



Figure 1: Our pick for `Mul` is natural and is a monoid at each n .

0.1 Base Category is Monoidal

We construct the base category, **Eff**, as follows:

- $U = E$, the set of ground effects in the non-polymorphic language.
- 1 is a singleton set.
- $U^n = E^n$, set of n -wide tuples of effects, \vec{e}

Hence when we treat effects that are well formed in Φ as morphisms, $E^n \rightarrow E$ in **Eff**, we should treat them as functions $f : E^n \rightarrow E$. Ground effects become point functions: $e : 1 \rightarrow E$, so the denotation of a ground effect is the constant value function: $\llbracket \Phi \vdash e : \text{Effect} \rrbracket = \vec{e} \mapsto e$. We extend the multiplication of ground effects to multiplication on effect functions, giving us our Mul_n operation $\text{Mul}_n(f, g)(\vec{e}) = (f\vec{e}) \cdot (g\vec{e})$. This Mul_n satisfies the naturality and monoidal requirements, as seen in figure 1. It is also trivially the case that $\text{Mul}_n(\llbracket \Phi \vdash e_1 : \text{Effect} \rrbracket, \llbracket \Phi \vdash e_2 : \text{Effect} \rrbracket) = \llbracket \Phi \vdash e_1 \cdot e_2 : \text{Effect} \rrbracket$.

0.2 S-Categories

The semantic category, $[E^0, \mathbf{Set}]$ of the effect-environment \diamond is isomorphic to \mathbf{Set} . Since each effect-environment is alpha equivalent to a natural number, the semantic category for Φ shall be represented as $\mathbb{C}(\Phi) = \mathbb{C}(n) = [E^n, \mathbf{Set}]$, the category of functions $E^n \rightarrow \mathbf{Set}$. Objects in $[E^n, \mathbf{Set}]$ are functions and we describe them by their actions on a generic vector of ground effects, $\vec{\epsilon}$. Morphisms in $[E^n, \mathbf{Set}]$ are natural transformations between the functions. So:

$$\begin{aligned} m &: A \rightarrow B \quad \text{In } [E^n, \mathbf{Set}] \\ m\vec{\epsilon} &: A\vec{\epsilon} \rightarrow B\vec{\epsilon} \quad \text{In } \mathbf{Set} \\ (f \circ g)\vec{\epsilon} &= (f\vec{\epsilon}) \circ (g\vec{\epsilon}) \\ \text{Id}_A(\vec{\epsilon}) &= \text{Id}_{A\vec{\epsilon}} \end{aligned}$$

So morphisms are dependently typed functions from a vector of ground effects to morphisms in \mathbf{Set} .

0.2.1 Each S-Category is a CCC

Since \mathbf{Set} is complete and a CCC, and E^n is small, since E is small, $[E^n, \mathbf{Set}]$ is a CCC.

$$\begin{aligned} (A \times B)\vec{\epsilon} &= (A\vec{\epsilon}) \times (B\vec{\epsilon}) \\ 1\vec{\epsilon} &= 1 \\ (B^A)\vec{\epsilon} &= (B\vec{\epsilon})^{(A\vec{\epsilon})} \\ \pi_1\vec{\epsilon} &= \pi_1 \\ \pi_2\vec{\epsilon} &= \pi_2 \\ \text{app}\vec{\epsilon} &= \text{app} \\ \text{cur}(f)\vec{\epsilon} &= \text{cur}(f\vec{\epsilon}) \\ \langle f, g \rangle\vec{\epsilon} &= \langle f\vec{\epsilon}, g\vec{\epsilon} \rangle \end{aligned}$$

0.2.2 The Terminal Co-Product

We can define the co-product point-wise.

$$\begin{aligned} (1 + 1)\vec{\epsilon} &= (1\vec{\epsilon} + 1\vec{\epsilon}) \\ &= (1 + 1) \\ \text{inl}\vec{\epsilon} &= \text{inl} \\ \text{inr}\vec{\epsilon} &= \text{inr} \\ [f, g]\vec{\epsilon} &= [f\vec{\epsilon}, g\vec{\epsilon}] \end{aligned}$$

This preserves the co-product diagram.

