0.1 Effect Weakening Definition

Introduce a relation $\omega : \Phi' \triangleright \Phi$ relating effect-environments.

0.1.1 Relation

- (Id) $\dfrac{\Phi \mathtt{Ok}}{\iota : \Phi \triangleright \Phi}$
- (Project) $\dfrac{\omega : \Phi' \triangleright \Phi}{\omega\pi : (\Phi', \alpha) \triangleright \Phi}$
- (Extend) $\dfrac{\omega : \Phi' \triangleright \Phi}{\omega\times : (\Phi', \alpha) \triangleright (\Phi, \alpha)}$

0.1.2 Weakening Properties

0.1.3 Effect Weakening Preserves Ok

$$\omega : \Phi' \triangleright \Phi \wedge \Phi \mathtt{Ok} \Leftarrow \Phi' \mathtt{Ok} \tag{1}$$

Proof

Case: ι

$$\Phi \mathtt{Ok} \wedge \iota : \Phi \triangleright \Phi \Leftarrow \Phi \mathtt{Ok}$$

Case: $\omega\pi$ By inversion,

$$\omega : \Phi' \triangleright \Phi \wedge \alpha \notin \Phi' \tag{2}$$

So, by induction, $\Phi' \mathtt{Ok}$ and hence $(\Phi', \alpha)\mathtt{Ok}$

Case: $\omega\times$ By inversion,

$$\omega : \Phi' \triangleright \Phi \wedge \alpha \notin \Phi' \tag{3}$$

So

$$\begin{aligned} (\Phi, \alpha)\mathtt{Ok} &\Rightarrow \Phi \mathtt{Ok} && (4) \\ &\Rightarrow \Phi' \mathtt{Ok} && (5) \\ &\Rightarrow (\Phi', \alpha)\mathtt{Ok} && (6) \\ & && (7) \end{aligned}$$

0.1.4 Domain Lemma

$$\omega : \Phi' \triangleright \Phi \Rightarrow (\alpha \notin \Phi \Rightarrow \alpha \notin \Phi')$$

Proof By trivial Induction.

0.1.5 Weakening Preserves Effect Well-Formed-Ness

If $\omega : \Phi' \triangleright \Phi$ then $\Phi \vdash \epsilon \implies \Phi' \vdash \epsilon$

Proof By induction over the well-formed-ness of effects

Case Ground By inversion, $\Phi \texttt{Ok} \wedge \epsilon \in E$. Hence by the ok-property, $\Phi' \texttt{Ok}$ So $\Phi' \vdash \epsilon$

Case Var $\Phi = \Phi'', \alpha$

So either:

Case: $\Phi' = \Phi''', \alpha$ So $\omega = \omega' \times$ So $\omega' : \Phi''' \triangleright \Phi''$, and hence:

$$(\text{Var})\frac{\Phi''', \alpha \texttt{Ok}}{\Phi''', \alpha \vdash \alpha} \tag{8}$$

Case: $\Phi' = \Phi''', \beta$ and $\beta \neq \alpha$

So $\omega = \omega' \pi$

By induction, $\omega' : \Phi''' \triangleright \Phi$ so

$$(\text{Weaken})\frac{\Phi''' \vdash \alpha}{\Phi' \vdash \alpha} \tag{9}$$

Case Weaken By inversion, $\Phi = \Phi'', \beta$.

So $\omega = \omega' \times$

And, $\Phi' = \Phi''', \beta$ So By inversion$\omega' : \Phi''' \triangleright \pi_1''$

So by induction

$$(\text{weak})\frac{\Phi''' \vdash \alpha}{\Phi' \vdash \alpha} \tag{10}$$

Case Monoid By inversion, $\Phi \vdash \epsilon_1$ and $\Phi \vdash \epsilon_2$. So by induction, $\Phi' \vdash \epsilon_1$ and $\Phi' \vdash \epsilon_2$, and so:

$$\Phi' \vdash \epsilon_1 \cdot \epsilon_2 \tag{11}$$

0.1.6 Weakening Preserves Type-Well-Formed-Ness

If $\omega : \Phi' \triangleright \Phi$ and $\Phi \vdash A$ then $\Phi' \vdash A$.

Proof:

Case Ground: By inversion, $\Phi \texttt{Ok}$, hence by property 1 of weakening, $\Phi' \texttt{Ok}$. Hence $\Phi' \vdash \gamma$.

