

## Ueli's Security Agency

This proof relies on the assumption that our "universe"  $U=\mathbb{Z}$ .

Shared secret: M

Public keys: a,b (coprine)

big prime: p

A = Rp(ma); B = Rp(mb)

We know a, b, p, A and B.

a and b are coprime => gcd(ab)=1 => =u,ve Z(u·a+v·b=1).

Evidenty, a must be positive and b must be negative or vice versa

Without loss of generality, we will assume v to be negative. We now

define u, v ∈ W as the positive integers satisfying u.a. - v.b = 1, where-

by our new v is just the absolute value of our old v.

From the definitions of A and B, are derive:

$$m^{\alpha} \equiv_{P} A$$

Hence, through algebraic manipulation, we get:

 $m^{\alpha} \equiv_{\rho} A \Rightarrow m^{\alpha} \equiv_{\rho} A^{\alpha} \Rightarrow m^{\alpha \cdot \alpha} \equiv_{\rho} A^{\alpha} \Rightarrow_{\rho} A^{\alpha} \Rightarrow_{\rho}$ 

(1)

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 $W_{\rho} \equiv^{b} \beta \Rightarrow (W_{\rho})_{\Lambda} \equiv^{b} \beta_{\Lambda} \Rightarrow W_{\Lambda \rho} \equiv^{b} \beta_{\Lambda}$ 

Pulting @ in O, we derive:  $m \cdot m^{Vb} = _{\rho} A^{u} \wedge m^{Vb} = _{\rho} 8^{V} \Rightarrow m \cdot 8^{V} = _{\rho} A^{u}$ (3) What is left to do is to compute the multiplicative inverse (B') of B' modulo p. As p is prime, B' definitely has such an inverse. Computing (BV) -1 can be done using the extended enclidion algorithm, as: Bv . (Bv)-1 =0 1 ⇒ B v· × = p 1 ( where x=(BV).1) => FLUE Z (W.P=1-B'.x) (we now define was the integer salishing  $\Rightarrow \omega \cdot p + x \cdot B^{V} = 1$ ⇒ gcd (ρ, β) - 1 Chere, the extended euc. alg. can be used to determine the linear combinations of p and Bu that add up to 1. The Pactor that By is multiplied with is equal to its multiplicative inverse) Now, we can use 3 to derive m. Bv. (Bv)-1 = Au. (Bv)-1 m = P Au. (Br)-1 In other words, m = Rp(Au.(BV)-1), as m < p. Answer: m= Rp (A4. (BV)-1), where a and v are positive integers such that wa - vb = 1 and (BV)-1. BV =p 1.