$$\begin{aligned}
([f, g] \circ \text{inl})\vec{e} &= [f\vec{e}, g\vec{e}] \circ \text{inl} \\
&= f\vec{e} \\
&\square \\
([f, g] \circ \text{inr})\vec{e} &= [f\vec{e}, g\vec{e}] \circ \text{inr} \\
&= g\vec{e} \\
&\square
\end{aligned}$$

$[f, g]$ is also unique in $[E^n, \mathbf{Set}]$. Suppose $l \circ \text{inl} = f$ and $l \circ \text{inr} = g$ in $[E^n, \mathbf{Set}]$. Then $l\vec{e} \circ \text{inl} = f\vec{e}$ and $l\vec{e} \circ \text{inr} = g\vec{e}$. Hence by the co-product in \mathbf{Set} , $l = [f\vec{e}, g\vec{e}]$ so $l = [f, g]$.

0.2.3 Ground Types and Terms

Each ground type in the non-polymorphic calculus has a fixed denotation $\llbracket \gamma \rrbracket \in \mathbf{obj} \ \mathbf{Set}$. The ground type in the polymorphic calculus hence has a denotation represented by the constant function.

$$\begin{aligned}
\llbracket \gamma \rrbracket : \quad E^n &\rightarrow \mathbf{obj} \ \mathbf{Set} \\
\vec{e} &\mapsto \llbracket \gamma \rrbracket
\end{aligned}$$

Each constant term C^A in the non-polymorphic calculus has a fixed denotation $\llbracket C^A \rrbracket \in \mathbf{Set}(1, A)$. So the morphism $\llbracket C^A \rrbracket$ in $[E^n, \mathbf{Set}]$ is the corresponding constant dependently typed morphism returning the $\llbracket C^A \rrbracket$ function in \mathbf{Set} .

$$\begin{aligned}
\llbracket C^A \rrbracket : \quad [E^n, \mathbf{Set}](1, A) \\
\vec{e} &\mapsto \llbracket C^A \rrbracket
\end{aligned}$$

0.2.4 Graded Monad

Given the strong graded monad $(T^0, \eta^0, \mu^0, \mathfrak{t}^0)$ on \mathbf{Set} , we can construct an appropriate graded monad $(T^n, \eta^n, \mu^n, \mathfrak{t}^n)$ on $[E^n, \mathbf{Set}]$. Through some mechanical proof and the naturality of the \mathbf{Set} strong graded monad, these morphisms are natural in their type parameters and form a strong graded monad in $[E^n, \mathbf{Set}]$.

$$\begin{aligned}
T^n : \quad (E^n, \cdot, \leq_n, 1_n) &\rightarrow [[E^n, \mathbf{Set}], [E^n, \mathbf{Set}]] \\
(T_f^n A)\vec{e} &= T_{(f\vec{e})}^0 A\vec{e} \\
(\eta_A^n)\vec{e} &= \eta_{A\vec{e}}^0 \\
(\mu_{f,g,A}^n)\vec{e} &= \mu_{(f\vec{e}), (g\vec{e}), (A\vec{e})}^0 \\
(\mathfrak{t}_{f,A,B}^n)\vec{e} &= \mathfrak{t}_{(f\vec{e}), (A\vec{e}), (B\vec{e})}^0
\end{aligned}$$

