Recitation 12

Warmup: How to deal with repeated elements in BucketSort and MergeSort?

/** Merge contents of arrays S1 and S2 into properly sized array S. */

```
2
      public static <K> void merge(K[] S1, K[] S2, K[] S, Comparator<K> comp) {
 3
        int i = 0, j = 0;
 4
        while (i + j < S.length) {
          if (j == S2.length \mid | (i < S1.length && comp.compare(S1[i], S2[j]) < 0))
 5
 6
            S[i+j] = S1[i++];
                                             // copy ith element of S1 and increment i
 7
          else
 8
            S[i+j] = S2[j++];
                                             // copy jth element of S2 and increment j
 9
      }
10
     Code Fragment 12.1: An implementation of the merge operation for a Java array.
      /** Merge-sort contents of array S. */
1
      public static <K> void mergeSort(K[] S, Comparator<K> comp) {
2
3
        int n = S.length;
        if (n < 2) return;
                                                             // array is trivially sorted
        // divide
5
6
        int mid = n/2;
        K[] S1 = Arrays.copyOfRange(S, 0, mid);
7
                                                             // copy of first half
        K[] S2 = Arrays.copyOfRange(S, mid, n);
                                                             // copy of second half
8
        // conquer (with recursion)
9
        mergeSort(S1, comp);
                                                             // sort copy of first half
10
11
        mergeSort(S2, comp);
                                                             // sort copy of second half
        // merge results
12
13
        merge(S1, S2, S, comp);
                                               // merge sorted halves back into original
14
```

Code Fragment 12.2: An implementation of the recursive merge-sort algorithm for a Java array (using the merge method defined in Code Fragment 12.1).

R-12.4 Is our array-based implementation of merge-sort given in Section 12.1.2 stable? Explain why or why not.

- C-12.47 Given two sets A and B represented as sorted sequences, describe an efficient algorithm for computing $A \oplus B$, which is the set of elements that are in A or B, but not in both.
- C-12.35 Suppose we are given an n-element sequence S such that each element in S represents a different vote for president, where each vote is given as an integer representing a particular candidate, yet the integers may be arbitrarily large (even if the number of candidates is not). Design an $O(n \log n)$ -time algorithm to see who wins the election S represents, assuming the candidate with the most votes wins.
- C-12.37 Consider the voting problem from Exercise C-12.35, but now suppose the integers 1 to k are used to identify k < n candidates. Design an O(n)-time algorithm to determine who wins the election.