

# Overview of last class

# Why layering?

dealing with complex systems:

- explicit structure allows identification, relationship of complex system's pieces
  - layered *reference model* for discussion
- modularization eases maintenance, updating of system
  - change of implementation of layer's service transparent to rest of system
  - e.g., change in gate procedure doesn't affect rest of system
- layering considered harmful? (view all the aspects)
  - Redundant functions
  - Cross-layer

# Network security

- **field of network security:**
  - how bad guys can attack computer networks
  - how we can defend networks against attacks
  - how to design architectures that are immune to attacks
- **Internet not originally designed with (much) security in mind**
  - *original vision*: “a group of mutually trusting users attached to a transparent network” 😊
  - Internet protocol designers playing “catch-up”
  - security considerations in all layers!

# Internet History

*1961-1972: Early packet-switching principles*

*1972-1980: Internetworking principles, new and proprietary nets*

*1980-1990: core protocols (TCP/IP/DNS), a proliferation of networks*

*1990, 2000's: commercialization, the Web, new apps*

*2000's-2010's: Quality of Service, wireless/mobile*

*2010's-present: Next generation Internet (Social Networking, IOT, SDN, CDN, Cloud, AI ...)*

# Introduction: summary

*covered a “ton” of material!*

- Internet overview
- what's a protocol?
- network edge, core, access network
  - packet-switching versus circuit-switching
  - Internet structure
- performance: loss, delay, throughput
- layering, service models
- security
- history

*you now have:*

- context, overview, “feel” of networking
- more depth, detail *to follow!*

**internetworking principles:**

- minimalism, autonomy - no internal changes required to interconnect networks
- best effort service model
- stateless routers
- decentralized control

# Class Today

# Chapter 2

## Application Layer

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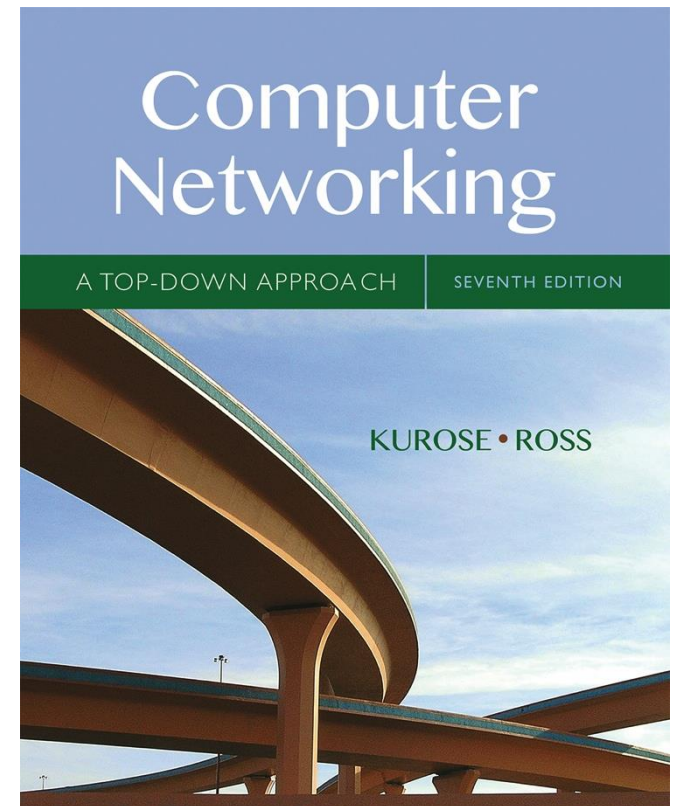
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## *Computer Networking: A Top Down Approach*

7<sup>th</sup> edition

Jim Kurose, Keith Ross

Pearson/Addison Wesley

April 2016

# Chapter 2: outline

## 2.1 principles of network applications

## 2.2 Web and HTTP

## 2.3 electronic mail

- SMTP, POP3, IMAP

## 2.4 DNS

## 2.5 P2P applications

## 2.6 video streaming and content distribution networks

## 2.7 socket programming with UDP and TCP



# Chapter 2: application layer

## our goals:

- conceptual, implementation aspects of network application protocols
  - transport-layer service models
  - client-server paradigm
  - peer-to-peer paradigm (**scalability**)
  - content distribution networks (**large scale with good QoS**)
- learn about protocols by examining popular application-level protocols
  - HTTP
  - FTP
  - SMTP / POP3 / IMAP
  - DNS
- creating network applications
  - socket API

# Some network apps

- e-mail
- web
- text messaging
- remote login
- P2P file sharing (BitTorrent, edonkey)
- multi-user network games
- streaming stored video (YouTube, Hulu, Netflix)
- voice over IP (e.g., Skype)
- real-time video conferencing
- social networking
- search
- ...
- ...