Naturality of the Graded Monad Transformations			
$ \begin{array}{ccc} A\vec{\epsilon} & \xrightarrow{\eta_{(A\vec{\epsilon})}^0} & T_1^0(A\vec{\epsilon}) \\ \downarrow f\vec{\epsilon} & & \downarrow T_1^0(f\vec{\epsilon}) \\ B\vec{\epsilon} & \xrightarrow{\eta_{(B\vec{\epsilon})}^0} & T_1^0(B\vec{\epsilon}) \end{array} $	$ \begin{array}{ccc} T_{(f\vec{\epsilon})}^0 T_{(g\vec{\epsilon})}^0 (A\vec{\epsilon}) & \xrightarrow{\mu_{f\vec{\epsilon}, g\vec{\epsilon}, (B\vec{\epsilon})}^0} & T_{(f\vec{\epsilon}) \cdot (g\vec{\epsilon})}^0 (A\vec{\epsilon}) \\ \downarrow T_{f\vec{\epsilon}}^0 T_{g\vec{\epsilon}}^0 m\vec{\epsilon} & & \downarrow T_{(f\vec{\epsilon}) \cdot (g\vec{\epsilon})}^0 m\vec{\epsilon} \\ T_{(f\vec{\epsilon})}^0 T_{(g\vec{\epsilon})}^0 (B\vec{\epsilon}) & \xrightarrow{\mu_{f\vec{\epsilon}, g\vec{\epsilon}, (B\vec{\epsilon})}^0} & T_{(f\vec{\epsilon}) \cdot (g\vec{\epsilon})}^0 (A\vec{\epsilon}) \end{array} $		
$ \begin{array}{ccc} A\vec{\epsilon} \times T_{f\vec{\epsilon}}^0(B\vec{\epsilon}) & \xrightarrow{\tau_{f\vec{\epsilon}, (A\vec{\epsilon}), (B\vec{\epsilon})}^0} & T_{f\vec{\epsilon}}^0(A\vec{\epsilon} \times B\vec{\epsilon}) \\ \downarrow (m\vec{\epsilon} \times \text{Id}_{T_{f\vec{\epsilon}}^0 B}) & & \downarrow T_{(f\vec{\epsilon})}^0(m\vec{\epsilon} \times \text{Id}_{B\vec{\epsilon}}) \\ A'\vec{\epsilon} \times T_{f\vec{\epsilon}}^0(B\vec{\epsilon}) & \xrightarrow{\tau_{f\vec{\epsilon}, (A'\vec{\epsilon}), (B\vec{\epsilon})}^0} & T_{f\vec{\epsilon}}^0(A'\vec{\epsilon} \times B\vec{\epsilon}) \end{array} $	$ \begin{array}{ccc} A\vec{\epsilon} \times T_{f\vec{\epsilon}}^0(B\vec{\epsilon}) & \xrightarrow{\tau_{f\vec{\epsilon}, (A\vec{\epsilon}), (B\vec{\epsilon})}^0} & T_{f\vec{\epsilon}}^0(A\vec{\epsilon} \times B\vec{\epsilon}) \\ \downarrow (\text{Id}_{A\vec{\epsilon}} \times T_{f\vec{\epsilon}}^0(m\vec{\epsilon})) & & \downarrow T_{(f\vec{\epsilon})}^0(\text{Id}_{A\vec{\epsilon}} \times m\vec{\epsilon}) \\ A\vec{\epsilon} \times T_{f\vec{\epsilon}}^0(B'\vec{\epsilon}) & \xrightarrow{\tau_{f\vec{\epsilon}, (A\vec{\epsilon}), (B'\vec{\epsilon})}^0} & T_{f\vec{\epsilon}}^0(A\vec{\epsilon} \times B'\vec{\epsilon}) \end{array} $		

Figure 2: Naturality Squares for the graded monad natural transformations. The naturality of each transformation α in $[E^n, \mathbf{Set}]$ derives from the naturality of each of its component transformations $\alpha\vec{\epsilon}$ in \mathbf{Set}

