

## 0.1 Beta and Eta Equivalence

### 0.1.1 Beta-Eta conversions

- (Lambda-Beta) 
$$\frac{\Phi \mid \Gamma, x : A \vdash v_2 : B \quad \Phi \mid \Gamma \vdash v_1 : A}{\Phi \mid \Gamma \vdash (\lambda x : A. v_1) v_2 \approx v_1 [v_2/x] : B}$$
- (Lambda-Eta) 
$$\frac{\Phi \mid \Gamma \vdash v : A \rightarrow B}{\Phi \mid \Gamma \vdash \lambda x : A. (v x) \approx v : A \rightarrow B}$$
- (Left Unit) 
$$\frac{\Phi \mid \Gamma \vdash v_1 : A \quad \Phi \mid \Gamma, x : A \vdash v_2 : \mathbf{M}_\epsilon B}{\Phi \mid \Gamma \vdash \text{do } x \leftarrow \text{return } v_1 \text{ in } v_2 \approx v_2 [v_1/x] : \mathbf{M}_\epsilon B}$$
- (Right Unit) 
$$\frac{\Phi \mid \Gamma \vdash v : \mathbf{M}_\epsilon A}{\Phi \mid \Gamma \vdash \text{do } x \leftarrow v \text{ in return } x \approx v : \mathbf{M}_\epsilon A}$$
- (Associativity) 
$$\frac{\Phi \mid \Gamma \vdash v_1 : \mathbf{M}_{\epsilon_1} A \quad \Phi \mid \Gamma, x : A \vdash v_2 : \mathbf{M}_{\epsilon_2} B \quad \Phi \mid \Gamma, y : B \vdash v_3 : \mathbf{M}_{\epsilon_3} C}{\Phi \mid \Gamma \vdash \text{do } x \leftarrow v_1 \text{ in } (\text{do } y \leftarrow v_2 \text{ in } v_3) \approx \text{do } y \leftarrow (\text{do } x \leftarrow v_1 \text{ in } v_2) \text{ in } v_3 : \mathbf{M}_{\epsilon_1 \cdot \epsilon_2 \cdot \epsilon_3} C}$$
- (Unit) 
$$\frac{\Phi \mid \Gamma \vdash v : \mathbf{Unit}}{\Phi \mid \Gamma \vdash v \approx () : \mathbf{Unit}}$$
- (if-true) 
$$\frac{\Phi \mid \Gamma \vdash v_1 : A \quad \Phi \mid \Gamma \vdash v_2 : A}{\Phi \mid \Gamma \vdash \text{if}_A \text{ true then } v_1 \text{ else } v_2 \approx v_1 : A}$$
- (if-false) 
$$\frac{\Phi \mid \Gamma \vdash v_2 : A \quad \Phi \mid \Gamma \vdash v_1 : A}{\Phi \mid \Gamma \vdash \text{if}_A \text{ false then } v_1 \text{ else } v_2 \approx v_2 : A}$$
- (If-Eta) 
$$\frac{\Phi \mid \Gamma, x : \mathbf{Bool} \vdash v_2 : A \quad \Phi \mid \Gamma \vdash v_1 : \mathbf{Bool}}{\Phi \mid \Gamma \vdash \text{if}_A v_1 \text{ then } v_2 [\text{true}/x] \text{ else } v_2 [\text{false}/x] \approx v_2 [v_1/x] : A}$$
- (Effect-beta) 
$$\frac{\Phi \vdash \epsilon \quad \Phi, \alpha \mid \Gamma \vdash v : A}{\Phi \mid \Gamma \vdash (\Lambda \alpha. v \epsilon) \approx v [\epsilon/\alpha] : A [\epsilon/\alpha]}$$
- (Effect-eta) 
$$\frac{\Phi \mid \Gamma \vdash v : \forall \alpha. A}{\Phi \mid \Gamma \vdash \Lambda \alpha. (v \alpha) \approx v : \forall \alpha. A}$$

