$[0] \le :$ 

0.1 CCC

The section should be a cartesian closed category. That is it should have:

• A Terminal object 1

• Binary products

• Exponentials

Further more, it should have a co-product of the terminal object 1. This is required for the beta-eta equivalence of if-then-else terms.

$$\mathbf{1} \xrightarrow{inl} A \xleftarrow{inr} \mathbf{1}$$

For each:

$$\mathbf{1} \stackrel{f}{\longrightarrow} A \stackrel{g}{\longleftarrow} \mathbf{1}$$

There exists unique  $[f,g]: 1+1 \to A$  such that:

$$\begin{array}{c}
A \\
f[f,g] \uparrow \\
1 \xrightarrow{\text{inl}} 1 + 1 \xleftarrow{\text{inr}} 1
\end{array}$$

### 0.2 Graded Pre-Monad

The category should have a graded pre-monad. That is:

• An endo-functor indexed by the po-monoid on effects:  $T:(\mathbb{E},\cdot 1,) \to \mathtt{Cat}(\mathbb{C},\mathbb{C})$ 

• A unit natural transformation:  $\eta: \mathtt{Id} \to T_1$ 

• A join natural transformation:  $\mu_{\epsilon_1,\epsilon_2}$ , :  $T_{\epsilon_1}T_{\epsilon_2} \to T_{\epsilon_1\cdot\epsilon_2}$ 

Subject to the following commutative diagrams:

### 0.2.1 Left Unit

$$T_{\epsilon}A \xrightarrow{T_{\epsilon}\eta_{A}} T_{\epsilon}T_{1}A$$

$$\downarrow^{\operatorname{Id}_{T_{\epsilon}A}} \downarrow^{\mu_{\epsilon,1,A}}$$

$$T_{\epsilon}A$$

## 0.2.2 Right Unit

$$T_{\epsilon}A \xrightarrow{\eta_{T_{\epsilon}A}} T_{1}T_{\epsilon}A$$

$$\downarrow^{\operatorname{Id}_{T_{\epsilon}A}} \downarrow^{\mu_{1,\epsilon,A}}$$

$$T_{\epsilon}A$$

### 0.2.3 Associativity

$$T_{\epsilon_{1}}T_{\epsilon_{2}}T_{\epsilon_{3}}A \xrightarrow{\mu_{\epsilon_{1},\epsilon_{2},T_{\epsilon_{3}}}A} T_{\epsilon_{1}\cdot\epsilon_{2}}T_{\epsilon_{3}}A \qquad 1$$

$$\downarrow T_{\epsilon_{1}}\mu_{\epsilon_{2},\epsilon_{3},A} \qquad \downarrow \mu_{\epsilon_{1}\cdot\epsilon_{2},\epsilon_{3},A}$$

$$T_{\epsilon_{1}}T_{\epsilon_{2}\cdot\epsilon_{3}}A \xrightarrow{\mu_{\epsilon_{1},\epsilon_{2}\cdot\epsilon_{3}}A} T_{\epsilon_{1}\cdot\epsilon_{2}\cdot\epsilon_{3}}A$$

#### 0.3 Tensor Strength

### 0.3.3 Unitor Law

$$1 \times T_{\epsilon} A \xrightarrow{\mathfrak{t}_{\epsilon,1,A}} T_{\epsilon}(1 \times A)$$

$$\downarrow^{\lambda_{T_{\epsilon}A}} \qquad \downarrow_{T_{\epsilon}(\lambda_{A})} \text{ Where } \lambda : 1 \times \text{Id} \to \text{Id is the left-unitor. } (\lambda = \pi_{2})$$

$$T_{\epsilon}A$$

**Tensor Strength and Projection** Due to the left-unitor law, we can develop a new law for the commutativity of  $\pi_2$  with  $t_{...}$ 

$$\pi_{2,A,B} = \pi_{2,1,B} \circ (\langle \rangle_A \times \mathrm{Id}_B)$$

And  $\pi_{2,\pmb{1}}$  is the left unitor, so by tensorial strength:

$$\begin{split} T_{\epsilon}\pi_{2} \circ \mathbf{t}_{\epsilon,A,B} &= T_{\epsilon}\pi_{2,1,B} \circ T_{\epsilon}(\langle \rangle_{A} \times \mathrm{Id}_{B}) \circ \mathbf{t}_{\epsilon,A,B} \\ &= T_{\epsilon}\pi_{2,1,B} \circ \mathbf{t}_{\epsilon,1,B} \circ (\langle \rangle_{A} \times \mathrm{Id}_{B}) \\ &= \pi_{2,1,B} \circ (\langle \rangle_{A} \times \mathrm{Id}_{B}) \\ &= \pi_{2} \end{split} \tag{1}$$

So the following commutes:

$$A \times T_{\epsilon}B \xrightarrow{\mathbf{t}_{\epsilon,A,B}} T_{\epsilon}(A \times B)$$

$$\uparrow^{T_{\epsilon}\pi_{2}}$$

$$\uparrow^{T_{\epsilon}\pi_{2}}$$

$$\uparrow^{T_{\epsilon}B}$$

### 0.3.4 Commutativity with Join

$$A \times T_{\epsilon_1} T_{\epsilon_2} B \xrightarrow{\mathbf{t}_{\epsilon_1,A,T_{\epsilon_2}} B} T_{\epsilon_1} (A \times T_{\epsilon_2} B) \xrightarrow{\mathbf{T}_{\epsilon_1} \mathbf{t}_{\epsilon_2,A,B}} T_{\epsilon_1} T_{\epsilon_2} (A \times B)$$

$$\downarrow \mu_{\epsilon_1,\epsilon_2,A \times B}$$

$$A \times T_{\epsilon_1 \cdot \epsilon_2} B \xrightarrow{\mathbf{t}_{\epsilon_1 \cdot \epsilon_2,A,B}} T_{\epsilon_1 \cdot \epsilon_2} (A \times B)$$

