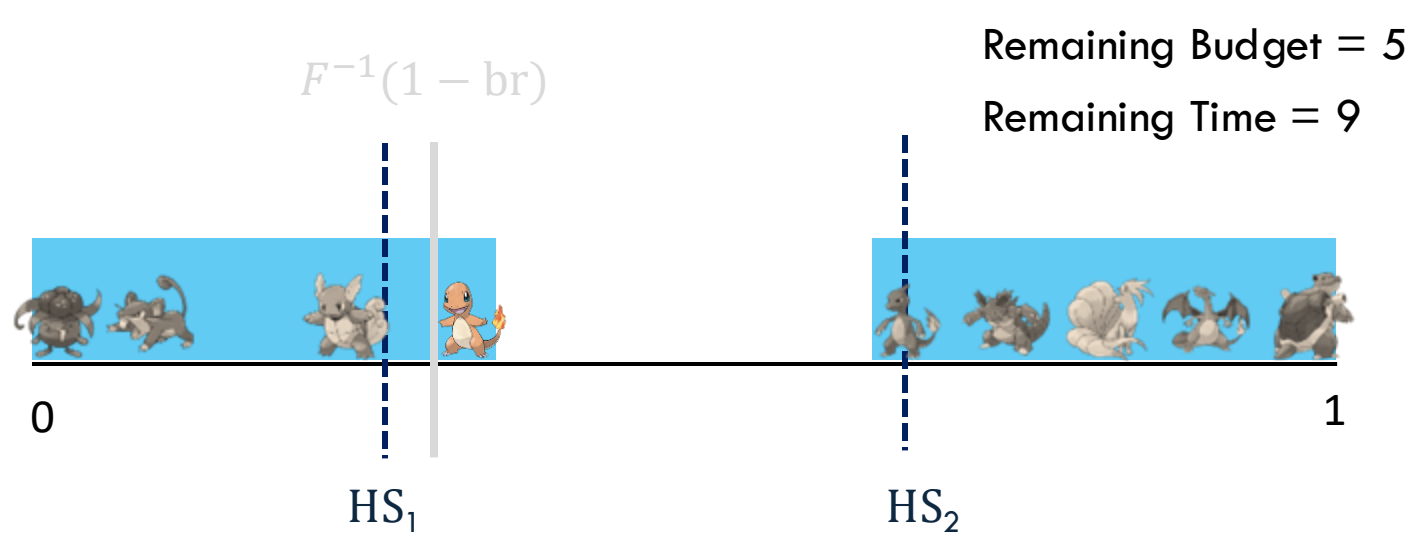
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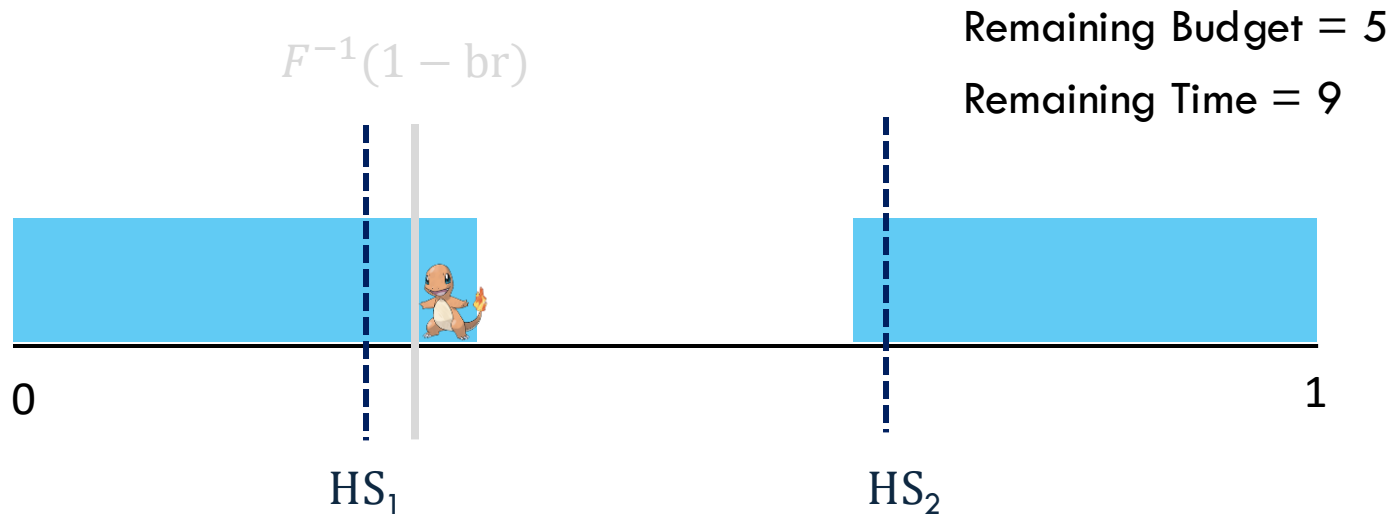


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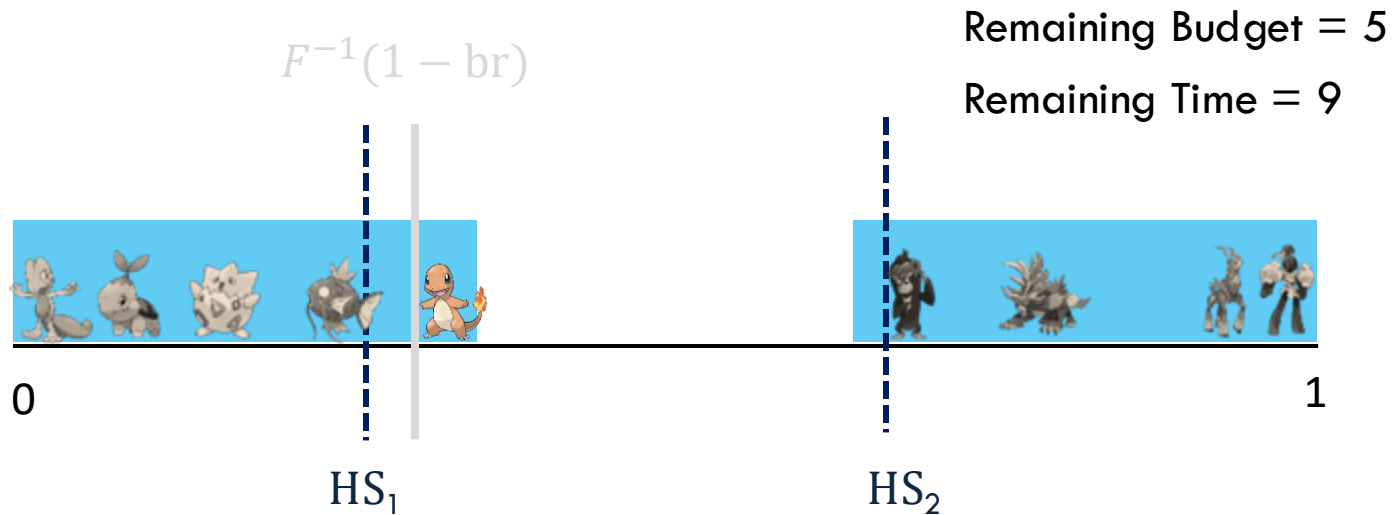




