Recursion Part 1

April 6, 2022

Administrative Notes

Project 1! Project 1! Project 1!

- Have I annoyed you yet? Get this done!!

Project 2 will likely be released this weekend; and will be due on Monday, April 25th. Look for announcements

Next week's lectures are virtual

The test is two weeks from today (April 20)

Recursion

Recursion

When a function calls itself, that's called "recursion" or "recursive programming." And it turns out to be a very useful way to solve certain problems. Like, those where certain things have to be done over and over, with only slight differences.

Up until now, when we've had things that needed to be done repeatedly, we used loops. Either for loops, or while loops.

Using loops is called "iteration" or "iterative programming."

Every programming problem that can be solved with recursion can also be solved with iteration!!

So why use recursion? Sometimes it just makes the problem easier to solve

So how does recursion work?

Somebody gives you a task to do.

If the task is easy enough, you just do the task yourself. This is called the "base case"

If the task is too hard to do directly, you offer to do a part of the task, then get somebody else to do the rest of it for you. *This is called the recursive case.*

This is only gonna work if you actually do part of the task yourself, and get other people to do a simpler version of the task. Otherwise we're in an infinite loop.

Recursion in the real world

(Outside of cyberspace that is)

All the real-world examples are lame; I admit that up front. That's one of the things that can make recursion difficult to understand

- If you're the kind of person who learns new concepts by analogies to things you already understand, recursion is a problem.
- There aren't any good analogies.

Nonetheless, we'll try. But I'll warn you the example is lame.

An example: bring me my new laptop!!!

My new laptop is in, but not here yet. Can you bring it to me? It's ten miles away. You have to go on foot and carry it.

Iterative solution:

- Go one mile. Then go another mile. Then go a third mile. Keep doing this until you've run all ten miles.
- Pick up my laptop. Come back one mile. Then come back another mile. Then come back a third mile. Keep doing this until you've brought my laptop back.

It works - I get my laptop! But it's a lot of work.

The recursive version

You offer to go one mile. Then you ask somebody else to go get the laptop, which is 9 miles away.

She goes one mile, then asks somebody else to go get a laptop, 8 miles away.

This repeats until somebody is asked to go get a laptop that's only one mile away. That seems reasonable, so the last person just goes and gets the laptop and brings it one mile back.

Then each other person brings the computer back the one mile they came.

I got my computer. Nobody had to run 20 miles round trip; nobody covered more than two miles!

Notes on this example:

The *base case* was the last person, who agreed that one mile was really not too far to go to get the computer, and she just did it.

Everybody else was part of the *recursive case*. Make the problem a little bit simpler, and then get somebody else to do it. Each person had to make the problem a little simpler; that is, get closer to the laptop. If I asked you to go 10 miles, and you asked your friend to go 10 miles, and she asked her friend to go 10 miles... we're never going to get that computer!!!

This seems fake

It depends on having friends who happen to be one mile closer to the computer and who are willing to help out. What are the odds of that?

Okay, it's kind of a hokey example, but:

- What if those 'friends' are clones copies of you?
- And what if by 'you' we mean a software function that can clone copies of itself at will?

Now the example's not all that hokey, is it?

A software example: computing a factorial

Mathematically, n factorial (written as n!) for any positive integer n is just the product n * (n-1) * (n-2)*...* 1.

We know how to solve it iteratively. Now let's do it recursively

```
\begin{array}{lll} & & & & & \\ \text{def fact(n):} & & & & \\ \text{prod} = 1 & & & \\ \text{for i in range(n,1,-1):} & & & \\ \text{prod *= i} & & & \\ \text{return prod} & & & \\ \end{array}
```