Case Function: By inversion, $\Phi \vdash A$, $\Phi \vdash B$. So by induction $\Phi' \vdash A$, $\Phi' \vdash B$, hence,

$$\Phi' \vdash A \to B$$

Case Computation: By inversion$\Phi \vdash A$, and $\Phi \vdash \epsilon$.

So by induction and the effect-well-formed-ness theorem,

$\Phi' \vdash A$ and $\Phi' \vdash \epsilon$

So

$$\Phi' \vdash \mathtt{M}_\epsilon A$$

Case For All: By inversion, $\Phi, \alpha \vdash A$ Picking $\alpha \notin \Phi'$ using α-conversion.

So $\omega \times : (\Phi', \alpha) \triangleright (\Phi, \alpha)$

So $(\Phi', \alpha) \vdash A$

So $\Phi \vdash \forall \alpha. A$

0.1.7 Corollary

$$\omega : \Phi' \triangleright \Phi \wedge \Phi \vdash \Gamma \mathtt{Ok} \implies \Phi' \vdash \Gamma \mathtt{Ok}$$

Case Nil: By inversion$\Phi \mathtt{Ok}$ so $\Phi \vdash \diamond \mathtt{Ok}$

Case Var: By inversion$\Phi \vdash \Gamma \mathtt{Ok}$, $x \in \mathtt{dom}(\Gamma)$, $\Phi \vdash A$

So by induction $\Phi' \vdash \Gamma \mathtt{Ok}$, and $\pi_1' \vdash \Gamma \mathtt{Ok}$

So $\Phi' \vdash (\Gamma, x : A) \mathtt{Ok}$

0.1.8 Effect Weakening preserves Type Relations

$$\Phi \mid \Gamma \vdash v \colon A \wedge \omega : \Phi' \triangleright \Phi \implies \Phi' \mid \Gamma \vdash v \colon A \tag{12}$$

Proof:

Case Constants: If $\Phi \vdash \Gamma \mathtt{Ok}$ then $\Phi' \vdash \Gamma \mathtt{Ok}$ so:

$$(\text{Const}) \frac{\Phi' \vdash \Gamma \mathtt{Ok}}{\Phi' \mid \Gamma \vdash \mathtt{C}^A \colon A} \tag{13}$$

Case Variables: If $\Phi \vdash \Gamma \mathtt{Ok}$ then $\Phi' \vdash \Gamma \mathtt{Ok}$ so: So, $\Phi' \mid G \vdash x \colon A$, if $\Phi \mid G \vdash x \colon A$

Case Lambda: By inversion, $\Phi \mid \Gamma, x : A \vdash v \colon B$, so by induction $\Phi' \mid \Gamma, x : A \vdash v \colon B$.

So,

$$\Phi' \mid \Gamma \vdash \lambda x : A.v : A \to B \tag{14}$$

Case Apply: By inversion$\Phi \mid \Gamma \vdash v_1 \colon A \to B$ and $\Phi \mid \Gamma \vdash v_2 \colon A$.

Hence by induction, $\Phi' \mid \Gamma \vdash v_1 \colon A \to B$ and $\Phi' \mid \Gamma \vdash v_2 \colon A$.

So

$$\Phi' \mid \Gamma \vdash \mathtt{app} v_1 v_2 \colon B$$

Case Return: By inversion$\Phi \mid \Gamma \vdash v\colon A$

So by induction $\Phi' \mid \Gamma \vdash v\colon A$

Hence $\Phi' \mid \Gamma \vdash \texttt{return}\ v : \texttt{M}_{\mathbf{1}} A$

Case Bind: By inversion$\Phi \mid \Gamma \vdash v_1\colon \texttt{M}_{\epsilon_1} A$ and $\Phi \mid \Gamma, x : A \vdash \epsilon_2\colon \texttt{M}_{\epsilon_2} A$.

Hence by induction $\Phi' \mid \Gamma \vdash v_1\colon \texttt{M}_{\epsilon_1} A$ and $\Phi' \mid \Gamma, x : A \vdash v_2\colon \texttt{M}_{\epsilon_2} A$.