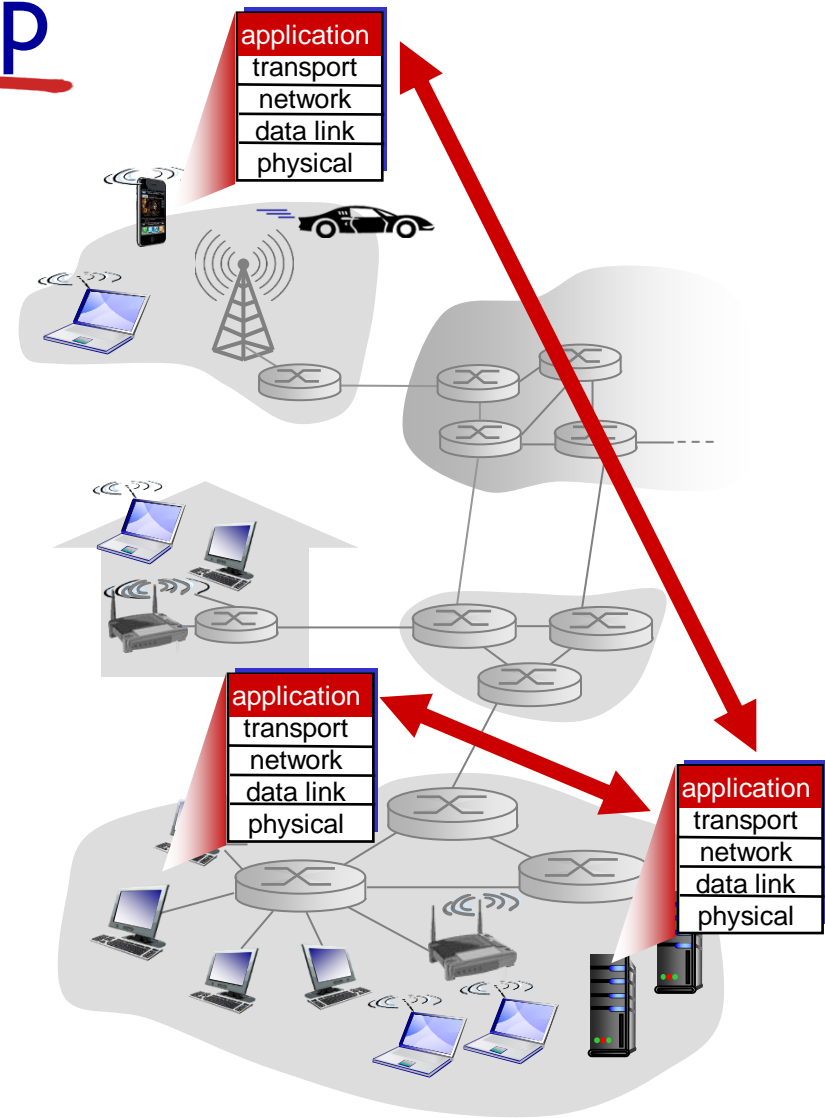
# Creating a network app

write programs that:

- run on (different) *end systems*
- communicate over network
- e.g., web server software communicates with browser software

no need to write software  
for network-core devices

- network-core devices do not run user applications
- applications on end systems allows for rapid app development, propagation



# Application architectures

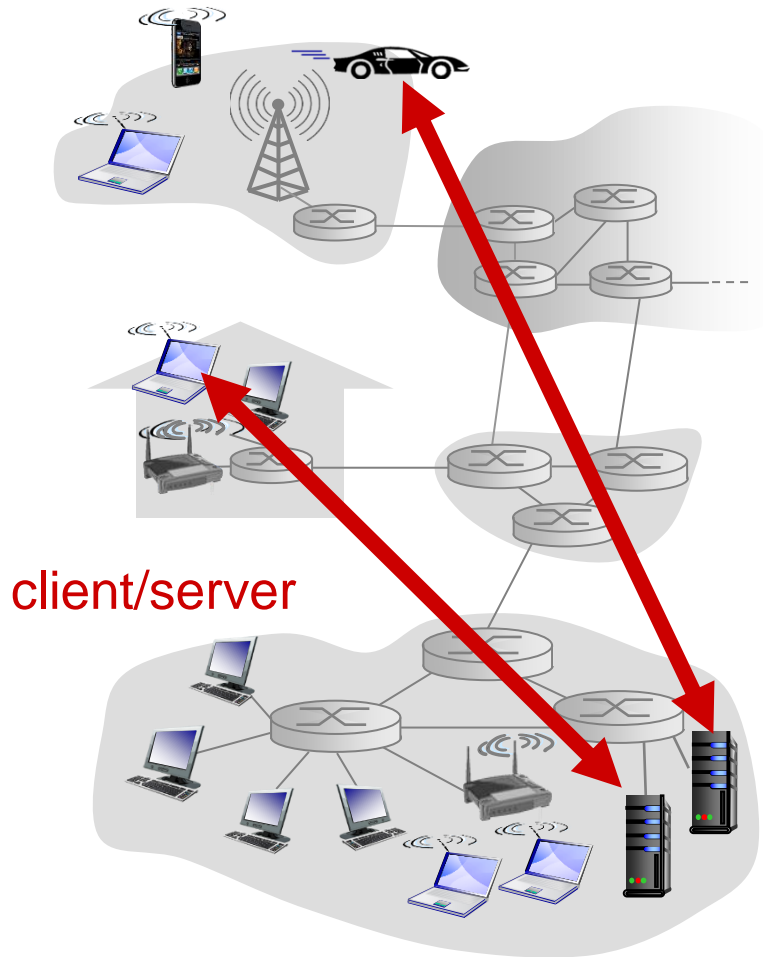
Two basic structure of applications:

- client-server
  - Bottleneck
  - Single point of failure
- peer-to-peer (P2P): scalability

Combine or use other techniques to enhance performance

No absolute winner forever: cloud computing

# Client-server architecture



## server:

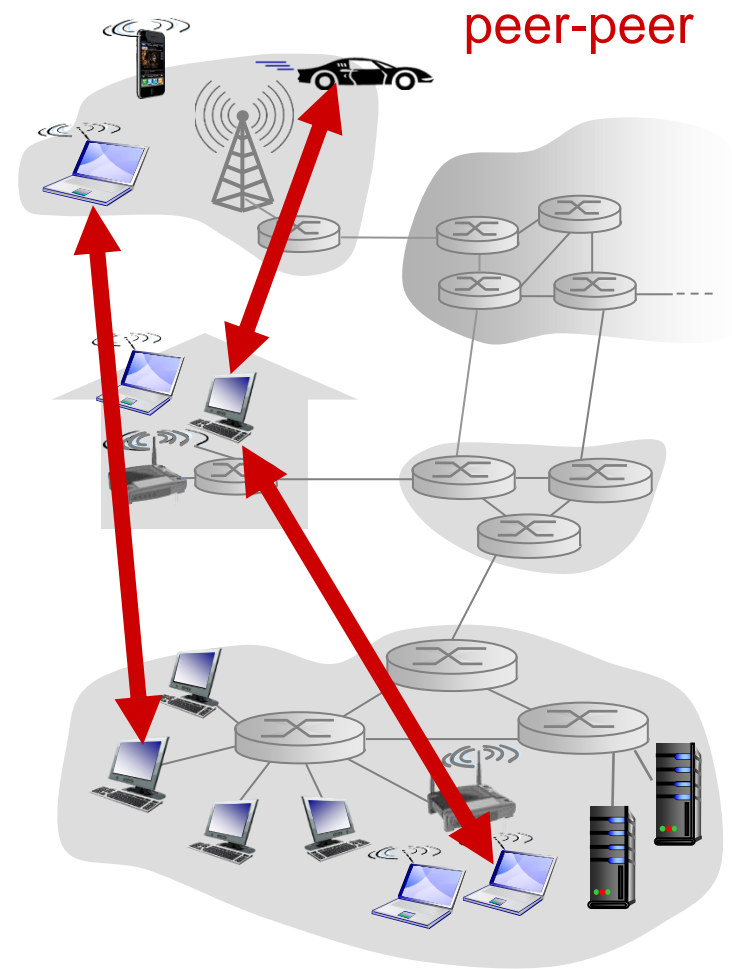
- always-on host
- permanent IP address
- data centers for scaling

## clients:

- communicate with server
- may be intermittently connected
- may have dynamic IP addresses
- do not communicate directly with each other