Visualizing recursion

Let's look visually at what's happening when the recursive function is executing:

http://www.pythontutor.com/visualize.html#mode=edit

Another example - the Fibonacci sequence

```
1, 1, 2, 3, 5, 8, 13,...
```

After the first two numbers, each number is the sum of the previous two numbers

```
That is, f(n) = f(n-1) + f(n-2)
```

```
\begin{array}{lll} \textit{Iterative} \\ \textit{def fib(n):} & \textit{Recursive} \\ \textit{def fib(n):} & \textit{def fib(n):} \\ \textit{return n} & \textit{if n < 3:} \\ \textit{Else:} & \textit{return 1} \\ \textit{fib = [1,1]} & \textit{else:} \\ \textit{for i in range(3,n+1)} & \textit{return (fib(n-1) + fib(n-2))} \\ & \textit{fib.append(fib[i-1] + fib[i-2])} \\ \textit{return (fib[n])} \end{array}
```

Tracing the recursive routine

http://www.pythontutor.com/visualize.html#mode=display

How Python Works

Python, like many (but not all) programming languages, uses a stack of frames to keep track of where it is in the program

- What instruction gets executed next
- What variables are in scope and can be accessed
- Where to go when this function finishes

Stack - stack of plates. Put a plate on top ('push'); take a plate off the top ('pop')

LIFO - last in, first out

This determines exactly how the program will work, and what statement will be executed after the current statement is executed

Only the frame on TOP of the stack can be executing!!

A stack is a data structure

A way of organizing data, with a defined set of rules about what can be done.

- Similar to a list or a dictionary, but with different operations

With a stack, you can only get the most recently added item from it.

- You can think of it more literally as a stack of papers where the most recently called functions information can be found on top.
- Here's an illustration

https://www.cs.usfca.edu/~galles/visualization/SimpleStack.html

Frames and the Python Stack

A **frame** is the set of all symbols (variables, constants, function names) currently in scope - currently known to the Python interpreter

When the program starts, it pushes the main program's frame onto the stack, and the main program executes.

When the main program calls a function, Python creates a new frame for that function and pushes that frame onto the stack

Since the function's frame is on top of the stack, that function is now executing

When a Function Ends

When a function ends, its frame is popped off the stack

- "Pop" in this sense means remove the top element from the stack
- All its parameters and local variables disappear
- Program Control returns to the frame that's now on top of the stack that is, whoever called that function!!

When does a function end?

- When a return statement is executed
- When there is no return statement, but all code in the function has been executed

When does it all end?

When the main program's code has been executed, the main program's frame is popped off the stack and that's the end of it. There's no place to return control, so the program is over

- This assumes there were no errors that ended your program prematurely

Back to Recursion - the Fibonacci sequence

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Tracing the recursive routine

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How do you solve a problem using recursion?

- 1. What is the base case? #there may be more than one
- 2. How do I describe a subproblem of my problem
 - a. If I repeat making that subproblem, do I get a base case?
 - b. At this point we should TRUST the recursive calls
- 3. Assuming I have the solution to the subproblem, <u>HOW DO I</u>

 <u>SOLVE MY PROBLEM WITH IT?</u>
 - I should be careful to make sure I'm returning the answer to MY PROBLEM

Palindromes

Calculate whether a string is a palindrome using recursion

What's the base case?

- A string that is zero characters long an empty string IS a palindrome
- A string that is one character long IS a palindrome
- A string where the first character is DIFFERENT from the last character is NOT a palindrome
 - "cat" is NOT a palindrome

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Palindromes - recursive case

What about the recursive case?

- IF the first character is the SAME as the last character, the string IS a palindrome if what's left when you throw away the first and last characters is a palindrome

yay - remove first character; remove last character; look at what's left

- a - IS a palindrome

Pseudocode

x='yay' IS a palindrome

First character equals last character

- Create the substring by throwing away the first and last character
 - x[1:-1] or (x[1:len(x)-1])
- We're left with the string 'a'

y = 'tt' IS a palindrome

Throw away first and last character

- we have an empty string left
- An empty string is a palindrome by definition (base case)

Now, let's write this code

Battab - atta - tt - empty string - IS a palindrome