Left Unit	Monad Laws	Right Unit
$ \begin{aligned} (\mu_{1,g,A}^n \circ \eta_{T_f^n A}^n)\vec{\epsilon} &= \mu_{1,(f\vec{\epsilon}), (A\vec{\epsilon})}^0 \circ (\eta_{T_{f\vec{\epsilon}}^0 A\vec{\epsilon}}^0) \\ &= \text{Id}_{T_{f\vec{\epsilon}}^0 A\vec{\epsilon}} \\ &= (\text{Id}_{T_f^n A})\vec{\epsilon} \end{aligned} $	$ \begin{aligned} (\mu_{f,1,A}^n \circ T_f^n \eta_A^n)\vec{\epsilon} &= \mu_{(f\vec{\epsilon}), 1, (A\vec{\epsilon})}^0 \circ T_{f\vec{\epsilon}}^0(\eta_{A\vec{\epsilon}}^0) \\ &= \text{Id}_{T_{f\vec{\epsilon}}^0 A\vec{\epsilon}} \\ &= (\text{Id}_{T_f^n A})\vec{\epsilon} \end{aligned} $	
Monad Associativity		
$ \begin{aligned} ((\mu_{f,(g\cdot h),A}^n) \circ T_f^n(\mu_{g,h,A}^n))\vec{\epsilon} &= \mu_{(f\vec{\epsilon}), ((g\vec{\epsilon}) \cdot (h\vec{\epsilon})), (A\vec{\epsilon})}^0 \circ T_{f\vec{\epsilon}}^0 \mu_{(h\vec{\epsilon}), (g\vec{\epsilon}), A\vec{\epsilon}}^0 \\ &= \mu_{((f\vec{\epsilon}) \cdot (g\vec{\epsilon})), (h\vec{\epsilon}), (A\vec{\epsilon})}^0 \circ \mu_{(f\vec{\epsilon}), (g\vec{\epsilon}), (T_{h\vec{\epsilon}}^0(A\vec{\epsilon}))}^0 \\ &= (\mu_{f \cdot g, h, A}^n \circ \mu_{f,g, T_h^n A}^n)\vec{\epsilon} \end{aligned} $		

Figure 3: The monad laws for the graded monad $(T^n, \eta^n, \mu^n, \tau^n)$ can be proved component-wise from the graded monad (T^0) on \mathbf{Set} .

Tensor Strength Laws

Bind Law

$$\begin{array}{ccc}
 A \times \mathsf{T}_f^n \mathsf{T}_g^n B & \xrightarrow{\mathfrak{t}_{f,A,\mathsf{T}_g^n B}} \mathsf{T}_f^n(A \times \mathsf{T}_g^n B) & \xrightarrow{\mathfrak{T}_f^n \mathfrak{t}_{g,A,B}} \mathsf{T}_f^n \mathsf{T}_g^n(A \times B) \\
 & \searrow \text{Id}_A \times \mu_{f,g,B}^n & \downarrow \mu_{f,g,A \times B}^n \\
 & A \times \mathsf{T}_{f \cdot g}^n B & \xrightarrow{\mathfrak{t}_{f \cdot g,A,B}} \mathsf{T}_{f \cdot g}^n(A \times B)
 \end{array}$$

$$\begin{aligned}
 & (\mathfrak{t}_{(f \cdot g),A,B}^n \circ (\text{Id}_A \times \mu_{f,g,B}^n)) \bar{\epsilon} \\
 &= (\mathfrak{t}_{((f \bar{\epsilon}), (g \bar{\epsilon})), (A \bar{\epsilon}), (B \bar{\epsilon})}^0 \circ (\text{Id}_{A \bar{\epsilon}} \times \mu_{(f \bar{\epsilon}), (g \bar{\epsilon}), (B \bar{\epsilon})}^n)) \\
 &= \mu_{(f \bar{\epsilon}), (g \bar{\epsilon}), (A \times B) \bar{\epsilon}}^0 \circ \mathsf{T}_{f \bar{\epsilon}}^0(\mathfrak{t}_{(g \bar{\epsilon}), (A \bar{\epsilon}), (B \bar{\epsilon})}^0) \circ \mathfrak{t}_{(f \bar{\epsilon}), (A \bar{\epsilon}), \mathsf{T}_{g \bar{\epsilon}}^0(B \bar{\epsilon})}^0 \\
 &= (\mu_{f,g,(A \times B)}^n \circ \mathsf{T}_f^n(\mathfrak{t}_{g,A,B}^n) \circ \mathfrak{t}_{f,A,\mathsf{T}_g^n(B)}^n) \bar{\epsilon}
 \end{aligned}$$

