

#### **Department of Information Technology**

National Institute of Technology Karnataka, Surathkal

# Distributed Memory Parallelism with MPI

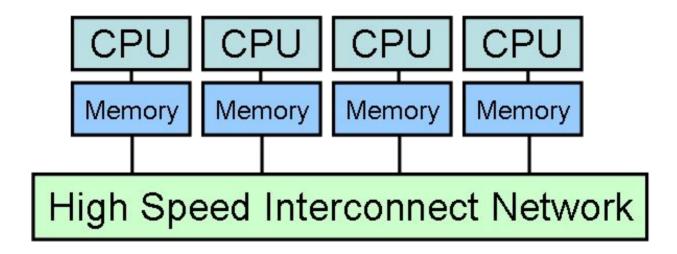
By,
Thanmayee,
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#### **Outline**

- Distributed Memory Architecture
- Introduction to MPI
- Structure of MPI program
- Types of Message Passing
- Basic Routines in Point to Point Communication
- Example programs on Point to Point Communication
- Basic Routines in Collective Communication
- Sample Programs on Collective Communication

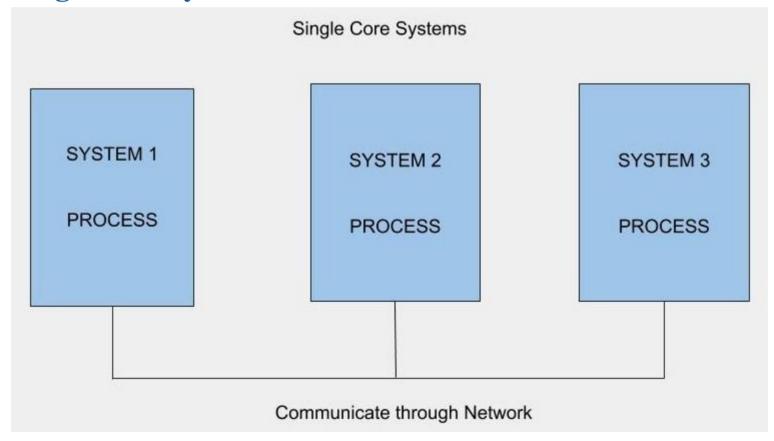
### Distributed Memory Architecture

- Each processor has its own memory
- They cannot access the memory of other processors.
- Any data that needs to be shared must be explicitly transmitted from one processor to another using <u>Message</u> <u>Passing.</u>



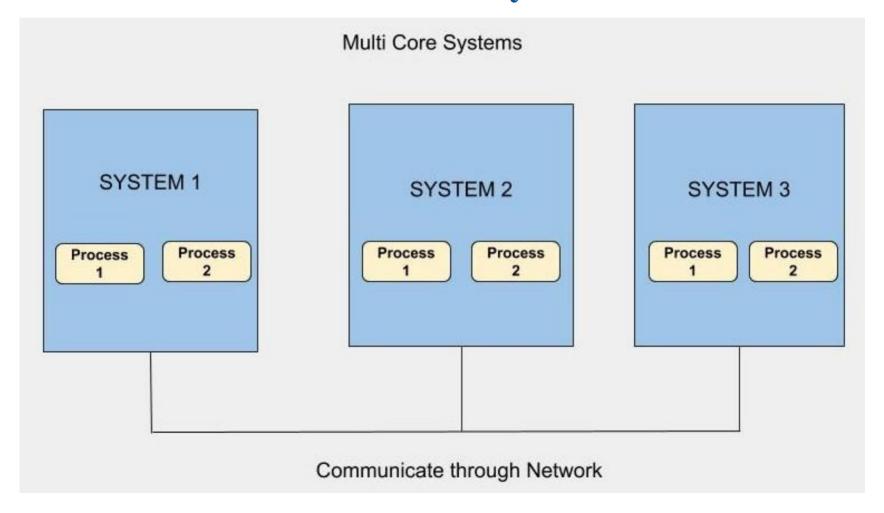
#### DISTRIBUTED MEMORY ARCHITECTURE

- Systems with single core communicating through distributed memory.
- Heterogeneous systems

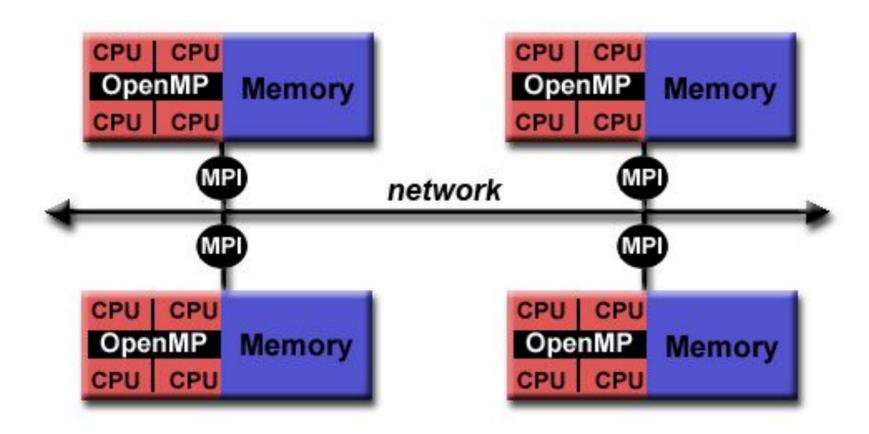


#### DISTRIBUTED MEMORY ARCHITECTURE

• Systems with multiple core communicating through shared and distributed memory

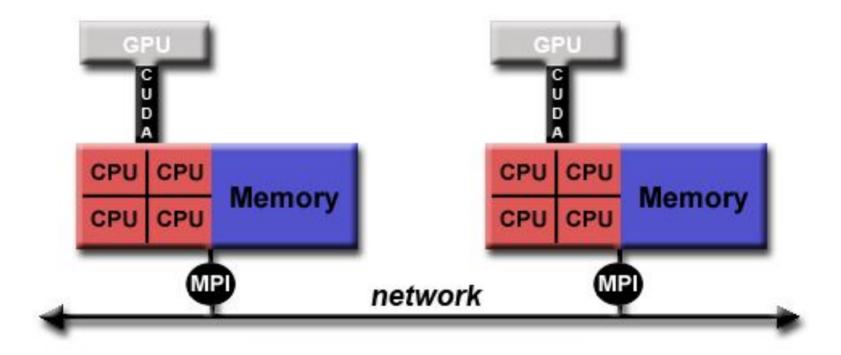


# **Hybrid Model**



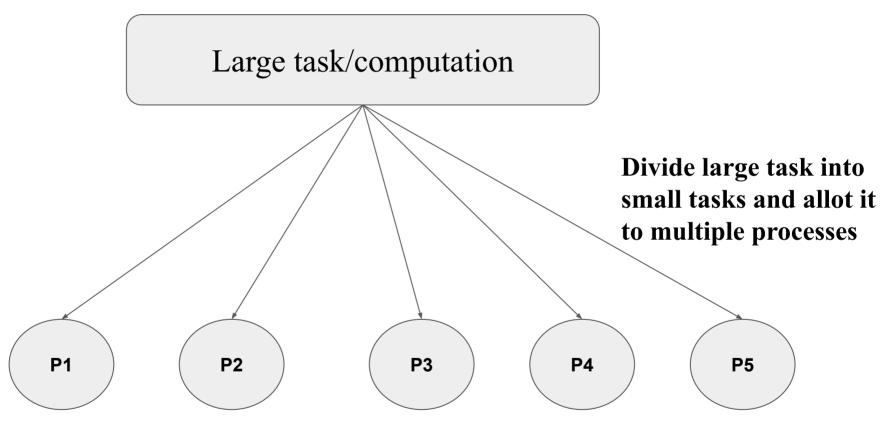
Reference: https://computing.llnl.gov/tutorials/parallel\_comp/#ModelsMessage

# **Hybrid Model**



Reference: https://computing.llnl.gov/tutorials/parallel\_comp/#ModelsMessage

## **Parallel Computation:**



For Example: N = 1,00,000 divided into P1=20,000, P2=20,000 .....