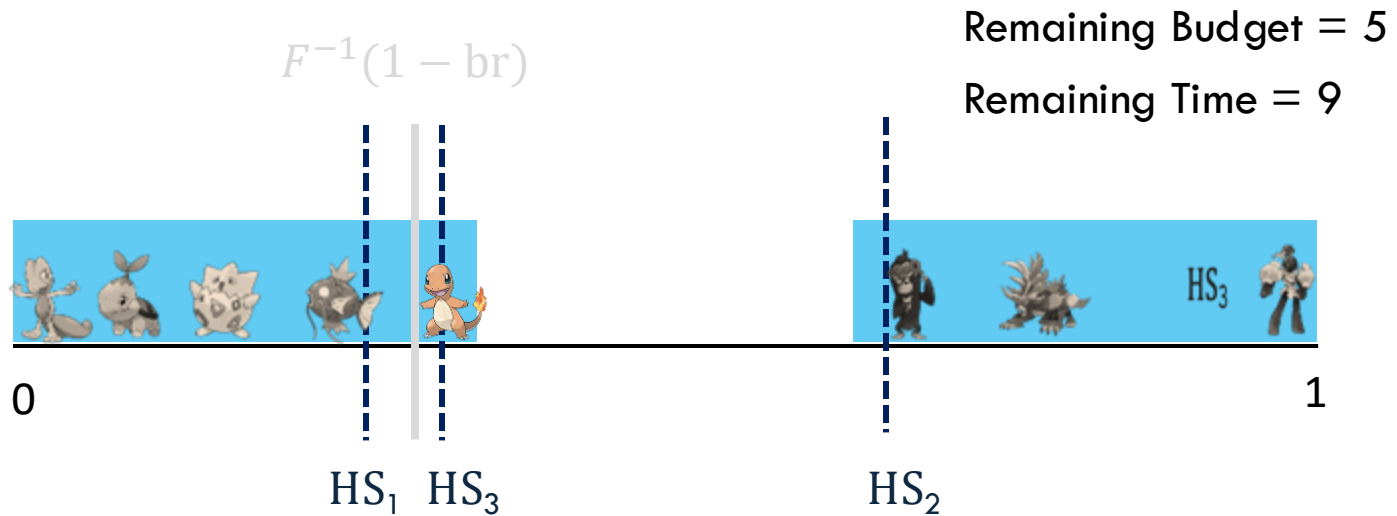
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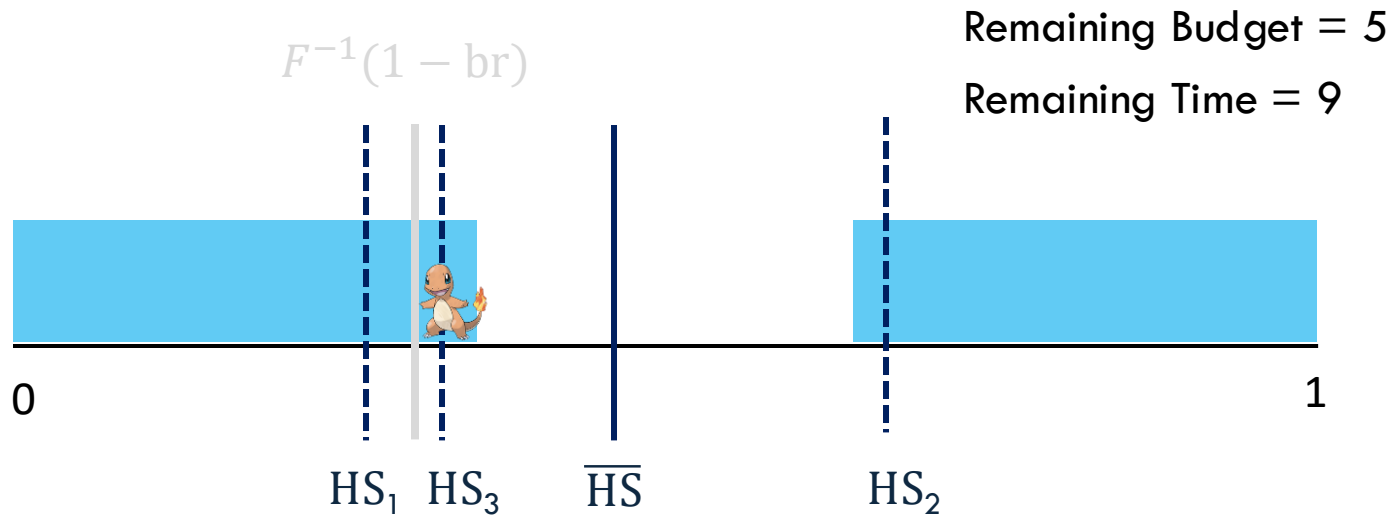
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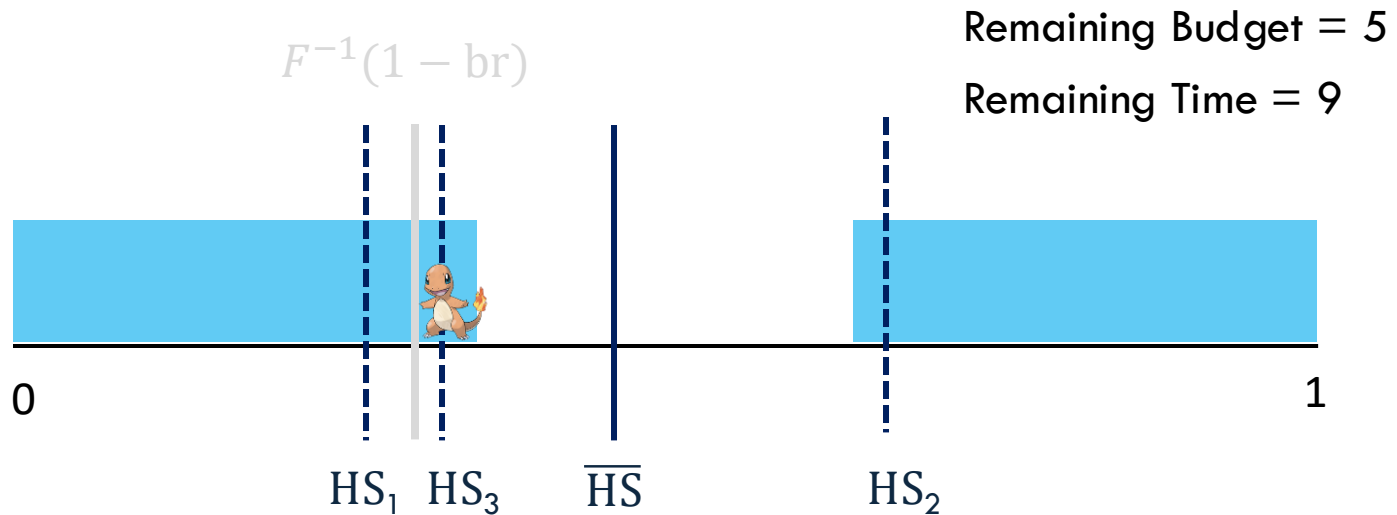
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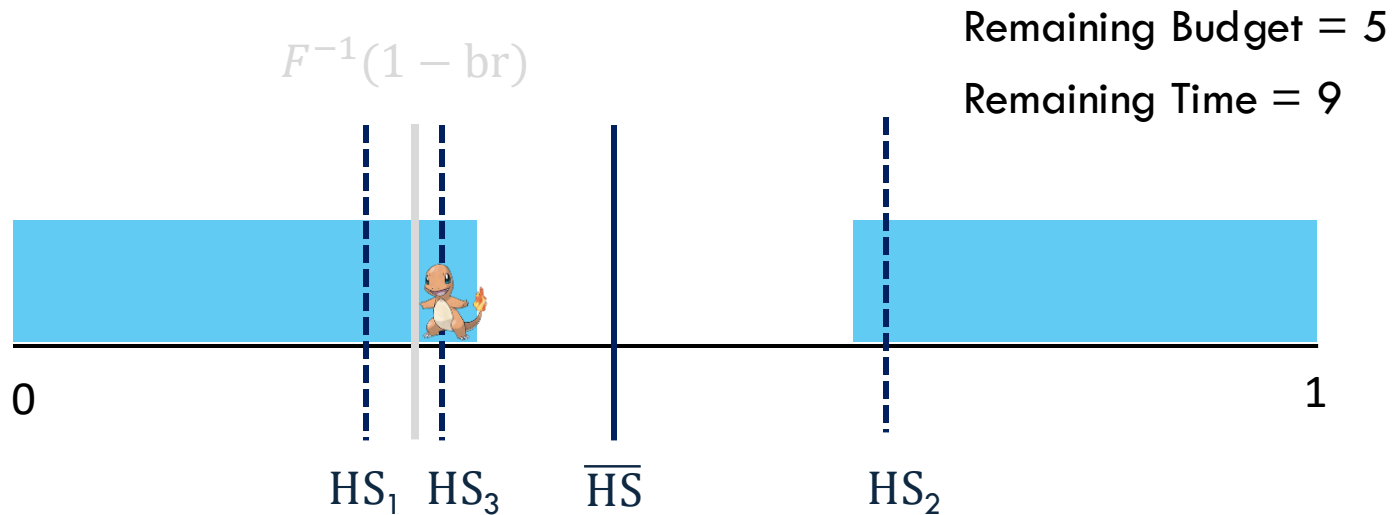
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Accept if the request type value is more than  $\overline{HS}$ , else reject the request



