Compressive Sensing

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Where Are We?

We ended up with following results

■ The optimal approach for sparse recovery is

 ℓ_0 -norm minimization

■ For unique sparse recovery via the optimal approach at least

$$s \leq M^{\star} \leq 2s$$

samples should be collected

 \bullet ℓ_0 -norm minimization is an NP-hard problem It cannot be implemented in practice

Where Are We?

Now, we want to check out

some sparse recovery algorithms

which can be used in practice

Algorithms for Sparse Recovery

There is a general simple rule

Unlike optimal approach, suboptimal approaches are not unique

So how shall we choose a suboptimal algorithm?

There is always a trade-off in between

- An algorithm might require less number of measurements
- It then probably is computationally more complex

So, it really depends on our budget!

What are the suboptimal approaches?

In general, anyone can come up with an approach!

However, there are popular algorithms with good trade-off

- \bullet ℓ_1 -Norm Minimization Algorithms
- Greedy Algorithms
- Iterative Thresholding Algorithms
- AMP-based Algorithms

We learn the last one in the second part of the lecture

 ℓ_1 -Norm Minimization Algorithms

What do we do in this approach?

In a nutshell,

We replace the ℓ_0 -norm with the ℓ_1 -norm

But why should it work fine?

We see it now!

Back to Our First Example

Remember the example in the introductory part

$$\begin{bmatrix} 1 & 1 & 1 & 0 \\ 1 & -1 & 3 & 0 \\ 0 & 0 & 1 & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix} = \begin{bmatrix} 3 \\ 1 \\ 1 \end{bmatrix}$$

We showed that the set of all solutions is

$$\mathbf{x}(\mathbf{u}) = \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix} = \begin{bmatrix} 2 - 2\mathbf{u} \\ 1 + \mathbf{u} \\ \mathbf{u} \\ 1 - \mathbf{u} \end{bmatrix}$$

Back to Our First Example

Remember the example in the introductory part

$$\begin{bmatrix} 1 & 1 & 1 & 0 \\ 1 & -1 & 3 & 0 \\ 0 & 0 & 1 & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix} = \begin{bmatrix} 3 \\ 1 \\ 1 \end{bmatrix}$$

The sparsest solution is the 2-sparse vector \mathbf{x} ($\mathbf{u} = 1$)

$$x(u=1) = \begin{bmatrix} 0 \\ 2 \\ 1 \\ 0 \end{bmatrix}$$

Remember the example in the introductory part

$$\begin{bmatrix} 1 & 1 & 1 & 0 \\ 1 & -1 & 3 & 0 \\ 0 & 0 & 1 & 1 \end{bmatrix} \begin{vmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{vmatrix} = \begin{bmatrix} 3 \\ 1 \\ 1 \end{bmatrix}$$

This means that x (u = 1)

$$x(u=1) = \operatorname{argmin} \|z\|_0$$
 subject to $Az = y$

Remember the example in the introductory part

$$\begin{bmatrix} 1 & 1 & 1 & 0 \\ 1 & -1 & 3 & 0 \\ 0 & 0 & 1 & 1 \end{bmatrix} \begin{vmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{vmatrix} = \begin{bmatrix} 3 \\ 1 \\ 1 \end{bmatrix}$$

Or equivalently,

$$x (u = 1) = \underset{x(u) \text{ for } u \in \mathbb{R}}{\operatorname{argmin}} ||x(u)||_{0}$$

Let's see how $\|\mathbf{x}(\mathbf{u})\|_0$ looks against \mathbf{u}

Back to Our First Example

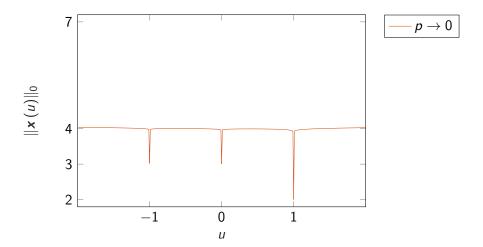
At u = -1, u = 0 and u = 1, x(u) is

$$\boldsymbol{x}(\boldsymbol{u}=-1) = \begin{bmatrix} 4\\0\\-1\\2 \end{bmatrix}, \ \boldsymbol{x}(\boldsymbol{u}=0) = \begin{bmatrix} 2\\1\\0\\1 \end{bmatrix}, \ \boldsymbol{x}(\boldsymbol{u}=1) = \begin{bmatrix} 0\\2\\1\\0 \end{bmatrix}$$

which means

$$\|\mathbf{x}(\mathbf{u}=-1)\|_0 = 3$$
, $\|\mathbf{x}(\mathbf{u}=0)\|_0 = 3$, $\|\mathbf{x}(\mathbf{u}=1)\|_0 = 2$

For any other choice of \mathbf{u} , $\|\mathbf{x}(\mathbf{u})\|_0 = 4$



We know that the key issue is with

 ℓ_0 -norm function

So, what if we relax it with an ℓ_p -norm?

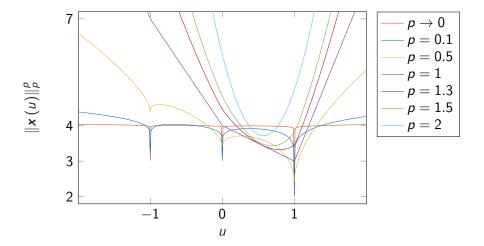
Let's do an experiment: We try to see how does

$$\|\mathbf{x}(\mathbf{u})\|_p^p$$

look like for different choices of p

Relaxation means we replace the problematic part with an approximation

Back to Our First Example



From the figure, one can see that

For $0 \le p \le 1$ the global minimum occur at the same u

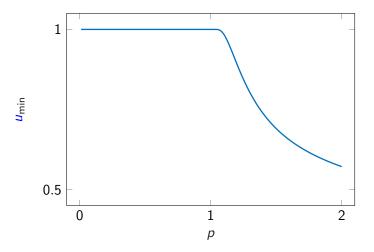
But . . .

As p > 1 the global minimizer starts to move

To see this clearly, let us plot

$$u_{\min} = \underset{u}{\operatorname{argmin}} \| \boldsymbol{x} (u) \|_{p}^{p}$$

against p



This is in fact a general property

 ℓ_p -norm minimization recovers a sparse vector for $0 \le p \le 1$

What would be then a good choice for relaxation?