So

$$\Phi' \mid \Gamma \vdash \texttt{do}\ x \leftarrow v_1\ \texttt{in}\ v_2 : \texttt{M}_{\epsilon_1 \cdot \epsilon_2} B \tag{15}$$

Case If: By inversion$\Phi \mid \Gamma \vdash v\colon \texttt{Bool}$, $\Phi \mid \Gamma \vdash v_1\colon A$, and $\Phi \mid \Gamma \vdash v_2\colon A$.

Hence by induction $\Phi' \mid \Gamma \vdash v\colon \texttt{Bool}$, $\Phi' \mid \Gamma \vdash v_1\colon A$, and $\Phi' \mid \Gamma \vdash v_2\colon A$.

So

$$\Phi' \mid \Gamma \vdash \texttt{if}_A\ v\ \texttt{then}\ v_1\ \texttt{else}\ v_2 : A \tag{16}$$

Case Subtype: By inversion$\Phi \mid \Gamma \vdash v\colon A$, and $A \leq: B$.

So by induction: $\Phi' \mid \Gamma \vdash v\colon A$, and $A \leq: B$.

So

$$\Phi' \mid \Gamma \vdash v\colon B \tag{17}$$

Case Effect-Lambda: By inversion$\Phi, \alpha \mid \Gamma \vdash v\colon A$

By picking $\alpha \notin \Phi'$ using α-conversion.

$$\omega\times : \Phi', \alpha \triangleright \Phi, \alpha \tag{18}$$

So by induction, $\Phi', \alpha \mid \Gamma \vdash v\colon A$

Hence,

$$\Phi' \mid \Gamma \vdash \Lambda\alpha.v\colon \forall a.A \tag{19}$$

Case Effect-Apply: By inversion, $\Phi \mid \Gamma \vdash v\colon \forall\alpha.A$, and $\Phi \vdash \epsilon$.

So by induction, $\Phi' \mid \Gamma \vdash v\colon \forall\alpha.A$

And by the well-formed-ness-theorem $\Phi' \vdash \epsilon$

Hence,

$$\Phi' \mid \Gamma \vdash v\ \epsilon\colon A\,[\epsilon/\alpha] \tag{20}$$

0.2 Type Environment Weakening

0.2.1 Relation

We define the ternary weakening relation $\Phi \vdash w : \Gamma' \triangleright \Gamma$ using the following rules.

- $(\mathrm{Id})\dfrac{\Phi \vdash \Gamma \texttt{Ok}}{\Phi \vdash \iota : \Gamma \triangleright \Gamma}$

- $(\mathrm{Project})\dfrac{\Phi \vdash \omega : \Gamma' \triangleright \Gamma \qquad x \notin \texttt{dom}(\Gamma')}{\Phi \vdash \omega\pi : \Gamma, x : A \triangleright \Gamma}$

- $(\mathrm{Extend})\dfrac{\Phi \vdash \omega : \Gamma' \triangleright \Gamma \qquad x \notin \texttt{dom}(\Gamma') \qquad A \leq: B}{\Phi \vdash w\times : \Gamma', x : A \triangleright \Gamma, x : B}$

0.2.2 Domain Lemma

If $\Phi \vdash \omega : \Gamma' \triangleright \Gamma$, then $\texttt{dom}(\Gamma) \subseteq \texttt{dom}(\Gamma')$.

Proof:

Case Id: Then $\Gamma' = \Gamma$ and so $\texttt{dom}(\Gamma') = \texttt{dom}(\Gamma)$.

Case Project: By inversion and induction, $\texttt{dom}(\Gamma) \subseteq \texttt{dom}(\Gamma') \subseteq \texttt{dom}(\Gamma' \cup \{x\})$

Case Extend: By inversion and induction, $\texttt{dom}(\Gamma) \subseteq \texttt{dom}(\Gamma')$ so

$$\texttt{dom}(\Gamma, x : A) = \texttt{dom}(\Gamma) \cup \{x\} \subseteq \texttt{dom}(\Gamma') \cup \{x\} = \texttt{dom}(\Gamma', x : A)$$

0.2.3 Theorem 1

If $\Phi \vdash \omega : \Gamma' \triangleright \Gamma$ and $\Phi \vdash \Gamma\texttt{Ok}$ then $\Phi \vdash \Gamma'\texttt{Ok}$

Proof:

Case Id:

$$(\mathrm{Id})\dfrac{\Phi \vdash \Gamma \texttt{Ok}}{\Phi \vdash \iota : \Gamma \triangleright \Gamma}$$

By inversion, $\Phi \vdash \Gamma\texttt{Ok}$.

