EMBEDED C

Preprocessor directive:

- 1-#include
- 2-#define
- 3-conditional directive
- 4-#error /#warning

5-stringification and concatenation

- Preprocessor >>> text replacement >>> # (only preprocessor understand it)
- Except: #pragma >>> compiler directive.

1-#include:

- To include header file (.h)
- How to include it ???
- a-#include <file.h>

standard library (built in) (هيروح يدور فيهم)

b-#include "file.h" or #include "path"

current directory if he does not find the file then he searches in standard lib.

#include "path"

Path has 2 types:

- a-absolute path: pc (don't use it)
- b-relative path: related to your working directory.
- **EX**: file.c >>>>> temp.h #include"new/new2/temp.h"

Main.c >>>>lib.h #include"..lib.h"

File.c

new

Lib.h New2

Temp.h,main.c

2-#define(macro):

has 4 syntax

- 1-#define identifier replacement list
- 2-#define ident(param) replacement list
- 3-#define ident(...) replacement list
- (...) >>> anything else
- #define karem(...) printf(__VA_ARG__)
- 4-#define ident(param,...) replacement list
- #define fun(a,...) printf((__VA_ARG__)

Types:

- a-object like macro.
- #define name value
 ex: #define x 10
- b-function like macro(parameterized macro)
- #define add(x,y) x+y int z=add(x,y); >>>> int z=x+y

Hint: #define add(space)(x,y) x+y

Int z=add(3,4)>>>> int z=(x,y)x+y(3,4)

Advantage:

1-fast.

disadvantage:

- 1-code size.
- 2-no type checking.
- 3-difficult to debug.
- When to use it ??? if it will not repeated
- In embedded: "BIT_MATH"

Macro name rule:

1-start with letter or underscore

2-without spaces or without special characters(except _)

- It is recommended that name is uppercase.
- Not recommended to put ';' at the end.

```
What will happen if ?????
```

```
#define x 10
```

1-int y=x >>> expansion >>> int y=10

2-int x=50>>>> expansion>>>>10=50 (compilation error)

3-int xy >>>>no expansion

```
#define z 10
```

#define x 10

1-int y=x >>>int y=10

Notes:

```
# define x 10
```

#define x 20

• Will give warning (any line of code after (# define x 10

)will operate that x=10 and any line of code after #define x=20

)will operate that x=20 and).

>>To avoid warning

define x 10

#undef x

#define x 20

- (.c)files=(.i)files
- concatenation operator '\' backslash used with function like micro

#define x \

===#define x 5

Macro:

- 1-Professional code
- 2-Difficult in debugging

❖ Different between #define and Enum

#define	enum
 Text replacement in preprocessor stage We can use float sentence. Not Used inside switch No memory space. Only one value and one name Any value 	 Text replacement in compiler stage Must be const int (not float) Use inside switch Int size Sane value to different names -2^n: 2^(n-1)-1

❖ Different between #define and TYPEDEF

typedef	#define(macro)
 Text replacement in compiler stage Char >>> u8 used for types only 	Text replacement in preprocessor stageUsed with values as well
 #define struct stud *ptr; ptr * m,n;(two pointer) 	 #define ptr struct stud* Ptr x,y;(will replace x by pointer to struct and y is a variable of type struct stud)

3-conditional directive:

a-#if cond1

#elif cond2

#else

#endif

We can use it to comment code

if	#if
• compiler	 preprocessor
Check variables.	 Check macros not variables.

Usage:

1-comment

#if 0

#endif

2-configuration

b-#ifdef, #ifndef

• use in header file guard.

c-#if defined ,#elif defined

4-#error(stop compilation,error message) , **#warning** (only warning)

Used with #if not if as if check in runtime

5-stringification (#) and concatenation(##)

• stringification (#):convert it to string

```
EX: #define printf(x) printf(#x)
Printf( ....... ) >>>> #==""
Concatenation:
EX: #define conc(x,y) x##y
Int z=conc(3,8) >>>> int z=38
```

Hint:

```
#define max(a,b) a>b?a:b

int main()
{
   int x=9,y=7,z;
   z=max(x,y)*2;
   printf("%d",z);
}
Output:9
```

To avoid that use brackets

```
#define max(a,b) ((a)>(b)?(a):(b))
int main()
{
int x=9,y=7,z;
z=max(x,y)*2;
printf("%d",z);
}
```

Usage of preprocessor directive:

- 1-configuration
- 2-readability
- 3-portability

Hint: #line 20 will consider the next line to be 20

#pragma

Usage:

- 1-#pragma optimize ("",off) >>>> no optimization
- 2-#pragma once >>>> replace file guard

Note: compiler dependent but file guard is independent

- 3-#pragma startup[priority], #pragma exit[priority]
- 4-add memory section
- #pragma region name=""

Tool chain:

Define: .c → tool chain → exe file

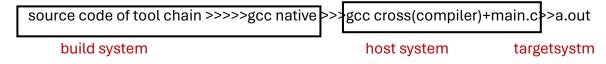
Consist of:

Preprocessor+compiler +assembler+linker+debugger+library +helper programmer (binary utilization) ex:objcopy,obj dumb,readelf

Types:

- a-cross: compile at pc and the output run on the target hardware.
- b-native: compile and run on the same machine.
- tool chain >>>> c.code

systems:

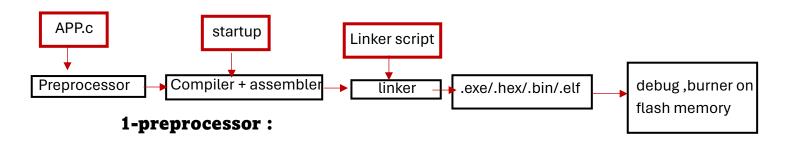


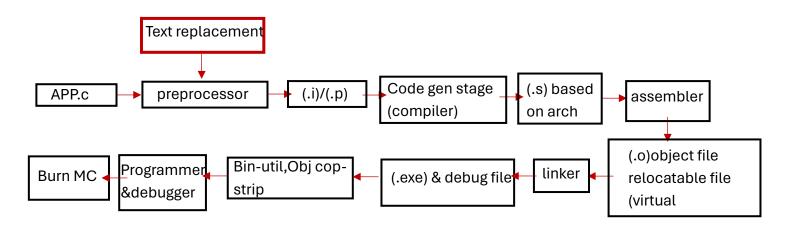
libraries:

- 1-static: we cannot see .c (object file(.a))
- 2-dynamic: we can see .c

how to ??? .c/.h >>>> .o >>> .a (static library)

Compilation process:





2-comiler:

- a-frontend stage: source code parsing + check syntax error
- syntax analysis /tokenization >>>> tokens
- o/p>>>> parse tree >>>> IR program

b-middle stage:

- Semantic analysis: check logical structure.
- optimization: why????

Reduce code size /exe-time/memory size.

How?? Has multilevel process

- 1-Remove Dead code.
- 2-Inline expansion of fun

- 3-Register allocation
- 4-Loop unrolling
- c-backend stage:
- code generation: convert into assembly from code language.
- memory allocation

note:

- compiler >>> symbol table (variable name/fun name)
- compiler >>> debug-info >>> debugger >>> digging code
- compiler >>> parse only one c file at a time

symbol has two sections:

- a-import: global & static variable which need from another file/fun (call)
- b-export: global & static variable/fun (implement)/debug info

3-linker:

- a-Symbol resolving
- b-section location: convert virtual address to physical address by using linker script it does section location by the help of locator(location counter) .operator

hint:

compilation flags

- -E >>> preprocessor stage (.i)
- -s >>>compilation (.s)
- -c>>> compile +ass (.o)

Booting sequence:

Power on /reset >>> entry point

Inst-life circle >>>(F,D,E)

Binary file contain:

a-startup code (.s/.c)

```
b-code (.text)
c-global & static >>> .data (init !=0)
d-global & static >>> .bss (uninit)
boot loader: software without os(startup code +main)
two cases:
case 1:
             baremetal SW
developer >>> EP >>> baremetal sw (startup+APP)
PC(program counter)>>> EP >>>reset of startup code
What startup code do ???
1-init_stack
2-copy (.data) from flash to RAM
3-reserve (.bss) in RAM
4-branch/jump >> main
case 2:
             bootloader
developer >>> EP >> bootloader (startup+main)
what boot loader do??
a-load source >>> dist at runtime
b-init modules
c-jump to startup (SW)>>>main
running moods:
a-ROM mode: case1
ram: stack/change data
rom: entry point/exe code +const data
b-RAM mode: case2
rom: entry point/code image
```

ram: stack/exe code /change da

ROM mood	RAM mood
 Very simple Fixed code address Relative small code Require smaller memory 	ComplexRelocatable modeFast why??(RAM > ROM)