

A Denotational Semantics for Polymorphic Effect Systems

Part III Project

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Motivating Polymorphic Effect Analysis

TODO: Syntax highlight

```
def logAction(  
    action: Unit => String  
): Unit {  
    log.info(action())  
}
```

```
logAction(() => FireMissiles(); "Launched Missiles)
```

```
logAction(() => throwError("My Error"))
```

```
logAction(() => readEnvironmentVariables)
```

What is a Denotational Semantics?

- A *compositional* mapping
 $\llbracket - \rrbracket : \text{Language Structure} \rightarrow \text{Mathematical Structure}$
- In particular want to define $\llbracket \Gamma \vdash t : A \rrbracket$
- Needs to be *sound* $t_1 \approx t_2 \implies \llbracket t_1 \rrbracket = \llbracket t_2 \rrbracket$
- And *adequate* for our needs $\llbracket t_1 \rrbracket = \llbracket d_2 \rrbracket \implies t_1 \approx t_2$

Contributions

- A sound set of requirements for denotational semantics of effect-polymorphic languages.
- A method to construct models for effect-polymorphic languages in Set.
- A proof of adequacy of such a model.

Denotational Semantics using Category Theory

- Interested in: Objects, Morphisms, and Functors
- Types and type environments are objects (e.g. sets) $\llbracket A \rrbracket, \llbracket \Gamma \rrbracket \in \text{obj } \mathbb{C}$
- Terms are morphisms $(\llbracket \Gamma \vdash t : A \rrbracket : \llbracket \Gamma \rrbracket \rightarrow \llbracket A \rrbracket)$

Language features (1) - Lambda Calculus

A cartesian closed category (CCC) consists of:

- Products $A \times B$ - models tuples
- A terminal object 1 - models the `Unit` type
- Exponential objects B^A - models functions as first-class objects

Language features (2.A) - Monads

A (strong) monad consists of:

- A functor $T : \mathbb{C} \rightarrow \mathbb{C}$
- Join and Unit natural transformations
 - $\mu_A : TTA \rightarrow TA$
 - $\eta_A : A \rightarrow TA$
- Tensor strength natural transformation $\tau_{A,B} : A \times TB \rightarrow T(A \times B)$

Language features (2.B) - Graded Monads

A (strong) graded monad consists of:

- An indexed functor $T_\epsilon : \mathbb{C} \rightarrow \mathbb{C}$
- Indexed Join and Unit natural transformations
 - $\mu_{\epsilon_1, \epsilon_2, A} : T_{\epsilon_1} T_{\epsilon_2} A \rightarrow T_{\epsilon_1 \cdot \epsilon_2} A$
 - $\eta_A : A \rightarrow T_1 A$
- Tensor strength natural transformation $\mathfrak{t}_{\epsilon, A, B} : A \times T_\epsilon B \rightarrow T_\epsilon(A \times B)$

An Effectful Language

$$v ::= k^A \mid x \mid \text{true} \mid \text{false} \mid () \mid \lambda x: A. v \mid v_1 v_2 \mid \text{return } v \\ \mid \text{do } x \leftarrow v_1 \text{ in } v_2 \mid \text{if}_A v \text{ then } v_1 \text{ else } v_2$$
$$A, B, C ::= \gamma \mid A \rightarrow B \mid M_e A$$
$$\text{(Return)} \frac{\Gamma \vdash v: A}{\Gamma \vdash \text{return } v: M_1 A} \quad \text{(Apply)} \frac{\Gamma \vdash v_1: A \rightarrow B \quad \Gamma \vdash v_2: A}{\Gamma \vdash v_1 v_2: B}$$

Semantics of EC

- Can build a model of EC when we have
 - CCC
 - Strong Graded Monad
 - Co-product and Subtyping (morphisms for if-statements)
- We'll call this an S-category

$$\begin{array}{c} \text{(Return)} \frac{f = \llbracket \Gamma \vdash v : A \rrbracket}{\llbracket \Gamma \vdash \text{return } v : M_1 A \rrbracket = \eta_A \circ f} \qquad \text{(Fn)} \frac{f = \llbracket \Gamma, x : A \vdash v : B \rrbracket : \Gamma \times A \rightarrow B}{\llbracket \Gamma \vdash \lambda x : A. v : A \rightarrow B \rrbracket = \text{cur}(f) : \Gamma \rightarrow B^A} \end{array}$$

An Ugly Example

TODO: syntax highlight this

```
let twiceIO =  $\lambda$  action:  $M_{IO}Unit$ . (  
  do _ <- action in action  
)
```

```
let twiceState =  $\lambda$  action:  $M_{State}Unit$ . (  
  do _ <- action in action  
)
```

```
do _ <- twiceState(increment) in twiceIO(writeLog)
```

Let's Add Polymorphism

$$v ::= .. \mid \Lambda \alpha. v \mid v \epsilon$$

$$A, B, C ::= ... \mid \forall \alpha. A$$

$$\epsilon ::= e \mid \alpha \mid \epsilon \cdot \epsilon$$

$$\text{(Effect-Gen)} \frac{\Phi, \alpha \mid \Gamma \vdash v : A}{\Phi \mid \Gamma \vdash \Lambda \alpha. v : \forall \alpha. A}$$

$$\text{(Effect-Spec)} \frac{\Phi \mid \Gamma \vdash v : \forall \alpha. A \quad \Phi \vdash \epsilon}{\Phi \mid \Gamma \vdash v \epsilon : A[\epsilon/\alpha]}$$