Case Project:

$$(\mathrm{Project})\dfrac{\Phi \vdash \omega : \Gamma' \triangleright \Gamma \qquad x \notin \texttt{dom}(\Gamma')}{\Phi \vdash \omega\pi : \Gamma, x : A \triangleright \Gamma}$$

By inversion, $\Phi \vdash \omega : \Gamma' \triangleright \Gamma$ and $x \notin \texttt{dom}(\Gamma')$.

Hence by induction $\Phi \vdash \Gamma'\texttt{Ok}$, $\Phi \vdash \Gamma\texttt{Ok}$. Since $x \notin \texttt{dom}(\Gamma')$, we have $\Phi \vdash \Gamma', x : A\texttt{Ok}$.

Case Extend: $(\mathrm{Extend})\dfrac{\Phi \vdash \omega : \Gamma' \triangleright \Gamma \qquad x \notin \texttt{dom}(\Gamma') \qquad A \leq: B}{\Phi \vdash w\times : \Gamma', x : A \triangleright \Gamma, x : B}$,

By inversion, we have

$\Phi \vdash \omega : \Gamma' \triangleright \Gamma$, $x \notin \texttt{dom}(\Gamma')$.

Hence we have $\Phi \vdash \Gamma\texttt{Ok}$, $\Phi \vdash \Gamma'\texttt{Ok}$, and by the domain Lemma, $\texttt{dom}(\Gamma) \subseteq \texttt{dom}(\Gamma')$, hence $x \notin \texttt{dom}(\Gamma)$. Hence, we have $\Phi \vdash \Gamma, x : A\texttt{Ok}$ and $\Phi \vdash \Gamma', x : A\texttt{Ok}$

0.2.4 Theorem 2

If $\Phi \mid \Gamma \vdash t\colon \tau$ and $\Phi \vdash \omega : \Gamma' \triangleright \Gamma$ then there is a derivation of $\Phi \mid \Gamma' \vdash t\colon \tau$

Proof: We induct over the structure of typing derivations of $\Phi \mid \Gamma \vdash t\colon \tau$, assuming $\Phi \vdash \omega : \Gamma' \triangleright \Gamma$ holds.

Case Var and Weaken: We case split on the weakening ω.

Case: $\omega = \iota$ Then $\Gamma' = \Gamma$, and so $\Phi \mid \Gamma' \vdash x\colon A$ holds and the derivation Δ' is the same as Δ

Case: $\omega = \omega'\pi$ Then $\Gamma' = (\Gamma'', x' : A')$ and $\Phi \vdash \omega' : \Gamma'' \triangleright \Gamma$. So by induction, there is a tree, Δ_1 deriving $\Phi \mid \Gamma'' \vdash x\colon A$, such that:

$$\text{(Weaken)}\frac{\dfrac{\Delta_1}{\Phi \mid \Gamma'' \vdash x\colon A}}{\Phi \mid \Gamma'', x' : A' \vdash x\colon A} \tag{21}$$

Case: $\omega = \omega'\times$ Then

$$\Gamma' = \Gamma''', x' : B \tag{22}$$
$$\Gamma = \Gamma'', x' : A' \tag{23}$$
$$B \leq: A \tag{24}$$

Case: $x = x'$ Then $A = A'$.

Then we derive the new derivation, Δ' as so:

$$\text{(Sub-type)}\frac{\text{(var)}\Phi \mid \Gamma''', x : B \vdash x\colon B \qquad B \leq: A}{\Phi \mid \Gamma' \vdash x\colon A} \tag{25}$$

Case: $x \neq x'$ Then

$$\Delta = \text{(Weaken)}\frac{\dfrac{\Delta_1}{\Phi \mid \Gamma'' \vdash x\colon A}}{\Phi \mid \Gamma \vdash x\colon A} \tag{26}$$

By induction with $\Phi \vdash \omega : \Gamma''' \triangleright \Gamma''$, we have a derivation Δ_1 of $\Phi \mid \Gamma''' \vdash x\colon A$

We have the weakened derivation:

$$\Delta' = \text{(Weaken)}\frac{\dfrac{\Delta_1'}{\Phi \mid \Gamma''' \vdash x\colon A}}{\Phi \mid \Gamma' \vdash x\colon A} \tag{27}$$

Case Constant: The constant typing rules, `()`, `true`, `false`, $\mathtt{C}^A$, all proceed by the same logic. Hence I shall only prove the theorems for the case $\mathtt{C}^A$.

$$\text{(Const)}\frac{\Gamma\,\mathtt{Ok}}{\Gamma \vdash \mathtt{C}^A\colon A} \tag{28}$$

By inversion, we have $\Phi \vdash \Gamma \texttt{Ok}$, so we have $\Phi \vdash \Gamma' \texttt{Ok}$.

