# Compressive Sensing

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Noisy Sparse Recovery

#### Where are We?

- Optimal recovery is given by  $\ell_0$ -norm minimization
- Best relaxation of optimal approach is  $\ell_1$ -norm minimization
  - We learned various forms of  $\ell_1$ -norm minimization
  - We also learned how to implement it
- We learned various iterative algorithms as well
  - Greedy algorithms: OMP, CoSaMP and Subspace Pursuit
  - Thresholding algorithms: IHT, HTP and IST

#### Where are We?

Up to this point, we have assumed noiseless sampling, i.e.,

Samples are modeled as y = Ax

However, in practice we usually have noisy sampling

Samples are modeled as  $\mathbf{y} = \mathbf{A}\mathbf{x} + \mathbf{w}$ 

We now want to modify the sparse recovery algorithms,

In order to deal with sampling noise

## Noisy Sparse Recovery

Noisy Sparse Recovery

## **Noisy Signal Sampling**

Remember the signal model

$$x(t) = \sum_{n=1}^{N} x_n \exp\left\{2\pi j f_n t\right\}$$

The noise-free sample at time instance  $t = t_m$  is

$$y_{m} = x(t_{m}) = \sum_{n=1}^{N} x_{n} e^{2\pi j f_{n} t_{m}} = \underbrace{\left[e^{2\pi j f_{1} t_{m}}, \dots, e^{2\pi j f_{N} t_{m}}\right]}_{\boldsymbol{a}_{m}^{T}} \underbrace{\left[\begin{matrix} x_{1} \\ \vdots \\ x_{N} \end{matrix}\right]}_{\boldsymbol{x}}$$
$$= \boldsymbol{a}_{m}^{T} \boldsymbol{x}$$

# **Noisy Signal Sampling**

In practice, sampling is noisy, i.e., we observe

$$y_m = \boldsymbol{a}_m^\mathsf{T} \boldsymbol{x} + \boldsymbol{w}_m$$

What is  $\mathbf{w_m}$ ?

It is an unknown number with following properties:

- $\mathbf{w}_{m}$  is positive or negative equal chance
- Almost all the time  $|\mathbf{w_m}| \le \epsilon$

## Noisy Signal Sampling

In practice, sampling is noisy, i.e., we observe

$$y_m = \boldsymbol{a}_m^\mathsf{T} \boldsymbol{x} + \boldsymbol{w}_m$$

So, how we could model it?

We could model it as a random process

- w<sub>m</sub> has zero mean
- It has a limited variance

This is enough for now! We later consider a more specific model

#### Sparse Recovery

$$\mathbf{A}\mathbf{z} = \mathbf{y}$$
 subject to  $\|\mathbf{z}\|_0 \le s$ 

How does the sparse recovery problem look like now?

Even for the true signal x, we have in this case

$$y \neq Ax$$

So, clearly we cannot use the earlier formulation!

#### Sparse Recovery

$$\mathbf{A}\mathbf{z} = \mathbf{y}$$
 subject to  $\|\mathbf{z}\|_0 \le s$ 

How does the sparse recovery problem look like now?

We however can say in this case that

$$\|\mathbf{y} - \mathbf{A}\mathbf{x}\|_{2}^{2} = \sum_{m=1}^{M} (\mathbf{y}_{m} - \mathbf{a}_{m}^{\mathsf{T}}\mathbf{x})^{2} = \sum_{m=1}^{M} \mathbf{w}_{m}^{2}$$

#### Sparse Recovery

$$\mathbf{A}\mathbf{z} = \mathbf{y}$$
 subject to  $\|\mathbf{z}\|_0 \le s$ 

How does the sparse recovery problem look like now?

Since  $|\mathbf{w_m}| \leq \epsilon$ , we could say

$$\|\mathbf{y} - \mathbf{A}\mathbf{x}\|_2^2 = \sum_{m=1}^M w_m^2 \le M\epsilon^2$$

So, the sparse solution satisfies  $\|\mathbf{x}\|_0 \leq s$ , and

$$\|\mathbf{y} - \mathbf{A}\mathbf{x}\|_2^2 \leq M\epsilon^2$$

We could hence formulate the sparse recovery problem as below:

#### Noisy Sparse Recovery

For sampling matrix **A** and samples **y** find **z**, such that

$$\|\mathbf{A}\mathbf{z} - \mathbf{y}\|_2^2 \le M\epsilon^2$$
 subject to  $\|\mathbf{z}\|_0 \le s$ 

#### Noisy Sparse Recovery

For sampling matrix A and samples y find z, such that

$$\|\mathbf{A}\mathbf{z} - \mathbf{y}\|_2^2 \le M\epsilon^2$$
 subject to  $\|\mathbf{z}\|_0 \le s$ 

This formulation also includes the earlier form

Set 
$$\epsilon \rightarrow 0$$
; then, we have

$$\|\mathbf{A}\mathbf{z} - \mathbf{y}\|_{2}^{2} \le 0 \implies \|\mathbf{A}\mathbf{z} - \mathbf{y}\|_{2}^{2} = 0$$
  
  $\implies \mathbf{A}\mathbf{z} = \mathbf{y} \implies \text{noise-free case}$ 

Is the solution unique in this case?

Assume that x is the true signal with Supp(x) = S. We now construct a new signal  $x^{\sharp}$  out of x as below:

- Choose  $i \in S$
- Set  $x_i^{\sharp} = x_j + \alpha$  and  $x_n^{\sharp} = x_n$  for or  $n \neq j$

We could hence write

$$\|\mathbf{A}\mathbf{x}^{\sharp} - \mathbf{y}\|_{2}^{2} = \|\mathbf{A}\mathbf{x} - \mathbf{y} + \alpha \mathbf{a}_{j}\|_{2}^{2}$$

$$\leq \|\mathbf{A}\mathbf{x} - \mathbf{y}\|_{2}^{2} + \|\alpha \mathbf{a}_{j}\|_{2}^{2}$$

$$= \|\mathbf{A}\mathbf{x} - \mathbf{y}\|_{2}^{2} + \alpha^{2}\|\mathbf{a}_{j}\|_{2}^{2}$$

Is the solution unique in this case?