Computation is same. Data is different. Single Program Multiple Data

#### INTRODUCTION TO MPI

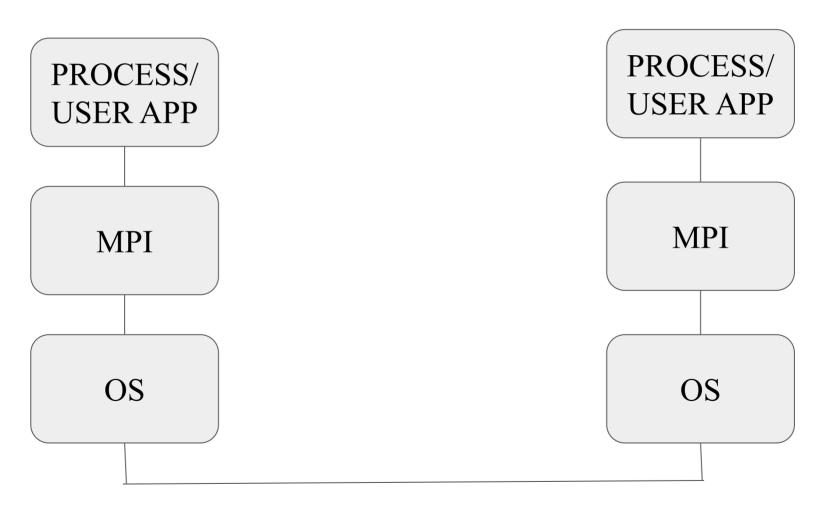
#### What is MPI?

- Message Passing Interface is a specification.
  - A standard for vendors to implement.
- It is a library, i.e. a set of subroutines, functions and constants
- Allows Message Passing between processes.
- It is based on Single Program, Multiple Data (SPMD)
  - Every process executes the same program
  - Each process performs computations on its local variables, then communicates with other processes, in order to get the final result.

# MPI: Major Goals

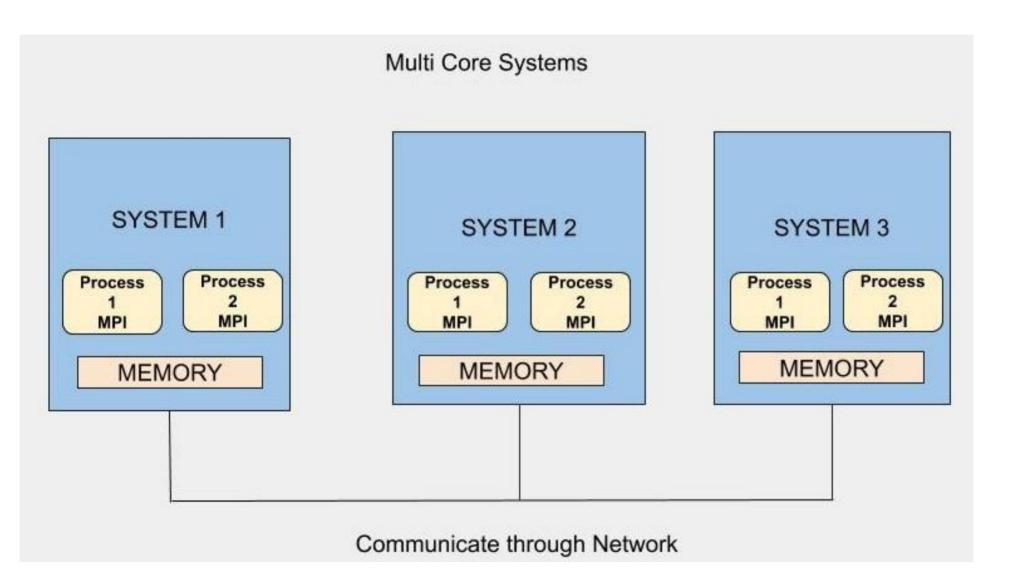
- Portability:
  - An MPI library exists on ALL parallel computing platforms so it is highly portable.
- Support heterogeneity
- High performance through efficient implementations
- Encourage overlap of communication and computations.
- Reliability

#### MPI is a Middleware



**NETWORK** 

#### MPI is a Middleware



# **MPI** Implementations

- OpenMPI (<u>www.open-mpi.org</u>)
- MPICH (<u>www.mpich.org</u>)
- HP MPI
- Intel MPI
- Scali MPI
- IBM MPI

# Thank You



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#### STRUCTURE OF MPI PROGRAM

**MPI Include File** 

**Initialize MPI Environment** 

**Computations and Message Passing** 

**Terminate MPI Environment** 

#### **MPI** Routines

#### . Start and terminate:

To initialize and terminate the MPI environment

#### . Communicators:

To identify the communication world (cluster of processes)

#### . Getting Information :

To get the number of processes and process ids

#### Sending and Receiving messages :

Actual computation and communication

#### STRUCTURE OF MPI PROGRAM

MPI Include File

#include<mpi.h>

Initialize MPI Environment

MPI\_Init(&argc,&argv);

**Computations and Message Passing** 

Terminate MPI Environment

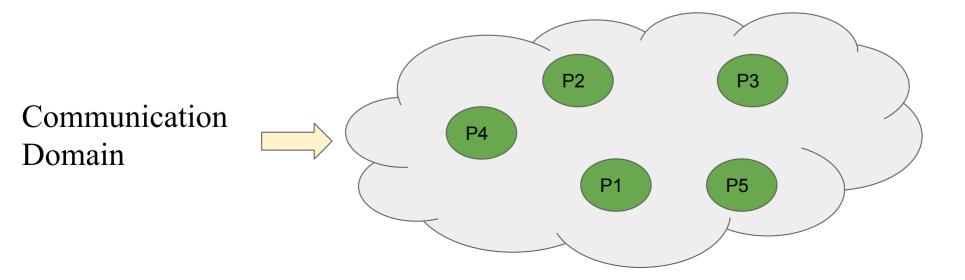
MPI\_Finalize();

#### **MPI Start and Terminate Routines**

```
#include<stdio.h>
int main(int argc,char **argv)
        MPI_Init(&argc,&argv);
        MPI_Finalize();
        return 0;
```

#### Communicators

- MPI defines communication domain set of processes that can communicate with each other.
- MPI\_comm : data type stores information about communication domains.
- Default communicator MPI COMM WORLD



# **Getting Information**

- MPI Comm size
- MPI\_Comm\_rank

- Syntax :
- int MPI\_Comm\_size(MPI\_Comm comm, int \*size)
- Int MPI\_Comm\_rank(MPI\_Comm comm, int \*rank)

# General MPI Program

```
#include<mpi.h>
       int main(int argc,char **argv)
       MPI_Init(&argc,&argv);
       MPI_Comm_size(MPI_COMM_WORLD,&size);
       MPI_Comm_rank(MPI_COMM_WORLD,&rank);
       MPI_Finalize();
       return 0;
```

