

# key derivation with easier measurable security

caveman

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hi — i propose *ciphart*, a memory-hard key derivation function that has a security gain that's measurable more objectively and more conveniently than anything in class known to date.

to nail this goal, *ciphart*'s security gain is measured in the unit of *relative entropy bits*. relative to what? relative to the encryption algorithm that's used later on. therefore, this *relative entropy bits* measure is guaranteed to be true when the encryption algorithm that's used with *ciphart* is also the same one that's used to encrypt the data afterwards.

my reference implementation is available here<sup>1</sup>.

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## 1 ciphart

### parameters:

- $W$  each task's size, at least 32 bytes.
- $M$  total memory in multiples of  $2W$ .
- $R$  number of rounds per task.
- $B$  added security in *relative entropy bits*.
- enc encryption function.
- $k$  initial key.

### input:

- $T$   $\leftarrow M/W$
- $P$   $\leftarrow \max(2, \lceil 2^B / (TR) \rceil)$
- $m_{l,s,t}$  memory for  $t^{th}$  task, in  $s^{th}$  segment, in  $l^{th}$  lane to work on.
- $n$   $\leftarrow 0$ , a variable with enough bytes to store nonces in.  $n[0]$  means first 64 bits.  $n[1]$  means second 64 bits.

### output:

- $\hat{k}$  better key, with  $B$ , or more, *relative entropy bits*.

### steps:

```
1 while true do
2   for  $s = 1, 2, \dots, S$  do
3     for  $l = 1, 2, \dots, L$  do
4       for  $t = 1, 2, \dots, T$  do
5         for  $r = 1, 2, \dots, R$  do
6           if  $t = 1$  then
7              $m_{l,s,t} \leftarrow \text{enc}(m_{l,s,T}, n, k)$ ;
8           else
9              $m_{l,s,t} \leftarrow \text{enc}(m_{l,s,t-1}, n, k)$ ;
10             $n[1] \leftarrow n[1] + 1$ ;
11             $k \leftarrow f(m_{l,s,t}[-64:], p, l, s, t)$ ;
12          if  $\log_2(n[0] \times SLTR + n[1]) \geq B$  then
13            go to step 15;
14           $n[0] \leftarrow n[0] + 1$ ;
15 while true do
16    $n[0] \leftarrow n[0] + 1$ ;
17   for  $l = 1, 2, \dots, L$  do
18     if  $\text{len}(\hat{k}) \geq K$  then
19       return  $\hat{k}[0 : K]$ 
20      $\hat{k} \leftarrow \hat{k} \parallel \text{enc}(m_{l,S,T}[1], n, k)$ ;
21      $n[1] \leftarrow n[1] + 1$ ;
```

<sup>1</sup><https://github.com/Al-Caveman/ciphart>

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