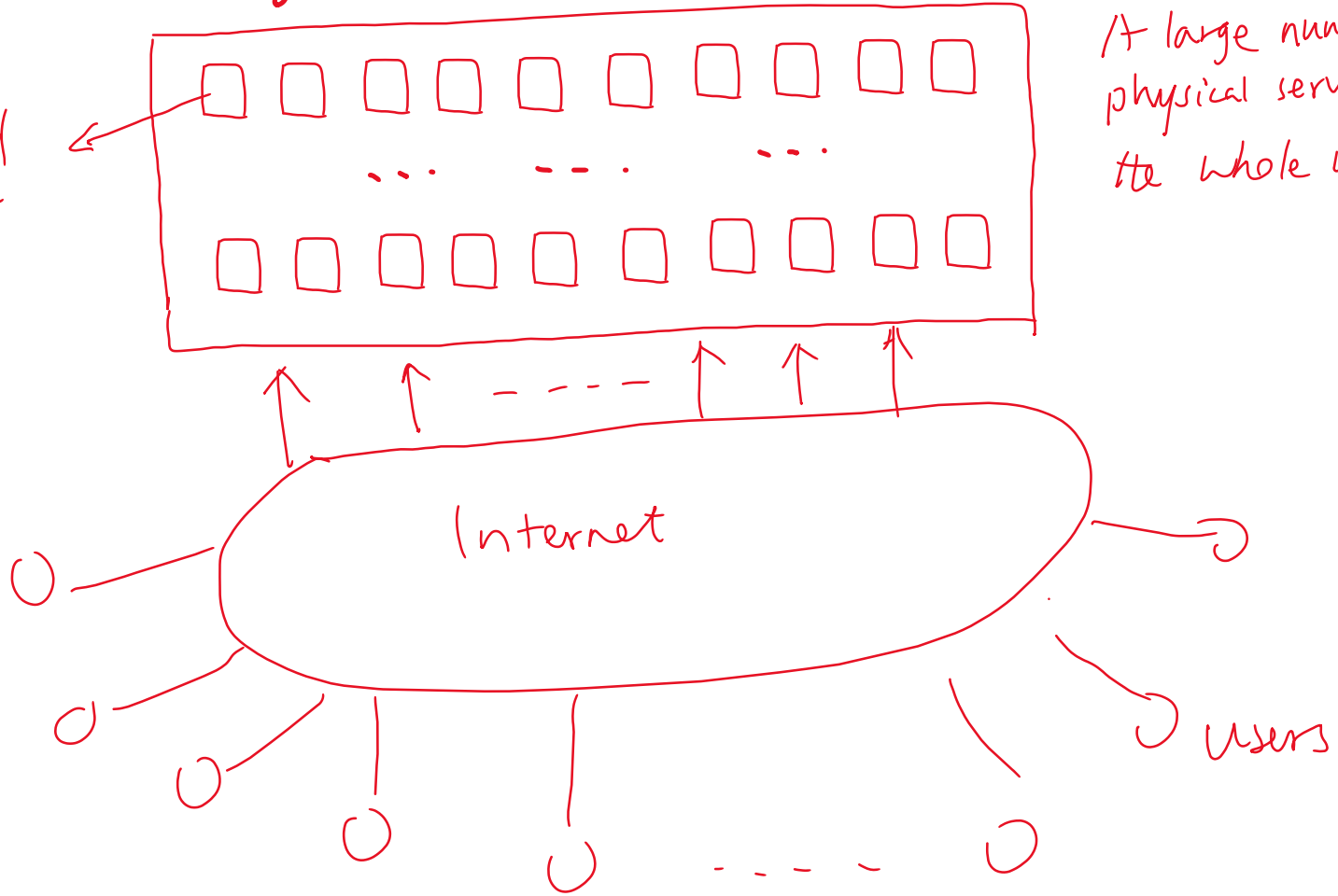
# P2P architecture

- no always-on server
- arbitrary end systems directly communicate
- peers request service from other peers, provide service in return to other peers
  - *self scalability* – new peers bring new service capacity, as well as new service demands
- peers are intermittently connected and change IP addresses
  - complex management



Google.com (a logic entity)

physical  
Server



A large number of physical servers over the whole world

# Processes communicating

*process*: program running within a host

- within same host, two processes communicate using **inter-process communication** (defined by OS)
- processes in different hosts communicate by exchanging **messages**

clients, servers

*client process*: process that initiates communication

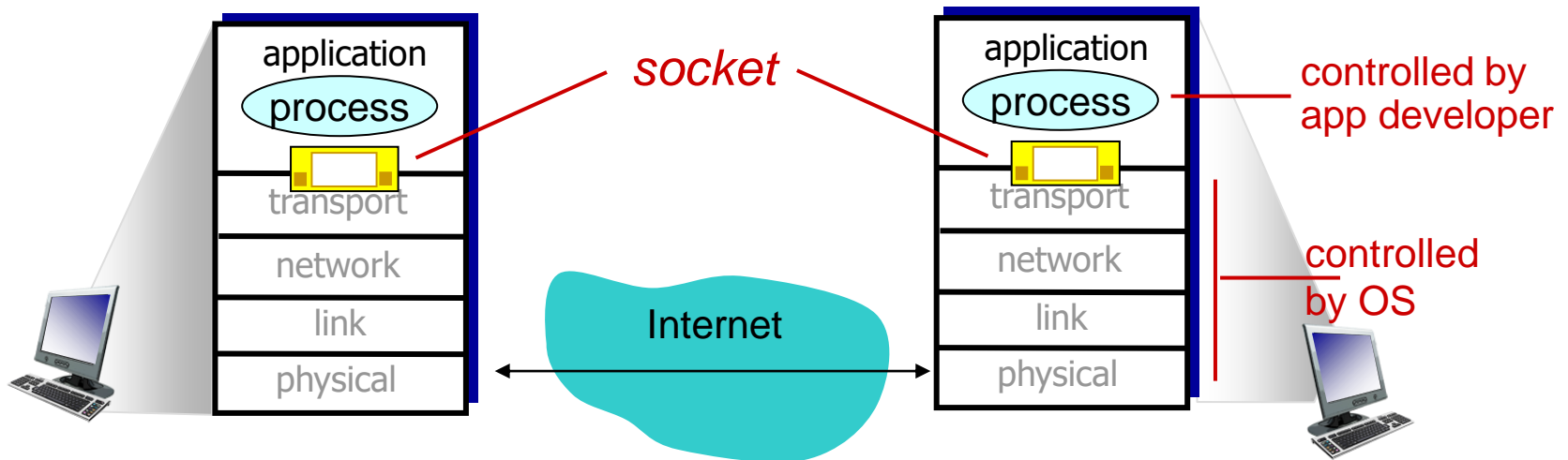
*server process*: process that waits to be contacted