### 0.1.2 Equivalence Relation

- (Reflexive) 
$$\frac{\Phi \mid \Gamma \vdash v : A}{\Phi \mid \Gamma \vdash v \approx v : A}$$
- (Symmetric) 
$$\frac{\Phi \mid \Gamma \vdash v_1 \approx v_2 : A}{\Phi \mid \Gamma \vdash v_2 \approx v_1 : A}$$
- (Transitive) 
$$\frac{\Phi \mid \Gamma \vdash v_1 \approx v_2 : A \quad \Phi \mid \Gamma \vdash v_2 \approx v_3 : A}{\Phi \mid \Gamma \vdash v_1 \approx v_3 : A}$$

### 0.1.3 Congruences

- (Effect-Abs) 
$$\frac{\Phi, \alpha \mid \Gamma \vdash v_1 \approx v_2 : A}{\Phi \mid \Gamma \vdash \Lambda \alpha. v_1 \approx \Lambda \alpha. v_2 : \forall \alpha. A}$$
- (Effect-Apply) 
$$\frac{\Phi \mid \Gamma \vdash v_1 \approx v_2 : \forall \alpha. A \quad \Phi \vdash \epsilon}{\Phi \mid \Gamma \vdash v_1 \epsilon \approx v_2 \epsilon : A[\epsilon/\alpha]}$$
- (Lambda) 
$$\frac{\Phi \mid \Gamma, x : A \vdash v_1 \approx v_2 : B}{\Phi \mid \Gamma \vdash \lambda x : A. v_1 \approx \lambda x : A. v_2 : A \rightarrow B}$$
- (Return) 
$$\frac{\Phi \mid \Gamma \vdash v_1 \approx v_2 : A}{\Phi \mid \Gamma \vdash \text{return } v_1 \approx \text{return } v_2 : \mathbf{M}_1 A}$$
- (Apply) 
$$\frac{\Phi \mid \Gamma \vdash v_1 \approx v'_1 : A \rightarrow B \quad \Phi \mid \Gamma \vdash v_2 \approx v'_2 : A}{\Phi \mid \Gamma \vdash v_1 v_2 \approx v'_1 v'_2 : B}$$
- (Bind) 
$$\frac{\Phi \mid \Gamma \vdash v_1 \approx v'_1 : \mathbf{M}_{\epsilon_1} A \quad \Phi \mid \Gamma, x : A \vdash v_2 \approx v'_2 : \mathbf{M}_{\epsilon_2} B}{\Phi \mid \Gamma \vdash \text{do } x \leftarrow v_1 \text{ in } v_2 \approx \text{do } c \leftarrow v'_1 \text{ in } v'_2 : \mathbf{M}_{\epsilon_1 \cdot \epsilon_2} B}$$
- (If) 
$$\frac{\Phi \mid \Gamma \vdash v \approx v' : \text{Bool} \quad \Phi \mid \Gamma \vdash v_1 \approx v'_1 : A \quad \Phi \mid \Gamma \vdash v_2 \approx v'_2 : A}{\Phi \mid \Gamma \vdash \text{if}_A v \text{ then } v_1 \text{ else } v_2 \approx \text{if}_A v \text{ then } v'_1 \text{ else } v'_2 : A}$$
- (Subtype) 
$$\frac{\Phi \mid \Gamma \vdash v \approx v' : A \quad A \leq_\Phi B}{\Phi \mid \Gamma \vdash v \approx v' : B}$$

## 0.2 Beta-Eta Equivalence Implies Both Sides Have the Same Type

If  $\Phi \mid \Gamma \vdash v \approx v' : A$  then each derivation of  $\Phi \mid \Gamma \vdash v \approx v' : A$  can be converted to a derivation of  $\Phi \mid \Gamma \vdash v : A$  and  $\Phi \mid \Gamma \vdash v' : A$  by induction over the beta-eta equivalence relation derivation.

### 0.2.1 Equivalence Relations

**Case Reflexive:** By inversion we have a derivation of  $\Phi \mid \Gamma \vdash v : A$ .

**Case Symmetric:** By inversion  $\Phi \mid \Gamma \vdash v' \approx v : A$ . Hence by induction, derivations of  $\Phi \mid \Gamma \vdash v' : A$  and  $\Phi \mid \Gamma \vdash v : A$  are given.