- For 0 , we can recover a sparse vector
- For p > 1, we deal with convex optimization

Thus, the best choice for relaxation is

 ℓ_1 -norm

Algorithms for Sparse Recovery \$\ell_1\$-Norm Minimization Algorithms | Implementing Basis Pursuit Algorithm | Final Points

Basis Pursuit Algorithm

ℓ_1 -Norm Minimization Algorithms

Section 1: Basis Pursuit Algorithm

Basis Pursuit

Basis Pursuit

Basis pursuit recovers the signal from samples in **y** by

$$\min \|\mathbf{z}\|_1$$
 subject to $\mathbf{A}\mathbf{z} = \mathbf{y}$

Although we have seen it intuitively, we still want to make sure

Basis pursuit recovers a sparse signal

and know what is the sparsity level

An Initial Claim

For $\mathbf{A} \in \mathbb{R}^{M \times N}$ basis pursuit recovers \mathbf{x}^*

$$\mathbf{x}^* = \operatorname{argmin} \|\mathbf{z}\|_1$$
 subject to $\mathbf{A}\mathbf{z} = \mathbf{y}$

Then, \mathbf{x}^* has at most M non-zero elements

To show that this claim is correct consider this reminder.

The kernel of a fat matrix has members other than $oldsymbol{0}$

ker fat
$$\mathbf{A} \neq \{\mathbf{0}\}$$

Let's start with the opposite assumption

 x^* has s > M non-zero elements

We show the support of \mathbf{x}^* with \mathcal{S} , where $|\mathcal{S}| = s$

Clearly, x* satisfies the linear system of equations

$$\mathbf{A}\mathbf{x}^{\star}=\mathbf{y}\rightsquigarrow\mathbf{A}_{\mathcal{S}}\mathbf{x}_{\mathcal{S}}^{\star}=\mathbf{y}$$

The submatrix \mathbf{A}_{S} is of size $M \times s$ which is fat

Let's start with the opposite assumption

 x^* has s > M non-zero elements

We show the support of \mathbf{x}^* with \mathcal{S} , where $|\mathcal{S}| = s$

Since \mathbf{A}_{S} is fat, we can find non-zero $\mathbf{v} \in \mathbb{R}^{s}$ such that

$$\mathbf{v} \in \ker \mathbf{A}_{\mathcal{S}} \leadsto \mathbf{A}_{\mathcal{S}} \mathbf{v} = 0$$

We now construct a new vector $\mathbf{x}^{\sharp} = \mathbf{x}^{\star} + t\tilde{\mathbf{v}}$

 $\tilde{\mathbf{v}} \in \mathbb{R}^N$ is s-sparse with support \mathcal{S} and $\tilde{\mathbf{v}}_{\mathcal{S}} = \mathbf{v}$

Let's start with the opposite assumption

 x^* has s > M non-zero elements

We show the support of \mathbf{x}^* with \mathcal{S} , where $|\mathcal{S}| = s$

Well, clearly we have

$$\mathbf{A} \mathbf{x}^{\sharp} = \mathbf{A} (\mathbf{x}^{\star} + t \tilde{\mathbf{v}})$$

= $\mathbf{A}_{\mathcal{S}} (\mathbf{x}_{\mathcal{S}}^{\star} + t \tilde{\mathbf{v}}_{\mathcal{S}}) = \mathbf{y} + t \mathbf{0} = \mathbf{y}$

Now, what about its ℓ_1 – norm?

Let's start with the opposite assumption

 x^* has s > M non-zero elements

We show the support of \mathbf{x}^* with \mathcal{S} , where $|\mathcal{S}| = s$

We now determine the ℓ_1 -norm of \mathbf{x}^{\sharp}

$$\|\mathbf{x}^{\sharp}\|_{1} = \|\mathbf{x}^{\star} + t\tilde{\mathbf{v}}\|_{1}$$
$$= \sum_{n=1}^{N} |\mathbf{x}_{n}^{\star} + t\tilde{\mathbf{v}}_{n}| = \sum_{n \in \mathcal{S}} |\mathbf{x}_{n}^{\star} + t\tilde{\mathbf{v}}_{n}|$$

Let's start with the opposite assumption

 x^* has s > M non-zero elements

We show the support of \mathbf{x}^* with \mathcal{S} , where $|\mathcal{S}| = s$

We first note that

If
$$|t| < \frac{|x_n^{\star}|}{|\tilde{v}_n|} \rightsquigarrow |t\tilde{v}_n| < |x_n^{\star}|$$

In this case, the sign of $x_n^* + t\tilde{v}_n$ is same as the sign of x_n^*

Let's start with the opposite assumption

 x^* has s > M non-zero elements

We show the support of \mathbf{x}^* with \mathcal{S} , where $|\mathcal{S}| = s$

Now. we focus on an interval of t in which

$$|t| < \min_{n \in \mathcal{S}} \frac{|x_n^*|}{|\tilde{v}_n|} := \theta$$

With $|t| < \theta$, we have $|x_n^* + t\tilde{v}_n| = \operatorname{Sgn}(x_n^*)(x_n^* + t\tilde{v}_n)$ for $n \in \mathcal{S}$

Let's start with the opposite assumption

 x^* has s > M non-zero elements

We show the support of \mathbf{x}^* with \mathcal{S} , where $|\mathcal{S}| = s$

The ℓ_1 -norm of \mathbf{x}^{\sharp} is hence written for $|t| < \theta$ as

$$\|\mathbf{x}^{\sharp}\|_{1} = \sum_{n \in \mathcal{S}} \operatorname{Sgn}(\mathbf{x}_{n}^{\star} + t\tilde{\mathbf{v}}_{n})(\mathbf{x}_{n}^{\star} + t\tilde{\mathbf{v}}_{n})$$

$$= \sum_{n \in \mathcal{S}} \operatorname{Sgn}(\mathbf{x}_{n}^{\star})(\mathbf{x}_{n}^{\star} + t\tilde{\mathbf{v}}_{n})$$

