

## 0.1 Introduce Substitutions

### 0.1.1 Substitutions as SNOC lists

Definition of  $\sigma$

### 0.1.2 Trivial Properties of substitutions

$\text{fv}(\sigma)$

$$\text{fv}(\diamond) = \emptyset \quad (1)$$

$$\text{fv}(\sigma, x := v) = \text{fv}(\sigma) \cup \text{fv}(v) \quad (2)$$

$\text{dom}(\sigma)$

$$\text{dom}(\diamond) = \emptyset \quad (3)$$

$$\text{dom}(\sigma, x := v) = \text{dom}(\sigma) \cup \{x\} \quad (4)$$

$x \# \sigma$

$$x \# \sigma \Leftrightarrow x \notin (\text{fv}(\sigma) \cup \text{dom}(\sigma)) \quad (5)$$

### 0.1.3 Effect of substitutions

### 0.1.4 Well Formedness

### 0.1.5 Simple Properties Of Substitution

If  $\Gamma' \vdash \sigma : \Gamma$  then: **TODO: Number these**

- $\Gamma'0k$  and  $\Gamma'0k$
- $\omega : \Gamma'' \triangleright \Gamma'$  implies  $\Gamma'' \vdash \sigma : \Gamma$

## 0.2 Substitution Preserves Typing

**TODO: State property** **TODO: Proof by induction over type relation**

## 0.3 Semantics of Substitution

### 0.3.1 Denotation of Substitutions

**TODO: Fill in from p98**

### 0.3.2 Lemma

**TODO: Fill in from p98**

### 0.3.3 Substitution Theorem

**TODO: There is Tikz code here to draw the Substitution Theorem diagram, but it compiles v slowly** If  $\Gamma \vdash t : \tau$  and  $\Gamma' \vdash \sigma : \Gamma$  then

## 0.4 Single Substitution