- aside: applications with P2P architectures have client processes & server processes



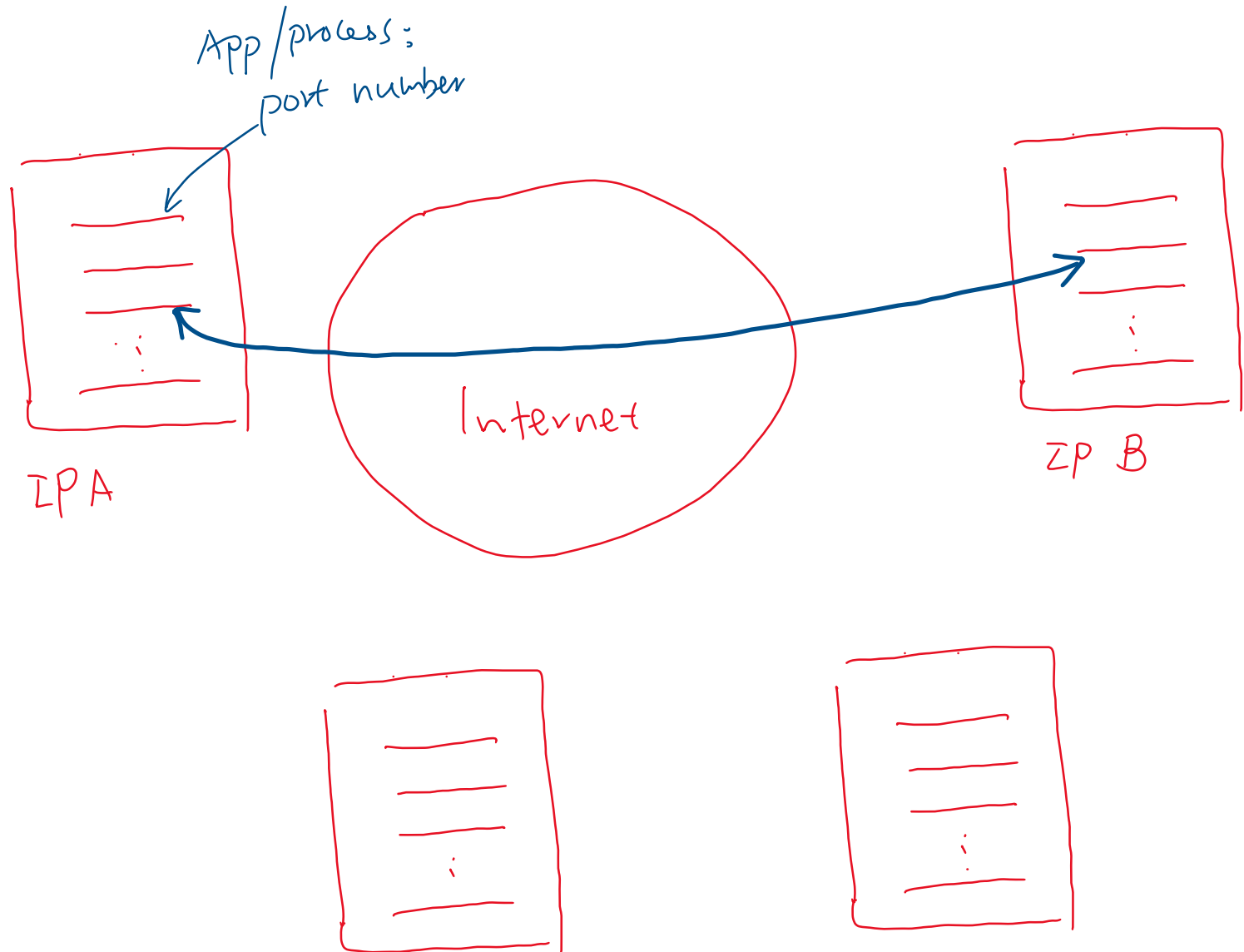
# Sockets

- process sends/receives messages to/from its **socket**
- socket analogous to door
  - sending process shoves message out door
  - sending process relies on transport infrastructure on other side of door to deliver message to socket at receiving process



# Addressing processes

- to receive messages, process must have *identifier*
- host device has unique 32-bit IP address
- Q: does IP address of host on which process runs suffice for identifying the process?
  - A: no, *many* processes can be running on same host
- *identifier* includes both **IP address** and **port numbers** associated with process on host.
- example port numbers:
  - HTTP server: 80
  - mail server: 25
- to send HTTP message to gaia.cs.umass.edu web server:
  - **IP address**: 128.119.245.12
  - **port number**: 80
- more shortly...



# App-layer protocol defines

- **types of messages exchanged,**
  - e.g., request, response
- **message syntax:**
  - what fields in messages & how fields are delineated
- **message semantics**
  - meaning of information in fields
- **rules** for when and how processes send & respond to messages

## **open protocols:**

- defined in RFCs
- allows for interoperability
- e.g., HTTP, SMTP

## **proprietary protocols:**

- e.g., Skype

# What transport service does an app need?

## data integrity/reliability

- some apps (e.g., file transfer, web transactions) require 100% reliable data transfer
- other apps (e.g., audio) can tolerate some loss

## timing

- some apps (e.g., Internet telephony, interactive games) require low delay to be “effective”

## throughput

- some apps (e.g., multimedia) require minimum amount of throughput to be “effective”
- other apps (“elastic apps”) make use of whatever throughput they get

## security

- encryption, data integrity, ...

# Transport service requirements: common apps

application	data loss	throughput	time sensitive
file transfer	no loss	elastic	no
e-mail	no loss	elastic	no
Web documents	no loss	elastic	no
real-time audio/video	loss-tolerant	audio: 5kbps-1Mbps video: 10kbps-5Mbps	yes, 100' s msec
stored audio/video	loss-tolerant	same as above	
interactive games	loss-tolerant	few kbps up	yes, few secs
text messaging	no loss	elastic	yes, 100' s msec yes and no

# Internet transport protocols services

## TCP service:

- *reliable transport* between sending and receiving process
- *flow control*: sender won't overwhelm receiver
- *congestion control*: throttle sender when network overloaded
- *does not provide*: timing, minimum throughput guarantee, security
- *connection-oriented*: setup required between client and server processes

## UDP service:

- *unreliable data transfer* between sending and receiving process
- *does not provide*: reliability, flow control, congestion control, timing, throughput guarantee, security, or connection setup,

Q: why bother? Why is there a UDP?