Let columns of **A** be normalized and remember  $\mathbf{w} = \mathbf{A}\mathbf{x} - \mathbf{y}$ 

$$\|\mathbf{A}\mathbf{x}^{\sharp} - \mathbf{y}\|_{2}^{2} \leq \|\mathbf{w}\|_{2}^{2} + \alpha^{2}$$

This means that if  $\|\mathbf{w}\|_2^2 < M\epsilon^2$ , for any  $\alpha$  with

$$0 < \alpha^2 \le M\epsilon^2 - \|\mathbf{w}\|_2^2$$

Thus,  $x^{\sharp}$  is also a solution to the sparse recovery problem!

Is the solution unique in this case?

What is the chance of having  $\|\mathbf{w}\|_2^2 < M\epsilon^2$ , i.e.,

$$\Pr\left\{\|\mathbf{w}\|_2^2 \neq M\epsilon^2\right\}$$

Well! It depends on distribution of  $\mathbf{w}$ , but for continuous  $\mathbf{w}$ 

$$\Pr\left\{\|\mathbf{w}\|_2^2 \neq M\epsilon^2\right\} = 1$$

This means that for almost all realizations.

The solution is not unique

If the solution is **not unique**, which one should be taken?

In fact, in this case, we are not recovering the true signal

We are estimating it by an estimate

The choice of this estimate however depends on

The metric by which we evaluate the quality

For now, we focus on the LS method

We get back to this point later on in the next chapters

## Sparse Recovery via LS Method

Noisy Sparse Recovery

One metric to evaluate the quality of our estimate is

Residual Sum of Squares (RSS)

#### Residual Sum of Squares

Let z with  $||z||_0 \le s$  be an estimate of the true signal x. The RSS for this estimate is defined as

$$\Delta\left(\boldsymbol{z}\right) = \|\mathbf{A}\boldsymbol{z} - \boldsymbol{y}\|_{2}^{2}$$

Why should RSS be a good metric?

We need few definitions to answer this question

#### Singular Values

For matrix  $\mathbf{A} \in \mathbb{R}^{M \times N}$ , the singular values are defined as

$$\sigma_m(\mathbf{A}) = \sqrt{\lambda_m(\mathbf{A}\mathbf{A}^\mathsf{T})}$$

for  $m \in \{1, ..., M\}$ , where  $\lambda_m(\mathbf{A}\mathbf{A}^T)$  is m-th eigenvalue of  $\mathbf{A}\mathbf{A}^T$ 

$$\mathbf{A} = \begin{bmatrix} 1 & 1 & 3 \\ 2 & 1 & 1 \end{bmatrix} \rightsquigarrow \mathbf{A}\mathbf{A}^\mathsf{T} = \begin{bmatrix} 11 & 6 \\ 6 & 6 \end{bmatrix} \rightsquigarrow \begin{cases} \lambda_1 = 2 \\ \lambda_2 = 15 \end{cases}$$

The singular values are hence  $\sigma_1 = \sqrt{2}$  and  $\sigma_2 = \sqrt{15}$ 

Singular values of a matrix bound the norm of its projections

#### **Bounding Norm of Projections**

For any  $\mathbf{x} \in \mathbb{R}^N$  and matrix  $\mathbf{A} \in \mathbb{R}^{M \times N}$ , we have

$$\sigma_{\min}^2(\mathbf{A}) \|\mathbf{x}\|_2^2 \le \|\mathbf{A}\mathbf{x}\|_2^2 \le \sigma_{\max}^2(\mathbf{A}) \|\mathbf{x}\|_2^2$$

and

$$\sigma_{\min}(\mathbf{A}) = \min_{m} \sigma_{m}(\mathbf{A})$$
 and  $\sigma_{\max}(\mathbf{A}) = \max_{m} \sigma_{m}(\mathbf{A})$ 

You prove these bounds using singular value decomposition of **A** 

This result can be represented alternatively as

#### **Bounding Norm of Projections**

For any  $\mathbf{x} \in \mathbb{R}^N$  and matrix  $\mathbf{A} \in \mathbb{R}^{M \times N}$ , there exist a  $\kappa$ 

$$\sigma_{\min}^2\left(\mathbf{A}\right) \le \kappa \le \sigma_{\max}^2\left(\mathbf{A}\right)$$

for which we have

$$\|\mathbf{A}\mathbf{x}\|_{2}^{2} = \kappa \|\mathbf{x}\|_{2}^{2}$$

Why should RSS be a good metric?

Let's assume w is a random vector whose entries have

zero mean and variance ξ<sup>2</sup>

This means that for a vector  $\mathbf{a} \in \mathbb{R}^M$ , we have

$$\mathcal{E}\left[\mathbf{a}^{\mathsf{T}}\mathbf{w}\right] = \mathcal{E}\left[\sum_{m=1}^{M} a_{m} \mathbf{w}_{m}\right] = \sum_{m=1}^{M} a_{m} \mathcal{E}\left[\mathbf{w}_{m}\right] = 0$$

$$\mathcal{E}\left[\|\mathbf{w}\|_{2}^{2}\right] = \mathcal{E}\left[\sum_{m=1}^{M}|\mathbf{w}_{m}|^{2}\right] = \sum_{m=1}^{M}\mathcal{E}\left[|\mathbf{w}_{m}|^{2}\right] = M\xi^{2}$$

Why should RSS be a good metric?