Let's start with the opposite assumption

 x^* has s > M non-zero elements

We show the support of \mathbf{x}^* with \mathcal{S} , where $|\mathcal{S}| = s$

The ℓ_1 -norm of \mathbf{x}^{\sharp} is hence written for $|t| < \theta$ as

$$\|\mathbf{x}^{\sharp}\|_{1} = \sum_{n \in \mathcal{S}} \operatorname{Sgn}(\mathbf{x}_{n}^{\star}) (\mathbf{x}_{n}^{\star} + t \tilde{\mathbf{v}}_{n})$$
$$= \sum_{n \in \mathcal{S}} \operatorname{Sgn}(\mathbf{x}_{n}^{\star}) \mathbf{x}_{n}^{\star} + t \sum_{n \in \mathcal{S}} \operatorname{Sgn}(\mathbf{x}_{n}^{\star}) \tilde{\mathbf{v}}_{n}$$

Let's start with the opposite assumption

 x^* has s > M non-zero elements

We show the support of \mathbf{x}^* with \mathcal{S} , where $|\mathcal{S}| = s$

We note that

$$\sum_{n \in \mathcal{S}} \operatorname{Sgn}(x_n^*) x_n^* = \sum_{n \in \mathcal{S}} |x_n^*|$$
$$= \|\mathbf{x}^*\|_1$$

Let's start with the opposite assumption

 x^* has s > M non-zero elements

We show the support of \mathbf{x}^* with \mathcal{S} , where $|\mathcal{S}| = s$

So, we can write

$$\|\mathbf{x}^{\sharp}\|_{1} = \|\mathbf{x}^{\star}\|_{1} + t \sum_{n \in \mathcal{S}} \operatorname{Sgn}\left(\mathbf{x}_{n}^{\star}\right) \tilde{v}_{n}$$

for all $|t| < \theta$

Let's start with the opposite assumption

 x^* has s > M non-zero elements

We show the support of \mathbf{x}^* with \mathcal{S} , where $|\mathcal{S}| = s$

Interval $|t| \leq \theta$ always include a real number t_1 such that

$$t_1 \sum_{n \in \mathcal{S}} \operatorname{Sgn}\left(x_n^{\star}\right) \tilde{v}_n < 0$$

This means that we can always find \mathbf{x}^{\sharp} , such that $\|\mathbf{x}^{\sharp}\|_{1} < \|\mathbf{x}^{\star}\|_{1}$

Let's start with the opposite assumption

 x^* has s > M non-zero elements

We show the support of \mathbf{x}^* with \mathcal{S} , where $|\mathcal{S}| = s$

This is now a contradiction with the fact that

 \mathbf{x}^{\star} is the global minimizer of ℓ_1 -norm on the feasible set

Our initial claim is thus proven by contradiction

To avoid the contradiction what do we need?

- First of all x^* be s-sparse with s < M
- Secondly, for A we need to have

$$\ker \mathbf{A}_{\mathcal{S}} = \{\mathbf{0}\}$$

for the recovery support S with |S| < s

Sparsity of Basis Pursuit Recovery

Basis pursuit always recovers a s-sparse signal with s < M and the columns of \mathbf{A}_{S} are linearly independent

Summary

Up to this point, we have learned that

- Basis pursuit $\equiv \ell_1$ -norm minimization gives Best relaxation of the optimal sparse recovery
- We can make sure that basis pursuit recovers A signal whose sparsity is at most as much as # of samples

Now the question is that

■ How can we implement the basis pursuit algorithm?

We study it in the next section

Implementing Basis Pursuit Algorithm

Getting Back to Basis Pursuit

Basis Pursuit

Basis pursuit recovers the signal from samples in **y** by

$$\min \|\mathbf{z}\|_1$$
 subject to $\mathbf{A}\mathbf{z} = \mathbf{y}$

One question which comes to mind

Is basis pursuit, the only relaxation via ℓ_1 -norm?

We see shortly that the answer is No!

Alternative Forms of ℓ_1 -Norm Minimization

Implementing Basis Pursuit Algorithm

Section 1: Alternative Forms of ℓ_1 -Norm Minimization

Alternative Forms of \(\ell_1 \)-Norm Minimization

Basis Pursuit with Quadratic Constraint

Basis Pursuit

$$\min \|\mathbf{z}\|_1$$
 subject to $\mathbf{A}\mathbf{z} = \mathbf{y}$

The linear constraint in basis pursuit can be translated as

$$\mathbf{A}\mathbf{z} = \mathbf{y} \iff \mathbf{A}\mathbf{z} - \mathbf{y} = \mathbf{0} \iff \|\mathbf{A}\mathbf{z} - \mathbf{y}\|_2^2 = 0$$

Basis Pursuit with Quadratic Constraint

Basis Pursuit

$$\min \|\mathbf{z}\|_1$$
 subject to $\mathbf{A}\mathbf{z} = \mathbf{y}$

Now, we consider the following quadratic constraint

$$\|\mathbf{A}\mathbf{z} - \mathbf{y}\|_2^2 \leq \eta$$

What happens if $\eta \to 0$?

$$\|\mathbf{A}\mathbf{z} - \mathbf{y}\|_2^2 \le \eta \iff \|\mathbf{A}\mathbf{z} - \mathbf{y}\|_2^2 = 0 \iff \mathbf{A}\mathbf{z} = \mathbf{y}$$

Basis Pursuit with Quadratic Constraint

Basis Pursuit

$$\min \|\mathbf{z}\|_1$$
 subject to $\mathbf{A}\mathbf{z} = \mathbf{y}$

So, we can alternatively solve the following problem and let $\eta \to 0$

Basis Pursuit with Quadratic Constraint

$$\min \|\mathbf{z}\|_1$$
 subject to $\|\mathbf{A}\mathbf{z} - \mathbf{y}\|_2^2 \le \eta$

Basis Pursuit with Quadratic Constraint

Basis Pursuit with Quadratic Constraint

$$\min \|\mathbf{z}\|_1$$
 subject to $\|\mathbf{A}\mathbf{z} - \mathbf{y}\|_2^2 \leq \eta$