An Ugly Example - With a Makeover

TODO: Syntax highlighting

```
let twice =  $\Lambda$  eff.(  
   $\lambda$  action:  $M_{eff}Unit$ . (  
    do _ <- action in action  
  )  
)
```

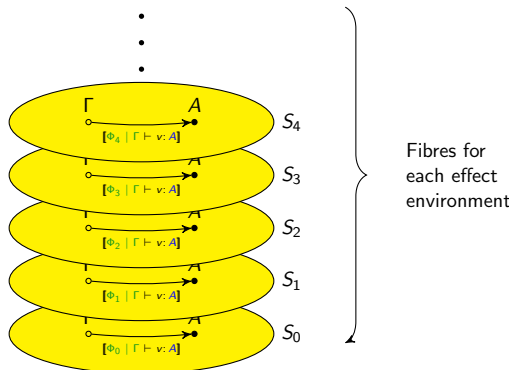
```
do _ <- (twice State increment) in (twice IO writeLog)
```

How do we Model the Semantics of a Polymorphic Language?

- For a fixed effect variable environment Φ and terms with no polymorphic sub-terms, we have EC
- Effect-variable environments of length n are isomorphic by α -equivalence

How do we Model the Semantics of a Polymorphic Language?

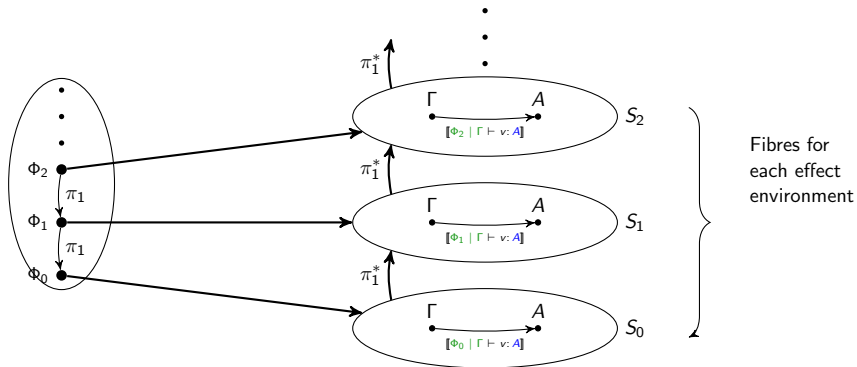
- So we instantiate an S-category for each environment.
- The type rule for quantification requires us to move between categories
TODO: Type rule here.
- Functors are required.



Base Category

- Need to be able to reason about the effect environment categorically
- Model effects and their environments using category \mathbf{Eff} :
 - Objects are products of the set of ground effects $\{*\}, E, E \times E, \dots E^n..$
 - Morphisms are monotone functions
- Represent Φ as an object E^n
- Transformations (e.g. substitutions) between environments become functions

Indexed Category



Quantification

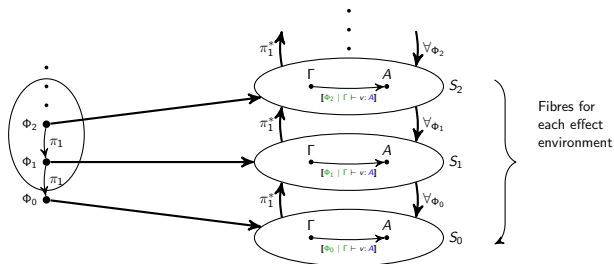
- What about effect-generalisation?

- (Effect-Gen)
$$\frac{\Phi, \alpha \mid \Gamma \vdash v : A}{\Phi \mid \Gamma \vdash \Lambda \alpha. v : \forall \alpha. A}$$

- Need to map $\llbracket \Phi, \alpha \mid \Gamma \vdash v : A \rrbracket$ to $\llbracket \Phi \mid \Gamma \vdash \Lambda \alpha. v : \forall \alpha. A \rrbracket$

- For specialisation to work, needs: $\pi_1^* \dashv \forall_I$

Instantiating a Model (1)



- Can we actually instantiate a category with the required structure?
- Starting point: a model of EC in Set

Instantiating a Model (2) - Fibres

- The fibre $\mathbb{C}(n)$ is the category of functors $[E^n, \text{Set}]$
- I.E. objects are functions that take a vector of ground effects and return a set $\llbracket \Phi \vdash A: \text{Type} \rrbracket : E^n \rightarrow \text{obj}(\text{Set})$.
- Morphisms are dependent functions that return functions in $f : (\vec{\epsilon} : E^n) \rightarrow A\vec{\epsilon} \rightarrow B\vec{\epsilon}$
- These fibres have S-Category features

Instantiating a Model (4) - Functors and Adjunctions

- Re-indexing functors act by pre-composition

$$\begin{aligned}A &\in [E^n, \text{Set}] \\ \theta^*(A)\epsilon_m^{\rightarrow} &= A(\theta(\epsilon_m^{\rightarrow})) \\ \theta^*(f)\epsilon_m^{\rightarrow} &= f(\theta(\epsilon_m^{\rightarrow})) : \theta^*(A) \rightarrow \theta^*(B)\end{aligned}$$

- The quantification functor takes a product over all ground effects

$$\forall_{E^n}(A)\epsilon_n^{\rightarrow} = \prod_{\epsilon \in E} A(\epsilon_n^{\rightarrow}, \epsilon)$$

The End

- Sound: Proved for all indexed S-Categories ✓
- Compositional: By the definition of denotations ✓
- Adequate: Proved for an instantiation in Set ✓

TODO: Dissertation and github links