#### **Example: Hello World**

```
#include<mpi.h>
int main(int argc,char *argv[])
int size, myrank;
MPI_Init(&argc,&argv);
MPI Comm size(MPI COMM WORLD,&size);
MPI Comm rank(MPI COMM WORLD,&myrank);
printf("Process %d of %d, Hello World",myrank,size);
MPI Finalize();
return 0;
```

#### MPI Hello World:

```
tans@tans-Inspiron-3542:~/PC$ mpiexec -n 5 ./a.out
Process 0 of 5, Hello World
Process 1 of 5, Hello World
Process 4 of 5, Hello World
Process 2 of 5, Hello World
Process 3 of 5, Hello World
```

#### **MPI Include File**



#### **Initialize MPI Environment**



#### **Computations and Message Passing**



**Terminate MPI Environment** 



# **Types of Message Passing:**

#### Point to Point

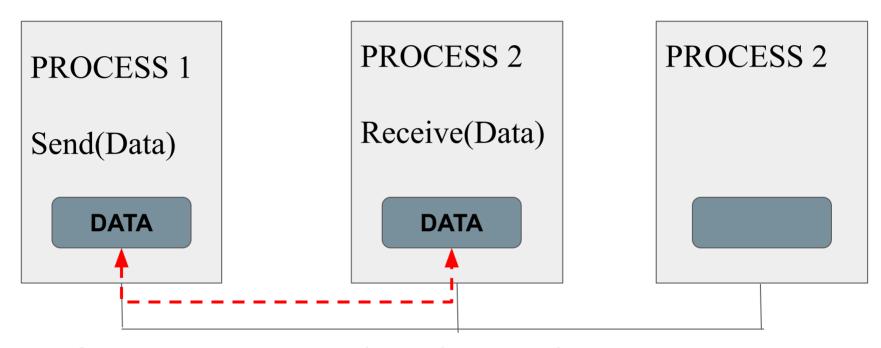
- Two processes
- Send and Receive are the basic functions

#### Collective messages

- Group of processes involved in communication
- Functions like Broadcast, Scatter, Gather, Parallel Reduction

#### **Point to Point Communication**

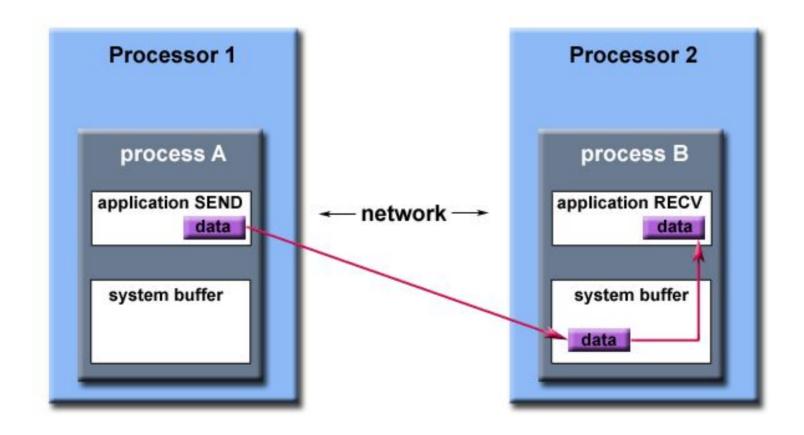
• Two processes involved in sending and receiving data.



- ID of sender and receiver is required.
- Specify what has to be sent and received.
- Communication needs to be synchronized.
- Communication makes use of buffers.

#### **Point to Point Communication**

• Data transfer from Sender Process to Receiver Process.



Path of a message buffered at the receiving process

#### Send and Receive Variants

- Blocking Send and Receive
- Non Blocking Send and Receive
- Based on modes of Communication:
  - Standard
  - Synchronous
  - Buffered
  - Ready

## **Blocking Send and Receive**

- Basic Send and Receive routine for point to point communication.
- MPI Routines:
  - MPI\_Send()
  - o MPI Recv()

# Thank You



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### **Blocking Send and Receive**

MPI\_Send()

MPI\_Send (void \*buf, int count, MPI\_Datatype type,int dest, int tag, MPI\_Comm comm)

#### **Parameters:**

**buf:** initial address of send buffer

**count :** number of elements in send buffer (nonnegative integer)

**datatype:** datatype of each send buffer element. Ex: MPI\_INT,

MPI\_CHAR

**dest:** rank of destination (integer)

tag: message tag (integer). For tagging send and receive.

**comm**: Communication domain of the communicating processes.

### **Blocking Send and Receive**

MPI\_Recv():

MPI\_Recv(void \*buf, int count, MPI\_Datatype datatype, int source, int tag, MPI\_Comm comm, MPI\_Status \*status)

#### **Parameters:**

**buf:** initial address of receive buffer

**count :** max number of elements in receive buffer (nonnegative integer)

**datatype:** datatype of each receive buffer element. Ex: MPI\_INT,

MPI\_CHAR

**source:** rank of source (integer)

tag: message tag (integer). For tagging send and receive.

**comm**: Communication domain of the communicating processes.

status: status object (Status). It is a structure containing information about

source, tag and error code.

### • MPI DATATYPES:

Table 1: Basic C datatypes in MPI

MPI Datatype	C data type
MP1_CHAR	signed char
MP1_SHORT	signed short int
MP1_INT	signed int
MP1_LONG	signed long int
MP1_UNS1GNED_CHAR	unsigned char
MP1_UNSIGNED_SHORT	unsigned short int
MP1_UNSIGNED	unsigned int
MP1_UNSIGNED_LONG	unsigned long int
MP1_FLOAT	float
MP1_DOUBLE	double
MP1_LONG_DOUBLE	long double
MP1_BYTE	
MP1_PA CKED	

# General MPI Program

```
#include<mpi.h>
    int main(int argc,char **argv)
{
        ...
        MPI_Init(&argc,&argv);
        ...
        MPI_Comm_size(MPI_COMM_WORLD,&size);
        MPI_Comm_rank(MPI_COMM_WORLD,&rank);
```

### COMPUTATIONS AND MESSAGE PASSING

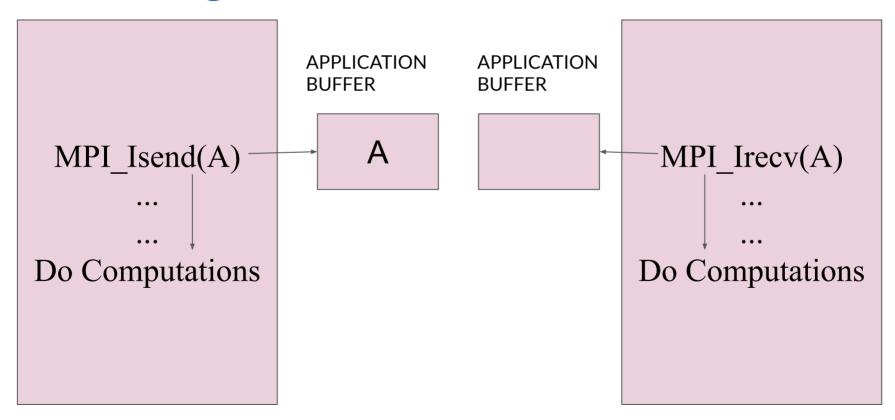
```
MPI_Finalize();
...
return 0;
```

# MPI Example - 1

```
for(i=0;i<50;i++) / Process 0 initializes array x
   x[i]=i+1;
if(myrank==0)
MPI_Send(x,10,MPI_INT,1,1,MPI_COMM_WORLD);
else if(myrank==1)
MPI_Recv(y,10,MPI_INT,0,1,MPI_COMM_WORLD,&status);
printf("Process %d Received Data from Process %d\n",
myrank, status. MPI_SOURCE);
   for(i=0;i<10;i++)
   printf("%d\t",y[i]);
          Process 1 Recieved data from Process 0
                                                                  10
```