Attention

You might think

 $\eta \to 0$ is the only meaningful choice

But, as we see later

Other choices of η are also useful in noisy settings

Basis Pursuit

$$\min \|\mathbf{z}\|_1$$
 subject to $\mathbf{A}\mathbf{z} = \mathbf{y}$

Let us construct a new objective function $F_{\lambda}(z)$ as below

$$F_{\lambda}(\mathbf{z}) = \frac{1}{2\lambda} \|\mathbf{A}\mathbf{z} - \mathbf{y}\|_{2}^{2} + \|\mathbf{z}\|_{1}$$

for some real $\lambda > 0$

Basis Pursuit

$$\min \|\mathbf{z}\|_1$$
 subject to $\mathbf{A}\mathbf{z} = \mathbf{y}$

What happens for this objective, if $\lambda \to 0$?

$$F_{\lambda}(z) = \frac{1}{2\lambda} \|\mathbf{A}z - \mathbf{y}\|_{2}^{2} + \|\mathbf{z}\|_{1} = \begin{cases} \|\mathbf{z}\|_{1} & \text{if } \mathbf{A}z = \mathbf{y} \\ +\infty & \text{if } \mathbf{A}z \neq \mathbf{y} \end{cases}$$

Basis Pursuit

$$\min \|\mathbf{z}\|_1$$
 subject to $\mathbf{A}\mathbf{z} = \mathbf{y}$

So, if the solution of basis pursuit has a finite ℓ_1 -norm, we have

$$\begin{cases} \lim_{\lambda \to 0} \min F_{\lambda}(z) \\ = \min \{ \|z\|_{1} \text{ subject to } \mathbf{A}z = y \} \cup \{+\infty\} \end{cases}$$
$$= \{ \min \|z\|_{1} \text{ subject to } \mathbf{A}z = y \}$$

Basis Pursuit

$$\min \|\mathbf{z}\|_1$$
 subject to $\mathbf{A}\mathbf{z} = \mathbf{y}$

We could in fact recover the basis pursuit solution by

Basis Pursuit Denoising

$$\min \frac{1}{2} \|\mathbf{A}\mathbf{z} - \mathbf{y}\|_2^2 + \lambda \|\mathbf{z}\|_1$$

From Denoising to Pure Basis Pursuit

Let \mathbf{x}^* and $\mathbf{x}(\lambda)$ be defined as

$$\mathbf{x}^* = \operatorname{argmin} \|\mathbf{z}\|_1$$
 subject to $\mathbf{A}\mathbf{z} = \mathbf{y}$

$$oldsymbol{x}(\lambda) = \operatorname{argmin} \frac{1}{2} \|\mathbf{A} \mathbf{z} - \mathbf{y}\|_2^2 + \lambda \|\mathbf{z}\|_1$$

Then, if \mathbf{x}^* is the unique solution, we have

$$\lim_{\lambda \to 0} \mathbf{x}(\lambda) = \mathbf{x}^{\star}$$

How can we show this equivalency?

Before, we start with the proof, remember that

 \mathbf{x}^{\star} is the solution of basis pursuit

So, for any solution x of the underdetermined equation, i.e.,

$$\mathbf{A}\mathbf{x}=\mathbf{y}$$

We have

$$\|oldsymbol{x}\|_1 \geq \|oldsymbol{x}^\star\|_1$$

How can we show this equivalency?

First of all, since x^* is the basis pursuit solution, we have

$$\mathbf{A}\mathbf{x}^{\star}=\mathbf{y}$$

This means that

$$\frac{1}{2} \underbrace{\|\mathbf{A}\mathbf{x}^{\star} - \mathbf{y}\|_{2}^{2}}_{0} + \lambda \|\mathbf{x}^{\star}\|_{1} = \lambda \|\mathbf{x}^{\star}\|_{1}$$

How can we show this equivalency?

If we define

$$G_{\lambda}(z) = \frac{1}{2} \|\mathbf{A}z - \mathbf{y}\|_{2}^{2} + \lambda \|\mathbf{z}\|_{1}$$

Our conclusion simply says that

$$G_{\lambda}\left(\mathbf{x}^{\star}\right) = \lambda \|\mathbf{x}^{\star}\|_{1}$$

How can we show this equivalency?

Secondly, we know that

$$\mathbf{x}(\lambda)$$
 is the minimizer of $G_{\lambda}(\mathbf{z})$

This means that for any z

$$G_{\lambda}\left(\mathbf{x}\left(\lambda\right)\right) \leq G_{\lambda}\left(\mathbf{z}\right)$$

This is also true for $\mathbf{z} = \mathbf{x}^*$, i.e., $G_{\lambda}(\mathbf{x}(\lambda)) < G_{\lambda}(\mathbf{x}^*)$ or

$$\frac{1}{2}\|\mathbf{A}\boldsymbol{x}\left(\lambda\right)-\boldsymbol{y}\|_{2}^{2}+\lambda\|\boldsymbol{x}\left(\lambda\right)\|_{1}\leq\lambda\|\boldsymbol{x}^{\star}\|_{1}$$

How can we show this equivalency?

Noting that

$$\|\mathbf{A}\mathbf{x}(\lambda) - \mathbf{y}\|_{2}^{2}$$
, $\lambda \|\mathbf{x}(\lambda)\|_{1}$ and $\lambda \|\mathbf{x}^{\star}\|_{1}$ are non-negative

We can conclude out of the last inequality that

$$\lambda \| \boldsymbol{x}(\lambda) \|_{1} \leq \lambda \| \boldsymbol{x}^{\star} \|_{1} \rightsquigarrow \| \boldsymbol{x}(\lambda) \|_{1} \leq \| \boldsymbol{x}^{\star} \|_{1}$$

$$\frac{1}{2}\|\mathbf{A}\mathbf{x}\left(\lambda\right)-\mathbf{y}\|_{2}^{2}\leq\lambda\|\mathbf{x}^{\star}\|_{1}$$

How can we show this equivalency?

We know at the moment that

$$\|\mathbf{x}(\lambda)\|_1 \leq \|\mathbf{x}^{\star}\|_1$$

which is also true if $\lambda \to 0$. Now, remember that

$$\frac{1}{2}\|\mathbf{A}\boldsymbol{x}\left(\lambda\right)-\boldsymbol{y}\|_{2}^{2}\leq\lambda\|\boldsymbol{x}^{\star}\|_{1}$$

If we send $\lambda \to 0$, we conclude that

$$\lim_{\lambda \to 0} \frac{1}{2} \|\mathbf{A}\mathbf{x}(\lambda) - \mathbf{y}\|_{2}^{2} \le 0 \leadsto \lim_{\lambda \to 0} \frac{1}{2} \|\mathbf{A}\mathbf{x}(\lambda) - \mathbf{y}\|_{2}^{2} = 0$$

How can we show this equivalency?