Hence

$$(\text{Const})\frac{\Phi \vdash \Gamma' \texttt{Ok}}{\Phi \mid \Gamma' \vdash \texttt{C}^A \colon A} \tag{29}$$

Holds.

Case Lambda: By inversion, we have a derivation Δ_1 giving

$$\Delta = (\text{Fn})\frac{\dfrac{\Delta_1}{\Phi \mid \Gamma, x : A \vdash v \colon B}}{\Phi \mid \Gamma \vdash \lambda x : A.v : A \to B} \tag{30}$$

Since $\Phi \vdash \omega : \Gamma' \triangleright \Gamma$, we have:

$$\Phi \vdash \omega \times : (\Gamma, x : A) \triangleright (\Gamma, x : A) \tag{31}$$

Hence, by induction, using $\Phi \vdash \omega \times : (\Gamma, x : A) \triangleright (\Gamma, x : A)$, we derive Δ_1':

$$\Delta' = (\text{Fn})\frac{\dfrac{\Delta_1'}{\Phi \mid \Gamma', x : A \vdash v \colon B}}{\Phi \mid \Gamma', x : A \vdash \lambda x : A.v : A \to B} \tag{32}$$

Case Sub-typing:

$$(\text{Sub-type})\frac{\Phi \mid \Gamma \vdash v \colon A \qquad A \leq\colon B}{\Phi \mid \Gamma \vdash v \colon B} \tag{33}$$

by inversion, we have a derivation Δ_1

$$\frac{\Delta_1}{\Phi \mid \Gamma \vdash v \colon A} \tag{34}$$

So by induction, we have a derivation Δ_1' such that:

$$(\text{Sub-type})\frac{\dfrac{\Delta_1'}{\Phi \mid \Gamma' \vdash v \colon a} \qquad A \leq\colon B}{\Phi \mid \Gamma' \vdash v \colon B} \tag{35}$$

Case Return: We have the sub-derivation Δ_1 such that

$$\Delta = (\text{Return})\frac{\dfrac{\Delta_1}{\Phi \mid \Gamma \vdash v \colon A}}{\Phi \mid \Gamma \vdash \texttt{return}\ v : \texttt{M}_{\texttt{1}} A} \tag{36}$$

Hence, by induction, with $\Phi \vdash \omega : \Gamma' \triangleright \Gamma$, we find the derivation Δ_1' such that:

$$\Delta' = (\text{Return})\frac{\dfrac{\Delta_1'}{\Phi \mid \Gamma' \vdash v \colon A}}{\Phi \mid \Gamma' \vdash \texttt{return}\ v : \texttt{M}_{\texttt{1}} A} \tag{37}$$

Case Apply: By inversion, we have derivations Δ_1, Δ_2 such that

$$\Delta = (\text{Apply})\dfrac{\dfrac{\Delta_1}{\Phi \mid \Gamma \vdash v_1 \colon A \to B} \qquad \dfrac{\Delta_2}{\Phi \mid \Gamma \vdash v_2 \colon A}}{\Phi \mid \Gamma \vdash v_1\ v_2 \colon B} \tag{38}$$

By induction, this gives us the respective derivations: Δ_1', Δ_2' such that

$$\Delta' = (\text{Apply})\dfrac{\dfrac{\Delta_1'}{\Phi \mid \Gamma' \vdash v_1 \colon A \to B} \qquad \dfrac{\Delta_2'}{\Phi \mid \Gamma' \vdash v_2 \colon A}}{\Phi \mid \Gamma' \vdash v_1\ v_2 \colon B} \tag{39}$$