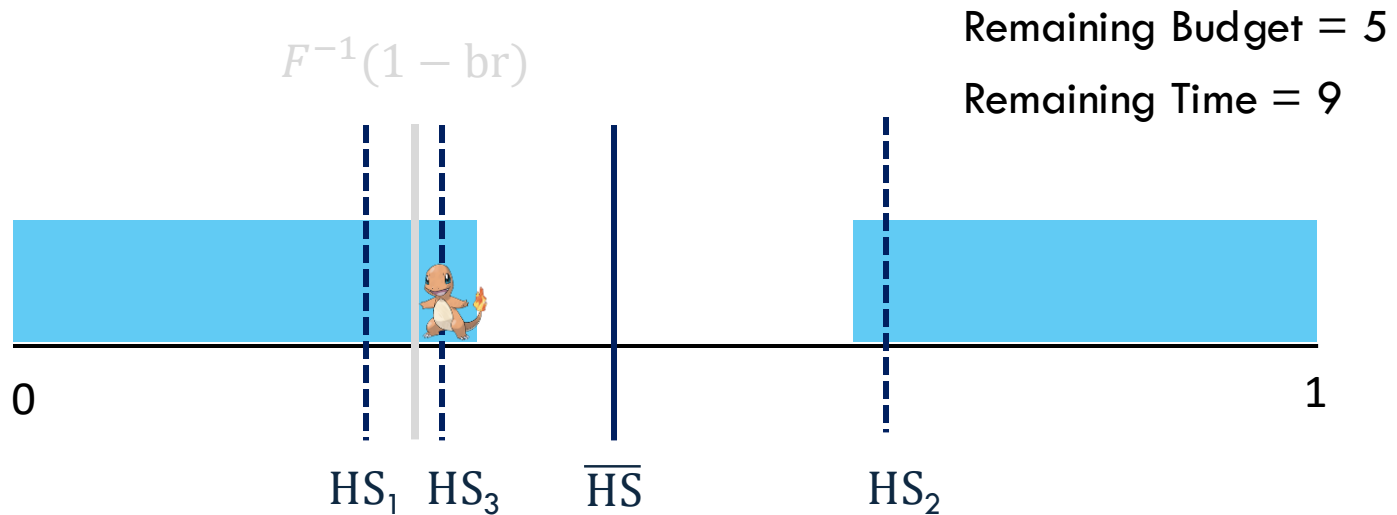
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$$\text{Regret(RAMS)} = O(\log^2 T)$$

# CwG via multiple simulations

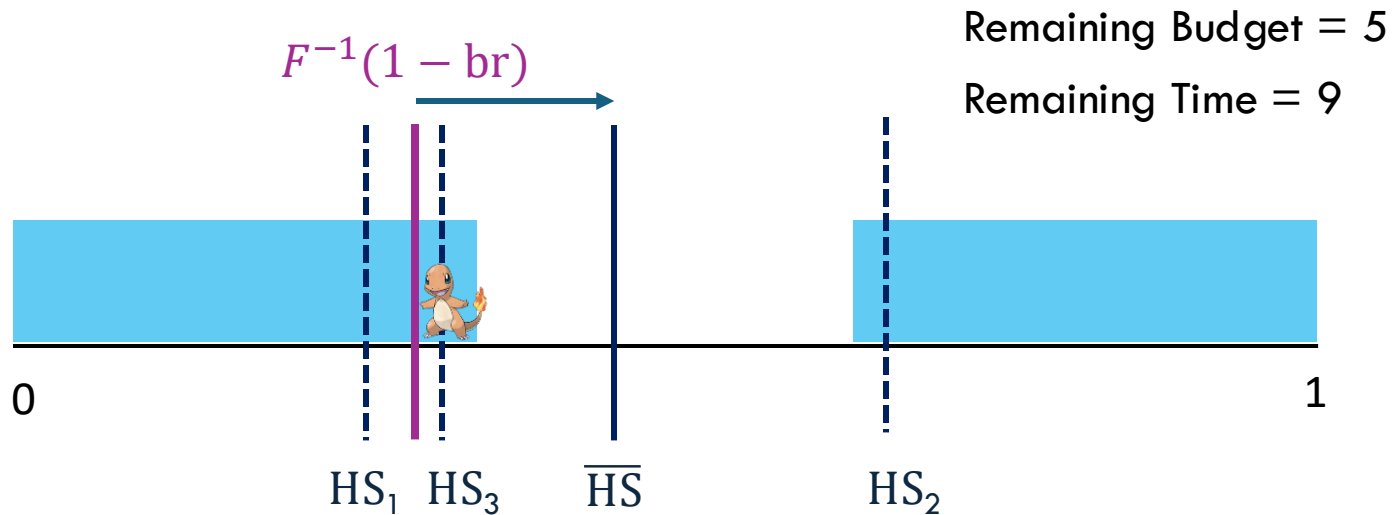


Accept if the request type value is more than  $\overline{HS}$ , else reject the request

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$$\text{Regret(CE)} = \Omega(\sqrt{T}) \text{ (highly sub-optimal regret scaling)}$$

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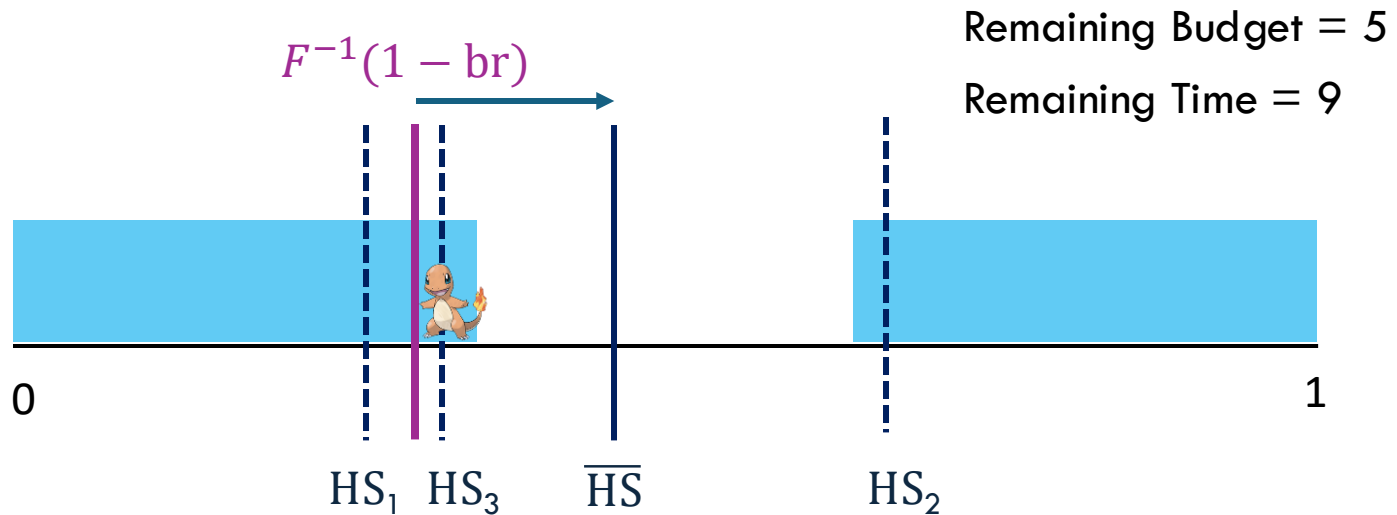
Accept if the request type value is more than  $\overline{HS}$ , else reject the request

## Conservativeness with respect to Gaps Principle

If the CE threshold  $F^{-1}(1 - br)$  is close to a gap, use the gap as the threshold. Otherwise continue using the CE threshold.