**Case Transitive:** By inversion, there exists  $v_2$  such that  $\Phi \mid \Gamma \vdash v_1 \approx v_2 : A$  and  $\Phi \mid \Gamma \vdash v_2 \approx v_3 : A$ . Hence by induction, we have derivations of  $\Phi \mid \Gamma \vdash v_1 : A$  and  $\Phi \mid \Gamma \vdash v_3 : A$ .

### 0.2.2 Beta-Eta conversions

**Case Lambda:** By inversion, we have  $\Phi \mid \Gamma, x : A \vdash v_1 : B$  and  $\Phi \mid \Gamma \vdash v_2 : A$ . Hence by the typing rules, we have:

$$\frac{\text{(Lambda)} \frac{\Phi \mid \Gamma, x : A \vdash v_1 : B}{\Phi \mid \Gamma \vdash \lambda x : A. v_1 : A \rightarrow B} \quad \Phi \mid \Gamma \vdash v_2 : A}{\text{(Apply)} \frac{}{\Phi \mid \Gamma \vdash (\lambda x : A. v_1) v_2 : A}}$$

By the substitution rule **TODO: which?**, we have

$$(\text{Substitution}) \frac{\Phi \mid \Gamma, x : A \vdash v_1 : B \quad \Phi \mid \Gamma \vdash v_2 : A}{\Phi \mid \Gamma \vdash v_1 [v_2/x] : B}$$

**Case Left Unit:** By inversion, we have  $\Phi \mid \Gamma \vdash v_1 : A$  and  $\Phi \mid \Gamma, x : A \vdash v_2 : M_\epsilon B$

Hence we have:

$$(\text{Bind}) \frac{(\text{Return}) \frac{\Phi \mid \Gamma \vdash v_1 : A}{\Phi \mid \Gamma \vdash \text{return } v_1 : M_1 A} \quad \Phi \mid \Gamma, x : A \vdash v_2 : M_\epsilon B}{\Phi \mid \Gamma \vdash \text{do } x \leftarrow \text{return } v_1 \text{ in } v_2 : M_{1 \cdot \epsilon} B = M_\epsilon B} \quad (1)$$

And by the substitution typing rule we have: **TODO: Which Rule?**

$$\Phi \mid \Gamma \vdash v_2 [v_1/x] : M_\epsilon B \quad (2)$$

**Case Right Unit:** By inversion, we have  $\Phi \mid \Gamma \vdash v : M_\epsilon A$ .

Hence we have:

$$(\text{Bind}) \frac{\Phi \mid \Gamma \vdash v : M_\epsilon A \quad (\text{Return}) \frac{(\text{var}) \frac{\Phi \vdash \Gamma, x : A \text{Ok}}{\Phi \mid \Gamma, x : A \vdash x : A}}{\Phi \mid \Gamma, x : A \vdash \text{return } v : M_1 A}}{\Phi \mid \Gamma \vdash \text{do } x \leftarrow v \text{ in return } x : M_{\epsilon \cdot 1} A = M_\epsilon A} \quad (3)$$

**Case Associativity:** By inversion, we have  $\Phi \mid \Gamma \vdash v_1 : M_{\epsilon_1} A$ ,  $\Phi \mid \Gamma, x : A \vdash v_2 : M_{\epsilon_2} B$ , and  $\Phi \mid \Gamma, y : B \vdash v_3 : M_{\epsilon_3} C$ .

$$\Phi \vdash (\iota\pi \times) : (\Gamma, x : A, y : B) \triangleright (\Gamma, y : B)$$

So by the weakening property **TODO: which?**,  $\Phi \mid \Gamma, x : A, y : B \vdash v_3 : M_{\epsilon_3} C$