# Non Blocking Send and Receive

- Allows overlapping of computation and communication
- Advantage is Performance Gain



# Non Blocking Send and Receive

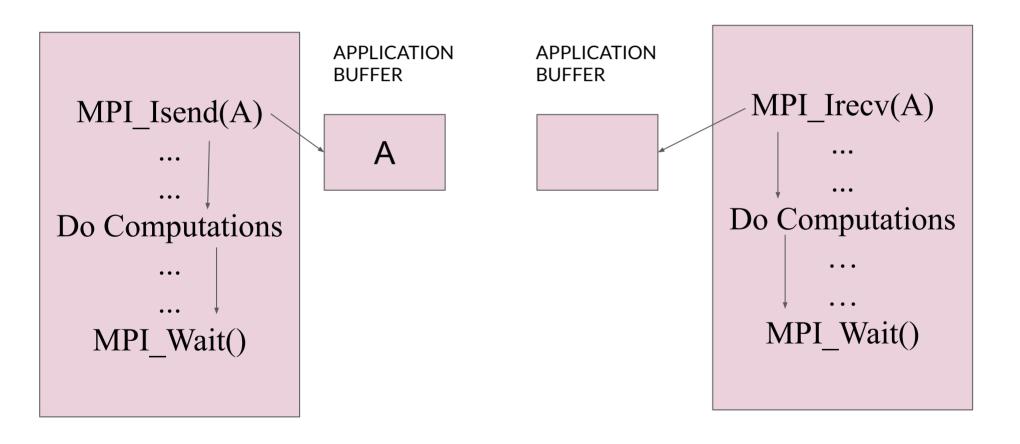
MPI\_Isend (&buf,count,datatype,dest,tag,comm,&request)

MPI\_Irecv (&buf,count,datatype,source,tag,comm,&request)

### **Parameters:**

- Same as Send() and Recv() except for request
- request: handle. This helps to get information about MPI\_Isend and MPI\_Irecv status.
- Used in routines : MPI\_Wait() and MPI\_Test()

# Non Blocking Send and Receive



# MPI\_Wait() and MPI\_Test()

### Syntax:

```
int MPI_Wait( MPI_Request *request, MPI_Status *status );
int MPI_Test( MPI_Request *request, int *flag, MPI_Status *status );
```

- If request is set to MPI\_REQUEST\_NULL (set if operation is completed) then:
  - MPI\_Wait returns immediately with an empty status.
  - MPI\_Test sets flag to true and returns an empty status.

# MPI Example - 2

```
if(myrank==0)
x=10;
MPI_Isend(&x,1,MPI_INT,1,20,MPI_COMM_WORLD,&request);
printf("Send returned immediately\n");
else if(myrank==1)
MPI_Irecv(&x,1,MPI_INT,0,25,MPI_COMM_WORLD,&request);
printf("Receive returned immediately\n");
printf("Process %d of %d, Value of x is %d\n",myrank,size,x);
      tans@tans-Inspiron-3542:~/PC$ mpiexec -n 2 ./a.out
      Send returned immediately
      Receive returned immediately
      Process 1 of 2, Value of x is 0
```

### What is the risk here?

```
if(myrank==0)
x=10;
MPI_Isend(&x,1,MPI_INT,1,20,MPI_COMM_WORLD,&request);
printf("Send returned immediately\n");
x = x + 10;
```

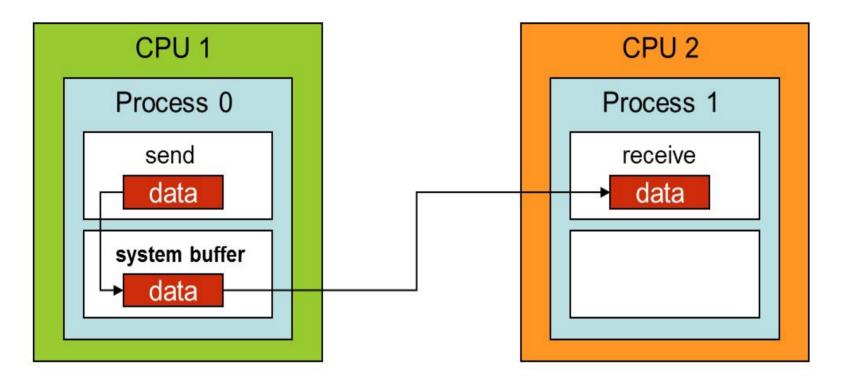
### Make sure that x is available for reuse:

```
if(myrank==0)
x=10;
MPI Isend(&x,1,MPI INT,1,20,MPI COMM WORLD,&request);
printf("Send returned immediately\n");
MPI Wait(request, status)
x=x+10;
```

### **Communication Modes**

- Standard Mode: Calls block until message has been either transferred or copied to an internal buffer for later delivery. Ex: MPI\_Send() and MPI\_Recv()
- Buffered Mode: Send may start and return before a matching receive. MPI\_Bsend()
- Synchronous Mode: Call blocks until matching receive has been posted and the message reception has started. MPI\_Ssend()
- Ready Mode: Requires that a matching receive is already posted. MPI\_Rsend().

### **Buffered Mode**



MPI\_BUFFER\_ATTACH( buffer, size)

buffer initial buffer address (choice)

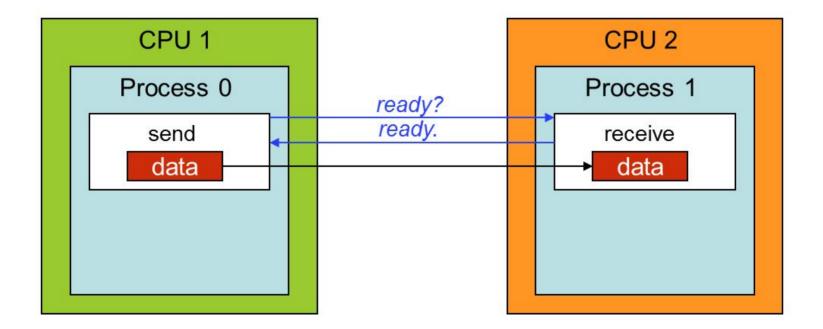
size buffer size, in bytes (integer)

NOTE: A user may specify a buffer to be used for buffering messages sent in buffered mode.