# CwG via multiple simulations



Accept if the request type value is more than  $\overline{HS}$ , else reject the request

## Connections to “Dual Averaging”

The different HS thresholds are the shadow prices of the budget for different scenarios, the bid price is computed by averaging the HS thresholds

Network  
Revenue  
Management

Repeated  
Auctions with  
Budgets



Multi-secretary



Order  
Fulfillment

Dynamic  
Matching

**Repeatedly Act using Multiple Simulations (RAMS)**

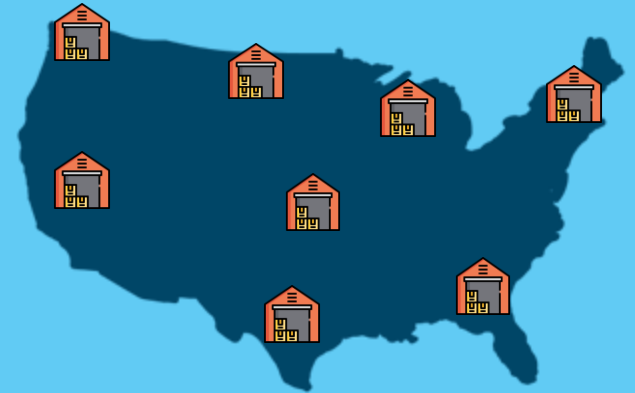
# RAMS beyond multi-secretary

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State (Budget)  $B_t$  and feasible set of actions  $A_t$

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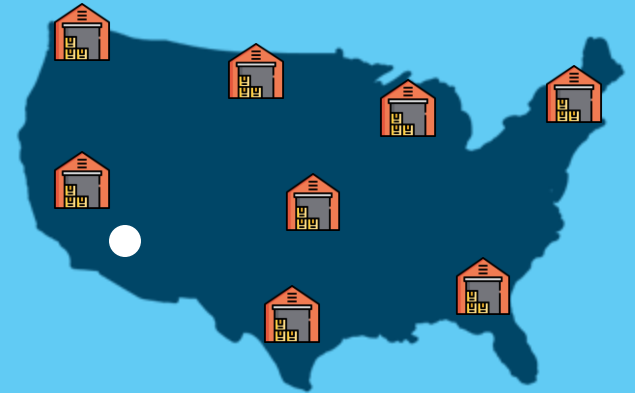
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# RAMS beyond multi-secretary

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Request  $\theta_t = (r_t, c_t)$  arrives at time  $t$

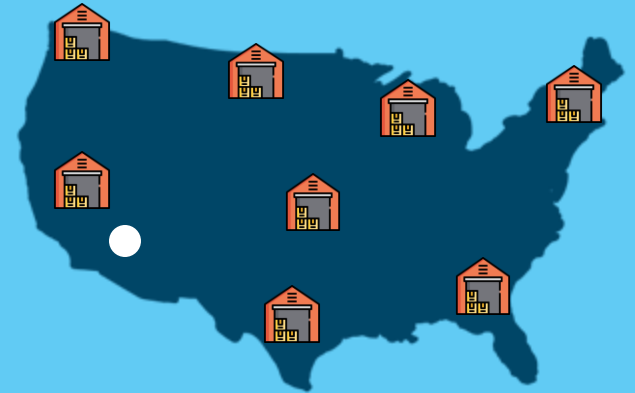


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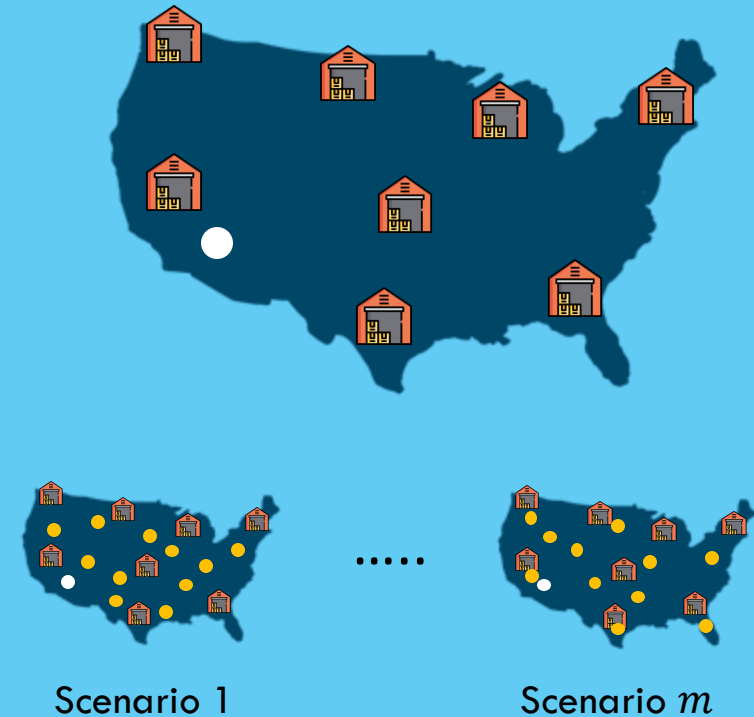
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$\theta_t, \theta_{t+1}^{(1)}, \theta_{t+2}^{(1)}, \dots, \theta_T^{(1)}$  Scenario 1

$\theta_t, \theta_{t+1}^{(2)}, \theta_{t+2}^{(2)}, \dots, \theta_T^{(2)}$  Scenario 2

$\vdots$

$\theta_t, \theta_{t+1}^{(m)}, \theta_{t+2}^{(m)}, \dots, \theta_T^{(m)}$  Scenario  $m$