## 0.4 Commutativity with Unit

$$\begin{array}{c} A\times B \xrightarrow{\operatorname{Id}_A\times \eta_B} A\times T_1B \\ & & \downarrow^{\operatorname{t}_{1,A,B}} \\ & & & \downarrow^{\operatorname{t}_{1,A,B}} \end{array}$$

## 0.5 Commutativity with $\alpha$

Let 
$$\alpha_{A,B,C} = \langle \pi_1 \circ \pi_1, \langle \pi_2 \circ \pi_1, \pi_2 \rangle \rangle : ((A \times B) \times C) \to (A \times (B \times C))$$

$$(A \times B) \times T_{\epsilon}C \xrightarrow{\mathbf{t}_{\epsilon,(A \times B),C}} T_{\epsilon}((A \times B) \times C)$$

$$\downarrow^{\alpha_{A,B,T_{\epsilon}C}} \downarrow^{T_{\epsilon}\alpha_{A,B,C}}$$

$$A \times (B \times T_{\epsilon}C) \xrightarrow{\mathbf{t}_{\epsilon,A},(B \times C)} A \times T_{\epsilon}(B \times C) \xrightarrow{\mathbf{t}_{\epsilon,A,(B \times C)}} T_{\epsilon}(A \times (B \times C))$$

## 0.6 Sub-Effecting

For each instance of the pre-order  $(\mathbb{E},)$ ,  $\epsilon_1\epsilon_2$ , there exists a natural transformation  $[\![\epsilon_1\epsilon_2]\!]:T_{\epsilon_1}\to T_{\epsilon_2}$  that commutes with  $t_{.,:}$ 

### 0.6.1 Sub-Effecting and Tensor Strength

$$\begin{array}{c} A \times T_{\epsilon_1} B^{\operatorname{\mathsf{Id}}_A \times \llbracket \epsilon_1 \epsilon_2 \rrbracket} A \times T_{\epsilon_2} B \\ \downarrow^{\operatorname{\mathsf{t}}_{\epsilon_1,A,B}} & \downarrow^{\operatorname{\mathsf{t}}_{\epsilon_2,A,B}} \\ T_{\epsilon_1} (A \times B) \xrightarrow{\llbracket \epsilon_1 \epsilon_2 \rrbracket} A \times B \end{array}$$

### 0.6.2 Sub-effecting and Monadic Join

Since the monoid operation on effects is monotone, we can introduce the following diagram.

$$\begin{split} T_{\epsilon_{1}}T_{\epsilon_{2}} & \xrightarrow{T_{\epsilon_{1}} \llbracket \epsilon_{2}\epsilon_{2}' \rrbracket} T_{\epsilon_{1}}T_{\epsilon_{2}'} \xrightarrow{\llbracket \epsilon_{1}\epsilon_{1}' \rrbracket_{M,T_{\epsilon_{2}'}}} T_{\epsilon_{1}'}T_{\epsilon_{2}'} \\ & \downarrow^{\mu_{\epsilon_{1},\epsilon_{2},}} & \downarrow^{\mu_{\epsilon_{1}',\epsilon_{2}',\epsilon_{2}'}} T_{\epsilon_{1}'\cdot\epsilon_{2}'} \\ T_{\epsilon_{1}\cdot\epsilon_{2}} & \xrightarrow{\llbracket \epsilon_{1}\cdot\epsilon_{2}\epsilon_{1}'\cdot\epsilon_{2}' \rrbracket} T_{\epsilon_{1}'\cdot\epsilon_{2}'} \end{split}$$

# 0.7 Sub-typing

The denotation of ground types  $\llbracket \_ \rrbracket$  is a functor from the pre-order category of ground types  $(\gamma, \gamma)$  to  $\mathbb C$ . This pre-ordered sub-category of  $\mathbb C$  is extended with the rule for function sub-typing to form a larger pre-ordered sub-category of  $\mathbb C$ .

$$(\text{Function Subtyping}) \frac{f = \llbracket A'A \rrbracket \ \ g = \llbracket BB' \rrbracket \ \ h = \llbracket \epsilon_1 \epsilon_2 \rrbracket}{rhs = \llbracket A \to \mathsf{M}_{\epsilon_1} BA' \to \mathsf{M}_{\epsilon_2} B' \rrbracket : (T_{\epsilon_1} B)^A \to (T_{\epsilon_2} B')^{A'}}$$

$$\begin{split} rhs = & (h_{B'} \circ T_{\epsilon_1} g)^{A'} \circ (T_{\epsilon_1} B)^f \\ = & \operatorname{cur}(h_{B'} \circ T_{\epsilon_1} g \circ \operatorname{app}) \circ \operatorname{cur}(\operatorname{app} \circ (\operatorname{Id}_{T_{\epsilon_1} B^{A'}} \times f)) \end{split} \tag{2}$$