Image Reference: https://www.codingame.com/playgrounds/47058/have-fun-with-mpi-in-c/communication-modes

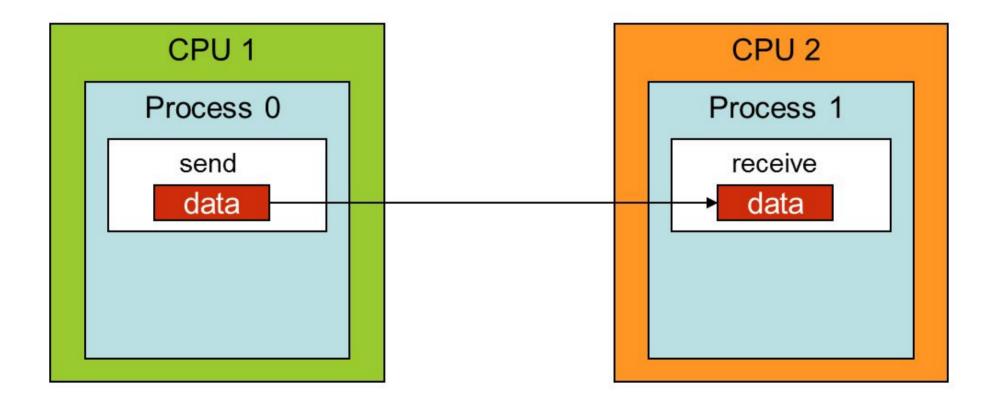
# Synchronous Mode

We see that the data is not copied to system buffer.



# Ready Mode

We make use of MPI\_Barrier() to wait for the receive to be posted. This will not result in error.



# MPI-Example - 3

```
if(myrank==0) {

// Blocking send will expect matching receive at the destination In Standard mode,Send will return after copying the data to the buffer

MPI_Send(x,10,MPI_INT,1,1,MPI_COMM_WORLD);

// This send will be initiated and matching receive is already there so the program will not lead to deadlock

MPI_Send(y,10,MPI_INT,1,2,MPI_COMM_WORLD);
}
```

```
else if(myrank==1)
{
//P1 will block as it has not received a matching send with tag 2

MPI_Recv(x,10,MPI_INT,0,2,MPI_COMM_WORLD,&status);

MPI_Recv(y,10,MPI_INT,0,1,MPI_COMM_WORLD,MPI_STATUS_IGNORE);
}
```

# MPI Example 3

#### **PROCESS 1**

 $MPI\_Send(x,10,...1,1,...);$ 

*MPI\_Send(y,10,..,1,2,..)*;

#### **PROCESS 2**

```
MPI_Recv(x,10,...,0,2,....);
\leq_{BLOCK}
MPI_Recv(y,10,...,0,1,....);
```

# MPI Example - 4

```
if(myrank==0) {
    MPI_Ssend(x,10,MPI_INT,1,1,MPI_COMM_WORLD);
    MPI_Send(y,10,MPI_INT,1,2,MPI_COMM_WORLD);
}
```

```
else if(myrank==1)
{
    MPI_Recv(x,10,MPI_INT,0,2,MPI_COMM_WORLD,&status);

MPI_Recv(y,10,MPI_INT,0,1,MPI_COMM_WORLD,MPI_STATUS_IGNORE);
}
```

# MPI Example - 4

```
if(myrank==0) {
    MPI_Ssend(x,10,MPI_INT,1,1,MPI_COMM_WORLD);

// Synchronous Blocking send will expect matching receive at the destination.
This results in deadlock.

MPI_Send(y,10,MPI_INT,1,2,MPI_COMM_WORLD); // This call will not be executed
}
```

```
else if(myrank==1)
{
    MPI_Recv(x,10,MPI_INT,0,2,MPI_COMM_WORLD,&status); // P1 will block
as it has not received a matching send with tag 2

MPI_Recv(y,10,MPI_INT,0,1,MPI_COMM_WORLD,MPI_STATUS_IGNORE);
}
```

# Thank You



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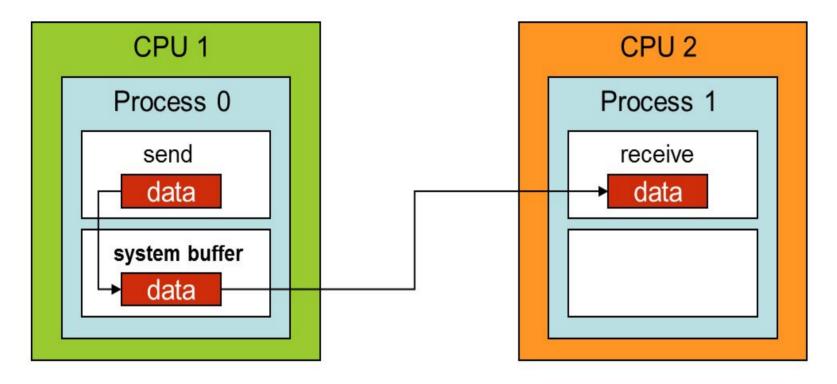
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- Synchronous Mode: Call blocks until matching receive has been posted and the message reception has started. MPI\_Ssend()
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### **Buffered Mode**



MPI\_BUFFER\_ATTACH( buffer, size)

buffer initial buffer address (choice)

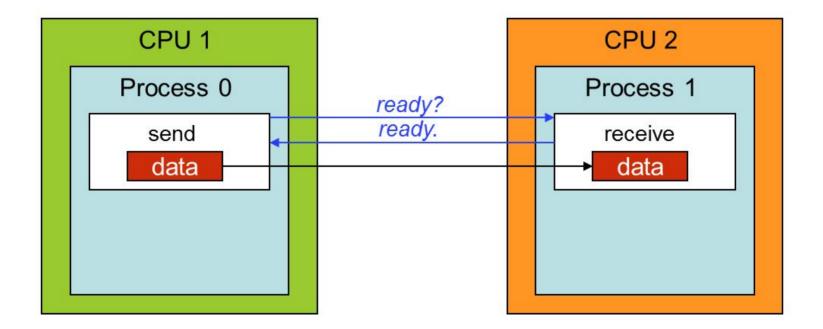
size buffer size, in bytes (integer)

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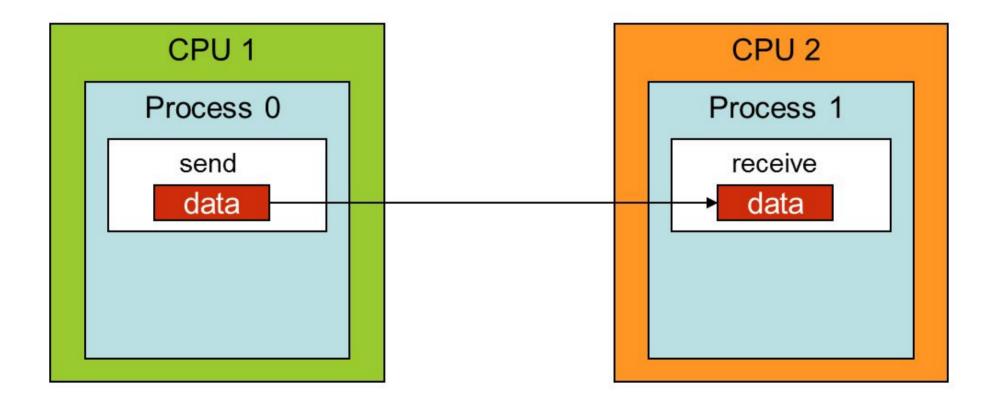
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We see that the data is not copied to system buffer.



# Ready Mode

We make use of MPI\_Barrier() to wait for the receive to be posted. This will not result in error.