# Internet apps: application, transport protocols

<b>application</b>	<b>application layer protocol</b>	<b>underlying transport protocol</b>
e-mail	SMTP [RFC 2821]	TCP
remote terminal access	Telnet [RFC 854]	TCP
Web	HTTP [RFC 2616]	TCP
file transfer	FTP [RFC 959]	TCP
streaming multimedia	HTTP (e.g., YouTube), RTP [RFC 1889]	TCP or UDP
Internet telephony	SIP, RTP, proprietary (e.g., Skype)	TCP or UDP



# Securing TCP

## TCP & UDP

- no encryption
- cleartext passwds sent into socket traverse Internet in cleartext

## SSL (secure sockets layer)

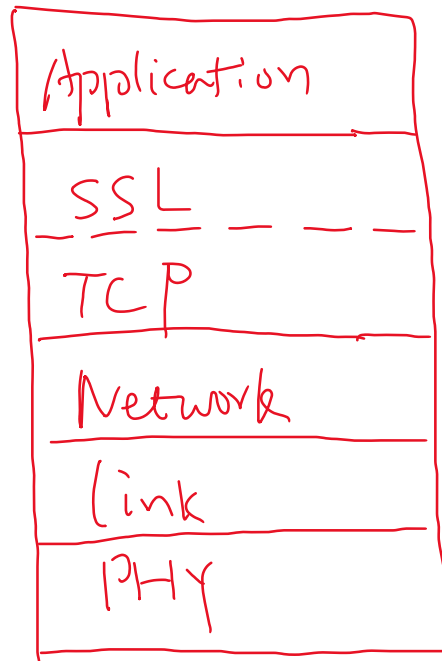
- provides encrypted TCP connection
- data integrity
- end-point authentication

## SSL is at app layer

- apps use SSL libraries, that “talk” to TCP

## SSL socket API

- cleartext passwords sent into socket traverse Internet encrypted
- see Chapter 8



Transport layer Security

TLS

# Chapter 2: outline

2.1 principles of network applications

2.2 Web and HTTP

2.3 electronic mail

- SMTP, POP3, IMAP

2.4 DNS

2.5 P2P applications

2.6 video streaming and content distribution networks

2.7 socket programming with UDP and TCP

# Web and HTTP

*First, a review...*

- *web page* consists of *objects*
- object can be HTML file, JPEG image, Java applet, audio file,...
- web page consists of *base HTML-file* which includes *several referenced objects*
- each object is addressable by a *URL*, e.g.,


`www.someschool.edu/someDept/pic.gif`

host name

path name

Yu Cheng's Homepage x New tab x | +

← ↻ ↗ Not secure | ece.iit.edu/~yucheng/

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**Open positions for PhD students**

Yu Cheng received the B.E. degree in Electrical Engineering from the University of Waterloo, Waterloo, Ontario, Canada, in 1995, and the M.Sc. degree in Electrical Engineering from the University of Toronto, Ontario, Canada, in July 2006. Since August 2006, he has been an Assistant Professor in the Department of Electrical and Computer Engineering, Illinois Institute of Technology, Chicago, Illinois, USA, and he is now a full Professor. His research interests include wireless network performance analysis, network optimization, and machine learning. He received a Postdoctoral Fellowship Award from the Natural Sciences and Engineering Research Council of Canada (NSERC) in 2007, a CAREER Award from the National Science Foundation (NSF) in 2008, and a Best Paper Award from the IEEE Communications Society in 2011. He has co-edited the book "Performance Analysis of Wireless Networks" (Springer, 2011). He is currently serving as the General Chair for the IEEE Conference on Communications, Computer and Security (CCS) 2017, and the General Chair for the IEEE Conference on Network and Information Security (NIS) 2017. He is an IEEE senior member.

**Research Interests**

- Performance analysis, protocol design, and energy efficiency of wireless networks
- Machine learning in wireless networks
- Age of information or information freshness optimization

**from 2021 Spring) in the areas of age of information or machine learning over wireless networks**

from Tsinghua University, Beijing, China, in 1995 and 1998, respectively, and the Ph.D. degree in Electrical and Computer Engineering from the University of Toronto, Ontario, Canada, in July 2006, he was a postdoctoral research fellow in the Department of Electrical and Computer Engineering, University of Toronto, Ontario, Canada. Since August 2006, he has been an Assistant Professor in the Department of Electrical and Computer Engineering, Illinois Institute of Technology, Chicago, Illinois, USA, and he is now a full Professor. His research interests include wireless network performance analysis, network optimization, and machine learning. He received a Postdoctoral Fellowship Award from the Natural Sciences and Engineering Research Council of Canada (NSERC) in 2007, a CAREER Award from the National Science Foundation (NSF) in 2008, and a Best Paper Award from the IEEE Communications Society in 2011. He has co-edited the book "Performance Analysis of Wireless Networks" (Springer, 2011). He is currently serving as the General Chair for the IEEE Conference on Communications, Computer and Security (CCS) 2017, and the General Chair for the IEEE Conference on Network and Information Security (NIS) 2017. He is an IEEE senior member.

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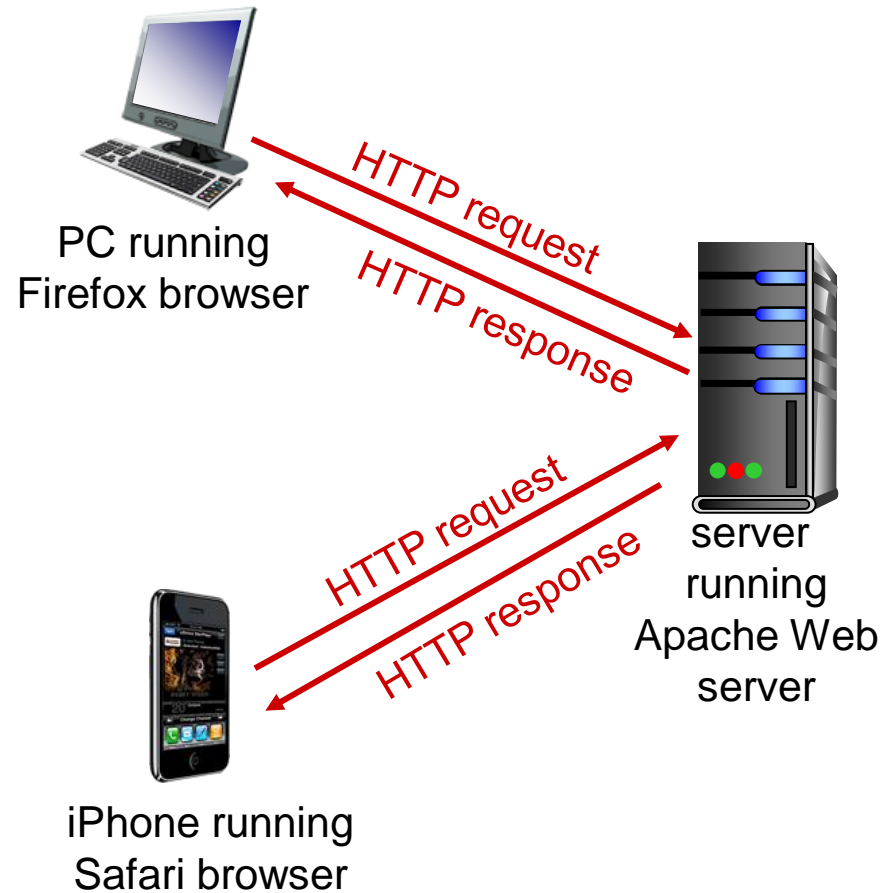
Search