Hence we can construct the type derivations:

$$(\text{Bind}) \frac{\Phi \mid \Gamma \vdash v_1 : M_{\epsilon_1} A \quad (\text{Bind}) \frac{\Phi \mid \Gamma, x : A \vdash v_2 : M_{\epsilon_2} B \quad \Phi \mid \Gamma, x : A, y : B \vdash v_3 : M_{\epsilon_3} C}{\Phi \mid \Gamma, x : A \vdash x v_2 v_3 : M_{\epsilon_2 \cdot \epsilon_3} C}}{\Phi \mid \Gamma \vdash \text{do } x \leftarrow v_1 \text{ in } (\text{do } y \leftarrow v_2 \text{ in } v_3) : M_{\epsilon_1 \cdot \epsilon_2 \cdot \epsilon_3} C} \quad (4)$$

and

$$(\text{Bind}) \frac{(\text{Bind}) \frac{\Phi \mid \Gamma \vdash v_1 : M_{\epsilon_1} A \quad \Phi \mid \Gamma, x : A \vdash v_2 : M_{\epsilon_2} B}{\Phi \mid \Gamma \vdash \text{do } x \leftarrow v_1 \text{ in } v_2 : M_{\epsilon_1 \cdot \epsilon_2} B} \quad \Phi \mid \Gamma, y : B \vdash v_3 : M_{\epsilon_3} C}{\Phi \mid \Gamma \vdash \text{do } y \leftarrow (\text{do } x \leftarrow v_1 \text{ in } v_2) \text{ in } v_3 : M_{\epsilon_1 \cdot \epsilon_2 \cdot \epsilon_3} C} \quad (5)$$

**Case Eta:** By inversion, we have  $\Phi \mid \Gamma \vdash v : A \rightarrow B$

By weakening, we have  $\Phi \vdash \iota\pi : (\Gamma, x : A) \triangleright \Gamma$  Hence, we have

$$\begin{array}{c}
\Phi \mid (\Gamma, x : A) \vdash x : A \quad (\text{weakening}) \frac{\Phi \mid \Gamma \vdash v : A \rightarrow B \quad \Phi \vdash \iota\pi : \Gamma, x : A \triangleright \Gamma}{\Phi \mid \Gamma, x : A \vdash v : A \rightarrow B} \\
(\text{App}) \frac{}{\Phi \mid \Gamma, x : A \vdash v x : B} \\
(\text{Fn}) \frac{}{\Phi \mid \Gamma \vdash \lambda x : A. (v x) : A \rightarrow B}
\end{array} \tag{6}$$

**Case If-True:** By inversion, we have  $\Phi \mid \Gamma \vdash v_1 : A$ ,  $\Phi \mid \Gamma \vdash v_2 : A$ . Hence by the typing lemma **TODO: Which?**, we have  $\Phi \vdash \Gamma \text{Ok}$  so  $\Phi \mid \Gamma \vdash \text{true} : \text{Bool}$  by the axiom typing rule.

Hence

$$(\text{If}) \frac{\Phi \mid \Gamma \vdash \text{true} : \text{Bool} \quad \Phi \mid \Gamma \vdash v_1 : A \quad \Phi \mid \Gamma \vdash v_2 : A}{\Phi \mid \Gamma \vdash \text{if}_A \text{ true then } v_1 \text{ else } v_2 : A} \tag{7}$$

**Case If-False:** As above,

Hence

$$(\text{If}) \frac{\Phi \mid \Gamma \vdash \text{false} : \text{Bool} \quad \Phi \mid \Gamma \vdash v_1 : A \quad \Phi \mid \Gamma \vdash v_2 : A}{\Phi \mid \Gamma \vdash \text{if}_A \text{ false then } v_1 \text{ else } v_2 : A} \tag{8}$$

**Case If-Eta:** By inversion, we have:

$$\Phi \mid \Gamma \vdash v_1 : \text{Bool} \tag{9}$$

and

$$\Phi \mid \Gamma, x : \text{Bool} \vdash v_2 : A \tag{10}$$

Hence we also have  $\Phi \vdash \Gamma \text{Ok}$ . Hence, the following also hold:

$\Phi \mid \Gamma \vdash \text{true} : \text{Bool}$ , and  $\Phi \mid \Gamma \vdash \text{false} : \text{Bool}$ .