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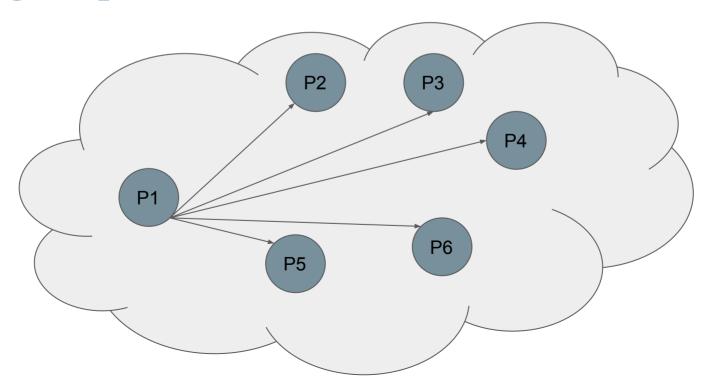
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### **Collective Communication**

- Multiple processes in same communicator involve in collective communication.
- They are blocking calls.
- No tags required.



### **Collective Communication**

- Barrier
- Broadcast
- Scatter
- Gather
- Reduce
- Scattery
- Gathery

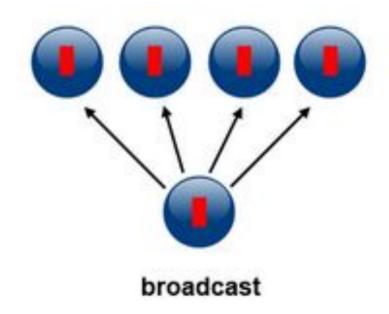
## **Collective communication: MPI\_Barrier**

- Mainly used for synchronization
- The call returns only after all the processes have called Barrier function.
- Uses:
  - Access to files
  - Achieve consistency

Syntax: MPI\_Barrier(MPI\_COMM\_WORLD)

# **Collective Communication: Broadcast**

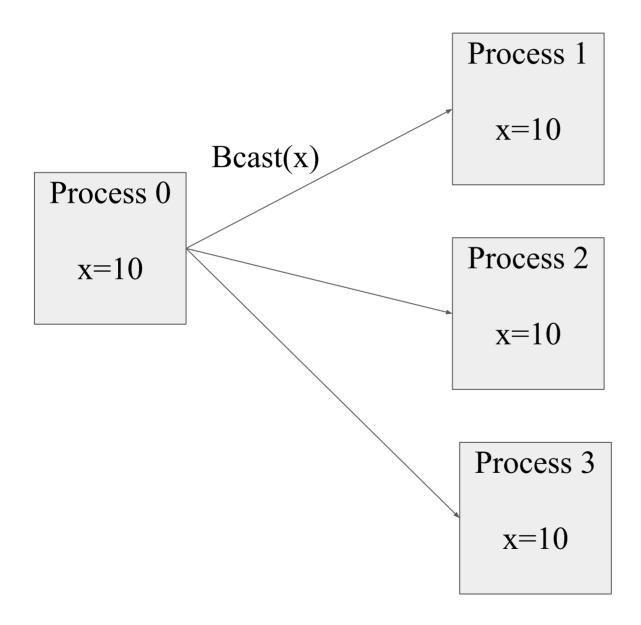
- MPI\_Bcast(buf, count, datatype, source, comm)
  - buf: send buffer of sender and receive buffer of receiver
  - o source: process which sends data to others



# MPI Example - 5

```
if(myrank==0)
scanf("%d",&x);
MPI_Bcast(&x,1,MPI_INT,0,MPI_COMM_WORLD);
printf("Value of x in process %d: %d\n",myrank,x);
MPI_Finalize();
return 0;
```

# Bcast():



# **Broadcast Output:**

```
tans@tans-Inspiron-3542:~/PC$ mpiexec -n 4 ./a.out

3

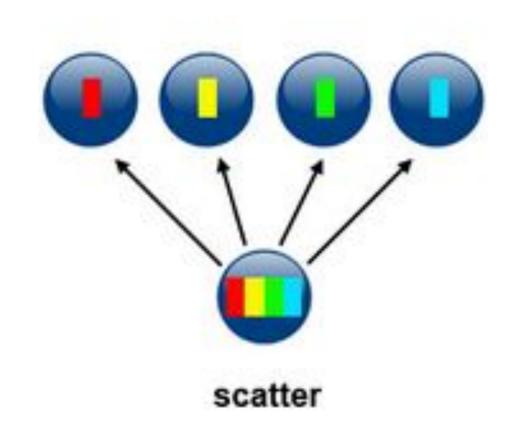
Value of x in process 0 : 3

Value of x in process 1 : 3

Value of x in process 2 : 3

Value of x in process 3 : 3
```

### **Collective Communication: Scatter**



### **Collective Communication: Scatter**

MPI\_Scatter(sendbuf, sendcount, datatype, recvbuf, recvcount, datatype, root, comm)

### **Parameters:**

sendbuf: sender buffer

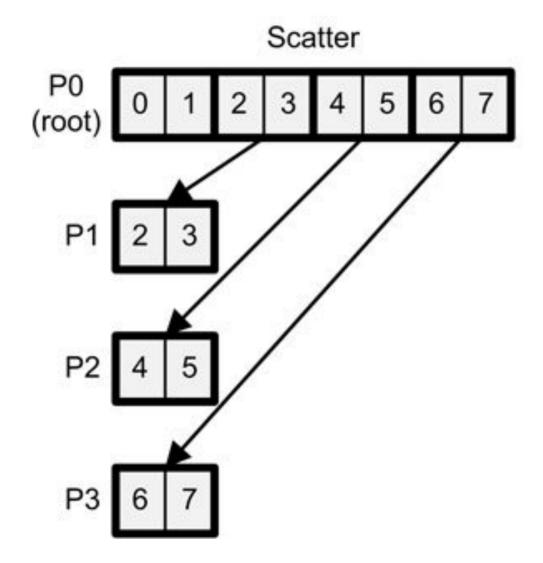
sendcount: specify the number of elements to be sent. recvcount should be same as sendcount

recvbuf: recv buffer

root: Sender

### **MPI\_Scatter**

**Example:** 



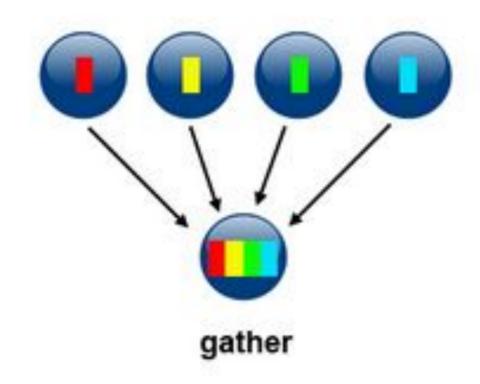
### MPI Example - 6

```
if(myrank==0)
printf("Enter values into array x:\n");
for(i=0;i<8;i++)
scanf("\%d",&x[i]);
MPI Scatter(x,2,MPI INT,y,2,MPI INT,0,MPI COMM WORLD);
for(i=0;i<2;i++)
printf("\nValue of y in process %d : %d\n",myrank,y[i]);
```

### Output

```
tans@tans-Inspiron-3542:~/PC$ mpiexec -n 4 ./a.out
Enter values into array x:
1 2 3 4 5 6 7 8
Value of y in process 0 : 1
Value of y in process 0 : 2
Value of y in process 1 : 3
Value of y in process 1 : 4
Value of y in process 2 : 5
Value of y in process 2 : 6
Value of y in process 3 : 7
Value of y in process 3 : 8
```

### **Collective Communication: Gather**



### **Collective Communication: Gather**

MPI\_Gather(sendbuf, sendcount, datatype, recvbuf, recvcount, datatype, root, comm)