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1/18/2023

# HTTP overview

## HTTP: hypertext transfer protocol

- Web's application layer protocol
- client/server model
  - **client**: browser that requests, receives, (using HTTP protocol) and "displays" Web objects
  - **server**: Web server sends (using HTTP protocol) objects in response to requests



# HTTP overview (continued)

## *uses TCP:*

- client initiates TCP connection (creates socket) to server, port 80
- server accepts TCP connection from client
- HTTP messages (application-layer protocol messages) exchanged between browser (HTTP client) and Web server (HTTP server)
- TCP connection closed

## *HTTP is “stateless”*

- server maintains no information about past client requests

*aside*

### protocols that maintain “state” are complex!

- past history (state) must be maintained
- if server/client crashes, their views of “state” may be inconsistent, must be reconciled

# HTTP connections

## *non-persistent HTTP*

- at most one object sent over TCP connection
  - connection then closed
- downloading multiple objects required multiple connections

## *persistent HTTP*

- multiple objects can be sent over single TCP connection between client, server



# Non-persistent HTTP

suppose user enters URL:

`www.someSchool.edu/someDepartment/home.index`

(contains text,  
references to 10  
jpeg images)

1a. HTTP client initiates TCP connection to HTTP server (process) at `www.someSchool.edu` on port 80

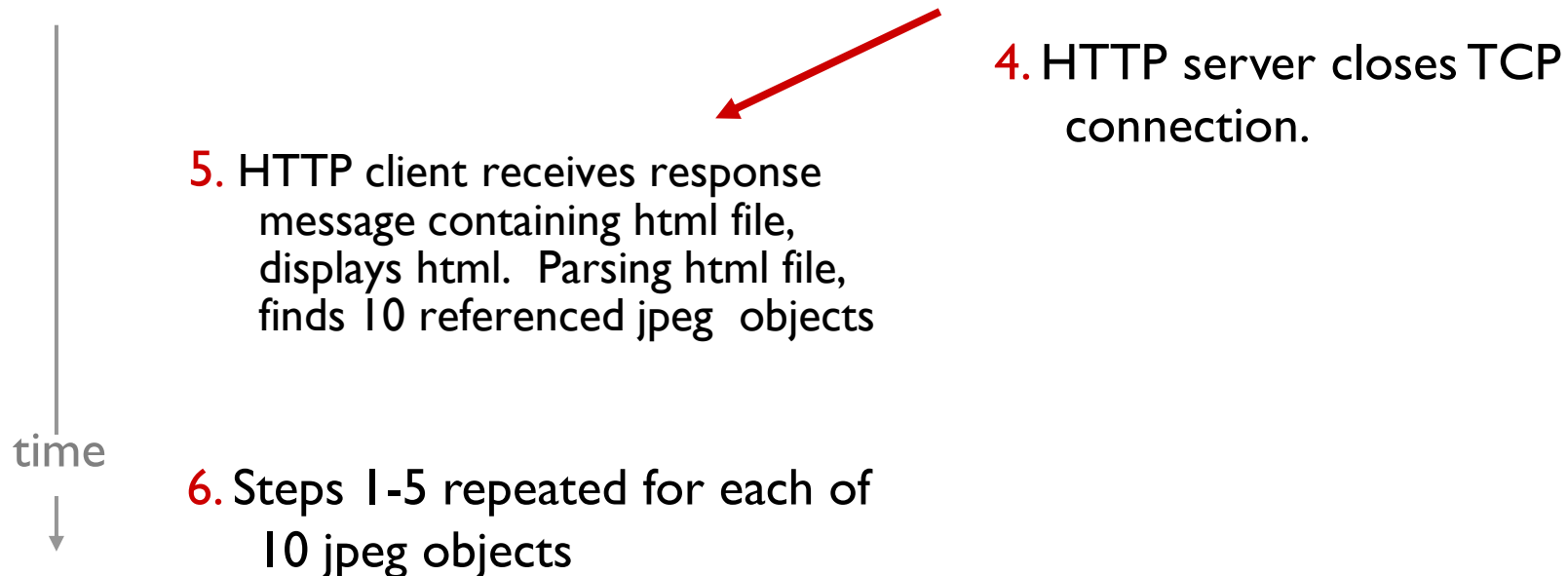
1b. HTTP server at host `www.someSchool.edu` waiting for TCP connection at port 80. “accepts” connection, notifying client

2. HTTP client sends HTTP *request message* (containing URL) into TCP connection socket. Message indicates that client wants object `someDepartment/home.index`

3. HTTP server receives request message, forms *response message* containing requested object, and sends message into its socket