The average RSS is then given by

$$\mathcal{E}\left[\Delta\left(\mathbf{z}\right)\right] = \mathcal{E}\left[\|\mathbf{A}\mathbf{z} - \mathbf{y}\|_{2}^{2}\right]$$

$$= \mathcal{E}\left[\|\mathbf{A}\mathbf{z} - \mathbf{A}\mathbf{x} - \mathbf{w}\|_{2}^{2}\right]$$

$$= \mathcal{E}\left[\|\mathbf{A}\left(\mathbf{z} - \mathbf{x}\right) - \mathbf{w}\|_{2}^{2}\right]$$

$$= \|\mathbf{A}\left(\mathbf{z} - \mathbf{x}\right)\|_{2}^{2} - 2\mathcal{E}\left[\left(\mathbf{z} - \mathbf{x}\right)^{\mathsf{T}}\mathbf{A}^{\mathsf{T}}\mathbf{w}\right] + \mathcal{E}\left[\|\mathbf{w}\|_{2}^{2}\right]$$

$$= \|\mathbf{A}\left(\mathbf{z} - \mathbf{x}\right)\|_{2}^{2} + M\xi^{2}$$

Why should RSS be a good metric?

Using, the bound on the projection norm, we could write

$$\mathcal{E}\left[\Delta\left(\mathbf{z}\right)\right] = \|\mathbf{A}\left(\mathbf{z} - \mathbf{x}\right)\|_{2}^{2} + M\xi^{2}$$
$$= \kappa \|\mathbf{z} - \mathbf{x}\|_{2}^{2} + M\xi^{2}$$

for a  $\kappa$  bounded by singular values of **A**. We could hence say

Average RSS 
$$\propto \|\mathbf{z} - \mathbf{x}\|_2^2$$

Thus, bounding RSS leads to restriction of the recovery error

Noisy Sparse Recovery

# Sparse Recovery via LS Method

Section 1: Regularizing LS via  $\ell_0$ -Norm

# Regularizing LS for Sparse Recovery

The LS method finds an estimate of the signal by minimizing RSS

This could however lead to non-sparse solution

To restrict the LS method to sparse recovery,

we regularize LS via the sparsity constraint

#### LS Regularized via $\ell_0$ -Norm

$$\min \|\mathbf{A}\mathbf{z} - \mathbf{y}\|_2^2 \qquad \text{subject to } \|\mathbf{z}\|_0 \le s$$

Regularizing LS via  $\ell_0$ -Norm

# Regularizing LS for Sparse Recovery

#### LS Regularized via $\ell_0$ -Norm

$$\min_{\mathbf{z}} \|\mathbf{A}\mathbf{z} - \mathbf{y}\|_2^2$$
 subject to  $\|\mathbf{z}\|_0 \leq s$ 

How could we get rid of s?

We could find an equivalent optimization

Constraint ← Objective

We have done it before in this lecture

Regularizing LS via  $\ell_0$ -Norm

## Regularized LS: $\ell_0$ -Norm Minimization

Assume LS regularized via  $\ell_0$ -norm has the unique solution  $\mathbf{x}^*$ 

Let's plot the RSS for all **z** which satisfy  $\|\mathbf{z}\|_0 \leq s$ 



- **x** is the unique solution
- **z** are all other points with  $\|\mathbf{z}\|_0 \leq s$

We could now think of one  $\varepsilon$ , for which

$$\|\mathbf{A}\mathbf{x}^{\star} - \mathbf{y}\|_{2}^{2} \leq \varepsilon$$
 and  $\|\mathbf{A}\mathbf{z} - \mathbf{y}\|_{2}^{2} > \varepsilon$ 

Regularizing LS via Lo-Norm

# Regularized LS: $\ell_0$ -Norm Minimization

Assume LS regularized via  $\ell_0$ -norm has the unique solution  $\mathbf{x}^*$ 

Since the solution is unique, we could alternatively say

Among all signals **z** which satisfy  $\|\mathbf{A}\mathbf{z} - \mathbf{y}\|_2^2 \leq \varepsilon$ ,

 $\mathbf{x}^{\star}$  is the only one with  $\|\mathbf{x}^{\star}\|_{0} < s$ 

This means that any  $\mathbf{z} \neq \mathbf{x}^*$  which satisfies  $\|\mathbf{A}\mathbf{z} - \mathbf{y}\|_2^2 \le \varepsilon$  has

$$\|z\|_0 > s$$

## Regularized LS: $\ell_0$ -Norm Minimization

Assume LS regularized via  $\ell_0$ -norm has the unique solution  $\mathbf{x}^*$ 

This concludes that

 $\mathbf{x}^{\star}$  has minimum sparsity among all  $\mathbf{z}$  with  $\|\mathbf{z}\|_{0} \leq \varepsilon$ 

#### Modified $\ell_0$ -Norm Minimization

$$\mathbf{x}^* = \operatorname{argmin} \|\mathbf{z}\|_0$$
 subject to  $\|\mathbf{A}\mathbf{z} - \mathbf{y}\|_2^2 \le \varepsilon$ 

Noisy Sparse Recovery

## Sparse Recovery via LS Method

Section 2: Regularizing LS via  $\ell_1$ -Norm

# LS with $\ell_1$ -Norm Regularization

Any form of  $\ell_0$ -norm minimization is NP-hard

Similar to noise-free case, we address the computational issue by

 $\ell_1$ -norm relaxation

Starting from the modified  $\ell_0$ -norm minimization, we end up with

Basis Pursuit with Quadratic Constraint

$$\mathbf{z}^* = \underset{\mathbf{z}}{\operatorname{argmin}} \|\mathbf{z}\|_1$$
 subject to  $\|\mathbf{A}\mathbf{z} - \mathbf{y}\|_2^2 \le \varepsilon$ 

Unlike the noise-free case, now  $\varepsilon$  does not tend to zero!

