## 0.1 CCC

The section should be a cartesian closed category. That is it should have:

- A Terminal object 1
- Binary products
- Exponentials

## 0.2 Graded Pre-Monad

The category should have a graded pre-monad. That is:

- An endofunctor indexed by the po-monad on effects:  $T:(\mathbb{E},\cdot\,\mathbf{1},\leq)\to\mathsf{Cat}(\mathbb{C},\mathbb{C})$
- A unit natural transformation:  $\eta: \mathtt{Id} \to T_{\mathtt{1}}$
- A join natural transformation:  $\mu_{\epsilon_1,\epsilon_2}$ ,  $:T_{\epsilon_1}T_{\epsilon_2}\to T_{\epsilon_1\cdot\epsilon_2}$

Subject to the following commutative diagrams:

#### 0.2.1 Left Unit

$$T_{\epsilon}A \xrightarrow{T_{\epsilon}\eta_{A}} T_{\epsilon}T_{1}A$$

$$\downarrow^{\operatorname{Id}_{T_{\epsilon}A}} \downarrow^{\mu_{\epsilon,1,A}}$$

$$T_{\epsilon}A$$

## 0.2.2 Right Unit

$$T_{\epsilon}A \underbrace{\begin{array}{c} \eta_{T_{\epsilon}A} \\ \text{Id}_{T_{\epsilon}A} \end{array}}_{T_{\epsilon}T_{\epsilon}A} I_{1}T_{1}A$$

### 0.2.3 Associativity

$$\begin{split} T_{\epsilon_{1}}T_{\epsilon_{2}}T_{\epsilon_{3}} \overset{\mu_{\epsilon_{1},\epsilon_{2},T_{\epsilon_{3}}}}{\longrightarrow} \overset{A}{T_{\epsilon_{1}\cdot\epsilon_{2}}}T_{\epsilon_{3}}A \\ \downarrow T_{\epsilon_{1}}\mu_{\epsilon_{2},\epsilon_{3},A} & \downarrow \mu_{\epsilon_{1}\cdot\epsilon_{2},\epsilon_{3},A} \\ T_{\epsilon_{1}}T_{\epsilon_{2}\cdot\epsilon_{3}} \overset{\mu_{\epsilon_{1},\epsilon_{2}\cdot\epsilon_{3}}}{\longrightarrow} T_{\epsilon_{1}\cdot\epsilon_{2}\cdot\epsilon_{3}}A \end{split}$$

# 0.3 Tensor Strength

The category should also have tensorial strength over its products and monads. That is, it should have a natural transformation

$$t_{\epsilon,A,B}: A \times T_{\epsilon}B \to T_{\epsilon}(A \times B)$$

Satisfying the following rules:

#### 0.3.1 Left Naturality

$$A \times T_{\epsilon}B \xrightarrow{\operatorname{Id}_{A} \times T_{\epsilon}f} A \times T_{\epsilon}B'$$

$$\downarrow^{\operatorname{t}_{\epsilon,A,B}} \qquad \downarrow^{\operatorname{t}_{\epsilon,A,B'}}$$

$$T_{\epsilon}(A \times B)^{T_{e}(\operatorname{Id}_{A} \times f)}T_{\epsilon}(A \times B')$$

## 0.3.2 Right Naturality

$$A \times T_{\epsilon}B \xrightarrow{f \times \operatorname{Id}_{T_{\epsilon}B}} A' \times T_{\epsilon}B$$

$$\downarrow \operatorname{t}_{\epsilon,A,B} \qquad \downarrow \operatorname{t}_{\epsilon,A',B}$$

$$T_{\epsilon}(A \times B) \xrightarrow{T_{\epsilon}(f \times \operatorname{Id}_{B})} T_{\epsilon}(A' \times B)$$

## 0.3.3 Unitor Law

$$1 \times T_{\epsilon} A \xrightarrow{\mathbf{t}_{\epsilon,1,A}} T_{\epsilon}(1 \times A)$$

$$\downarrow^{\lambda_{T_{\epsilon}A}} \qquad \downarrow_{T_{\epsilon}(\lambda_{A})} \text{ Where } \lambda : 1 \times \text{Id} \to \text{Id is the left-unitor. } (\lambda = \pi_{2})$$

$$T_{\epsilon}A$$

**Tensor Strength and Projection** Due to the left-unitor law, we can develop a new law for the commutivity of  $\pi_2$  with  $t_{..}$ 

$$\pi_{2,A,B} = \pi_{2,1,B} \circ (\langle \rangle_A \times \mathrm{Id}_B)$$

And  $\pi_{2,1}$  is the left unitor, so by tensorial strength:

$$T_{\epsilon}\pi_{2} \circ \mathsf{t}_{\epsilon,A,B} = T_{\epsilon}\pi_{2,1,B} \circ T_{\epsilon}(\langle \rangle_{A} \times \mathsf{Id}_{B}) \circ \mathsf{t}_{\epsilon,A,B}$$

$$= T_{\epsilon}\pi_{2,1,B} \circ \mathsf{t}_{\epsilon,1,B} \circ (\langle \rangle_{A} \times \mathsf{Id}_{B})$$

$$= \pi_{2,1,B} \circ (\langle \rangle_{A} \times \mathsf{Id}_{B})$$

$$= \pi_{2}$$

$$(1)$$

So the following commutes:

$$A \times T_{\epsilon}B \xrightarrow{\mathbf{t}_{\epsilon,A,B}} T_{\epsilon}(A \times B)$$

$$\downarrow^{T_{\epsilon}\pi_{2}}$$

$$\downarrow^{T_{\epsilon}\pi_{2}}$$

$$T_{\epsilon}B$$

#### 0.3.4 Commutativity with Join

$$A \times T_{\epsilon_1} T_{\epsilon_2} B \xrightarrow{\mathbf{t}_{\epsilon_1,A,T_{\epsilon_2}}} T_{\epsilon_1} (A \times T_{\epsilon_2} B) \xrightarrow{T_{\epsilon_1} \mathbf{t}_{\epsilon_2,A,B}} T_{\epsilon_1} T_{\epsilon_2} (A \times B) \xrightarrow{\mu_{\epsilon_1,\epsilon_2,A \times B}} A \times T_{\epsilon_1 \cdot \epsilon_2} B \xrightarrow{\mathbf{t}_{\epsilon_1 \cdot \epsilon_2,A,B}} T_{\epsilon_1 \cdot \epsilon_2} (A \times B)$$

## 0.4 Commutativity with Unit

$$\begin{array}{c} A \times B \xrightarrow{\operatorname{Id}_A \times \eta_B} A \times T_{\epsilon}B \\ & & \downarrow^{\operatorname{t}_{\epsilon,A,B}} \\ & & & T_{\epsilon}(A \times B) \end{array}$$

## **0.5** Commutativity with $\alpha$

Let 
$$\alpha_{A,B,C} = \langle \pi_1 \circ \pi_1, \langle \pi_2 \circ \pi_1, \pi_2 \rangle \rangle : ((A \times B) \times C) \to (A \times (B \times C))$$

$$(A \times B) \times T_{\epsilon}C \xrightarrow{\mathsf{t}_{\epsilon,(A \times B),C}} T_{\epsilon}((A \times B) \times C)$$

$$\downarrow^{\alpha_{A,B,T_{\epsilon}C}} \downarrow^{T_{\epsilon}\alpha_{A,B,C}} TODO: Needed?$$

$$A \times (B \times T_{\epsilon}C) \xrightarrow{\mathsf{d}_{A} \times \mathsf{t}_{\epsilon,B},C} A \times T_{\epsilon}(B \times C) \xrightarrow{\mathsf{t}_{\epsilon,A,(B \times C)}} T_{\epsilon}(A \times (B \times C))$$

# 0.6 Subeffecting

For each instance of the pre-order  $(\mathbb{E}, \leq)$ ,  $\epsilon_1 \leq \epsilon_2$ , there exists a natural transformation  $[\epsilon_1 \leq \epsilon_2]: T_{\epsilon_1} \to T_{\epsilon_2}$  that commutes with  $t_{\cdot,:}$ :

#### 0.6.1 Subeffecting and Tensor Strength

$$\begin{array}{c} A \times T_{\epsilon_1} B \overset{\operatorname{Id}_A \times \llbracket \epsilon_1 \leq \epsilon_2 \rrbracket}{\longrightarrow} A \times T_{\epsilon_2} B \\ \qquad \qquad \qquad \downarrow^{\operatorname{t}_{\epsilon_1,A,B}} \qquad \qquad \downarrow^{\operatorname{t}_{\epsilon_2,A,B}} \\ T_{\epsilon_1} (A \times B) \overset{\llbracket \epsilon_1 \leq \epsilon_2 \rrbracket}{\longrightarrow} T_{\epsilon_2} (A \times B) \end{array}$$

# 0.7 Subtyping

The denotation of ground types  $\llbracket \_ \rrbracket_M$  is a functor from the pre-order category of ground types  $(\gamma, \leq :_{\gamma})$  to  $\mathbb{C}$ . This pre-ordered sub-category of  $\mathbb{C}$  is extended with the rule for function subtyping to form a larger pre-ordered sub-category of  $\mathbb{C}$ .

(Function Subtyping) 
$$\frac{f = [\![A' \le : A]\!]_M}{rhs = [\![A \to M_{\epsilon_1} B \le : A' \to M_{\epsilon_2} B']\!]_M : (T_{\epsilon_1} B)^A \to (T_{\epsilon_2} B')^{A'}}$$

$$rhs = (h_{B'} \circ T_{\epsilon_1} g)^A \circ (T_{\epsilon_1} B)^f$$

$$= \operatorname{cur}(h_{B'} \circ T_{\epsilon_1} g \circ \operatorname{app}) \circ \operatorname{cur}(\operatorname{app} \circ (\operatorname{Id} \times f))$$
(2)

#### 0.8 If natural transformation

There exists a natural transformation  $\mathbf{If}_A : (\mathtt{Bool} \times (A \times A)) \to A$  Satisfying the following:

• If 
$$A \circ \langle \llbracket \mathsf{true} \rrbracket_M \circ \langle \rangle_{\Gamma}, \langle t, f \rangle \rangle = t$$

• If 
$$_{A}\circ\langle \llbracket \mathtt{false} \rrbracket_{M}\circ\langle \rangle_{\Gamma}\,, \langle t,f \rangle \rangle = f$$