## Explain that we implicitly carry around a derivation in the denotation

• Denotation for each typing relation derivation

$$- (\text{Unit}) \frac{}{\llbracket \Gamma \vdash () : \text{Unit} \rrbracket_M = \llbracket () \rrbracket_M \circ \langle \rangle_{\Gamma} : \Gamma \to \llbracket \text{Unit} \rrbracket_M} \\ - (\text{Const}) \frac{}{\llbracket \Gamma \vdash \text{C}^A : A \rrbracket_M = \llbracket \text{C}^A \rrbracket_M \circ \langle \rangle_{\Gamma} : \Gamma \to \llbracket \text{A} \rrbracket_M} \\ - (\text{True}) \frac{}{\llbracket \Gamma \vdash \text{true} : \text{Bool} \rrbracket_M = \llbracket \text{true} \rrbracket_M \circ \langle \rangle_{\Gamma} : \Gamma \to \llbracket \text{Bool} \rrbracket_M} \\ - (\text{False}) \frac{}{\llbracket \Gamma \vdash \text{true} : \text{Bool} \rrbracket_M = \llbracket \text{false} \rrbracket_M \circ \langle \rangle_{\Gamma} : \Gamma \to \llbracket \text{Bool} \rrbracket_M} \\ - (\text{Lambda}) \frac{}{\llbracket \Gamma \vdash \text{true} : A \to M_c B : \Gamma \times A \to T_c B} \\ \frac{}{\llbracket \Gamma \vdash \lambda x : A . C : A \to M_c B \rrbracket_M = \text{cur}(f) : \Gamma \to (T_c B)^A} \\ - (\text{Return}) \frac{}{\llbracket \Gamma \vdash \text{return} v : M_1 A \rrbracket_M = \eta_A \circ f} \\ - (\text{Subtype}) \frac{}{\llbracket \Gamma \vdash v : A \rrbracket_M : \Gamma \to Ag = \llbracket A \le : B \rrbracket_M} \\ \frac{}{\llbracket \Gamma \vdash v : B \rrbracket_M = g \circ f : \Gamma \to B} \\ - (\text{Subeffect}) \frac{}{\llbracket \Gamma \vdash v : \text{Bool} \rrbracket_M g = \llbracket \Gamma \vdash C : M_c A \rrbracket_M h = \llbracket \Gamma \vdash C_2 : M_c A \rrbracket_M} \\ \frac{}{\llbracket \Gamma \vdash \text{if}_{e,A} v \text{then} C_1 \text{else} C_2 : M_c A \rrbracket_M = \text{if}_{e,A} \circ \langle f, \langle g, h \rangle \rangle : \Gamma \to T_c A} \\ - (\text{Bind}) \frac{}{\llbracket \Gamma \vdash \text{dox} \leftarrow C_1 \text{in} C_2 : M_{c_1} A : \Gamma \to T_{c_1} A g = \llbracket \Gamma, x : A \vdash C_2 : M_{c_2} B \rrbracket_M \rrbracket_M : \Gamma \times A \to T_{c_2} B} \\ \frac{}{\llbracket \Gamma \vdash \text{dox} \leftarrow C_1 \text{in} C_2 : M_{c_1} A : \Gamma \to T_{c_1} A g = \llbracket \Gamma, x : A \vdash C_2 : M_{c_2} B \rrbracket_M \rrbracket_M : \Gamma \times A \to T_{c_2} B} \\ \frac{}{\llbracket \Gamma \vdash \text{dox} \leftarrow C_1 \text{in} C_2 : M_{c_1} A : \Gamma \to T_{c_1} A g = \llbracket \Gamma, x : A \vdash C_2 : M_{c_2} B \rrbracket_M \rrbracket_M : \Gamma \times A \to T_{c_2} B} \\ \frac{}{\llbracket \Gamma \vdash \text{dox} \leftarrow C_1 \text{in} C_2 : M_{c_1} A : \Gamma \to T_{c_1} A g = \llbracket \Gamma, x : A \vdash C_2 : M_{c_2} B \rrbracket_M \rrbracket_M : \Gamma \times A \to T_{c_2} B} \\ \frac{}{\llbracket \Gamma \vdash \text{dox} \leftarrow C_1 \text{in} C_2 : M_{c_1} A : \Gamma \to T_{c_1} A g = \llbracket \Gamma, x : A \vdash C_2 : M_{c_2} B \rrbracket_M \rrbracket_M : \Gamma \times A \to T_{c_2} B} \\ \frac{}{\llbracket \Gamma \vdash \text{dox} \leftarrow C_1 \text{in} C_2 : M_{c_1} A : \Gamma \to T_{c_1} A g = \llbracket \Gamma, x : A \vdash C_2 : M_{c_2} B \rrbracket_M \rrbracket_M : \Gamma \times A \to T_{c_1} B B} \\ \frac{}{\llbracket \Gamma \vdash \text{dox} \leftarrow C_1 \text{in} C_2 : M_{c_1} A : \Gamma \to T_{c_1} A g = \llbracket \Gamma, x : A \vdash C_2 : M_{c_2} B \rrbracket_M \rrbracket_M : \Gamma \times A \to T_{c_1} B B} \\ \frac{}{\llbracket \Gamma \vdash \text{dox} \leftarrow C_1 \text{in} C_2 : M_{c_1} A : \Gamma \to T_{c_1} A g = \llbracket \Gamma, x : A \vdash C_2 : M_{c_2} B \rrbracket_M : \Gamma \to T_{c_1} C : \Gamma \to T_{c_$$

- Denotations of Types
- - Fill in from notebook
  - For each ground type  $g \in \gamma$
  - morphism  $[\![A\leq:B]\!]_M:[\![A]\!]_M\to [\![B]\!]_M$  for each  $A\leq:B$
  - Natural Transformation  $[\![\epsilon_1 \leq \epsilon_2]\!]_M : T_{\epsilon_1} \to T_{\epsilon_2}$  for each  $\epsilon_1 \leq \epsilon_2$