Regularizing LS via  $\ell_1$ -Norm

# LS with $\ell_1$ -Norm Regularization

As we showed in Part 2, we could alternatively solve the following:

#### Basis Pursuit Denoising

$$\mathbf{x}^{\star} = \underset{\mathbf{z}}{\operatorname{argmin}} \frac{1}{2} \|\mathbf{A}\mathbf{z} - \mathbf{y}\|_{2}^{2} + \lambda \|\mathbf{z}\|_{1}$$

Again we note that, in noisy scenarios  $\lambda$  does not tend to zero!

Remember that optimal  $\lambda$  varies proportional to  $\varepsilon$ 

$$\lambda \uparrow \leadsto \varepsilon \uparrow \qquad \lambda \downarrow \leadsto \varepsilon \downarrow$$

Regularizing LS via  $\ell_1$ -Norm

# LS with $\ell_1$ -Norm Regularization

LASSO also describes an alternative form:

#### **LASSO**

$$\mathbf{x}^* = \underset{\mathbf{z}}{\operatorname{argmin}} \|\mathbf{A}\mathbf{z} - \mathbf{y}\|_2^2$$
 subject to  $\|\mathbf{z}\|_1 \leq \tau$ 

Remember the initial form of regularized LS

#### LS Regularized via $\ell_0$ -Norm

$$\min \|\mathbf{A}\mathbf{z} - \mathbf{y}\|_2^2$$
 subject to  $\|\mathbf{z}\|_0 \leq s$ 

So, here  $\tau$  is proportional to the sparsity

## LS with $\ell_1$ -Norm Regularization

LASSO also describes an alternative form:

#### **LASSO**

$$\mathbf{x}^* = \underset{\mathbf{z}}{\operatorname{argmin}} \|\mathbf{A}\mathbf{z} - \mathbf{y}\|_2^2$$
 subject to  $\|\mathbf{z}\|_1 \leq \tau$ 

One may note that . . .

LASSO is in fact the original form of the LS method which is regularized by  $\ell_1$ -norm

This is in fact how LASSO was developed initially by Tibshirani!

## Summary

In a nutshell, up to this point we have learned that

With noisy samples,

Exact recovery → Signal Estimation

- The optimal approach is not unique anymore
   It depends on the metric which quantifies estimation quality
- LS method considers RSS as the metric
  - Noisy sparse recovery in this case reduces to modified version of  $\ell_0$ -norm minimization
  - The relaxed version is given by  $\ell_1$ -norm minimization

We now try another approach for noisy sparse recovery

#### Dantzig Approach

Noisy Sparse Recovery

#### Maximum Error Correlation

A less popular, but yet known, metric for quality of our estimate is

Maximum Error Correlation (MEC)

#### Maximum Error Correlation

Let z with  $||z||_0 \le s$  be an estimate of the true signal x. The MEC for this estimate is defined as

$$\Delta(z) = \|\mathbf{A}^{\mathsf{T}}(\mathbf{A}z - \mathbf{y})\|_{\infty}$$

#### Maximum Error Correlation

MEC was taken initially as a metric for noisy sparse recovery by

Emmanuel Candes and Terence Tao

Why it is then called Dantzig approach?

Well! The so-called "father of linear programming"

Georg Dantzig

passed away in 2005 while

Candes and Tao were working on this approach

So, they paid tribute to him

But, why should MEC be a good metric?

Say, a genie has told us that

support of the sparse signal is S

We hence want to find an estimation z whose support is S

Now the question is . . .

Is there an alternative way rather than the LS method?

Let us stick to our earlier definition of residual

$$r = Ax - y$$

MEC is the  $\ell_{\infty}$ -norm of the inner product

$$\mathbf{A}^{\mathsf{T}}$$

But how does the entries of this inner product look like?

How does the entries of this inner product look like?

Consider the noisy sampling model

$$\mathbf{y} = \mathbf{A}\mathbf{x} + \mathbf{w}$$

Residual is written as

$$r = Az - y = Az - Ax - w$$
  
=  $A(z - x) - w$ 

How does the entries of this inner product look like?

Therefore, the inner product looks like

$$\mathbf{A}^{\mathsf{T}} = \mathbf{A}^{\mathsf{T}} \mathbf{A} (z - x) - \mathbf{A}^{\mathsf{T}} w$$

There are two main components here

- **E**stimation error. i.e., z x
- Sampling noise

How does the entries of this inner product look like?

Let's for a moment forget about the sampling noise; then,

- If the genie has told us the right support; then, we could recover  $\mathbf{x}$  exactly

  The good estimation  $\mathbf{z}$  leads to  $\mathbf{A}^\mathsf{T}\mathbf{r} = \mathbf{0}$
- If the genie has told us the wrong support; then, our estimation is always inexact
  - **A**<sup>T</sup>**r** contains always a non-zero entry

How does the entries of this inner product look like?

Let's for a moment forget about the sampling noise; then,

If the genie has told us the right support; then,

$$\|\mathbf{A}^\mathsf{T}_{\mathbf{r}}\|_{\infty} = 0$$

If the genie has told us the wrong support; then,

$$\|\mathbf{A}^\mathsf{T}_{\mathbf{r}}\|_{\infty} \neq 0$$

What does the noise-free case illustrate?

It indicates that

MEC is proportional to the quality of estimation

More precisely, we could say

The most accurate estimation leads to minimum MEC

What about the noisy case?