Commutativity with Unit

$$\begin{array}{ccc}
 A \times B & \xrightarrow{\text{Id}_A \times \eta_B} A \times \mathsf{T}_1 B & \\
 & \searrow \eta_{A \times B} & \downarrow \mathfrak{t}_{1,A,B} \\
 & \mathsf{T}_1^n(A \times B) &
 \end{array}$$

$$\begin{aligned}
 (\mathfrak{t}_{1,A,B}^n \circ (\text{Id}_A \times \eta_A^n)) \bar{\epsilon} &= \mathfrak{t}_{1,(A \bar{\epsilon}), (B \bar{\epsilon})}^0 \circ (\text{Id}_{A \bar{\epsilon}} \times \eta_{A \bar{\epsilon}}^0) \\
 &= \eta_{A \bar{\epsilon} \times B \bar{\epsilon}}^0 \\
 &= (\eta_{A \times B}^n) \bar{\epsilon}
 \end{aligned}$$

Commutativity with α

Let $\alpha_{A,B,C} = \langle \pi_1 \circ \pi_1, \langle \pi_2 \circ \pi_1, \pi_2 \rangle \rangle : ((A \times B) \times C) \rightarrow (A \times (B \times C))$

$$\begin{array}{ccc}
 (A \times B) \times \mathsf{T}_\epsilon^n C & \xrightarrow{\mathfrak{t}_{\epsilon,(A \times B),C}} \mathsf{T}_\epsilon^n((A \times B) \times C) & \\
 \downarrow \alpha_{A,B,\mathsf{T}_\epsilon^n C} & & \downarrow \mathsf{T}_\epsilon^n \alpha_{A,B,C} \\
 A \times (B \times \mathsf{T}_\epsilon^n C) & \xrightarrow{\text{Id}_A \times \mathfrak{t}_{\epsilon,B,C}} A \times \mathsf{T}_\epsilon^n(B \times C) & \xrightarrow{\mathfrak{t}_{\epsilon,A,(B \times C)}} \mathsf{T}_\epsilon^n(A \times (B \times C))
 \end{array}$$

$$\begin{aligned}
 & (\mathsf{T}_f^n \alpha_{A,B,C} \circ \mathfrak{t}_{f,A \times B,C}^n) \bar{\epsilon} \\
 &= \mathsf{T}_{f \bar{\epsilon}}^0 \alpha_{A \bar{\epsilon}, B \bar{\epsilon}, C \bar{\epsilon}} \circ \mathfrak{t}_{(f \bar{\epsilon}), (A \times B) \bar{\epsilon}, (C \bar{\epsilon})}^0 \\
 &= \mathfrak{t}_{(f \bar{\epsilon}), (A \bar{\epsilon}), (B \bar{\epsilon} \times C \bar{\epsilon})}^0 \circ (\text{Id}_{A \bar{\epsilon}} \times \mathfrak{t}_{(f \bar{\epsilon}), (B \bar{\epsilon}), (C \bar{\epsilon})}^0) \circ \alpha_{A \bar{\epsilon}, B \bar{\epsilon}, C \bar{\epsilon}} \\
 &= (\mathfrak{t}_{f,A,(B \times C)}^n \circ (\text{Id}_A \times \mathfrak{t}_{f,B,C}^n) \circ \alpha_{A,B,C}) \bar{\epsilon}
 \end{aligned}$$

Figure 4: The tensor strength laws can be proved component wise from the strength of the monad T^0

Subeffect Natural Transformation

$$\begin{aligned} \llbracket f \leq_n g \rrbracket &: \mathbf{T}_f^n \rightarrow \mathbf{T}_g^n \\ \llbracket f \leq_n g \rrbracket A\vec{\epsilon} &: \mathbf{T}_{f\vec{\epsilon}}^n(A\vec{\epsilon}) \rightarrow \mathbf{T}_{g\vec{\epsilon}}^n(B\vec{\epsilon}) \\ &= \llbracket f\vec{\epsilon} \leq_0 g\vec{\epsilon} \rrbracket A\vec{\epsilon} \end{aligned}$$

Figure 5: Definition of subeffecting natural transformations

0.2.5 Subeffecting

Given a collection of subeffecting natural transformation in \mathbf{Set} , $\llbracket \epsilon_1 \leq_0 \epsilon_2 \rrbracket : \mathbf{T}_{\epsilon_1}^0 \rightarrow \mathbf{T}_{\epsilon_2}^0$ we can form subeffect natural transformations in $[E^n, \mathbf{Set}]$ as seen in figure 5. This natural transformation has all the required properties. These are proved in figure 6.

