

# Design and Analysis of Algorithms

NAME : AKSHIT SINGH

SECTION : F

ROLL No: 42

COURSE: B.Tech CSE

Tutorial-4

1.  $T(n) = 3T(n/2) + n^2$

Ans.  $a = 3, b = 2$

$$n^{\log_b a} = n^{\log_2 3}$$

Comparing  $n^{\log_2 3}$  and  $n^2$

$$n^{\log_2 3} < n^2 \quad (\text{Case 3})$$

$\therefore$  according to master's Theorem,

$$T(n) = O(n^2)$$

2.  $T(n) = 4T(n/2) + n^2$

$$a = 4, b = 2$$

$$n^{\log_b a} = n^{\log_2 4} = n^2 = f(n) \quad (\text{Case 2})$$

$\therefore$  according to master's theorem  $T(n) = O(n^2 \log n)$

3.  $T(n) = T(n/2) + 2^n$

$$a = 1, b = 2$$

$$n^{\log_2 1} = n^0 = 1$$

$$1 < 2^n \quad (\text{Case 3})$$

$\therefore$  According to master's theorem  $T(n) = O(2^n)$

4.  $T(n) = 2^n T(n/2) + n^n$

$\therefore$  Master's theorem is not applicable as  $a$  is function.

5.  $T(n) = 16T(n/4) + n$

$a = 16, b = 4, f(n) = n$

$n^{\log_b a} = n^{\log_4 16} = n^2$

$n^2 > P(n)$  (Case 1)

$\therefore T(n) = O(n^2)$

6.  $T(n) = 2T(n/2) + n \log n$

$a = 2, b = 2, f(n) = n \log n$

$n^{\log_b a} = n^{\log_2 2} = n$

Now,  $f(n) > n$

$\therefore$  According to master's  $T(n) = O(n \log n)$

7.  $T(n) = 2T(n/2) + \frac{n}{\log n}$

$a = 2, b = 2, f(n) = \frac{n}{\log n}$

$n^{\log_b a} = n^{\log_2 2} = n$

$n > f(n)$

$\therefore$  According to master's theorem  $T(n) = O(n)$

8.  $T(n) = 2T\left(\frac{n}{4}\right) + n^{0.51}$

$a=2, b=4, f(n) = n^{0.51}$

$n^{\log_b a} = n^{\log_4 2} = n^{0.5}$

$n^{0.5} < P(n)$

$\therefore$  According to master's theorem  $T(n) = O(n^{0.51})$

9.  $T(n) = 0.5T\left(\frac{n}{2}\right) + \frac{1}{n}$

$\therefore$  Master's not applicable as  $a < 1$

10.  $T(n) = 16T\left(\frac{n}{4}\right) + n!$

$a=16, b=4, f(n) = n!$

$n^{\log_b a} = n^{\log_4 16} = n^2$

$n^2 < n!$

$\therefore$  According to master's method,  $T(n) = O(n!)$ .

11.  $T(n) = 4T\left(\frac{n}{2}\right) + \log n$

Ans -  $a=4, b=2$

$f(n) = \log n$

$n^{\log_b a} = n^{\log_2 4} = n^2$

$n^2 \gg f(n)$

$\therefore$  According to master's method,  $T(n) = O(n^2)$

12.  $T(n) = \text{Sqrt}(n) + (n/2) + \log n$

∴ Master's Not applicable as  $a$  is not constant

13.  $T(n) = 3T(n/2) + n$

$a=3, b=2, f(n)=n$

$$n^{\log_b a} = n^{\log_2 3} = n^{1.58}$$

$$n^{\log_b a} > f(n)$$

∴ According to master's Theorem,  $T(n) = O(n^{\log_2 3})$

14.  $T(n) = 3T(n/3) + \sqrt{n}$

Ans  $\Rightarrow a=3, b=3, f(n)=\sqrt{n}$

$$n^{\log_b a} = n^{\log_3 3} = n$$

$$n > \sqrt{n}$$

∴ According to master's theorem,  $T(n) = O(n)$

15.  $T(n) = 4T(n/2) + Cn$

Ans  $\Rightarrow a=4, b=2, f(n)=C \cdot n$

$$n^{\log_b a} = n^{\log_2 4} = n^2$$

$$n^2 > C \cdot n$$

∴ According to master's theorem,  $T(n) = O(n^2)$



16.  $T(n) = 3T(n/4) + n \log n$

Ans  $a=3, b=4, f(n) = n \log n$

$$n^{\log_b a} = n^{\log_4 3} = n^{0.79}$$

$$n^{0.79} < n \log n$$

$\therefore$  According to master's theorem,  $T(n) = O(n \log n)$

17.  $T(n) = 3T(n/3) + n/2$

Ans  $\rightarrow a=3, b=3, f(n) = n/2$

$$n^{\log_b a} = n^{\log_3 3} = n$$

$$O(n) = O(n/2)$$

$\therefore$  According to master's theorem  
 $T(n) = O(n \log n)$

18.  $T(n) = 6T(n/3) + n^2 \log n$

Ans  $\rightarrow a=6, b=3, f(n) = n^2 \log n$

$$n^{\log_b a} = n^{\log_3 6} = n^{1.63}$$

$$n^{1.63} < n^2 \log n$$

$\therefore$  According to master's theorem  $T(n) = O(n^2 \log n)$

19.  $T(n) = 4T(n/2) + n/\log n$

Ans  $a=4, b=2, f(n) = n/\log n$

$$n^{\log_b a} = n^{\log_2 4} = n^2$$

$$n^2 > n/\log n$$

$\therefore$  According to Master's theorem  $T(n) = O(n^2)$

20.  $T(n) = 64T(n/8) - n^2 \log n$

Ans Master's theorem is not applicable as  $f(n)$  is not increasing function.

21.  $T(n) = 7T(n/3) + n^2$

Ans

$$a=7, b=3, f(n) = n^2$$

$$n^{\log_b a} = n^{\log_3 7} = n^{1.7}$$

$$n^{1.7} < n^2$$

$\therefore$  According to master's,  $T(n) = O(n^2)$

22.  $T(n) = T(n/2) + n(2 - \cos n)$

Ans

Master's method isn't applicable since regular condition is violated in Case 3.