For noisy sampling, we could extend everything using few tricks

Let us start with a simple inequality:

Consider a vector  $\mathbf{a} \in \mathbb{R}^N$ : then.

$$\|\mathbf{a}\|_{2} = \sqrt{\sum_{n=1}^{N} |a_{n}|^{2}} \ge \sqrt{\max_{n} |a_{n}|^{2}}$$

$$= \max_{n} |a_{n}| = \|\mathbf{a}\|_{\infty}$$

So the bottom line is that  $\|\mathbf{a}\|_{\infty} \leq \|\mathbf{a}\|_{2}$ 

Another simple inequality is

Consider a vector  $\mathbf{a} \in \mathbb{R}^N$ : then.

$$\|\boldsymbol{a}\|_{2} = \sqrt{\sum_{n=1}^{N} |a_{n}|^{2}} \leq \sqrt{N \max_{n} |a_{n}|^{2}}$$
$$= \sqrt{N} \max_{n} |a_{n}| = \sqrt{N} \|\boldsymbol{a}\|_{\infty}$$

So the bottom line is that  $\|\mathbf{a}\|_{\infty} \geq \|\mathbf{a}\|_{2}/\sqrt{N}$ 

Also remember this result that

#### Bounding Norm of Projections

For any  $\mathbf{x} \in \mathbb{R}^N$  and matrix  $\mathbf{A} \in \mathbb{R}^{M \times N}$ , there exist a  $\kappa$ 

$$\sigma_{\min}^2(\mathbf{A}) \le \kappa \le \sigma_{\max}^2(\mathbf{A})$$

for which we have

$$\|\mathbf{A}\mathbf{x}\|_{2}^{2} = \kappa \|\mathbf{x}\|_{2}^{2}$$

Here, bottom line is  $\|\mathbf{A}\mathbf{x}\|_2 = \kappa \|\mathbf{x}\|_2$  where  $\kappa$  mainly depends on  $\mathbf{A}$ 

What about the noisy case?

Let's bound the MEC in the noisy case from above first

$$\|\mathbf{A}^{\mathsf{T}}\mathbf{r}\|_{\infty} = \|\mathbf{A}^{\mathsf{T}}\mathbf{A}(z-x) - \mathbf{A}^{\mathsf{T}}\mathbf{w}\|_{\infty}$$

$$\leq \|\mathbf{A}^{\mathsf{T}}\mathbf{A}(z-x)\|_{\infty} + \|\mathbf{A}^{\mathsf{T}}\mathbf{w}\|_{\infty}$$

$$\leq \|\mathbf{A}^{\mathsf{T}}\mathbf{A}(z-x)\|_{2} + \|\mathbf{A}^{\mathsf{T}}\mathbf{w}\|_{2}$$

$$= \kappa_{1}\|z-x\|_{2} + \kappa_{2}\|\mathbf{w}\|_{2}$$

What about the noisy case?

Similarly, we can bound the MEC in the noisy case from below

$$\|\mathbf{A}^{\mathsf{T}}\mathbf{r}\|_{\infty} = \|\mathbf{A}^{\mathsf{T}}\mathbf{A}(z-\mathbf{x}) - \mathbf{A}^{\mathsf{T}}\mathbf{w}\|_{\infty}$$

$$\geq \|\mathbf{A}^{\mathsf{T}}\mathbf{A}(z-\mathbf{x})\|_{\infty} - \|\mathbf{A}^{\mathsf{T}}\mathbf{w}\|_{\infty}$$

$$\geq \frac{1}{\sqrt{N}}\|\mathbf{A}^{\mathsf{T}}\mathbf{A}(z-\mathbf{x})\|_{2} - \|\mathbf{A}^{\mathsf{T}}\mathbf{w}\|_{2}$$

$$= \hat{\kappa}_{1}\|z-\mathbf{x}\|_{2} - \kappa_{2}\|\mathbf{w}\|_{2}$$

What do the bounds say?

#### We cannot

- $\blacksquare$  play with the  $\kappa$  values, as they depend on the matrix
- play with  $\|\mathbf{w}\|_2$  as is given by noise

So, the key factor in the estimation which change the MEC is

estimation error 
$$\|\mathbf{z} - \mathbf{x}\|_2$$

Sparse Recovery with Dantzig Approach

Noisy Sparse Recovery

#### Dantzig Approach

Section 1: Sparse Recovery with Dantzig Approach

Sparse Recovery with Dantzig Approach

# Optimal Dantzig Approach

We can repeat the same approach, as we took with LS Now, we replace RSS with MEC

#### Optimal Dantzig Approach

$$\min_{\mathbf{z}} \|\mathbf{A}^{\mathsf{T}} (\mathbf{A}\mathbf{z} - \mathbf{y})\|_{\infty}$$
 subject to  $\|\mathbf{z}\|_{0} \leq s$ 

Well, we have the same issue as before

We need to get rid of s

We know very well how to do it

Constraint <-->
Objective

Sparse Recovery with Dantzig Approach

# Optimal Dantzig Approach

We hence can find some  $\varepsilon$ , and rewrite

#### Optimal Dantzig Approach

$$\min_{\mathbf{z}} \|\mathbf{z}\|_{0}$$
 subject to  $\|\mathbf{A}^{\mathsf{T}}(\mathbf{A}\mathbf{z} - \mathbf{y})\|_{\infty} \leq \varepsilon$ 

As always, this approach suffers from

Infeasible Computational Complexity

And as always, we deal with it by

 $\ell_1$ -norm relaxation

Dantzig Selector

Noisy Sparse Recovery

### Dantzig Approach

Section 2: Dantzig Selector

# Dantzig Selector

The relaxed version of optimal Dantzig is called Dantzig selector

#### **Dantzig Selector**

$$\min \| \mathbf{z} \|_1$$
 subject to  $\| \mathbf{A}^\mathsf{T} (\mathbf{A} \mathbf{z} - \mathbf{y}) \|_\infty \le \varepsilon$ 

#### Attention!

- Although not used, one could consider noiseless case by sending  $\varepsilon$  to zero
- Main reason of choosing  $\ell_{\infty}$ -norm is more complicated

#### Summary

Dantzig Selector

In a nutshell, up to this point we have learned that

With noisy samples.