time  
↓

# Non-persistent HTTP (cont.)

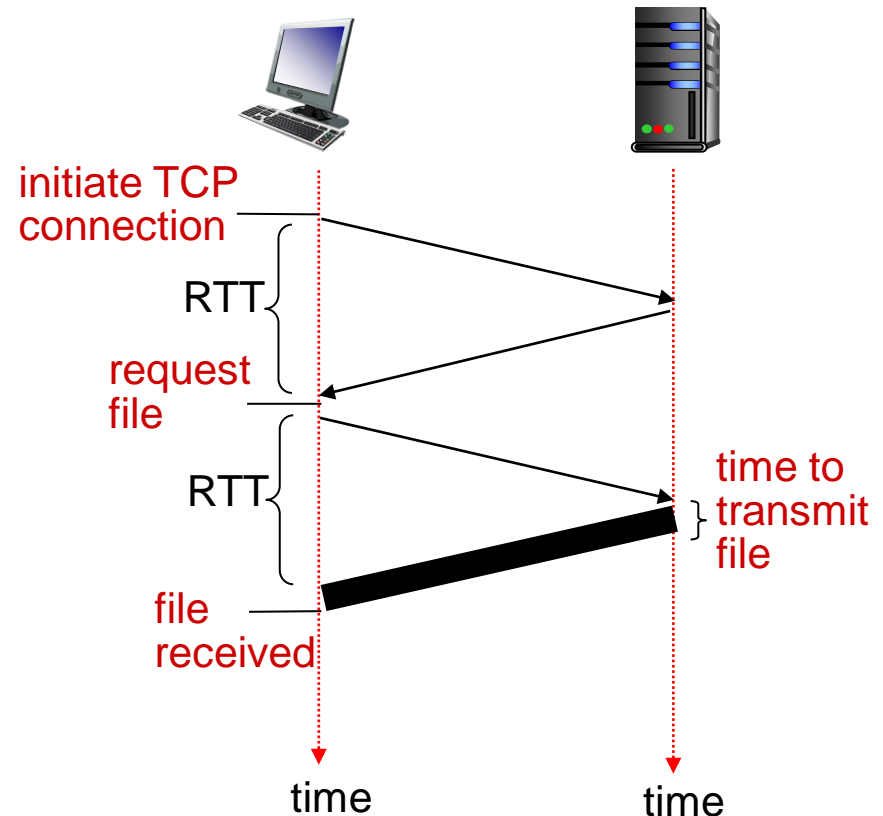


# Non-persistent HTTP: response time

**RTT (definition):** time for a small packet to travel from client to server and back

**HTTP response time:**

- one RTT to initiate TCP connection
- one RTT for HTTP request and first few bytes of HTTP response to return
- file transmission time
- non-persistent HTTP response time =  
 $2\text{RTT} + \text{file transmission time}$



# Persistent HTTP

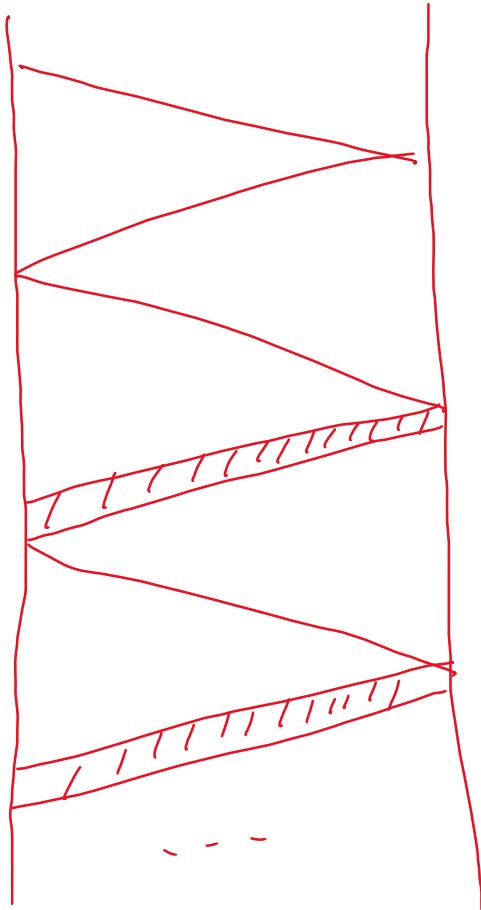
## *non-persistent HTTP issues:*

- requires 2 RTTs per object
- OS overhead for *each* TCP connection
- browsers often open parallel TCP connections to fetch referenced objects

## *persistent HTTP:*

- server leaves connection open after sending response
- subsequent HTTP messages between same client/server sent over open connection
- client sends requests as soon as it encounters a referenced object
- as little as one RTT for all the referenced objects

## Persistent HTTP



For  $N$  objects:

Persistent :

$$RTT + N(RTT + \text{object transmission time})$$

---

Non-persistent

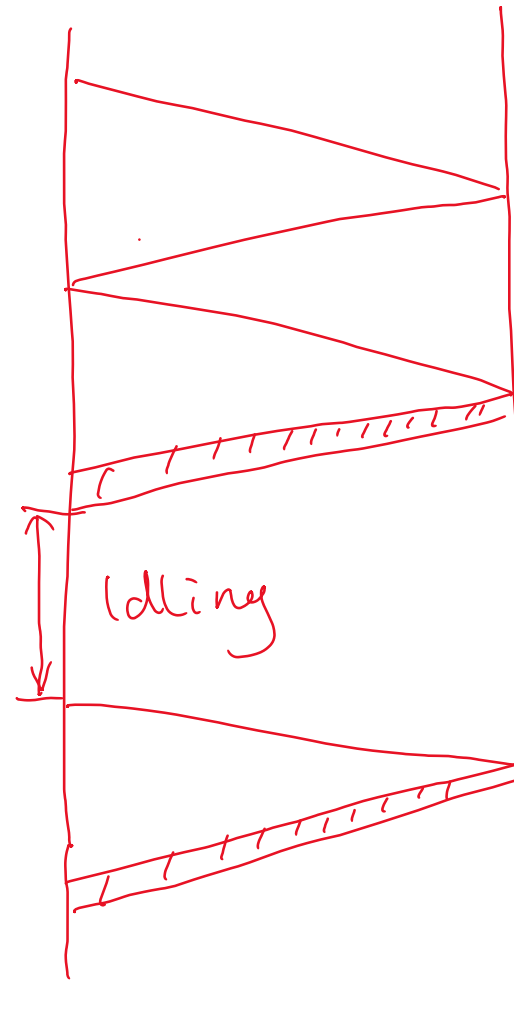
$$N(2RTT + \text{object transmission time})$$

but can be just

$2RTT + \text{object transmission time}$   
with  $N$  parallel TCP connections  
(i.e.  $N$  times bandwidth)