0.3 Re-indexing Functors

For a function $\theta : E^m \rightarrow E^n$, the re-indexing functor θ^* is defined as follows:

$$\begin{aligned} \theta^* &: [E^n, \mathbf{Set}] \rightarrow [E^m, \mathbf{Set}] \\ \theta^*(A)\vec{\epsilon}_m &= A(\theta(\vec{\epsilon}_m)) \\ f &: A \rightarrow B \in [E^n, \mathbf{Set}] \\ \theta^*(f)\vec{\epsilon}_m &= f(\theta(\vec{\epsilon}_m)) : A(\theta(\vec{\epsilon}_m)) \rightarrow B(\theta(\vec{\epsilon}_m)) \end{aligned}$$

This functor preserves all the S-category properties. These can be seen in figures 7-10

0.3.1 Quantification

We need to define $\forall_{E^n} : [E^{n+1}, \mathbf{Set}] \rightarrow [E^n, \mathbf{Set}]$

So

$$\begin{aligned} (\forall_{E^n} A)\vec{\epsilon}_n &= \prod_{\epsilon \in E} A(\vec{\epsilon}_n, \epsilon) \\ m &: A \rightarrow B \\ (\forall_{E^n} m) &: \forall_{E^n} A \rightarrow \forall_{E^n} B \\ (\forall_{E^n} m)\vec{\epsilon}_n &= \prod_{\epsilon \in E} m(\vec{\epsilon}_n, \epsilon) \end{aligned}$$

Subeffect Natural Transformation Properties

Naturality

$$\begin{array}{ccc} T_{f\vec{e}}^0 A\vec{e} & \xrightarrow{\llbracket f\vec{e} \leq_0 g\vec{e} \rrbracket A\vec{e}} & T_{g\vec{e}}^0 A\vec{e} \\ \downarrow T_{f\vec{e}}^0 m\vec{e} & & \downarrow T_{g\vec{e}}^0 m\vec{e} \\ T_{f\vec{e}}^0 B\vec{e} & \xrightarrow{\llbracket f\vec{e} \leq_0 g\vec{e} \rrbracket B\vec{e}} & T_{g\vec{e}}^0 B\vec{e} \end{array}$$

Commutes With Tensor Strength

$$\begin{array}{ccc} A \times T_f^n B & \xrightarrow{\text{Id}_A \times \llbracket f \leq_n g \rrbracket_B} & A \times T_g^n B \\ \downarrow \tau_{f,A,B}^n & & \downarrow \tau_{g,A,B}^n \\ T_f^n(A \times B) & \xrightarrow{\llbracket f \leq_n g \rrbracket_{A \times B}} & T_g^n(A \times B) \end{array}$$

$$\begin{aligned} & (\tau_{g,A,B}^n \circ (\text{Id}_A \times \llbracket f \leq_n g \rrbracket_B))\vec{e} \\ &= \tau_{(g\vec{e}), (A\vec{e}), (B\vec{e})}^0 \circ (\text{Id}_{A\vec{e}} \times \llbracket f\vec{e} \leq_0 g\vec{e} \rrbracket_{B\vec{e}}) \\ &= \llbracket f\vec{e} \leq_0 g\vec{e} \rrbracket_{(A \times B)\vec{e}} \circ \tau_{(f\vec{e}), (A\vec{e}), (B\vec{e})}^0 \\ &= (\llbracket f \leq_n g \rrbracket_{(A \times B)} \circ \tau_{f,A,B}^n)\vec{e} \end{aligned}$$