The identity

$$\lim_{\lambda \to 0} \frac{1}{2} \|\mathbf{A} \mathbf{x} (\lambda) - \mathbf{y}\|_2^2 = 0$$

means that if $\mathbf{x}(\lambda)$ converges as $\lambda \to 0$, we have

$$\mathbf{A}\lim_{\lambda\to 0}\mathbf{x}\left(\lambda\right)=\mathbf{y}$$

In other words.

 $\lim_{\lambda \to 0} \mathbf{x}(\lambda)$ is a solution of the underdetermined equation

How can we show this equivalency?

Well, we could say that

$$\|\lim_{\lambda \to 0} \mathbf{x}(\lambda)\|_1 \ge \|\mathbf{x}^{\star}\|_1$$

Put it beside the fact that

$$\|\lim_{\lambda \to 0} \mathbf{x}(\lambda)\|_1 \le \|\mathbf{x}^{\star}\|_1$$

We need to have

$$\|\lim_{\lambda \to 0} \mathbf{x}(\lambda)\|_1 = \|\mathbf{x}^{\star}\|_1$$

How can we show this equivalency?

The uniqueness of x* finally concludes that

$$\lim_{\lambda \to 0} \mathbf{x}(\lambda) = \mathbf{x}^*$$

Attention!

- Again, you can think of $\lambda \neq 0$ in noisy settings
- This alternative form is a key form, we use for implementation

Basis Pursuit with Quadratic Constraint

$$\min \|\mathbf{z}\|_1$$
 subject to $\|\mathbf{A}\mathbf{z} - \mathbf{y}\|_2^2 \le \eta$

Assume that we have the unique solution x^*

Assume that Z is the following set

$$\mathcal{Z} = \left\{ all \ \mathbf{z} \in \mathbb{R}^{N} : \|\mathbf{A}\mathbf{z} - \mathbf{y}\|_{2}^{2} \leq \eta \right\}$$

Alternative Forms of ℓ_1 -Norm Minimization

LASSO

Basis Pursuit with Quadratic Constraint

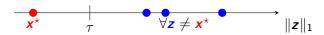
$$\min \|\mathbf{z}\|_1$$
 subject to $\|\mathbf{A}\mathbf{z} - \mathbf{y}\|_2^2 \le \eta$

Assume that we have the unique solution x^*

We thus have

$$\mathbf{z}^{\star} = \underset{\mathbf{z} \in \mathcal{Z}}{\operatorname{argmin}} \|\mathbf{z}\|_{1}$$

What happens if we plot the ℓ_1 -norm of all $\mathbf{z} \in \mathcal{Z}$?



- x* is the unique solution
- **z** are all other points in \mathcal{Z}

We could now think of one τ , for which

$$\|\mathbf{x}^{\star}\|_{1} \leq \tau$$
$$\|\mathbf{z}\|_{1} > \tau$$

We could hence say.

$$\{\mathbf{x}^{\star}\} = \mathcal{Z} \cap \{\text{all } \mathbf{z} : \|\mathbf{z}\|_1 \leq \tau\}$$

What about the other points $\mathbf{z} \neq \mathbf{x}^*$ in $\{all \ \mathbf{z} : \|\mathbf{z}\|_1 \leq \tau\}$?

Well, since \mathbf{x}^* is the unique solution, we have

$$\forall \{all \ \mathbf{z} : \|\mathbf{z}\|_1 \leq \tau\} \ni \mathbf{z} \neq \mathbf{x}^* : \mathbf{z} \notin \mathcal{Z}$$

This means

$$\forall \left\{ all \ \mathbf{z} : \|\mathbf{z}\|_1 \leq \tau \right\} \ni \mathbf{z} \neq \mathbf{x}^* : \|\mathbf{A}\mathbf{z} - \mathbf{y}\|_2^2 > \eta$$

So, we could break $\{all \ z : \|z\|_1 \le \tau\}$ into two cases

$$\{all \ \mathbf{z} : \|\mathbf{z}\|_1 \leq \tau\} : \begin{cases} \|\mathbf{A}\mathbf{z} - \mathbf{y}\|_2^2 > \eta & \mathbf{z} \neq \mathbf{x}^* \\ \|\mathbf{A}\mathbf{z} - \mathbf{y}\|_2^2 \leq \eta & \mathbf{z} = \mathbf{x}^* \end{cases}$$

This means that

$$\mathbf{x}^{\star} = \operatorname{argmin} \|\mathbf{A}\mathbf{z} - \mathbf{y}\|_{2}^{2} \text{ subject to } \|\mathbf{z}\|_{1} \leq \tau$$

This was discovered by Robert Tibshirani in 1996 under the name

Least Absolute Shrinkage and Selection Operator (LASSO)

LASSO

$$\min \|\mathbf{A}\mathbf{z} - \mathbf{y}\|_2^2$$
 subject to $\|\mathbf{z}\|_1 \le \tau$

Attention!