Exact recovery \simple Signal Estimation

- The optimal approach is not unique anymore
- I.S method considers RSS as the metric
- Dantzig approach considers MEC as the metric
  - Optimal Dantzig approach deal with

 $\ell_0$ -norm minimization

The  $\ell_1$ -norm relaxation is called Dantzig selector

We now go for iterative approaches

#### Iterative Algorithms

Noisy Sparse Recovery

#### Noisy Sparse Recovery via Iterative Algorithms

Is there any specific change in the iterative approaches?

Basically No!

Only the stopping criteria is changed

How do we change that?

We consider a quality metric, e.g., RSS or MEC

We stop when the metric is less than a threshold level

Noisy Sparse Recovery

#### Iterative Algorithms

Section 1: Greedy Algorithms

# Orthogonal Matching Pursuit

#### We had

- Start with  $\mathbf{x}^{(0)} = \mathbf{0}$  and  $S^{(0)} = \emptyset$
- Update the support for t > 1 as

$$\mathcal{S}^{(t)} = \mathcal{S}^{(t-1)} \cup \underset{n}{\operatorname{argmax}} \left| \left[ \mathbf{A}^{\mathsf{T}} \left( \mathbf{y} - \mathbf{A} \mathbf{x}^{(t-1)} \right) \right]_{n} \right|$$

■ Set  $\mathbf{z} = \mathbf{A}_{S(t)}^{\dagger} \mathbf{y}$  and update  $\mathbf{x}^{(t)}$  as

$$\mathbf{x}_{S^{(t)}}^{(t)} = \mathbf{z}$$

• Stop at iteration T, at which  $\mathbf{A}\mathbf{x}^{(T)} = \mathbf{y}$ 

# Orthogonal Matching Pursuit

#### Now, we have

- Start with  $\mathbf{x}^{(0)} = \mathbf{0}$  and  $\mathcal{S}^{(0)} = \emptyset$
- Update the support for t > 1 as

$$\mathcal{S}^{(t)} = \mathcal{S}^{(t-1)} \cup \underset{n}{\operatorname{argmax}} \left| \left[ \mathbf{A}^{\mathsf{T}} \left( \mathbf{y} - \mathbf{A} \mathbf{x}^{(t-1)} \right) \right]_{n} \right|$$

■ Set  $\mathbf{z} = \mathbf{A}_{S(t)}^{\dagger} \mathbf{y}$  and update  $\mathbf{x}^{(t)}$  as

$$\mathbf{x}_{S^{(t)}}^{(t)} = \mathbf{z}$$

• Stop at iteration T, at which  $\|\mathbf{A}\mathbf{x}^{(T)} - \mathbf{y}\|_2^2 < \varepsilon$ 

# Orthogonal Matching Pursuit

#### Or, we have

- Start with  $\mathbf{x}^{(0)} = \mathbf{0}$  and  $\mathcal{S}^{(0)} = \emptyset$
- Update the support for  $t \ge 1$  as

$$\mathcal{S}^{(t)} = \mathcal{S}^{(t-1)} \cup \underset{n}{\operatorname{argmax}} \left| \left[ \mathbf{A}^{\mathsf{T}} \left( \mathbf{y} - \mathbf{A} \mathbf{x}^{(t-1)} \right) \right]_{n} \right|$$

lacksquare Set  $oldsymbol{z} = oldsymbol{\mathsf{A}}_{\mathcal{S}^{(t)}}^{\dagger} oldsymbol{y}$  and update  $oldsymbol{x}^{(t)}$  as

$$\mathbf{x}_{S^{(t)}}^{(t)} = \mathbf{z}$$

• Stop at iteration T, at which  $\|\mathbf{A}^{\mathsf{T}} (\mathbf{A} \mathbf{x}^{(\mathsf{T})} - \mathbf{y})\|_{\infty} \leq \varepsilon$ 

# CoSaMP Algorithm

#### For CoSaMP, we had

- Start with  $\mathbf{x}^{(0)} = \mathbf{0}$  and  $S^{(0)} = \emptyset$ , and choose a K
- Update the support for t > 1 as

$$\hat{\mathcal{S}}^{(t)} = \mathcal{S}^{(t-1)} \cup \mathcal{D}_{\mathcal{K}} \left( \mathbf{A}^{\mathsf{T}} \left( \mathbf{y} - \mathbf{A} \mathbf{x}^{(t-1)} 
ight) 
ight)$$

- lacksquare Set  $oldsymbol{z} = oldsymbol{\mathsf{A}}_{\hat{\mathcal{S}}(t)}^{\dagger} oldsymbol{y}$  and update  $oldsymbol{u}^{(t)}$  as  $oldsymbol{u}_{\hat{\mathcal{S}}(t)}^{(t)} = oldsymbol{z}$
- Update once again  $S^{(t)} = D_s(\mathbf{u}^{(t)})$  and  $\mathbf{x}^{(t)} = T_s^H(\mathbf{u}^{(t)})$
- Stop at iteration T, at which  $\mathbf{A}\mathbf{x}^{(T)} = \mathbf{y}$

# CoSaMP Algorithm

Now, we have

- Start with  $\mathbf{x}^{(0)} = \mathbf{0}$  and  $\mathcal{S}^{(0)} = \emptyset$ , and choose a K
- Update the support for  $t \ge 1$  as

$$\hat{\mathcal{S}}^{(t)} = \mathcal{S}^{(t-1)} \cup \mathcal{D}_{\mathcal{K}} \left( \mathbf{A}^{\mathsf{T}} \left( \mathbf{y} - \mathbf{A} \mathbf{x}^{(t-1)} 
ight) 
ight)$$

