



University of
Pittsburgh

Algorithms and Data Structures 2

CS 1501

Spring 2022

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(Slides are adapted from Dr. Ramirez's and Dr. Farnan's CS1501 slides.)

Announcements

- Upcoming deadlines:
 - Assignment 2 late due date on 3/30
 - Lab 9 due on 4/1
 - Homework 10 due on 4/4
 - Assignment 3 and 4 due on 4/18
 - Used to be one assignment

Previous lecture ...

- Dynamic Connectivity Problem
 - Union/Find data structure

CourseMIRROR Reflections (most confusing)

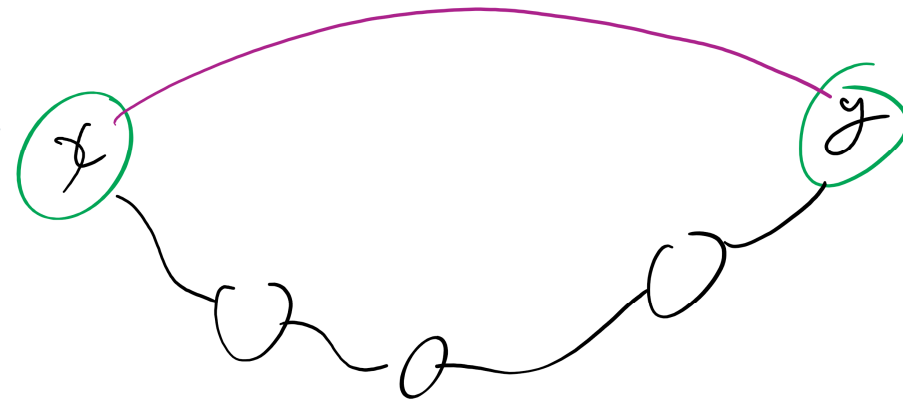
- I still found the heap sort confusing. Going through another example would help
- at the end of making the heap array, why do we switch the first item and the item at the boundary?
- I am still confused why we have to sort the Heap in descending order
- Personally, I feel like the class has gotten very fast lately, it's become harder to keep up with the new material
- The difference between storing UF as an array or as a tree
- using Union/find with Kruskals
- The union part was a bit confusing.
- I was confused on the runtime breakdowns for different prims algorithms
- The most confusing part of class today was applying a priority queue for Lazy Prim's and Eager Prim's algorithms.
- The middle part of lecture. The PQ and lazy/eager implementations were gone over very quickly. Going through an example would help

CourseMIRROR Reflections (most interesting)