### **Parameters:**

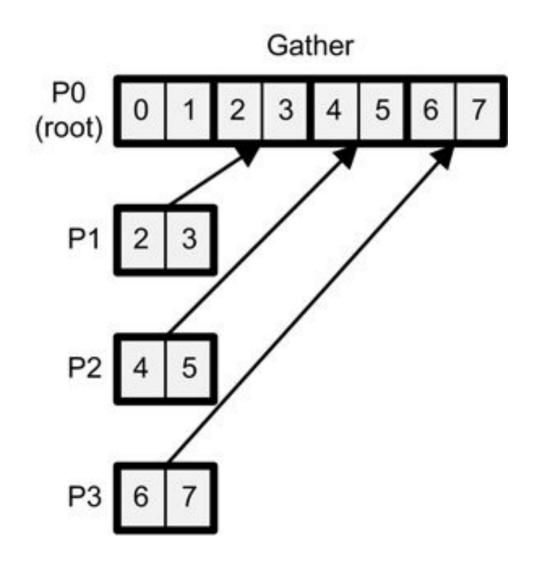
sendbuf: buffer of sending processes

sendcount and recvcount value is same

recvbuf: root process's buffer

root: process where the data is gathered

### **MPI\_Gather**



### MPI-Example 7

```
x=10, y[50]

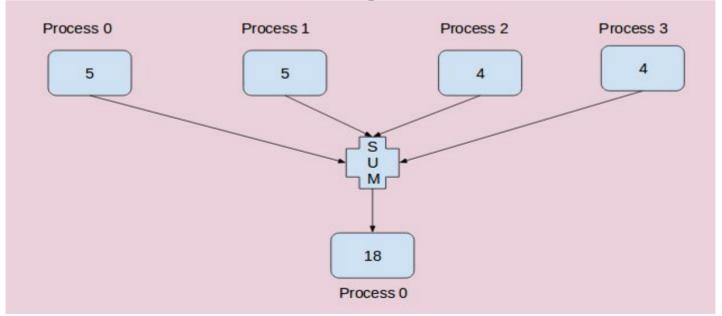
MPI_Gather(&x,1,MPI_INT,y,1,MPI_INT,0,MPI_COMM_WORLD);
// Value of x at each process is copied to array y in Process 0
if(myrank==0)
{
for(i=0;i<size;i++)
printf("\nValue of y[%d] in process %d: %d\n",i,myrank,y[i]);
}</pre>
```

### Output

```
tans@tans-Inspiron-3542:~/PC$ mpiexec -n 4 ./a.out
Value of y[0] in process 0:10
Value of y[1] in process 0 : 10
Value of y[2] in process 0:10
Value of y[3] in process 0:10
tans@tans-Inspiron-3542:~/PC$ mpiexec -n 6 ./a.out
Value of y[0] in process 0:10
Value of y[1] in process 0 : 10
Value of y[2] in process 0:10
Value of y[3] in process 0 : 10
Value of y[4] in process 0:10
Value of y[5] in process 0 : 10
```

### **Collective Communication: Reduce**

- Allows to perform computations on data present at multiple processes.
- Computations like: Sum, Product, Maximum, Minimum
- Stores the result in one process.



### **Collective Communication: Reduce**

MPI\_Reduce(sendbuf, recvbuf, count, datatype, operation, dest, comm)

### **Parameters:**

count: size of receive buffer

operation:

MPI name	Operation
MPI_MAX	Maximum
MPI_MIN	Minimum
MPI_SUM	Summation
MPI_PROD	Product
MPI_LAND	Logical AND
MPI_LOR	Logical OR
MPI_LXOR	Logical XOR

### MPI Example - 8

```
x=myrank;
MPI_Reduce(&x,&y,1,MPI_INT,MPI_SUM,0,MPI_COMM_WORLD);
if(myrank==0)
{
printf("Value of y after reduce : %d\n",y);
}
```

### Output

```
tans@tans-Inspiron-3542:~/PC$ mpiexec -n 3 ./a.out
Value of y after reduce : 3
tans@tans-Inspiron-3542:~/PC$ mpiexec -n 4 ./a.out
Value of y after reduce : 6
```

## Thank You



### **Department of Information Technology**

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# Distributed Memory Parallelism with MPI

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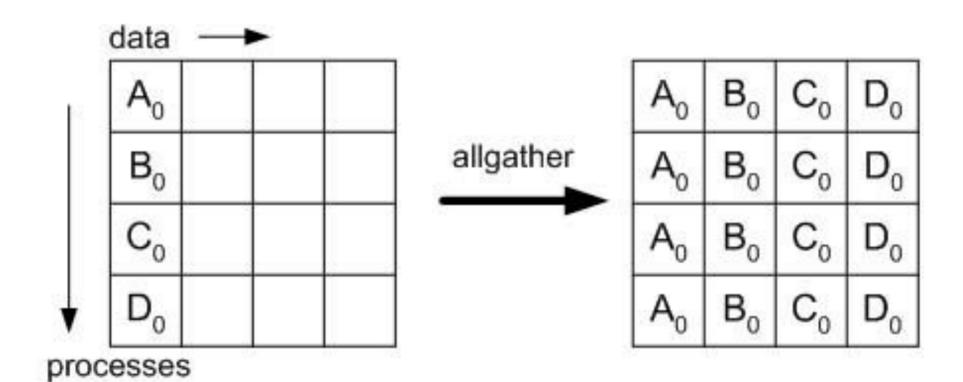
### **Outline**

- Distributed Memory Architecture
- Introduction to MPI
- Structure of MPI program
- Types of Message Passing
- Basic Routines in Point to Point Communication
- Example programs on Point to Point Communication
- Basic Routines in Collective Communication
- Sample Programs on Collective Communication

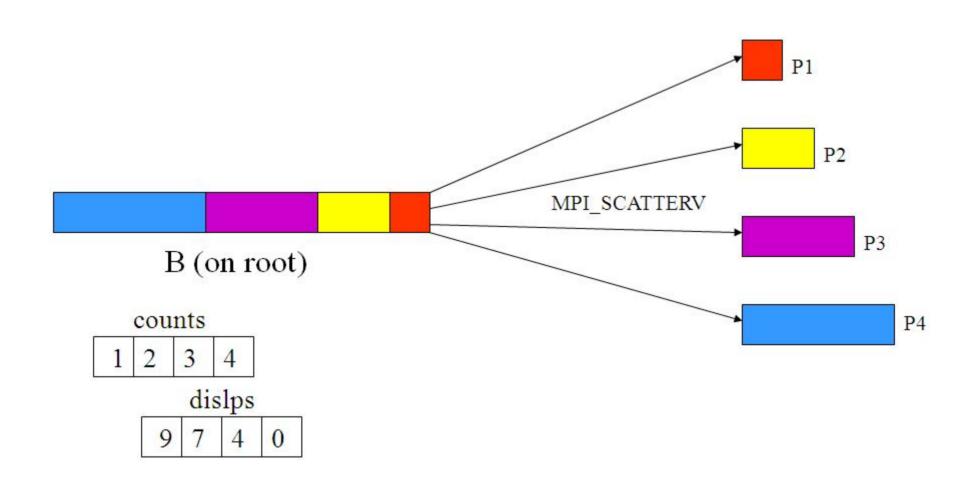
# **MORE Collective Communication Routines:**

- MPI\_Gatherv()
- MPI\_Scatterv()
- MPI\_Allgather
- MPI\_AllReduce()
- MPI\_Scan()
- MPI\_Comm\_Split()

### MPI\_Allgather



### **MPI\_Scatterv**



### **MPI\_Scatterv():**

MPI\_Scatterv(sendbuf, sendcounts, displacement, datatype, recvbuf, recvcount, datatype, root, comm)

### **Parameters:**

sendcounts: array with number of elements to be sent to each process. ex: sendcount[0]=10 means send 10 elements to Process zero. sendcount[1]=20 means send 20 elements to Process one.