- lacksquare Set  $oldsymbol{z} = oldsymbol{\mathsf{A}}_{\hat{\mathcal{S}}^{(t)}}^{\dagger} oldsymbol{y}$  and update  $oldsymbol{u}^{(t)}$  as  $oldsymbol{u}_{\hat{\mathcal{S}}^{(t)}}^{(t)} = oldsymbol{z}$
- Update once again  $S^{(t)} = D_s\left( \boldsymbol{u}^{(t)} \right)$  and  $\boldsymbol{x}^{(t)} = T_s^{\mathsf{H}}\left( \boldsymbol{u}^{(t)} \right)$
- Stop at iteration T, at which  $\|\mathbf{A}\mathbf{x}^{(T)} \mathbf{y}\|_2^2 \le \varepsilon$

Iterative Algorithms

# CoSaMP Algorithm

Or, we have

- Start with  $\mathbf{x}^{(0)} = \mathbf{0}$  and  $\mathcal{S}^{(0)} = \emptyset$ , and choose a K
- Update the support for  $t \ge 1$  as

$$\hat{\mathcal{S}}^{(t)} = \mathcal{S}^{(t-1)} \cup \mathcal{D}_{\mathcal{K}} \left( \mathbf{A}^{\mathsf{T}} \left( \mathbf{y} - \mathbf{A} \mathbf{x}^{(t-1)} 
ight) 
ight)$$

- lacksquare Set  $oldsymbol{z} = oldsymbol{\mathsf{A}}_{\hat{\mathcal{S}}^{(t)}}^{\dagger} oldsymbol{y}$  and update  $oldsymbol{u}^{(t)}$  as  $oldsymbol{u}_{\hat{\mathcal{S}}^{(t)}}^{(t)} = oldsymbol{z}$
- lacksquare Update once again  $\mathcal{S}^{(t)} = D_s\left(oldsymbol{u}^{(t)}
  ight)$  and  $oldsymbol{x}^{(t)} = T_s^{\mathsf{H}}\left(oldsymbol{u}^{(t)}
  ight)$
- Stop at iteration T, at which  $\|\mathbf{A}^{\mathsf{T}} (\mathbf{A} \mathbf{x}^{(\mathsf{T})} \mathbf{y})\|_{\infty} \leq \varepsilon$

# Subspace Pursuit

For subspace pursuit, we had

- Start with  $\mathbf{x}^{(0)} = \mathbf{0}$  and  $\mathcal{S}^{(0)} = \emptyset$ , and choose a K
- Update initially the support for  $t \ge 1$  as

$$\hat{\mathcal{S}}^{(t)} = \mathcal{S}^{(t-1)} \cup \mathcal{D}_{\mathcal{K}} \left( \mathbf{A}^{\mathsf{T}} \left( \mathbf{y} - \mathbf{A} \mathbf{x}^{(t-1)} 
ight) 
ight)$$

- lacksquare Set  $oldsymbol{z} = oldsymbol{\mathsf{A}}_{\hat{\mathcal{S}}(t)}^{\dagger} oldsymbol{y}$  and update  $oldsymbol{u}^{(t)}$  as  $oldsymbol{u}_{\hat{\mathcal{S}}(t)}^{(t)} = oldsymbol{z}$
- Update  $S^{(t)} = D_s(\mathbf{u}^{(t)})$
- lacksquare Set  $oldsymbol{z} = oldsymbol{\mathsf{A}}_{\mathcal{S}^{(t)}}^{\dagger} oldsymbol{y}$  and update  $oldsymbol{x}^{(t)}$  as  $oldsymbol{x}_{\mathcal{S}^{(t)}}^{(t)} = oldsymbol{z}$
- Stop at iteration T, at which  $\mathbf{A}\mathbf{x}^{(T)} = \mathbf{y}$

# Subspace Pursuit

#### Now, we have

- Start with  $\mathbf{x}^{(0)} = \mathbf{0}$  and  $\mathcal{S}^{(0)} = \emptyset$ , and choose a K
- Update initially the support for  $t \ge 1$  as

$$\hat{\mathcal{S}}^{(t)} = \mathcal{S}^{(t-1)} \cup \mathcal{D}_{\mathcal{K}} \left( \mathbf{A}^{\mathsf{T}} \left( \mathbf{y} - \mathbf{A} \mathbf{x}^{(t-1)} 
ight) 
ight)$$

- lacksquare Set  $oldsymbol{z} = oldsymbol{\mathsf{A}}_{\hat{\mathcal{S}}^{(t)}}^{\dagger} oldsymbol{y}$  and update  $oldsymbol{u}^{(t)}$  as  $oldsymbol{u}_{\hat{\mathcal{S}}^{(t)}}^{(t)} = oldsymbol{z}$
- Update  $S^{(t)} = D_s(\mathbf{u}^{(t)})$
- lacksquare Set  $oldsymbol{z} = oldsymbol{\mathsf{A}}_{\mathcal{S}^{(t)}}^{\dagger} oldsymbol{y}$  and update  $oldsymbol{x}^{(t)}$  as  $oldsymbol{x}_{\mathcal{S}^{(t)}}^{(t)} = oldsymbol{z}$
- Stop at iteration T, at which  $\|\mathbf{A}\mathbf{x}^{(T)} \mathbf{y}\|_2^2 \leq \varepsilon$