- The value of τ is determined in terms of η
- LASSO is the oldest ℓ_1 -norm minimization technique

The Homotopy Method

Implementing Basis Pursuit Algorithm

Section 2: The Homotopy Method

Homotopy in Nutshell

Remember that we can perform basis pursuit by denoising as

$$oldsymbol{x}\left(oldsymbol{\lambda}
ight) = \operatorname{argmin} rac{1}{2} \|oldsymbol{\mathsf{A}}oldsymbol{z} - oldsymbol{y}\|_2^2 + \lambda \|oldsymbol{z}\|_1$$

and then taking the limit $\lambda \to 0$

The homotopy method uses this property

- It first finds $\lambda^{(0)}$ such that $\mathbf{x}(\lambda^{(0)}) = \mathbf{0}$
- It then step by step updates $\mathbf{x}(\lambda^{(t)})$ till it gets to $\lambda^{(t)} = 0$

The Homotopy Method

The Homotopy Method: Derivation

Noting that the function

$$G_{\lambda}(\mathbf{z}) = \frac{1}{2} \|\mathbf{A}\mathbf{z} - \mathbf{y}\|_{2}^{2} + \lambda \|\mathbf{z}\|_{1}$$

is convex, we can find its minimizer by checking the extreme points

Let $\mathbf{x}(\lambda)$ be an extreme point; then, we find for $n \in \{1, ..., N\}$

$$\left[\mathbf{A}^{\mathsf{T}}\mathbf{A}\mathbf{x}\left(\lambda\right)-\mathbf{A}^{\mathsf{T}}\mathbf{y}\right]_{n}+\lambda\mathsf{Sgn}\left(x_{n}\left(\lambda\right)\right)=0\quad\forall n:x_{n}\left(\lambda\right)\neq0$$

$$\left| \left[\mathbf{A}^{\mathsf{T}} \mathbf{A} \mathbf{x} \left(\lambda \right) - \mathbf{A}^{\mathsf{T}} \mathbf{y} \right]_{n} \right| \leq \lambda$$
 $\forall n : x_{n} \left(\lambda \right) = 0$

[a] is the n-th entry of a and Sgn(x) is the sign of x

To have $\mathbf{x}(\lambda^{(0)}) = \mathbf{0}$, we need to choose $\lambda^{(0)}$, such that

$$\left| \left[\mathbf{A}^\mathsf{T} \mathbf{A} \mathbf{x} \left(\lambda^{(0)} \right) - \mathbf{A}^\mathsf{T} \mathbf{y} \right]_n \right| \leq \lambda^{(0)} \leadsto \left| \left[\mathbf{A}^\mathsf{T} \mathbf{y} \right]_n \right| \leq \lambda^{(0)}$$

for all $n \in \{1, ..., N\}$. One good choice is

$$\lambda^{(0)} = \left| \left[\mathbf{A}^\mathsf{T} \mathbf{y} \right]_{n_1} \right|$$

where n_1 is

$$n_1 = \underset{n}{\operatorname{argmax}} \left| \left[\mathbf{A}^\mathsf{T} \mathbf{y} \right]_n \right|$$

Let us now initiate our algorithm

We start at point $\mathbf{x}^{(0)} = \mathbf{0}$. and set

$$\lambda^{(0)} = \max_{n} \left| \left[\mathbf{A}^{\mathsf{T}} \mathbf{y} \right]_{n} \right|$$

For compactness, we further define the residual in iteration t

$$\mathbf{r}^{(t)} = \mathbf{A}^{\mathsf{T}} \left(\mathbf{A} \mathbf{x}^{(t)} - \mathbf{y} \right)$$

For instance at t = 0, we have $\mathbf{r}^{(0)} = -\mathbf{A}^{\mathsf{T}}\mathbf{v}$

So, we could say

We start at point $\mathbf{x}^{(0)} = \mathbf{0}$, and set

$$\lambda^{(0)} = \max_{n} \left| r_n^{(0)} \right|$$

and denote the index of largest residual element with n_1

where we consider

$$\mathbf{r}^{(t)} = \begin{bmatrix} r_1^{(t)} \\ \vdots \\ r_N^{(t)} \end{bmatrix}$$

Remember the extreme point equations

$$[\mathbf{A}^{\mathsf{T}}\mathbf{A}\mathbf{x}(\lambda) - \mathbf{A}^{\mathsf{T}}\mathbf{y}]_{n} + \lambda \operatorname{Sgn}(x_{n}(\lambda)) = 0 \quad \forall n : x_{n}(\lambda) \neq 0$$
$$|[\mathbf{A}^{\mathsf{T}}\mathbf{A}\mathbf{x}(\lambda) - \mathbf{A}^{\mathsf{T}}\mathbf{y}]_{n}| \leq \lambda \qquad \forall n : x_{n}(\lambda) = 0$$

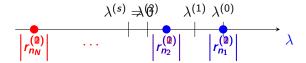
Note that at iteration t with $\lambda^{(t)}$ and $\mathbf{x}^{(t)} = \mathbf{x} (\lambda^{(t)})$, we have

$$\mathbf{r}^{(t)} = \mathbf{A}^{\mathsf{T}} \left(\mathbf{A} \mathbf{x}^{(t)} - \mathbf{y} \right)$$

So, in iteration t the equations read

$$\begin{vmatrix} r_n^{(t)} + \lambda^{(t)} \operatorname{Sgn}\left(x_n^{(t)}\right) = 0 & \forall n : x_n^{(t)} \neq 0 \\ \left| r_n^{(t)} \right| \leq \lambda^{(t)} & \forall n : x_n^{(t)} = 0 \end{aligned}$$

How does it look like? $x_n = 0 \ \forall n$, but $n = n_1, n_2, \ldots, n_s$



We now want to move in a direction which leads to $\lambda \to 0$

We set
$$\lambda^{(1)} = \lambda^{(0)} - \beta^{(1)}$$
 and update $\mathbf{x}^{(1)}$ as

$$\mathbf{x}^{(1)} = \mathbf{x}^{(0)} + \beta^{(1)} \mathbf{d}^{(1)}$$

 $d^{(1)}$ is a 1-sparse vector whose non-zero element is at n_1

$$d_{n_1}^{(1)} = -\frac{\operatorname{Sgn}\left(r_{n_1}^{(0)}\right)}{\boldsymbol{a}_{n_1}^{\mathsf{T}}\boldsymbol{a}_{n_1}}$$

with a_n being the n_1 -th column of A

We now want to move in a direction which leads to $\lambda \to 0$

We set
$$\lambda^{(1)} = \lambda^{(0)} - \beta^{(1)}$$
 and update $\mathbf{x}^{(1)}$ as

$$\mathbf{x}^{(1)} = \mathbf{x}^{(0)} + \beta^{(1)} \mathbf{d}^{(1)}$$

One could check that for any $0 < \beta^{(1)} \le \lambda^{(0)}$ we have

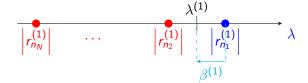
$$r_{n_1}^{(1)} + \lambda^{(1)} \operatorname{Sgn}\left(x_{n_1}^{(1)}\right) = 0$$