Commutes with Join

$$\begin{array}{ccc} T_f^n T_g^n & \xrightarrow{\tau_f^n \llbracket g \leq_n g' \rrbracket} T_f^n T_{g'}^n & \xrightarrow{\llbracket f \leq_n f' \rrbracket_{M, T_{g'}^n}} T_{f'}^n T_{g'}^n \\ \downarrow \mu_{f,g}^n & & \downarrow \mu_{f',g'}^n \\ T_{f \cdot g}^n & \xrightarrow{\llbracket f \cdot g \leq_n f' \cdot g' \rrbracket} & T_{f' \cdot g'}^n \end{array}$$

$$\begin{aligned} & (\llbracket f \cdot g \leq_n f' \cdot g' \rrbracket_A \circ \mu_{f,g,A}^n)\vec{e} \\ &= \llbracket (f\vec{e}) \cdot (g\vec{e}) \leq_0 (f'\vec{e}) \cdot (g'\vec{e}) \rrbracket_{A\vec{e}} \circ \mu_{(f\vec{e}), (g\vec{e}), (A\vec{e})}^0 \\ &= \mu_{(f\vec{e}), (g\vec{e}), (A\vec{e})}^0 \circ \llbracket f\vec{e} \leq_0 f'\vec{e} \rrbracket_{T_{g\vec{e}}^0(A\vec{e})} \circ T_{f\vec{e}}^0 \llbracket g\vec{e} \leq_0 g'\vec{e} \rrbracket_{(A\vec{e})} \\ &= \mu_{f,g,A}^n \circ \llbracket f \leq_n f' \rrbracket_{T_{g,A}^n} \circ T_f^n \llbracket g \leq_n g' \rrbracket_A \end{aligned}$$

Figure 6: The required properties of the subeffect natural transformations can be proved component-wise from the appropriate of the subeffect natural transformation on **Set**.

θ^* is Cartesian Closed	
$ \begin{aligned} (\theta^*(A \times B))\vec{\epsilon} &= (A \times B)(\theta\vec{\epsilon}) \\ &= (A(\theta\vec{\epsilon}) \times B(\theta\vec{\epsilon})) \\ &= (\theta^*A \times \theta^*B)\vec{\epsilon} \end{aligned} $	$ \begin{aligned} (\theta^*\pi_1)\vec{\epsilon} &= \pi_1(\theta\vec{\epsilon}) \\ &= \pi_1 \quad \text{Constant function} \\ &= \pi_1\vec{\epsilon} \end{aligned} $
$ \begin{aligned} (\theta^*\pi_2)\vec{\epsilon} &= \pi_2(\theta\vec{\epsilon}) \\ &= \pi_2 \quad \text{Constant function} \\ &= \pi_2\vec{\epsilon} \end{aligned} $	$ \begin{aligned} (\theta^*\langle f, g \rangle)\vec{\epsilon} &= (\langle f, g \rangle)(\theta\vec{\epsilon}) \\ &= \langle f(\theta\vec{\epsilon}), g(\theta\vec{\epsilon}) \rangle \\ &= \langle \theta^*f, \theta^*g \rangle\vec{\epsilon} \end{aligned} $
$ \begin{aligned} (\theta^*(A^B))\vec{\epsilon} &= (A^B)(\theta\vec{\epsilon}) \\ &= (A(\theta\vec{\epsilon}))^{(B(\theta\vec{\epsilon}))} \\ &= (\theta^*A)^{(\theta^*B)}\vec{\epsilon} \end{aligned} $	$ \begin{aligned} (\theta^*\text{app})\vec{\epsilon} &= \text{app}(\theta\vec{\epsilon}) \\ &= \text{app} \quad \text{Constant fn} \\ &= \text{app}\vec{\epsilon} \end{aligned} $
$ \begin{aligned} (\theta^*\text{cur}(f))\vec{\epsilon} &= \text{cur}(f)(\theta\vec{\epsilon}) \\ &= \text{cur}(f(\theta\vec{\epsilon})) \\ &= \text{cur}(\theta^*f) \end{aligned} $	$ \begin{aligned} (\theta^*1)\vec{\epsilon} &= 1(\theta\vec{\epsilon}) \\ &= 1 \\ &= 1\vec{\epsilon} \end{aligned} $
$ \begin{aligned} (\theta^*\langle \rangle_A)\vec{\epsilon} &= \langle \rangle_A(\theta\vec{\epsilon}) \\ &= \langle \rangle_{A(\theta\vec{\epsilon})} \\ &= \langle \rangle_{\theta^*A}\vec{\epsilon} \end{aligned} $	

Figure 7: Proof of the CCC-preserving property of re-indexing functors.