# Subspace Pursuit

#### Or, we have

- Start with  $\mathbf{x}^{(0)} = \mathbf{0}$  and  $\mathcal{S}^{(0)} = \emptyset$ , and choose a K
- Update initially the support for  $t \ge 1$  as

$$\hat{\mathcal{S}}^{(t)} = \mathcal{S}^{(t-1)} \cup \mathcal{D}_{\mathcal{K}} \left( \mathbf{A}^{\mathsf{T}} \left( \mathbf{y} - \mathbf{A} \mathbf{x}^{(t-1)} 
ight) 
ight)$$

- lacksquare Set  $oldsymbol{z} = oldsymbol{\mathsf{A}}_{\hat{\mathcal{S}}^{(t)}}^{\dagger} oldsymbol{y}$  and update  $oldsymbol{u}^{(t)}$  as  $oldsymbol{u}_{\hat{\mathcal{S}}^{(t)}}^{(t)} = oldsymbol{z}$
- Update  $S^{(t)} = D_s(\mathbf{u}^{(t)})$
- lacksquare Set  $oldsymbol{z} = oldsymbol{\mathsf{A}}_{\mathcal{S}^{(t)}}^{\dagger} oldsymbol{y}$  and update  $oldsymbol{x}^{(t)}$  as  $oldsymbol{x}_{\mathcal{S}^{(t)}}^{(t)} = oldsymbol{z}$
- Stop at iteration T, at which  $\|\mathbf{A}^{\mathsf{T}} (\mathbf{A} \mathbf{x}^{(\mathsf{T})} \mathbf{y})\|_{\infty} \leq \varepsilon$

Noisy Sparse Recovery

Thresholding Algorithms

#### Iterative Algorithms

Section 2: Thresholding Algorithms

#### Thresholding Algorithms

# Iterative Hard Thresholding

Similarly for IHT, we could modify the algorithm as

- It starts with  $\mathbf{x}^{(0)} = \mathbf{0}$  and  $\mathcal{S}^{(0)} = \emptyset$
- It updates the approximation of the signal as

$$\mathbf{x}^{(t)} = T_s^{\mathsf{H}} \left( \mathbf{x}^{(t-1)} + \mathbf{A}^{\mathsf{T}} \left( \mathbf{y} - \mathbf{A} \mathbf{x}^{(t-1)} 
ight) 
ight)$$

It stops at iteration T, at which  $\|\mathbf{A}\mathbf{x}^{(T)} - \mathbf{y}\|_2^2 \leq \varepsilon$ 

Or alternatively,

It stops at iteration T, at which  $\|\mathbf{A}^{\mathsf{T}}(\mathbf{A}\mathbf{x}^{(\mathsf{T})}-\mathbf{y})\|_{\infty}\leq \varepsilon$ 

# Hard Thresholding Pursuit

Similarly, we could modify HTP as

- It starts with  $\mathbf{x}^{(0)} = \mathbf{0}$  and  $\mathcal{S}^{(0)} = \emptyset$
- It updates the approximation of the support as

$$\mathcal{S}^{(t)} = D_{s} \left( \mathbf{x}^{(t-1)} + \mathbf{A}^{\mathsf{T}} \left( \mathbf{y} - \mathbf{A} \mathbf{x}^{(t-1)} \right) \right)$$

- $lacksymbol{I}$  It sets  $oldsymbol{z} = oldsymbol{\mathsf{A}}_{\mathcal{S}^{(t)}}^{\dagger} oldsymbol{y}$  and updates  $oldsymbol{x}_{\mathcal{S}^{(t)}}^{(t)} = oldsymbol{z}$
- It stops at iteration T, at which  $\|\mathbf{A}\mathbf{x}^{(T)} \mathbf{y}\|_2^2 \leq \varepsilon$

Or we could use MEC and say

• It stops at iteration T, at which  $\|\mathbf{A}^{\mathsf{T}} (\mathbf{A} \mathbf{x}^{(\mathsf{T})} - \mathbf{y})\|_{\infty} \leq \varepsilon$ 

# Iterative Soft Thresholding

Finally, we modify IST as

- It starts with  $\mathbf{x}^{(0)} = \mathbf{0}$  and  $\mathcal{S}^{(0)} = \emptyset$  and a fixed  $\beta$
- It updates the approximation of the signal as

$$\mathbf{x}^{(t)} = T_{\lambda}^{\mathsf{S}} \left( \mathbf{x}^{(t-1)} + \beta \mathbf{A}^{\mathsf{T}} \left( \mathbf{y} - \mathbf{A} \mathbf{x}^{(t-1)} \right) \right)$$

• It stops at iteration T, at which  $\|\mathbf{A}\mathbf{x}^{(T)} - \mathbf{y}\|_2^2 \leq \varepsilon$ 

Or alternatively,

• It stops at iteration T, at which  $\|\mathbf{A}^{\mathsf{T}} \left(\mathbf{A} \mathbf{x}^{(T)} - \mathbf{y}\right)\|_{\infty} \leq \varepsilon$ 

**Final Points** 

Noisy Sparse Recovery

#### Summary

In a nutshell, we have learned that

With noisy samples,

Exact recovery → Signal Estimation

- The optimal approach is not unique anymore
- LS method considers RSS as the metric
- Dantzig approach considers MEC as the metric
- Iterative algorithms remain the same

Only the stop criteria changes: RSS or MEC

#### What We Learn Next?

We now know most important sparse recovery algorithms, but

How do they perform?

Answering to this question is quite hard and long!

We go through this long way in the next chapter

#### Which Parts of Textbooks?

I would suggest to go over

Statistical Mechanics of Regularized Least Squares A. Bereyhi, 2020

and study the following parts:

■ Chapter 1: Till end of Section 1.3 and Chapter 2: Section 2.3

You can find the manuscript online at this link

For more details on Dantzig approach, you could check out

The Dantzig Selector: Statistical estimation when p is much larger than n, E. Candes and T. Tao, 2005

You can find the manuscript online at this link