

## Algorithms: CSE 202 — Homework 2

### **Problem 1: Maximum Length Chain of Subwords**

Input: We are given a set of  $n$  distinct strings of length at most  $k$  over a finite alphabet  $\sigma$ .

Output: A sequence of strings that form a chain under the (consecutive) subword relation; i.e., if the output is  $w_1, w_2, \dots, w_t$  then we can write  $w_{i+1} = uw_iv$  for some strings  $u, v$ .

Find a chain of maximum length.

### **Problem 2: Business plan**

Consider the following problem. You are designing the business plan for a start-up company. You have identified  $n$  possible projects for your company, and for,  $1 \leq i \leq n$ , let  $c_i > 0$  be the minimum capital required to start the project  $i$  and  $p_i > 0$  be the profit after the project is completed. You also know your initial capital  $C_0 > 0$ . You want to perform at most  $k$ ,  $1 \leq k \leq n$ , projects before the IPO and want to maximize your total capital at the IPO. Your company cannot perform the same project twice.

In other words, you want to pick a list of up to  $k$  distinct projects,  $i_1, \dots, i_{k'}$  with  $k' \leq k$ . Your *accumulated capital* after completing the project  $i_j$  will be  $C_j = C_0 + \sum_{h=1}^{j-1} p_{i_h}$ . The sequence must satisfy the constraint that you have sufficient capital to start the project  $i_{j+1}$  after completing the first  $j$  projects, i.e.,  $C_j \geq c_{i_{j+1}}$  for each  $j = 0, \dots, k' - 1$ . You want to maximize the final amount of capital,  $C_{k'}$ .

### **Problem 3: Minimum Cost Sum**

You are given a sequence  $a_1, a_2, \dots, a_n$  of nonnegative integers, where  $n \geq 1$ . You are allowed to take any two numbers and add them to produce their sum. However, each such addition has a cost which is equal to the sum. The goal is to find the sum of all the numbers in the sequence with minimum total cost. Describe an algorithm for finding the sum of the numbers in the sequence with minimum total cost. Argue the correctness of your algorithm.

### **Problem 4: Speech recognition**

We can use dynamic programming on a directed graph  $G = (V, E)$  for speech recognition. Each edge  $(u, v) \in E$  is labeled with a sound  $\sigma(u, v)$  from a finite set  $\Sigma$  of sounds. The labeled graph is a formal model of a person speaking a restricted language. Each path in the graph starting from a distinguished vertex  $v_0 \in V$  corresponds to a possible sequence of sounds produced by the model. The label of a directed path is defined to be the concatenation of the labels of the edges on that path.

Describe an efficient algorithm that, given an edge-labeled graph  $G$  with distinguished vertex  $v_0$  and a sequence  $s = (\sigma_1, \dots, \sigma_k)$  of characters from  $\Sigma$ , returns a path in  $G$  that begins at  $v_0$  and has  $s$  as its label, if any such path exists. Otherwise, the algorithm should return NO-SUCH-PATH. Analyze the running time of the algorithm. Clearly write any dynamic programming formulation you may use to solve this problem.