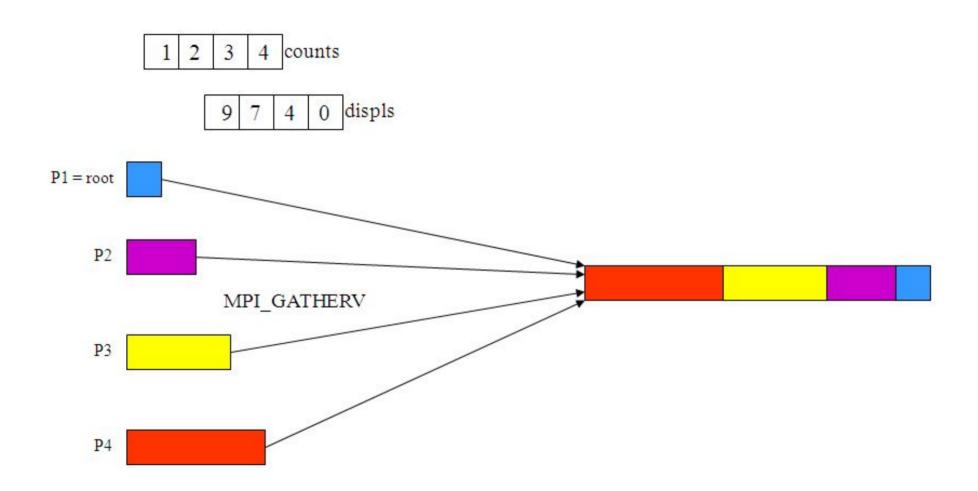
displacement: array which holds the index from where the data is to be sent to each process. Ex: disp[0]=0 means Process zero gets elements starting with index zero. disp[1]=10 means Process 1 will get elements starting from index 10.

### **MPI\_Scatterv**

```
9 if(myrank==0)
           printf("Enter the number of elements:\n");
12
13
14
15
16
17
18
19
20
21
           scanf("%d",&m);
           printf("Initializing array x:\n");
           for(i=0;i<m;i++)
           x[i]=i+1;
           disp[0]=0;
           for(i=0;i<size;i++)</pre>
           z[i]=i+2; // Process θ gets 2 elements, Process 1 gets 3, Process 2 gets 4 and so on
           disp[i+1]=disp[i]+z[i]; // disp[1]=2, disp[2]=5, disp[3]=9 and so on
  MPI_Scatterv(x,z,disp,MPI_INT,y,myrank+2,MPI_INT,0,MPI_COMM_WORLD);
```

```
Enter the number of elements:
                20
                Value of v in process 1:3
OUTPUT:
               Value of y in process 1 : 4
                Value of y in process 1 : 5
                Value of y in process 2:6
                Value of y in process 2 : 7
                Value of y in process 2 : 8
                Value of y in process 2 : 9
                Value of y in process 3 : 10
                Value of v in process 3: 11
                Value of y in process 3 : 12
                Value of y in process 3 : 13
                Value of y in process 3 : 14
                Value of y in process 4: 15
                Value of y in process 4: 16
                Value of y in process 4: 17
                Value of y in process 4: 18
                Value of y in process 4: 19
                Value of y in process 4: 20
                Initializing array x:
                Value of v in process 0 : 1
                Value of y in process 0 : 2
```

### **MPI\_Gatherv**



### **MPI\_Gatherv():**

MPI\_Gatherv(sendbuf, sendcount, datatype, recvbuf, recvcounts, displacements, datatype, root, comm)

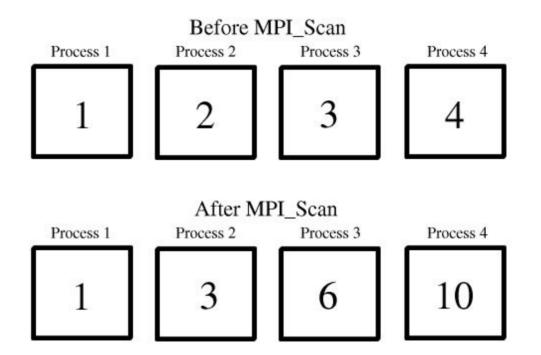
### **Parameters:**

recvcounts: array with number of elements to be received from each process.

displacement: array which holds the beginning index where the data is to be received from each process.

### MPI\_Scan

int MPI\_Scan(sendbuf, recvbuf, int count, datatype, MPI\_Op, comm)



# MPI\_Comm\_Split : Split the communication Domain

MPI\_Comm\_Split(MPI\_Comm comm, int color, int key, MPI\_Comm \*newcomm);

color: controls subset assignment

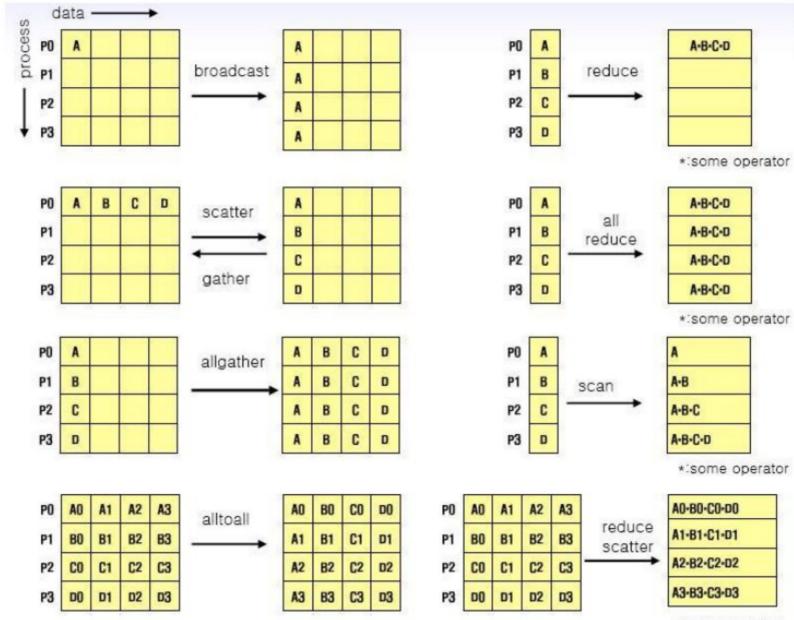
key: controls rank assignment of processes in different group

Ex: MPI\_Comm\_split(MPI\_COMM\_WORLD,0,0,&comm1);

MPI\_Comm\_split(MPI\_COMM\_WORLD,1,0,&comm2);

MPI\_Comm\_split(MPI\_COMM\_WORLD,2,0,&comm3);

### **MPI Collective Routine**



\*:some operator

Reference: Introduction to MPI and OpenMP (with Labs) Brandon Barker Computational Scientist Cornell University Center for Advanced Computing (CAC). https://www.cac.cornell.edu/

### Summary

- MPI provides a simplified way for sending and receiving messages
- MPI rich set of collective functions
- MPI helps for developing Scalable and Portable Parallel Programs
- MPI is the defacto standard for Distributed Memory Parallelism

## Thank You