Internet Infrastructure and Protocols (COMP5311)

Assignment Three (due on 30 Nov. 2011) Each question carries 8 marks, unless stated otherwise.

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1) Consider the IP network in Figure 1 that is subnetted with a fixed-length subnet mask. The numbers next to the LAN segments indicate their subnet numbers. The routers use RIP-I to share the routing information with split horizon and poisonous reverse. A hop count of 16 is used to represent infinity.

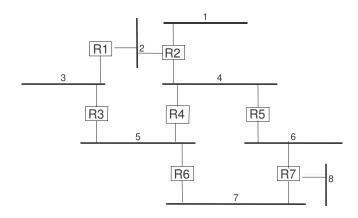


Fig. 1. An IP network running RIP-1.

A hacker gets hold of R2 and sends out a false route for subnet 8. Specifically, R2 sends out the following distance vector on subnet 4 and a similar distance vector on subnet 2.

Destination (subnet number)	Number of hops
1	1
2	1
3	2
4	(not included)
5	16
6	16
7	16
8	1 (a false value)

Which subnets (1-7) will be affected by this false route when the hosts on these subnets send packets to a destination on subnet 8, and why?

2) Consider case (a) in Figure 2. Routers A to D use a distance vector routing protocol with source tracing capability. You may use the hop count as the routing metric. Shortly after the routing protocol converges,

the link A-B breaks, describe how this routing protocol prevents routing loops from forming. Repeat it for case (b).

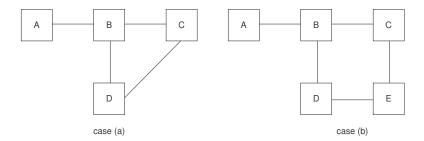


Fig. 2. Two IP networks running a distance vector routing protocol with source tracing capability.

- 3) To verify the answers to question 2 of assignment 2, we configured the hosts H1 (Ibm_b3:f0:cf) and H2 (Ibm_16:8b:06) with different subnet addresses, and the hosts communicated with each other through router R (Tp-LinkT_c9:41:05), as depicted in Figure 3. Their configurations and forwarding tables were:
 - R(eth1.0): 123.123.1.10; (eth1.1): 123.123.2.10; MAC Addr: 00:21:27:C9:41:05
 - H1 (eth0): 123.123.2.1/24; gateway 123.123.2.10; MAC Addr: 00:11:25:B3:F0:CF
 - H2 (eth0): 123.123.1.1/24; gateway 123.123.1.10; MAC Addr: 00:11:25:16:8B:06.

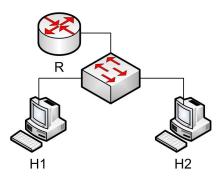


Fig. 3. Router R connected to hosts H1 and H2 through a switch.

	Destination	Gateway	Net mask	Interface
H1's routing table:	table: 123.123.2.0 123.123.2.1 2		255.255.255.0	eth0
	0.0.0.0	123.123.2.10	0.0.0.0	eth0
	Dantination	Catarra	Not most	Interface
	Destination	Gateway	Net mask	Interrace
H2's routing table:		123.123.1.1	255.255.255.0	eth0

	Destination	Gateway	Net mask	Interface
R's routing table:	123.123.2.0	123.123.2.10	255.255.255.0	eth1
	123.123.1.0	123.123.1.10	255.255.255.0	eth1

Moreover, their ARP caches were initially empty. Then H1 sent ping requests (ICMP echo requests) to H2. The following list records what actually happened.

- H1 broadcasted an ARP request for the MAC addr of 123.123.2.10.
- R replied the ARP request with its MAC addr (00:21:27:C9:41:05).
- H1 sent out a ping request to H2 (123.123.1.1).
- R received the ping request, sent an ICMP Redirect message to H1, and forwarded the ping request to H2.
- H2 replied the ping request by sending an ICMP echo reply to R. Similar to before, R sent an ICMP Redirect message to H2.
- H2 broadcasted an ARP request for the MAC address of 123.123.2.1, and H1 replied.
- H1 sent a second ping request to H2 via R, and ICMP Redirect messages were again issued by R to H1 and H2.
- Starting from the third ping request and on, H1 and H2 were able to exchange ping requests and replies directly, without going through R.

Figure 4 shows the Wireshark capture at H1 with the promiscuous mode off (i.e., captures only packets destined to H1 and sent by H1). Note that since this trace was captured at H1, it does not contain the communication between R and H2.