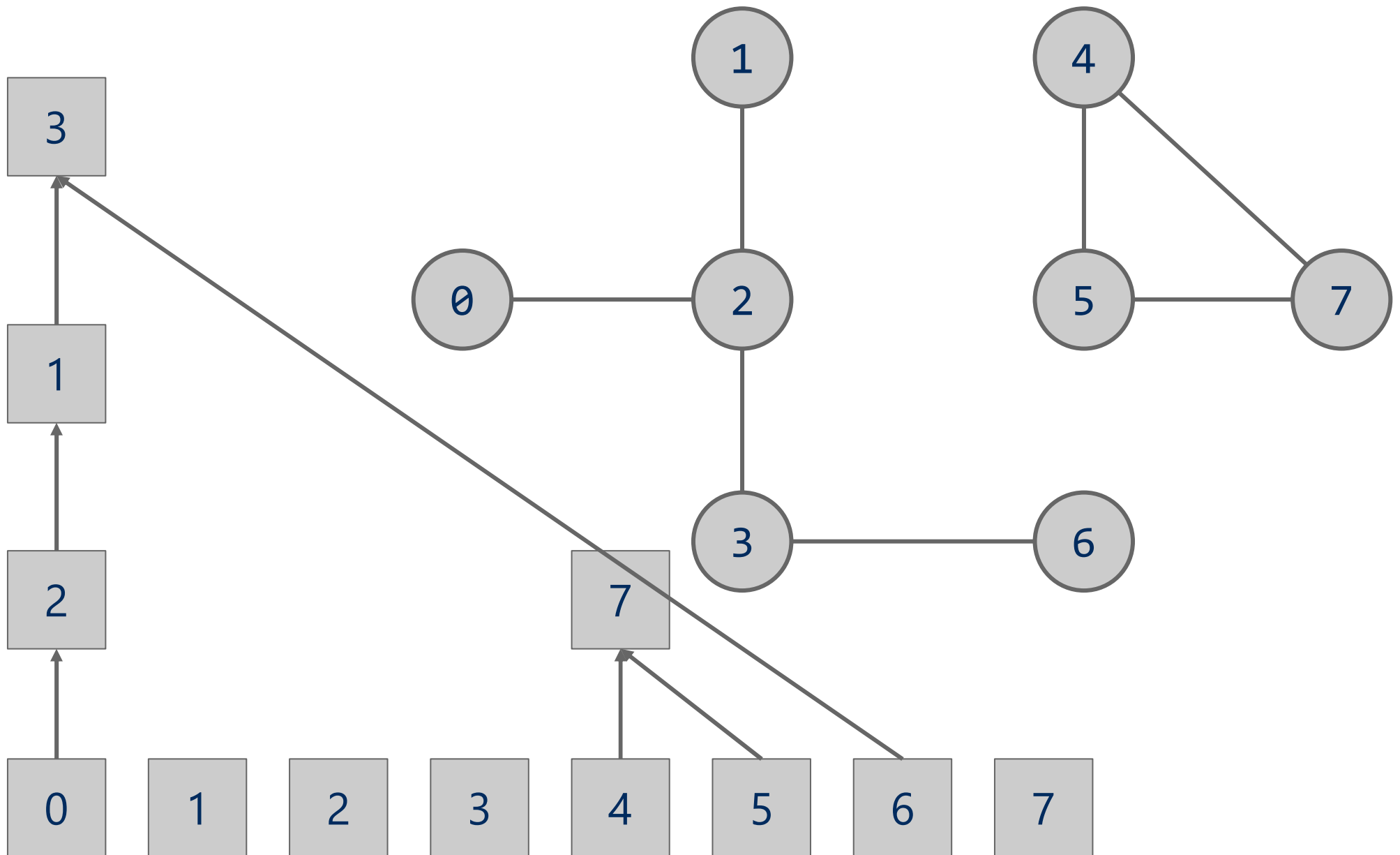
- Applicability of more complicated algorithms to help improve prim's and kruskal's
- I found it interesting how a priority queue could improve Prim's Algorithm
- Eager prim for improved runtime
- The runtime efficiency.
- Dynamic connectivity problem
- The connectivity problem and how a graph is a good approach to solving
- Connections of groups and the array of IDs
- The most interesting part of today's class was Union/Find ADT.
- The union examples were fun to walk through
- The new problem of the day, iterative graph building via unions/merges, seems interesting
- code for union find API