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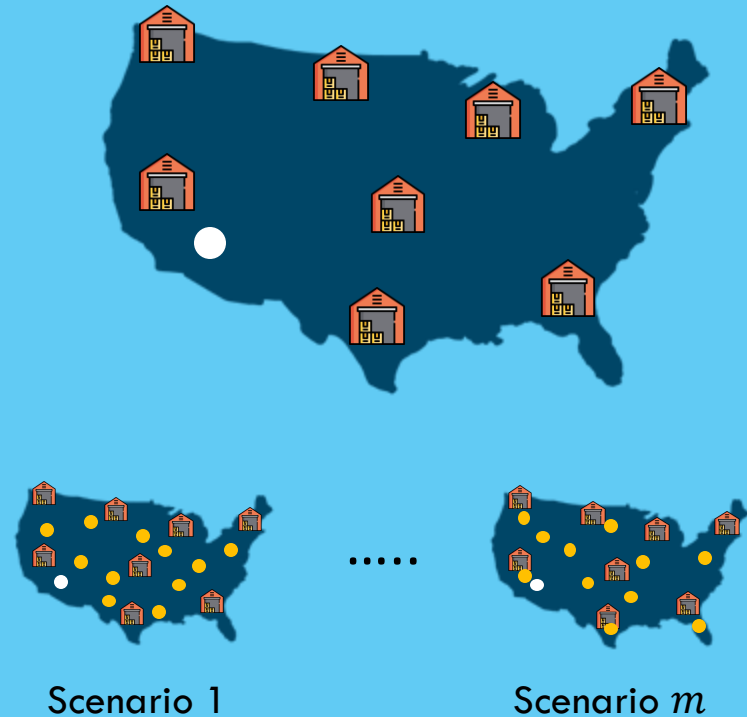
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For each scenario  $k$ , compute the **compensation**  
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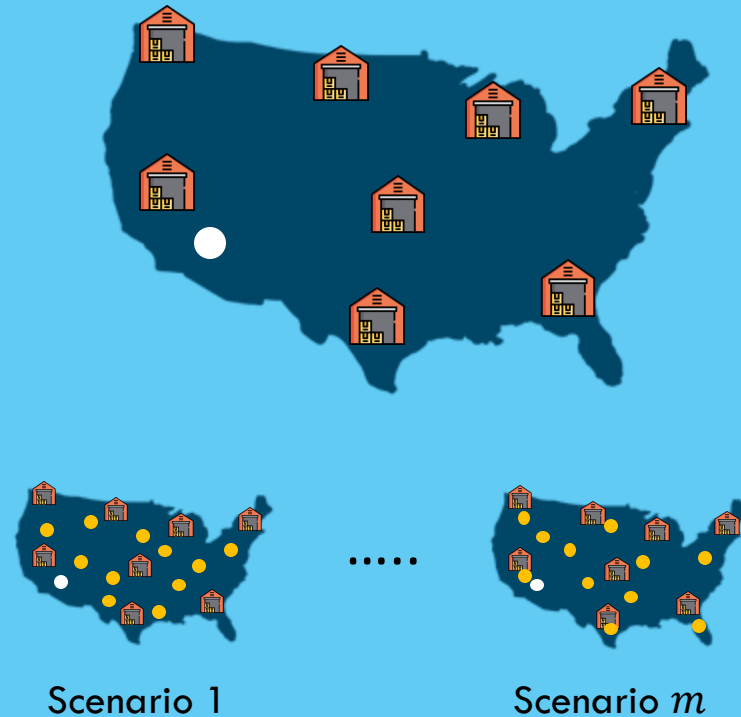
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**Compensation(scenario  $k$ ,  $a$ ) =**  
(Max total reward in scenario  $k$ ) –  
(Max total reward in scenario  $k$  if action  
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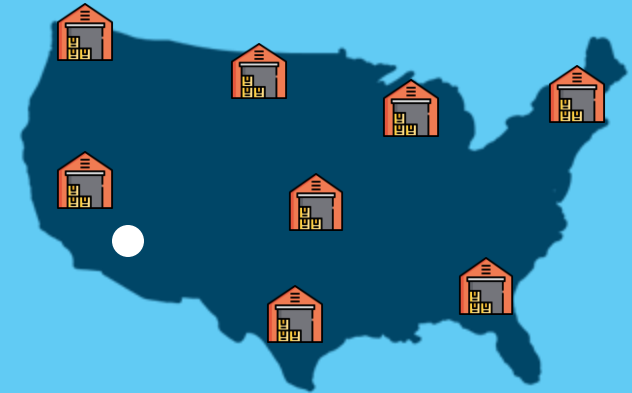
$\theta_t, \theta_{t+1}^{(2)}, \theta_{t+2}^{(2)}, \dots, \theta_T^{(2)}$  Scenario 2

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For each scenario  $k$ , compute the **compensation**  
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Take the **action** with the **minimum compensation**  
averaged over  $m$  scenarios



Scenario 1



Scenario  $m$

**Compensation(scenario  $k$ ,  $a$ ) =**  
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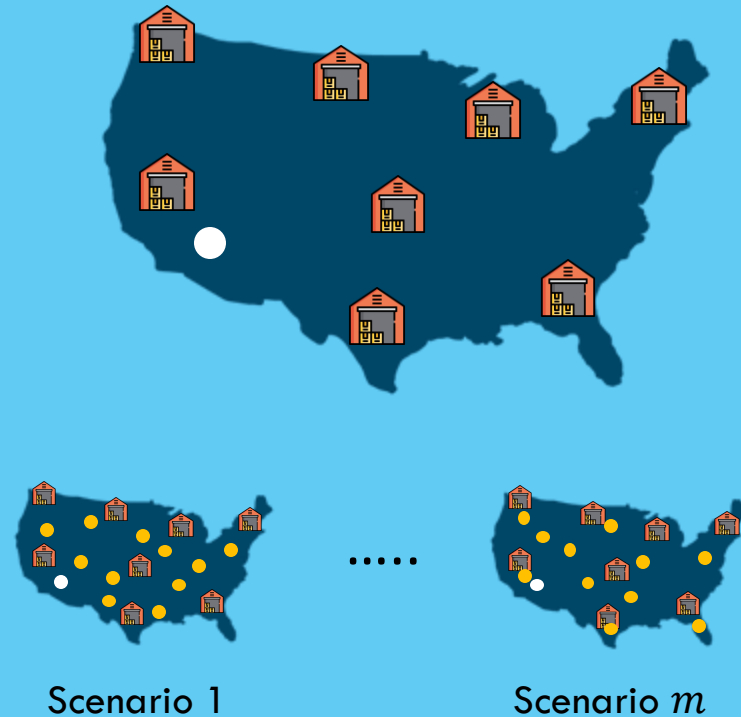
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Repeat the process



**Compensation(scenario  $k$ ,  $a$ ) =**  
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RAMS minimizes hindsight-based  
regret

# RAMS minimizes hindsight-based regret

Informal Meta Theorem [RAMS inherits guarantees of near-optimal algos].

Given a dynamic resource allocation setting, if there exists an algorithm **ALG** satisfying certain technical conditions, then

$$\text{Regret(RAMS)} \leq \text{Regret Upper Bound of } \mathbf{ALG} + \text{Sampling Error}$$

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Proof of the Informal Meta Theorem.

$$\text{Regret}(\text{RAMS}) = \sum_{t=1}^T \mathbb{E}[\text{Comp}_t(a_t^{\text{RAMS}})] \leq \sum_{t=1}^T \mathbb{E}[\text{Comp}_t(a_t^{\text{ALG}})]$$

Compensated Coupling or  
Performance Diff. Lemma

RAMS chooses the action with the  
minimum compensation

RAMS is on-par with SOTA



# RAMS is on-par with SOTA

Corollary of the Meta Theorem.