No	Time	Source	Destination	Protocol	Info
	1 0.000000	Ibm_b3:f0:cf	Broadcast	ARP	Who has 123.123.2.10? Tell 123.123.2.1
	2 0.000145	Tp-LinkT_c9:41:05	Ibm_b3:f0:cf	ARP	123.123.2.10 is at 00:21:27:c9:41:05
	3 0.000156	123.123.2.1	123.123.1.1	ICMP	Echo (ping) request
	4 0.000314	123.123.2.10	123.123.2.1	ICMP	Redirect (Redirect for host)
	5 0.000585	Tp-LinkT_c9:41:05	Broadcast	ARP	Who has 123.123.1.1? Tell 123.123.1.10
	6 0.000940	123.123.1.1	123.123.2.1	ICMP	Echo (ping) reply
	7 0.001103	Ibm_16:8b:06	Broadcast	ARP	Who has 123.123.2.1? Tell 123.123.1.1
	8 0.001117	Ibm_b3:f0:cf	Ibm_16:8b:06	ARP	123.123.2.1 is at 00:11:25:b3:f0:cf
	9 0.993191	123.123.2.1	123.123.1.1	ICMP	Echo (ping) request
	10 0.993334	123.123.2.10	123.123.2.1	ICMP	Redirect (Redirect for host)
	11 0.993587	123.123.1.1	123.123.2.1	ICMP	Echo (ping) reply
	12 1.992194	123.123.2.1	123.123.1.1	ICMP	Echo (ping) request
	13 1.992443	123.123.1.1	123.123.2.1	ICMP	Echo (ping) reply
	14 2.994116	123.123.2.1	123.123.1.1	ICMP	Echo (ping) request
	15 2.994339	123.123.1.1	123.123.2.1	ICMP	Echo (ping) reply

Fig. 4. A partial Wireshark capture of the packets when H1 pings H2 at H1 with the promiscuous mode off.

Given the information above, answer the questions below with succinct explanation.

- a) (2 marks) What is the destination MAC address in the frame of packet 3 in the Wireshark trace?
- b) (2 marks) What is the value in the "Gateway address" field in the ICMP Redirect message (i.e., packets 4 and 10 in the Wireshark trace)?
- c) (2 marks) How many ICMP header(s) is (are) included in the IP packet that carries the ICMP Redirect message?

- d) (2 marks) The Wireshark at H1 could capture R's ARP requests for H2's MAC address but not the ARP reply. What is the reason for that?
- e) (2 marks) H1 could send ping requests directly to H2 without first sending out an ARP request for H2's MAC address. What is the reason for that?
- f) (4 marks) If we examine the Wireshark trace captured at R with the promiscuous mode off, which packets in Figure 4 will *not* show up in the trace?
- g) (4 marks) If we examine the Wireshark trace captured at H2 with the promiscuous mode off, which packets in Figure 4 will *not* show up in the trace?
- h) (2 marks) Are the ICMP headers in packets 3 and 12 in Figure 4 identical?
- i) (2 marks) If the router's two interfaces ethl.0 and ethl.1 are now configured with two separate physical interfaces each of which is connected directly to a host, will R still send the ICMP Redirect messages? Note that the routing tables remain the same.
- 4) Figure 5 shows a similar plot (which was generated using a different simulator) as problem 1 of assignment 3. In this trace, we label the lost data segment (the one with a cross) to be segment 14, and other data segments are labeled consecutively. The only main difference as compared with the plot in the assignment is that the ACK values here refer to the next expected segment. Recall that ack-every-segment strategy is used, and assume that the rwnd value is always equivalent to 30 data segments.

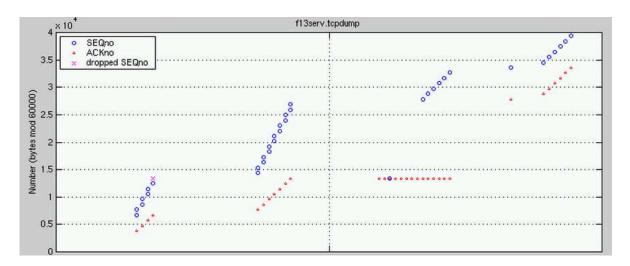


Fig. 5. TCP segment 14 is lost and is being fast retransmitted.

Answer the following questions concerning the trace in Figure 5 with succinct explanation. You do not need to know the ssthresh values to answer these questions.

- a) (2 marks) What is the value of cwnd by the time of transmitting data segment 14?
- b) (2 marks) Why is there an idle period between sending segment 14 and segment 15?

- c) (2 marks) What is the value of cwnd by the time of transmitting data segment 28?
- d) (2 marks) Why is there an idle period between sending segment 28 and segment 29?
- e) (2 marks) What is the value of cwnd by the time of transmitting data segment 29? Explain your answer *without* using ssthresh.
- f) (2 marks) Why is there an idle period between sending segment 34 and segment 35?
- g) (2 marks) What is the value of cwnd by the time of transmitting data segment 35? Explain your answer *without* using ssthresh.
- h) (2 marks) Assume that just after sending data segment 28, $snd_una = n$, and the MSS is given by m bytes. What are the values of snd_max and snd_nxt in terms of m and n?
- i) (2 marks) By the time of retransmitting segment 14, what are the values of snd_max and snd_nxt in terms of m and n?
- j) (4 marks) If the rwnd value is changed to 15 segments, what would be different in the plot?