So,
$$x_{n_1}^{(1)} \neq 0$$

We now want to move in a direction which leads to $\lambda \to 0$

We set
$$\lambda^{(1)}=\lambda^{(0)}-\beta^{(1)}$$
 and update $m{x}^{(1)}$ as $m{x}^{(1)}=m{x}^{(0)}+eta^{(1)}m{d}^{(1)}$

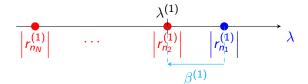
But we want to keep the others elements zero! So, we bound $\beta^{(1)}$



We now want to move in a direction which leads to $\lambda \to 0$

We set
$$\lambda^{(1)}=\lambda^{(0)}-\beta^{(1)}$$
 and update $m{x}^{(1)}$ as $m{x}^{(1)}=m{x}^{(0)}+\beta^{(1)}m{d}^{(1)}$

We try to find the largest possible choice



We now want to move in a direction which leads to $\lambda \to 0$

We set
$$\lambda^{(1)}=\lambda^{(0)}-\beta^{(1)}$$
 and update $m{x}^{(1)}$ as $m{x}^{(1)}=m{x}^{(0)}+eta^{(1)}m{d}^{(1)}$

One can verify that the largest choice of $\beta^{(1)}$ is

$$\beta^{(1)} = \min \bigcup_{n \neq n_1} Sub^{+} \left\{ \frac{\lambda^{(0)} + r_n^{(0)}}{1 - \left[\mathbf{A}^{\mathsf{T}} \mathbf{A} d^{(1)} \right]_n}, \frac{\lambda^{(0)} - r_n^{(0)}}{1 + \left[\mathbf{A}^{\mathsf{T}} \mathbf{A} d^{(1)} \right]_n} \right\}$$

with $Sub^+ \{\cdot\}$ calculating the subset positive elements

We now want to move in a direction which leads to $\lambda \to 0$

We set
$$\lambda^{(1)} = \lambda^{(0)} - \beta^{(1)}$$
 and update $\mathbf{x}^{(1)}$ as

$$\mathbf{x}^{(1)} = \mathbf{x}^{(0)} + \beta^{(1)} \mathbf{d}^{(1)}$$

The residual is then updated as

$$\mathbf{r}^{(1)} = \mathbf{A}^{\mathsf{T}} \left(\mathbf{A} \mathbf{x}^{(1)} - \mathbf{y} \right)$$

We now want to move in a direction which leads to $\lambda \to 0$

We set
$$\lambda^{(1)}=\lambda^{(0)}-\beta^{(1)}$$
 and update $m{x}^{(1)}$ as $m{x}^{(1)}=m{x}^{(0)}+\beta^{(1)}m{d}^{(1)}$

We now update the support $S^{(2)} = \{n_1, n_2\}$ with

$$\underline{n_2} = \operatorname{argmin} \bigcup_{n \neq n_1} Sub^+ \left\{ \frac{\lambda^{(0)} + r_n^{(0)}}{1 - \left[\mathbf{A}^\mathsf{T} \mathbf{A} \boldsymbol{d}^{(1)}\right]_n}, \frac{\lambda^{(0)} - r_n^{(0)}}{1 + \left[\mathbf{A}^\mathsf{T} \mathbf{A} \boldsymbol{d}^{(1)}\right]_n} \right\}$$

and exclude it in step 2

We keep repeating this strategy: We are at iteration $t \geq 1$

We set
$$\lambda^{(t)}=\lambda^{(t-1)}-\beta^{(t)}$$
 and update $\mathbf{x}^{(t)}$ as $\mathbf{x}^{(t)}=\mathbf{x}^{(t-1)}+\beta^{(t)}\mathbf{d}^{(t)}$

We have a $|S^{(t)}|$ -sparse path $\mathbf{d}^{(t)}$ with non-zero elements

$$\mathbf{z} = -\left(\mathbf{A}_{\mathcal{S}^{(t)}}^\mathsf{T} \mathbf{A}_{\mathcal{S}^{(t)}}\right)^{-1} \mathsf{Sgn}\left(\mathbf{r}_{\mathcal{S}^{(t)}}^{(t)}\right)$$

This means that we set $\mathbf{d}_{S(t)}^{(t)} = \mathbf{z}$ and the other entries of $\mathbf{d}^{(t)}$ zero

We keep repeating this strategy: We are at iteration $t \geq 1$

We set
$$\lambda^{(t)}=\lambda^{(t-1)}-\beta^{(t)}$$
 and update $\mathbf{x}^{(t)}$ as $\mathbf{x}^{(t)}=\mathbf{x}^{(t-1)}+\beta^{(t)}\mathbf{d}^{(t)}$

We now find the largest choice of $\beta^{(t)}$ which is

$$\beta_{+}^{(t)} = \min \bigcup_{n \notin \mathcal{S}^{(t)}} Sub^{+} \left\{ \frac{\lambda^{(t-1)} + r_{n}^{(t-1)}}{1 - \left[\mathbf{A}^{\mathsf{T}} \mathbf{A} \boldsymbol{d}^{(t)}\right]_{n}}, \frac{\lambda^{(t-1)} - r_{n}^{(t-1)}}{1 + \left[\mathbf{A}^{\mathsf{T}} \mathbf{A} \boldsymbol{d}^{(t)}\right]_{n}} \right\}$$

We keep repeating this strategy: We are at iteration $t \geq 1$

We set
$$\lambda^{(t)}=\lambda^{(t-1)}-\beta^{(t)}$$
 and update $\mathbf{x}^{(t)}$ as $\mathbf{x}^{(t)}=\mathbf{x}^{(t-1)}+\beta^{(t)}\mathbf{d}^{(t)}$

This means that the next non-zero entry holds at

$$n_{t+1}^{+} = \operatorname{argmin} \bigcup_{n \notin \mathcal{S}^{(t)}} Sub^{+} \left\{ \frac{\lambda^{(t-1)} + r_{n}^{(t-1)}}{1 - \left[\mathbf{A}^{\mathsf{T}}\mathbf{A}\boldsymbol{d}^{(t)}\right]_{n}}, \frac{\lambda^{(t-1)} - r_{n}^{(t-1)}}{1 + \left[\mathbf{A}^{\mathsf{T}}\mathbf{A}\boldsymbol{d}^{(t)}\right]_{n}} \right\}$$

We keep repeating this strategy: We are at iteration $t \geq 1$

We set
$$\lambda^{(t)}=\lambda^{(t-1)}-\beta^{(t)}$$
 and update $m{x}^{(t)}$ as $m{x}^{(t)}=m{x}^{(t-1)}+eta^{(t)}m{d}^{(t)}$

But remember that the residual gets updated in each step

It might be also the case that we get rid of one non-zero entry

So, we should also check this!