**Polynomial regret** for multi-secretary problem under different type distributions  
[this work]

**Bounded regret** for Network Revenue Management and Online Matching for a **few types** [Vera and Banerjee '21]

**Logarithmic regret** for Network Revenue Management with many types and non-degeneracy assumps. [Bray '23]

**Log-Squared regret** for Network Revenue Management with many types and w/o non-degeneracy assumps. [Jiang et. al '22]

# RAMS is on-par with SOTA

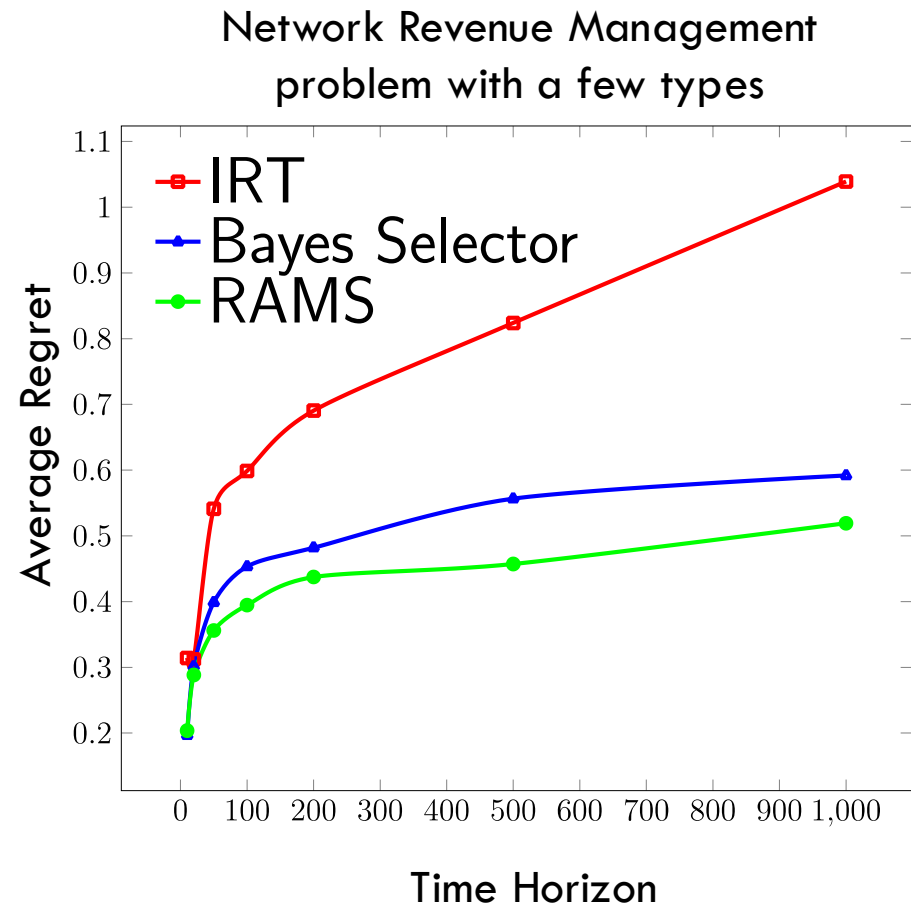
## Corollary of the Meta Theorem.

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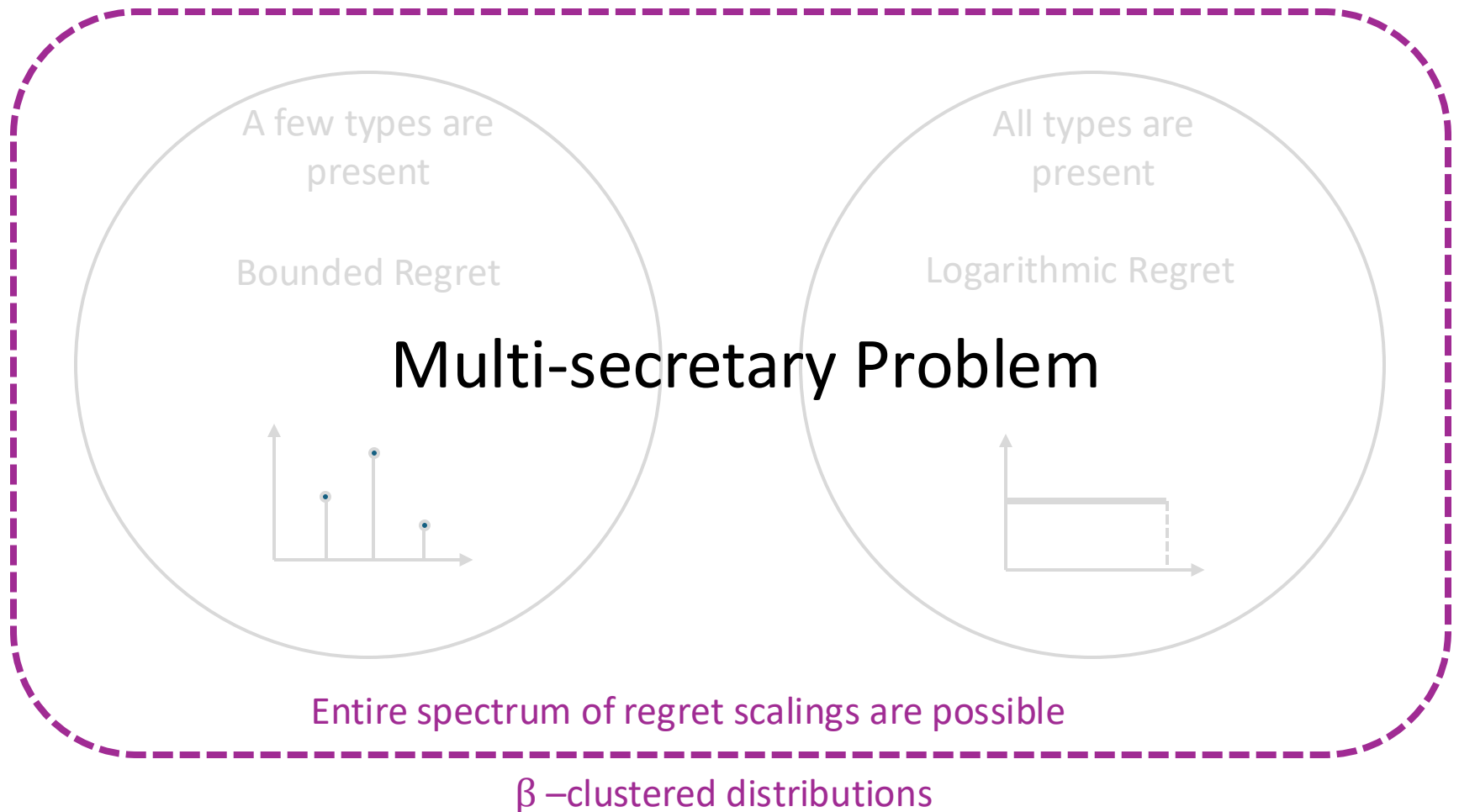
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# What is the interplay between the distribution of request types and achievable algorithmic performance?



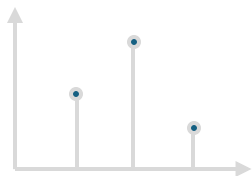
Can we design a **unified, simple and near-optimal** algorithms which works for all type distributions?

one policy to solve them all

Repeatedly **A**ct using **M**ultiple **S**imulations (**RAMS**)

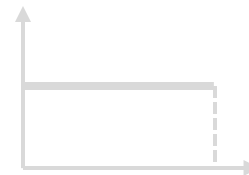
A few types are present

Bounded Regret



All types are present

Logarithmic Regret



Multi-secretary Problem

Entire spectrum of regret scalings are possible

$\beta$ -clustered distributions

Network  
Revenue  
Management

Repeated  
Auctions with  
Budgets



Multi-secretary



Order  
Fulfillment

Dynamic  
Pricing

**Repeatedly Act using Multiple Simulations (RAMS)**

Network  
Revenue  
Management



Repeated  
Auctions with  
Budgets



End of Part I



Order  
Fulfillment



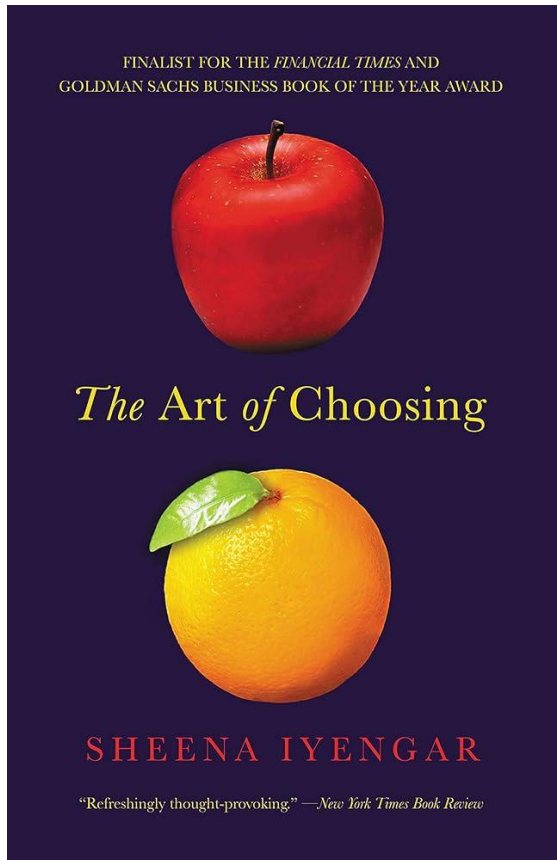
Dynamic  
Pricing

Repeatedly Act using Multiple Simulations (RAMS)