θ^* Preserves Co-products	
$ \begin{aligned} (\theta^*(1 + 1))\vec{\epsilon} &= (1 + 1)(\theta\vec{\epsilon}) \\ &= (1 + 1) \quad \text{Constant function} \\ &= (1 + 1)\vec{\epsilon} \end{aligned} $	$ \begin{aligned} (\theta^*\text{inl})\vec{\epsilon} &= \text{inl}(\theta\vec{\epsilon}) \\ &= \text{inl} \quad \text{Constant Fn} \\ &= \text{inl}\vec{\epsilon} \end{aligned} $
$ \begin{aligned} (\theta^*\text{inr})\vec{\epsilon} &= \text{inr}(\theta\vec{\epsilon}) \\ &= \text{inr} \quad \text{Constant Fn} \\ &= \text{inr}\vec{\epsilon} \end{aligned} $	$ \begin{aligned} (\theta^*[f, g])\vec{\epsilon} &= [f, g](\theta\vec{\epsilon}) \\ &= [f(\theta\vec{\epsilon}), g(\theta\vec{\epsilon})] \\ &= [\theta^*f, \theta^*g]\vec{\epsilon} \end{aligned} $

Figure 8: Proof that re-indexing functors preserve the $1 + 1$ co-product.

θ^* Preserves the Graded Monad

$$\begin{aligned} (\theta^* T_f^n A) \vec{\epsilon} &= T_f^n A(\theta \vec{\epsilon}) \\ &= T_{(f(\theta \vec{\epsilon}))}^0(A(\theta \vec{\epsilon})) \\ &= (T_{(f \circ \theta)}^m \theta^* A) \vec{\epsilon} \end{aligned}$$

$$\begin{aligned} (\theta^* \eta_A^n) \vec{\epsilon} &= \eta_A^n(\theta \vec{\epsilon}) \\ &= \eta_{A(\theta \vec{\epsilon})}^0 \\ &= \eta_{\theta^* A}^m \vec{\epsilon} \end{aligned}$$

$$\begin{aligned} (\theta^* \mu_{f,g,A}^n) \vec{\epsilon} &= \mu_{f,g,A}^n(\theta \vec{\epsilon}) \\ &= \mu_{f(\theta \vec{\epsilon}), g(\theta \vec{\epsilon}), A(\theta \vec{\epsilon})}^0 \\ &= \mu_{f \circ \theta, g \circ \theta, \theta^*(A)}^m(\vec{\epsilon}) \end{aligned}$$

$$\begin{aligned} (\theta^* \mathfrak{t}_{f,A,B}^n) \vec{\epsilon} &= \mathfrak{t}_{f,A,B}^n(\theta \vec{\epsilon}) \\ &= \mathfrak{t}_{(f(\theta \vec{\epsilon}), (A(\theta \vec{\epsilon})), (B(\theta \vec{\epsilon})))}^0 \\ &= \mathfrak{t}_{f \circ \theta, \theta^* A, \theta^* B}^m \vec{\epsilon} \end{aligned}$$

Figure 9: Re-indexing functors preserve the graded monad structure

θ^* preserves Ground Subtyping and Subeffecting

$$\begin{aligned} \theta^* (\llbracket A \leq_\gamma B \rrbracket) \vec{\epsilon} &= \llbracket A \leq_\gamma B \rrbracket(\theta \vec{\epsilon}) \\ &= \llbracket A \leq B \rrbracket \quad \text{Constant Function} \\ &= \llbracket A \leq B \rrbracket \vec{\epsilon} \end{aligned}$$

$$\begin{aligned} (\theta^* (\llbracket f \leq_n g \rrbracket A)) \vec{\epsilon} &= (\llbracket f \leq_n g \rrbracket A)(\theta \vec{\epsilon}) \\ &= (\llbracket f(\theta \vec{\epsilon}) \leq_n g(\theta \vec{\epsilon}) \rrbracket (A(\theta \vec{\epsilon}))) \\ &= (\llbracket \theta^* f \leq_m \theta^* g \rrbracket (\theta^* A)) \vec{\epsilon} \end{aligned}$$

Figure 10: Re-indexing functors preserve the ground subtyping and subeffecting morphisms.