Case If: By inversion, we have the sub-derivations $\Delta_1, \Delta_2, \Delta_3$, such that:

$$\Delta = (\text{If})\dfrac{\dfrac{\Delta_1}{\Phi \mid \Gamma \vdash v \colon \texttt{Bool}} \qquad \dfrac{\Delta_2}{\Phi \mid \Gamma \vdash v_1 \colon A} \qquad \dfrac{\Delta_3}{\Phi \mid \Gamma \vdash v_2 \colon A}}{\Phi \mid \Gamma \vdash \texttt{if}_A\ v\ \texttt{then}\ v_1\ \texttt{else}\ v_2 : A} \tag{40}$$

By induction, this gives us the sub-derivations $\Delta_1', \Delta_2', \Delta_3'$ such that

$$\Delta' = (\text{If})\dfrac{\dfrac{\Delta_1'}{\Phi \mid \Gamma' \vdash v \colon \texttt{Bool}} \qquad \dfrac{\Delta_2'}{\Phi \mid \Gamma' \vdash v_1 \colon A} \qquad \dfrac{\Delta_3'}{\Phi \mid \Gamma' \vdash v_2 \colon A}}{\Phi \mid \Gamma' \vdash \texttt{if}_A\ v\ \texttt{then}\ v_1\ \texttt{else}\ v_2 : A} \tag{41}$$

Case Bind: By inversion, we have derivations Δ_1, Δ_2 such that:

$$\Delta = (\text{Bind})\dfrac{\dfrac{\Delta_1}{\Phi \mid \Gamma \vdash v_1 \colon \texttt{M}_{\mathbb{E}_1} A} \qquad \dfrac{\Delta_2}{\Phi \mid \Gamma, x : A \vdash v_2 \colon \texttt{M}_{\epsilon_2} B}}{\Phi \mid \Gamma \vdash \texttt{do}\ x \leftarrow v_1\ \texttt{in}\ v_2 : \texttt{M}_{\epsilon_1 \cdot \epsilon_2} B} \tag{42}$$

If $\Phi \vdash \omega : \Gamma' \triangleright \Gamma$ then $\Phi \vdash \omega \times : \Gamma', x : A \triangleright \Gamma, x : A$, so by induction, we can derive Δ_1', Δ_2' such that:

$$\Delta' = (\text{Bind})\dfrac{\dfrac{\Delta_1'}{\Phi \mid \Gamma' \vdash v_1 \colon \texttt{M}_{\mathbb{E}_1} A} \qquad \dfrac{\Delta_2'}{\Phi \mid \Gamma', x : A \vdash v_2 \colon \texttt{M}_{\epsilon_2} B}}{\Phi \mid \Gamma' \vdash \texttt{do}\ x \leftarrow v_1\ \texttt{in}\ v_2 : \texttt{M}_{\epsilon_1 \cdot \epsilon_2} B} \tag{43}$$

Case Effect-Abstraction: By inversion,we have derivation Δ_1 deriving

$$(\text{Effect-Abs})\dfrac{\dfrac{\Delta_1}{\Phi, \alpha \mid \Gamma \vdash v \colon A}}{\Phi \mid \Gamma \vdash \Lambda\alpha.v \colon \forall\alpha.A} \tag{44}$$

By α-conversion, we have $\iota\pi : \Phi, \alpha \triangleright \Phi$, So we have $\Phi, \alpha \vdash \omega : \Gamma' \triangleright \Gamma$ so by induction, there exists Δ_1 deriving:

$$\Delta' = (\text{Effect-Abs})\dfrac{\dfrac{\Delta_1}{\Phi, \alpha \mid \Gamma' \vdash v \colon A}}{\Phi \mid \Gamma' \vdash \Lambda\alpha.v \colon \forall\alpha.A} \tag{45}$$

Case Effect-Application: By inversionwe have derivation Δ_1 deriving

$$\text{(Effect-App)}\,\frac{\dfrac{\Delta_1}{\Phi \mid \Gamma \vdash v\colon \forall\alpha.A} \qquad \Phi \vdash \epsilon}{\Phi \mid \Gamma \vdash v\ \epsilon\colon A\,[\epsilon/\alpha]} \tag{46}$$

So by induction, we have Δ_1' deriving

$$\text{(Effect-App)}\,\frac{\dfrac{\Delta_1'}{\Phi \mid \Gamma' \vdash v\colon \forall\alpha.A} \qquad \Phi \vdash \epsilon}{\Phi \mid \Gamma' \vdash v\ \epsilon\colon A\,[\epsilon/\alpha]} \tag{47}$$