# Part II

## The Impact of Rankings and Personalized Recommendations in Marketplaces

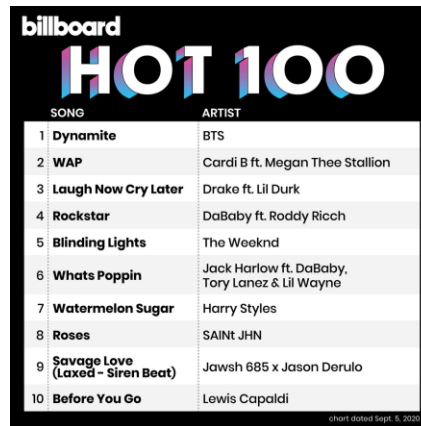
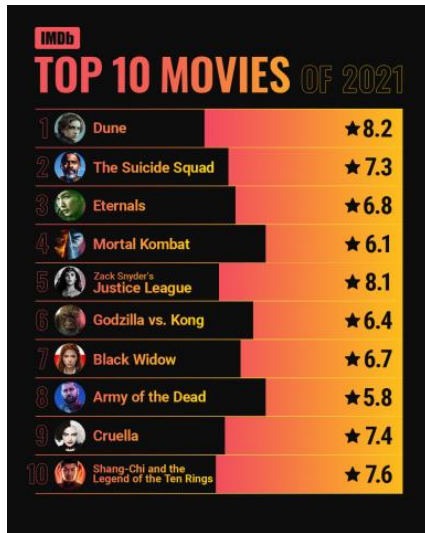
# People have ill-formed preferences



*"Majority of Americans would pursue a different degree, institution or major if given a chance to do it all over again"*



# Rankings offer a plausible solution to aid decision-making

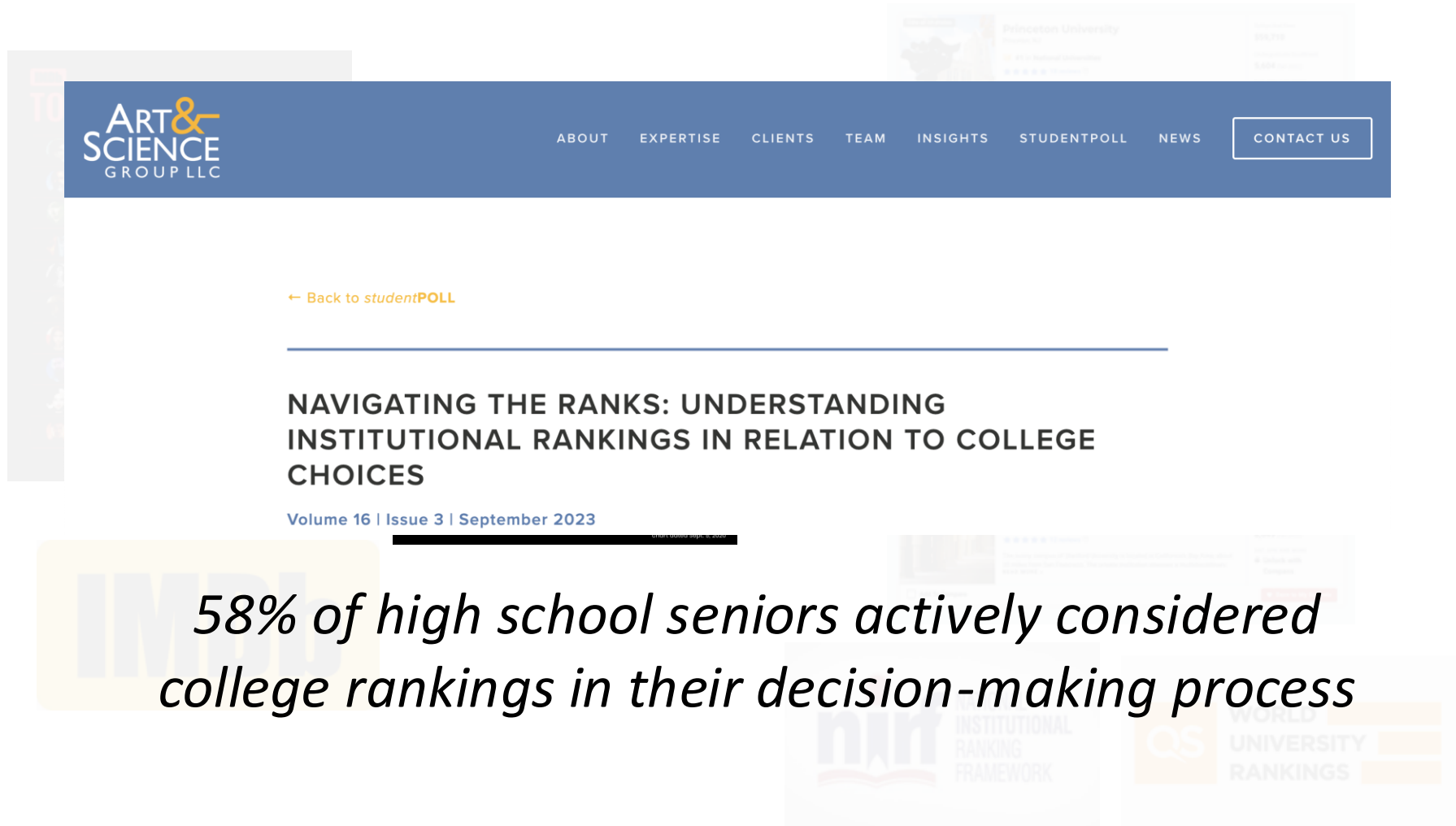


The screenshot displays a website with university profiles. Each profile includes a photo, name, location, ranking, reviews, tuition and fees, undergraduate enrollment, and a 'Save to My Schools' button.

University	Location	Ranking	Reviews	Tuition And Fees	Undergraduate Enrollment
Princeton University	Princeton, NJ	#1 in National Universities	18 reviews	\$59,710	5,604 (fall 2022)
Massachusetts Institute of Technology	Cambridge, MA	#2 in National Universities	17 reviews	\$60,156	4,657 (fall 2022)
Harvard University	Cambridge, MA	#3 in National Universities (tie)	17 reviews	\$59,076	7,240 (fall 2022)
Stanford University	Stanford, CA	#3 in National Universities (tie)	12 reviews	\$62,484	8,049 (fall 2022)



# Rankings offer a plausible solution to aid decision-making



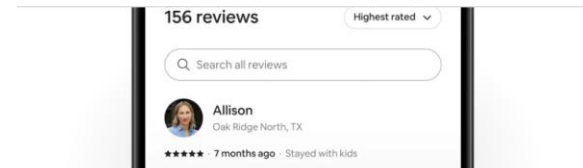
*58% of high school seniors actively considered college rankings in their decision-making process*

# Rankings are not personalized to individual tastes

NETFLIX



Instagram



## See the rating distribution

A new overall rating chart shows you a breakdown of a home's 1 to 5 star reviews.

## Sort by recency and rating

Now you can sort reviews to read the latest and greatest from prior guests.

## Find relevant reviews

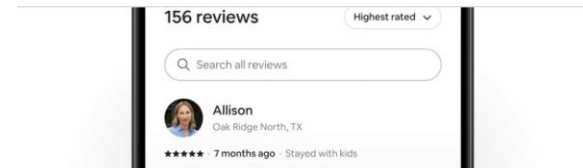
New details, like type of trip or length of stay, make it simple to pick out relevant reviews.

# Significant work on personalized recos in uncapacitated settings

NETFLIX



Instagram



## See the rating distribution

A new overall rating chart shows you a breakdown of a home's 1 to 5 star reviews.

## Sort by recency and rating

Now you can sort reviews to read the latest and greatest from prior guests.