Hence by the substitution theorem, we have:

$$(\text{If}) \frac{\Phi \mid \Gamma \vdash v_1 : \text{Bool} \quad \Phi \mid \Gamma \vdash v_2 [\text{true}/x] : A \quad \Phi \mid \Gamma \vdash v_2 [\text{false}/x] : A}{\Phi \mid \Gamma \vdash \text{if}_A v_1 \text{ then } v_2 [\text{true}/x] \text{ else } v_2 [\text{false}/x] : A} \tag{11}$$

and

$$\Phi \mid \Gamma \vdash v_2 [v_1/x] : A \tag{12}$$

**Case Effect-Beta:** By inversion,  $\Phi, \alpha \mid \Gamma \vdash v : A$  and  $\Phi \vdash \epsilon$ .

Then we have the following type derivation:

$$\begin{array}{c}
(\text{Effect-Fn}) \frac{\Phi, \alpha \mid \Gamma \vdash v : A}{\Phi \mid \Gamma \vdash \Lambda \alpha. v : \forall \alpha. A} \quad \Phi \vdash \epsilon \\
(\text{Effect-App}) \frac{}{\Phi \mid \Gamma \vdash \Lambda \alpha. v \epsilon : A [\epsilon/\alpha]}
\end{array} \tag{13}$$

And we can construct the single-effect-substitution:

$$(\text{Single Substitution}) \frac{\Phi \vdash \epsilon}{\Phi \vdash [\epsilon/\alpha] : (\Phi, \alpha)} \tag{14}$$

Hence by the substitution theorem,

$$\Phi \mid \Gamma \vdash v [\epsilon/\alpha] : A [\epsilon/\alpha] \tag{15}$$

**Case Effect-Eta:** By inversion  $\Phi \mid \Gamma \vdash v : \forall \alpha. A$

So the following derivation holds:

$$\begin{array}{c} \text{(Effect-weakening)} \frac{\Phi \mid \Gamma \vdash v : \forall \alpha. A}{\Phi, \alpha \mid \Gamma \vdash v : \forall \alpha. A} \quad \Phi, \alpha \vdash \alpha \\ \text{(Effect-App)} \frac{\quad}{\Phi, \alpha \mid \Gamma \vdash v \ \alpha : A[\alpha/\alpha] = A} \\ \text{(Effect-Fn)} \frac{\quad}{\Phi \mid \Gamma \vdash \Lambda \alpha. (v \ \alpha) : \forall \alpha. A} \end{array} \quad (16)$$

And

$$\Phi \mid \Gamma \vdash v : \forall \alpha. A \quad (17)$$

### 0.2.3 Congruences

Each congruence rule corresponds exactly to a type derivation rule. To convert to a type derivation, convert all preconditions, then use the equivalent type derivation rule.

**Case Lambda:** By inversion,  $\Phi \mid \Gamma, x : A \vdash v_1 \approx v_2 : B$ . Hence by induction  $\Phi \mid \Gamma, x : A \vdash v_1 : B$ , and  $\Phi \mid \Gamma, x : A \vdash v_2 : B$ .

So

$$\Phi \mid \Gamma \vdash \lambda x : A. v_1 : A \rightarrow B \quad (18)$$

and

$$\Phi \mid \Gamma \vdash \lambda x : A. v_2 : A \rightarrow B \quad (19)$$

Hold.

**Case Return:** By inversion,  $\Phi \mid \Gamma \vdash v_1 \approx v_2 : A$ , so by induction

$$\Phi \mid \Gamma \vdash v_1 : A$$

and

$$\Phi \mid \Gamma \vdash v_2 : A$$

Hence we have

$$\Phi \mid \Gamma \vdash \text{return } v_1 : \mathsf{M}_1 A$$

and

$$\Phi \mid \Gamma \vdash \text{return } v_2 : \mathsf{M}_1 A$$

**Case Apply:** By inversion, we have  $\Phi \mid \Gamma \vdash v_1 \approx v'_1 : A \rightarrow B$  and  $\Phi \mid \Gamma \vdash v_2 \approx v'_2 : A$ . Hence we have by induction  $\Phi \mid \Gamma \vdash v_1 : A \rightarrow B$ ,  $\Phi \mid \Gamma \vdash v_2 : A$ ,  $\Phi \mid \Gamma \vdash v'_1 : A \rightarrow B$ , and  $\Phi \mid \Gamma \vdash v'_2 : A$ .