We keep repeating this strategy: We are at iteration $t \geq 1$

We set
$$\lambda^{(t)}=\lambda^{(t-1)}-\beta^{(t)}$$
 and update $\mathbf{x}^{(t)}$ as $\mathbf{x}^{(t)}=\mathbf{x}^{(t-1)}+\beta^{(t)}\mathbf{d}^{(t)}$

This happen if step-size

$$\beta_{-}^{(t)} = \min \bigcup_{n \in \mathcal{S}^{(t)}: d_n^{(t)} \neq 0} Sub^{+} \left\{ -\frac{x_n^{(t-1)}}{d_n^{(t)}} \right\}$$

exists and is smaller than $\beta_{\perp}^{(t)}$

We keep repeating this strategy: We are at iteration $t \geq 1$

We set
$$\lambda^{(t)}=\lambda^{(t-1)}-\beta^{(t)}$$
 and update $\mathbf{x}^{(t)}$ as $\mathbf{x}^{(t)}=\mathbf{x}^{(t-1)}+\beta^{(t)}\mathbf{d}^{(t)}$

Could it be the case that $\beta^{(t)}$ does not exist?

$$\forall n \in \mathcal{S}^{(t)}: d_n^{(t)} \neq 0: -\frac{x_n^{(t-1)}}{d_n^{(t)}} < 0$$

We keep repeating this strategy: We are at iteration $t \geq 1$

We set
$$\lambda^{(t)}=\lambda^{(t-1)}-\beta^{(t)}$$
 and update $\mathbf{x}^{(t)}$ as $\mathbf{x}^{(t)}=\mathbf{x}^{(t-1)}+\beta^{(t)}\mathbf{d}^{(t)}$

If this happens, we should get rid of the entry at

$$n_{t+1}^- = \operatorname{argmin} \bigcup_{n \in S^{(t)}: d_n^{(t)} \neq 0} Sub^+ \left\{ -\frac{x_n^{(t-1)}}{d_n^{(t)}} \right\}$$

We keep repeating this strategy: We are at iteration $t \geq 1$

We set
$$\lambda^{(t)}=\lambda^{(t-1)}-\beta^{(t)}$$
 and update $\mathbf{x}^{(t)}$ as $\mathbf{x}^{(t)}=\mathbf{x}^{(t-1)}+\beta^{(t)}\mathbf{d}^{(t)}$

How should we choose $\beta^{(t)}$?

If
$$\beta_{-}^{(t)}$$
 exists and $\beta_{-}^{(t)}<\beta_{+}^{(t)}$; then,

- We set $\beta^{(t)} = \beta_{-}^{(t)}$
- We update $S^{(t+1)} = S^{(t)} \{ n_{t+1}^- \}$

We keep repeating this strategy: We are at iteration $t \geq 1$

We set
$$\lambda^{(t)}=\lambda^{(t-1)}-\beta^{(t)}$$
 and update $\mathbf{x}^{(t)}$ as $\mathbf{x}^{(t)}=\mathbf{x}^{(t-1)}+\beta^{(t)}\mathbf{d}^{(t)}$

How should we choose $\beta^{(t)}$?

Otherwise . . .

- We set $\beta^{(t)} = \beta_+^{(t)}$
- We update $S^{(t+1)} = S^{(t)} \cup \{n_{t+1}^+\}$

We keep repeating this strategy: We are at iteration $t \geq 1$

We set
$$\lambda^{(t)}=\lambda^{(t-1)}-\beta^{(t)}$$
 and update $\mathbf{x}^{(t)}$ as $\mathbf{x}^{(t)}=\mathbf{x}^{(t-1)}+\beta^{(t)}\mathbf{d}^{(t)}$

We finally update the residual as

$$\mathbf{r}^{(t)} = \mathbf{A}^{\mathsf{T}} \left(\mathbf{A} \mathbf{x}^{(t)} - \mathbf{y} \right)$$

and use it together with $S^{(t+1)}$ in the next step

We keep repeating this strategy: We stop at iteration s

At iteration s, we should see

$$\lambda^{(s)} = 0$$

Attention!

There might be the case that n_{t+1}^+ and n_{t+1}^- are not unique!

- In this case, some extra searches should be performed
- If interested, check the textbook: Page 480, Remark 15.3

We keep repeating this strategy: We stop at iteration s

At iteration s, we should see

$$\lambda^{(s)} = 0$$

Side Note

Why we call it the homotopy method? Because

 $\mathbf{d}^{(1)}, \dots, \mathbf{d}^{(s)}$ describes a homotopy path from $\mathbf{0}$ to \mathbf{x}^*

where \mathbf{x}^{\star} is the basis pursuit solution

Final Points

Summary

- ℓ_1 -norm minimization gives
 - Best relaxation of the optimal recovery
- \bullet ℓ_1 -norm minimization recovers a sparse signal whose
 - # of non-zero entries < # of samples
- ℓ_1 -norm minimization has various forms
 - Basis pursuit
 - Basis pursuit with quadratic constraint
 - Basis pursuit denoising
 - I ASSO
- We know how to implement it via the homotopy method

What We Learn Next?

For some applications,

 ℓ_1 -norm minimization is still computationally complex

what we can do in such applications?

We could go for much lighter approaches

- Greedy algorithms
- Thresholding algorithms

We are going to learn them in the next part

Which Parts of Textbooks?

We are over with this part

I would suggest to go over the textbook

A Mathematical Introduction to Compressive Sensing

S. Foucart and H. Rauhut, Book, 2013

and study the following parts:

Chapter 3: Sections 3.1 and Chapter 15