## Find relevant reviews

New details, like type of trip or length of stay, make it simple to pick out relevant reviews.

# My vision for the future

## CORE, The Academic Explorer

### Course Recommendation Engine

Enter your name:

Select your gender:

Male



### Provide your preference order for colleges and courses:

Select College 1:

Select College 2:

Select College 3:

Select College 4:

Recommender Systems for College Recommendations

## Electives Chatbot

This is based on the starter kit with ReactJS + NextJS + TypeScript. You can [download the source code](#) for this Starter Kit from GitHub.

Contact Us



I am a chat assistant to help you navigate classes offered at CBS. How can I help you?

Enter your message...

Send

Chatbots for Course Recommendations

# Research Questions

What are the implications of providing personalized recommendations in marketplaces with capacity constraints?

How does this contrast with marketplaces with no capacity constraints?

# In a Nutshell



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We study a prototypical utility model which comprises of public and private utility components





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Identify subtle interplay of personalized recommendations and capacity constraints on social welfare



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We study a prototypical utility model which comprises of public and private utility components



Identify subtle interplay of personalized recommendations and capacity constraints on social welfare



Personalized recos may have *little to no* benefit in markets with uncapacitated supply (think movies, songs)



# In a Nutshell



We study a prototypical utility model which comprises of public and private utility components



Identify subtle interplay of personalized recommendations and capacity constraints on social welfare



Personalized recos may have *little to no* benefit in markets with uncapacitated supply (think movies, songs)



Personalized recos unlock *significant* social welfare in capacitated supply settings (think AirBnB, college admits)



# Model

$n$  individuals



$n$  items



$$u(\text{person}, \text{item}) = \rho \times \text{megaphone} + (1 - \rho) \times \text{eye}$$

Common term

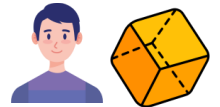
Depends only on



Known through  
publicly available  
rankings

Private term

Depends on both



Apriori unknown, can  
be known using  
personalized recos

# Model

$n$  individuals



$n$  items



$$u(\text{person}, \text{cube}) = \rho \times \text{megaphone} + (1 - \rho) \times \text{no eye}$$

Uncapacitated supply setting



Many-to-One matching

Individuals simply choose the item they want to consume

Capacitated supply setting



One-to-One matching

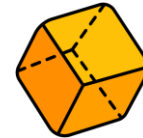
Individuals are matched by centralized clearinghouse using some matching mechanism

# Model

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Uncapacitated supply setting

Capacitated supply setting

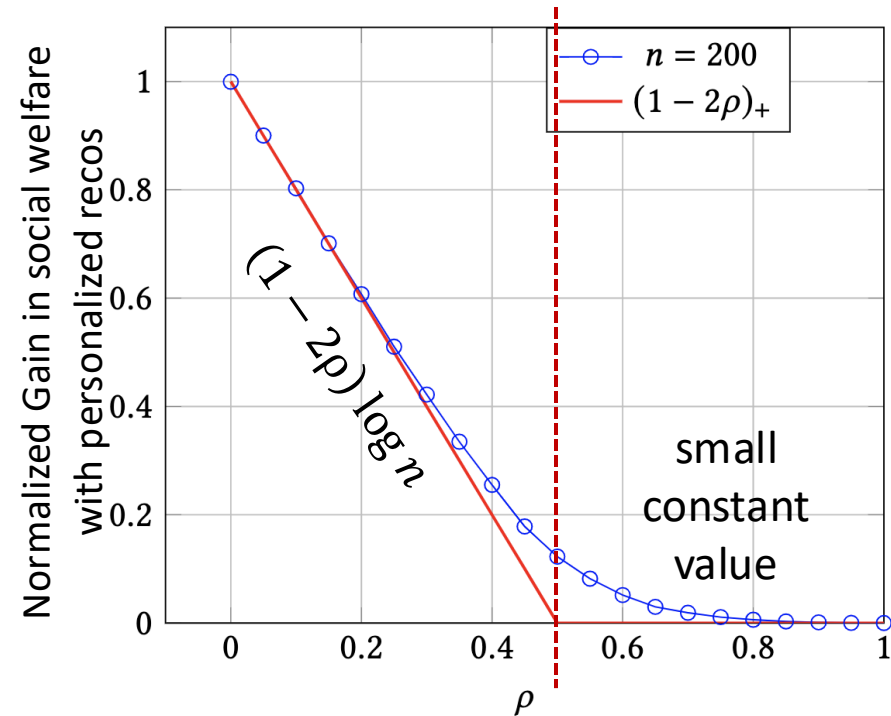
We care about the social welfare which is the average (across the  $n$  individuals) utility obtained by the individuals

# Results

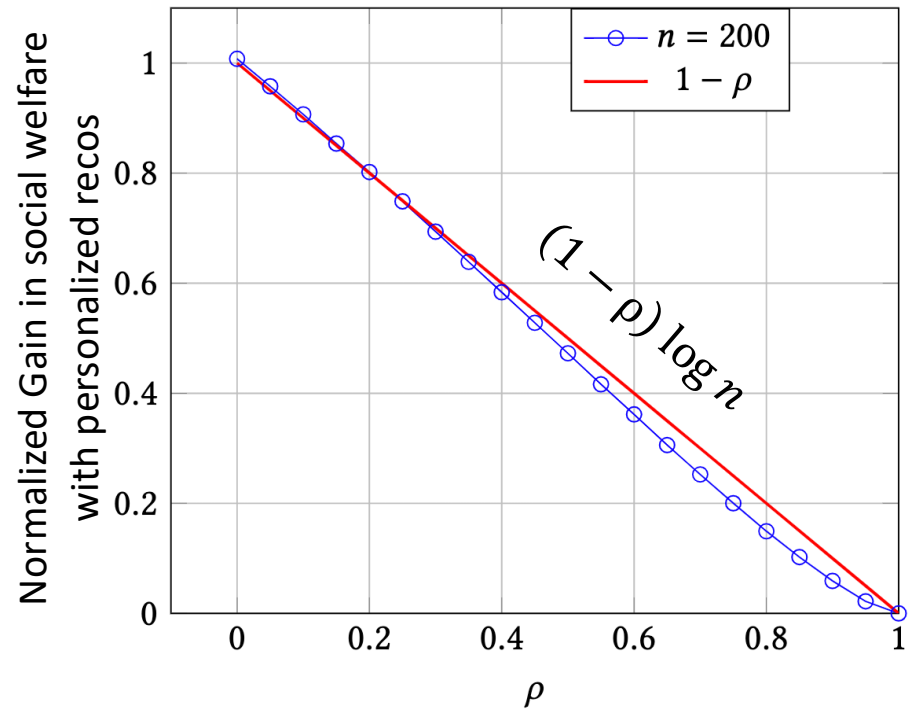
Common term is Exponentially distributed with rate 1

Private idiosyncratic term is Exponentially distributed with rate 1

Uncapacitated Supply Setting



Capacitated Supply Setting



# Part I in a Nutshell





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We study prototypical models of dynamic resource allocation problem (focus on the multisecretary prob)



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Shape of the reward distribution is a fundamental driver of algorithmic performance



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Workhorse policies can be highly suboptimal and near-optimal algorithms are overly specified



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Design a simple and near-optimal policy called Repeatedly Act using Multiple Simulations (**RAMS**)



# Part II in a Nutshell



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We study a prototypical utility model which comprises of public and private utility components



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Identify subtle interplay of personalized recommendations and capacity constraints on social welfare



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Personalized recos may have *little to no* benefit in markets with uncapacitated supply (think movies, songs)





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Identify subtle interplay of personalized recommendations and capacity constraints on social welfare



Personalized recos may have *little to no* benefit in markets with uncapacitated supply (think movies, songs)



Personalized recos unlock *significant* social welfare in capacitated supply settings (think AirBnB, college admits)



So long and **Thanks** for all the fish