So we have:

$$\Phi \mid \Gamma \vdash v_1 \ v_2 : B \quad (20)$$

and

$$\Phi \mid \Gamma \vdash v'_1 \ v'_2 : B \quad (21)$$

**Case Bind:** By inversion, we have:  $\Phi \mid \Gamma \vdash v_1 \approx v'_1 : \mathbb{M}_{\epsilon_1} A$  and  $\Phi \mid \Gamma, x : A \vdash v_2 \approx v'_2 : \mathbb{M}_{\epsilon_2} B$ . Hence by induction, we have  $\Phi \mid \Gamma \vdash v_1 : \mathbb{M}_{\epsilon_1} A$ ,  $\Phi \mid \Gamma \vdash v'_1 : \mathbb{M}_{\epsilon_1} A$ ,  $\Phi \mid \Gamma, x : A \vdash v_2 : \mathbb{M}_{\epsilon_2} B$ , and  $\Phi \mid \Gamma, x : A \vdash v'_2 : \mathbb{M}_{\epsilon_2} B$

Hence we have

$$\Phi \mid \Gamma \vdash \text{do } x \leftarrow v_1 \text{ in } v_2 : \mathbb{M}_{\epsilon_1, \epsilon_2} A \quad (22)$$

$$\Phi \mid \Gamma \vdash \text{do } x \leftarrow v'_1 \text{ in } v'_2 : \mathbb{M}_{\epsilon_1, \epsilon_2} A \quad (23)$$

**Case If:** By inversion, we have:  $\Phi \mid \Gamma \vdash v \approx v' : \text{Bool}$ ,  $\Phi \mid \Gamma \vdash v_1 \approx v'_1 : A$ , and  $\Phi \mid \Gamma \vdash v_2 \approx v'_2 : A$ .

Hence by induction, we have:

$$\Phi \mid \Gamma \vdash v : \text{Bool}, \Phi \mid \Gamma \vdash v' : \text{Bool},$$

$$\Phi \mid \Gamma \vdash v_1 : A, \Phi \mid \Gamma \vdash v'_1 : A,$$

$$\Phi \mid \Gamma \vdash v_2 : A, \text{ and } \Phi \mid \Gamma \vdash v'_2 : A.$$

So

$$\Phi \mid \Gamma \vdash \text{if}_A v \text{ then } v_1 \text{ else } v_2 : A \quad (24)$$

and

$$\Phi \mid \Gamma \vdash \text{if}_A v' \text{ then } v'_1 \text{ else } v'_2 : A \quad (25)$$

hold.

**Case Subtype:** By inversion, we have  $A \leq_{\Phi} B$  and  $\Phi \mid \Gamma \vdash v \approx v' : A$ . By induction, we therefore have  $\Phi \mid \Gamma \vdash v : A$  and  $\Phi \mid \Gamma \vdash v' : A$ .

Hence we have

$$\Phi \mid \Gamma \vdash v : B \quad (26)$$

$$\Phi \mid \Gamma \vdash v' : B \quad (27)$$

**Case Effect-Lambda:** By inversion,  $\Phi, \alpha \mid \Gamma \vdash v_1 \approx v_2 : A$ . So

$$\text{(Effect-Lambda)} \frac{\Phi, \alpha \mid \Gamma \vdash v_1 : A}{\Phi \mid \Gamma \vdash \Lambda \alpha. v_2 : \forall \alpha. A} \quad (28)$$

and

$$\text{(Effect-Lambda)} \frac{\Phi, \alpha \mid \Gamma \vdash v_2 : A}{\Phi \mid \Gamma \vdash \Lambda \alpha. v_2 : \forall \alpha. A} \quad (29)$$

**Case Effect-Apply:** By inversion,  $\Phi \mid \Gamma \vdash v_1 \approx v_2 : \forall \alpha. A$  and  $\Phi \vdash \epsilon$ .

So

$$\text{(Effect-App)} \frac{\Phi \mid \Gamma \vdash v_1 : \forall \alpha. A \quad \Phi \vdash \epsilon}{\Phi \mid \Gamma \vdash v_1 \epsilon : A [\alpha/\epsilon]} \quad (30)$$

and

$$\text{(Effect-App)} \frac{\Phi \mid \Gamma \vdash v_2 : \forall \alpha. A \quad \Phi \vdash \epsilon}{\Phi \mid \Gamma \vdash v_2 \epsilon : A [\alpha/\epsilon]} \quad (31)$$