Problem of the Day: Dynamic connectivity problem

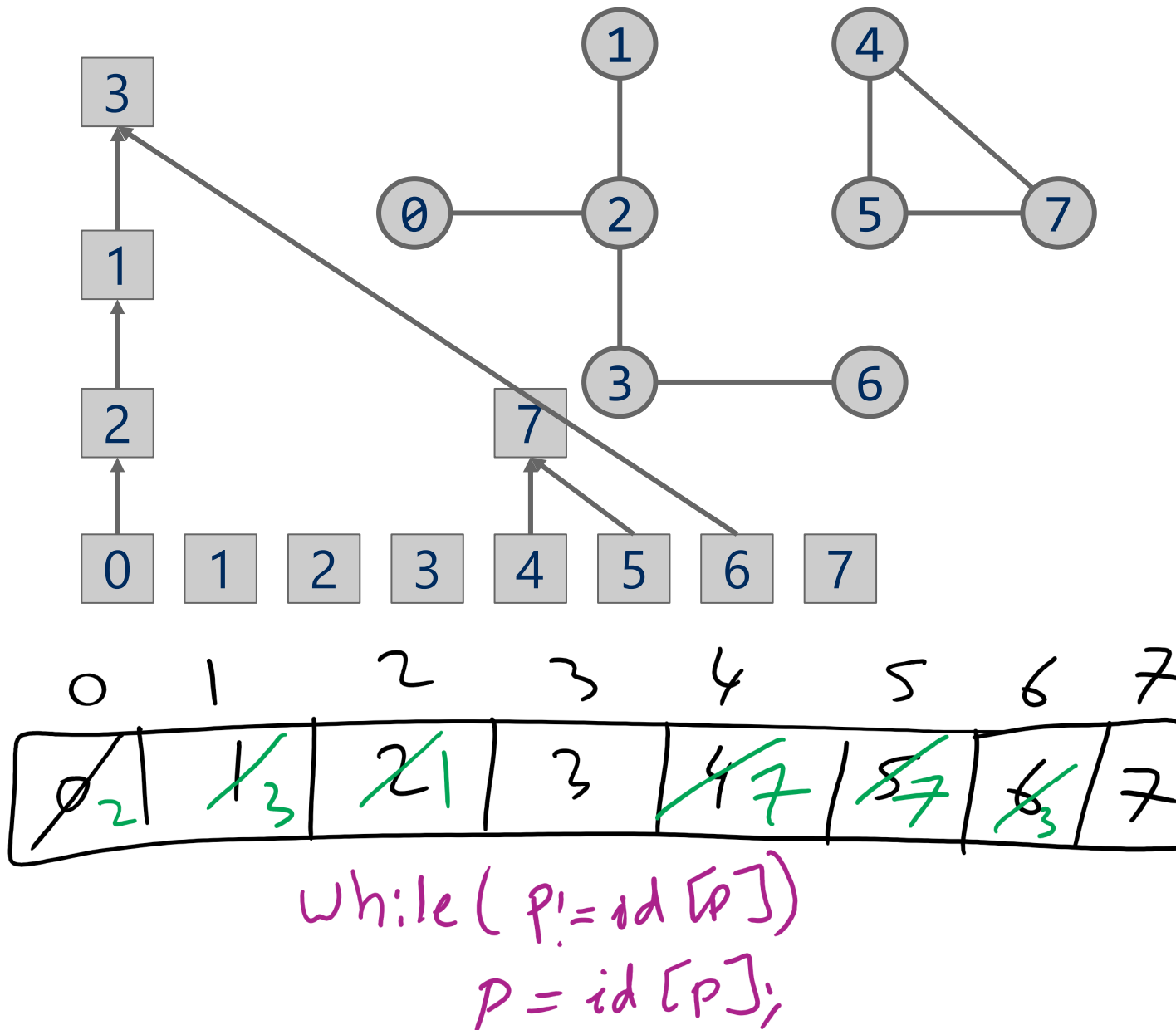
- Input:
 - A set of items initially in separate groups and
 - a sequence of merge/union operations, each operation merging two items
- Output:
 - At any point of time, we can be asked if two items are in the same group
 - Initially, the answer will be NO for any two items because they start in separate groups
- For a given graph G , can we determine whether or
- Can also be viewed as checking subset membership
- Important for many practical applications



Tree example



Forest of Trees Implementation



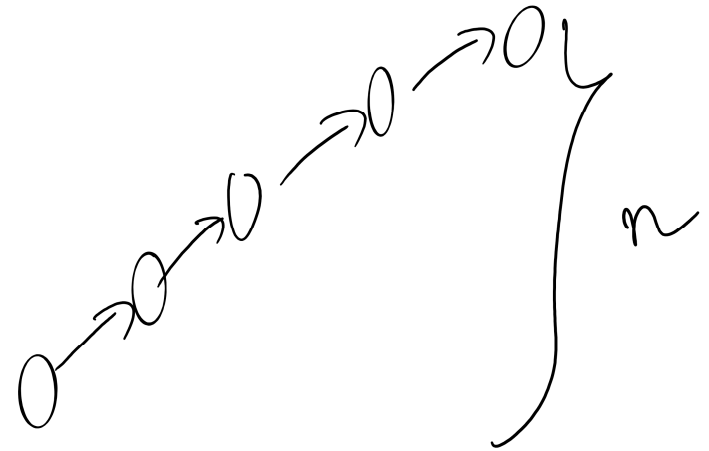
Implementation using the same id array

```
public int find(int p) {  
    while (p != id[p]) p = id[p];  
    return p;  
}
```

```
public void union(int p, int q) {  
    int i = find(p);  
    int j = find(q);  
    if (i == j) return;  
    id[i] = j;  
    count--;  
}
```