## *persistent connection management*

- Maximum lifetime
- Timeout



# HTTP request message

- two types of HTTP messages: *request, response*
- **HTTP request message:**
  - ASCII (human-readable format)
  - <https://developer.mozilla.org/en-US/docs/Web/HTTP/Headers>

request line  
(GET, POST,  
HEAD commands)

header  
lines

carriage return,  
line feed at start  
of line indicates  
end of header lines

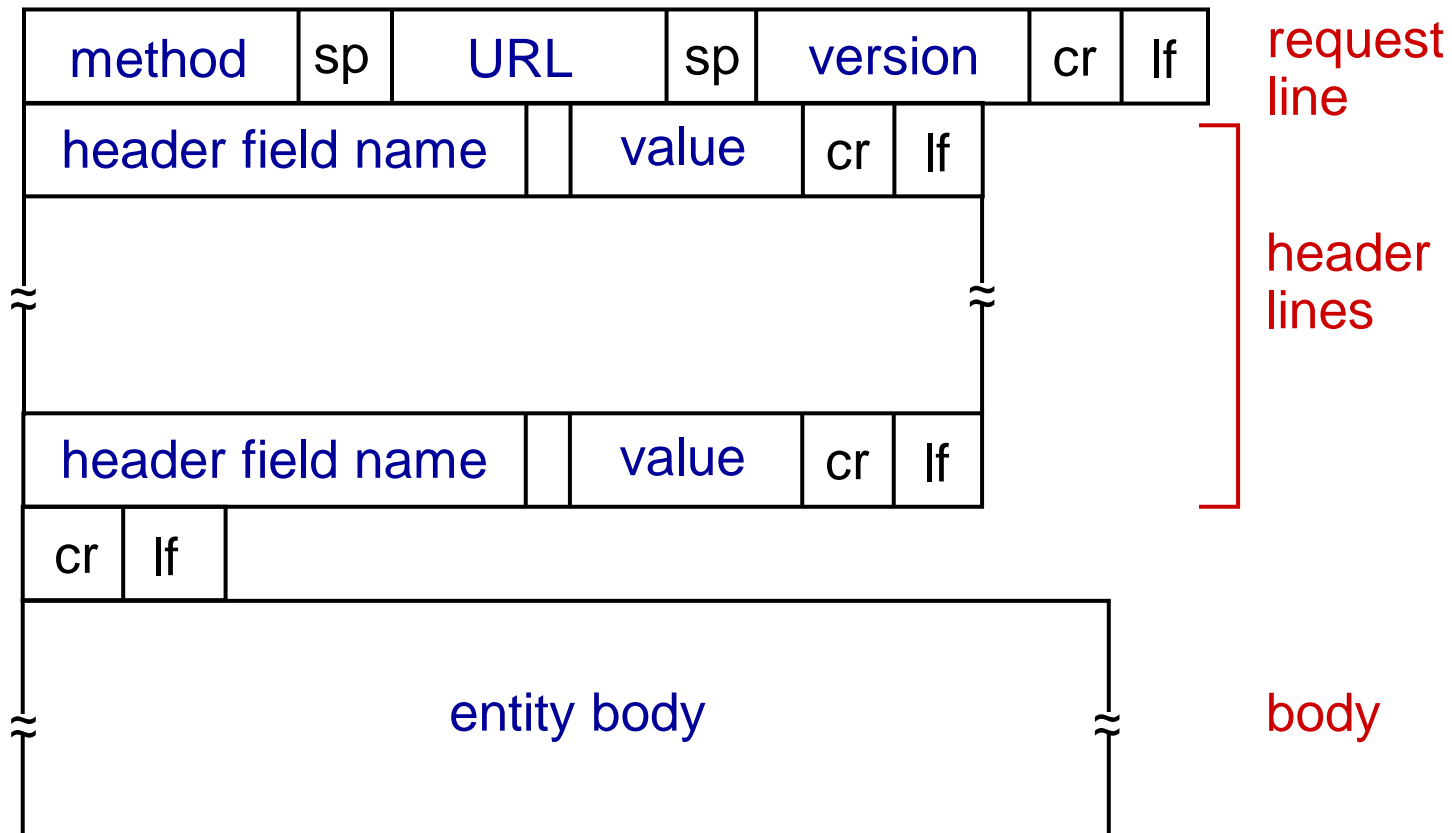
```
GET /index.html HTTP/1.1\r\n
Host: www-net.cs.umass.edu\r\n
User-Agent: Firefox/3.6.10\r\n
Accept: text/html,application/xhtml+xml\r\n
Accept-Language: en-us,en;q=0.5\r\n
Accept-Encoding: gzip,deflate\r\n
Accept-Charset: ISO-8859-1,utf-8;q=0.7\r\n
Keep-Alive: 115\r\n
Connection: keep-alive\r\n
\r\n
```

carriage return character

line-feed character

\* Check out the online interactive exercises for more  
examples: [http://gaia.cs.umass.edu/kurose\\_ross/interactive/](http://gaia.cs.umass.edu/kurose_ross/interactive/)

# HTTP request message: general format





# Uploading form input

## POST method:

- web page often includes form input
- input is uploaded to server in entity body

## URL method:

- uses GET method
- input is uploaded in URL field of request line:

`www.somesite.com/animalsearch?monkeys&banana`

# Method types

## HTTP/1.0:

- GET
- POST
- HEAD
  - asks server to leave requested object out of response

## HTTP/1.1:

- GET, POST, HEAD
- PUT
  - uploads file in entity body to path specified in URL field
- DELETE
  - deletes file specified in the URL field

# HTTP response message

status line

(protocol

status code

status phrase)

header  
lines

data, e.g.,  
requested  
HTML file

```
HTTP/1.1 200 OK\r\n
Date: Sun, 26 Sep 2010 20:09:20 GMT\r\n
Server: Apache/2.0.52 (CentOS)\r\n
Last-Modified: Tue, 30 Oct 2007 17:00:02
GMT\r\n
ETag: "17dc6-a5c-bf716880"\r\n
Accept-Ranges: bytes\r\n
Content-Length: 2652\r\n
Keep-Alive: timeout=10, max=100\r\n
Connection: Keep-Alive\r\n
Content-Type: text/html; charset=ISO-8859-
1\r\n
\r\n
data data data data data ...
```

\* Check out the online interactive exercises for more  
examples: [http://gaia.cs.umass.edu/kurose\\_ross/interactive/](http://gaia.cs.umass.edu/kurose_ross/interactive/)

# HTTP response status codes

- status code appears in 1st line in server-to-client response message.
- some sample codes:

## **200 OK**

- request succeeded, requested object later in this msg

## **301 Moved Permanently**

- requested object moved, new location specified later in this msg (Location:)

## **400 Bad Request**

- request msg not understood by server

## **404 Not Found**

- requested document not found on this server

## **505 HTTP Version Not Supported**

# Trying out HTTP (client side) for yourself

1. Telnet to your favorite Web server:

```
telnet gaia.cs.umass.edu 80
```

{ opens TCP connection to port 80  
(default HTTP server port)  
at gaia.cs.umass.edu.  
anything typed in will be sent  
to port 80 at gaia.cs.umass.edu

2. type in a GET HTTP request:

```
GET /kurose_ross/interactive/index.php HTTP/1.1  
Host: gaia.cs.umass.edu
```

{ by typing this in (hit carriage  
return twice), you send  
this minimal (but complete)  
GET request to HTTP server

3. look at response message sent by HTTP server!  
(or use Wireshark to look at captured HTTP request/response)