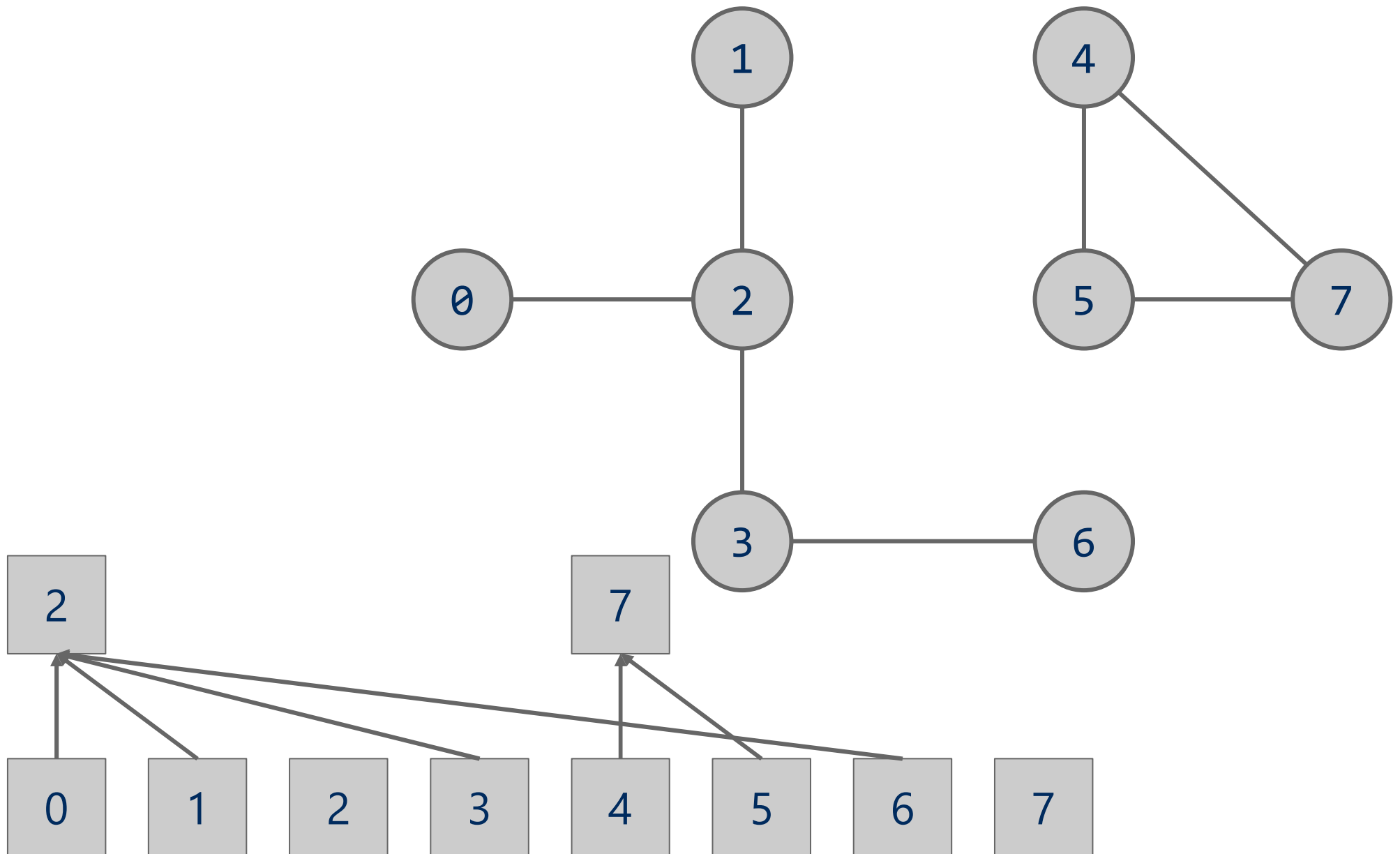
Forest of trees implementation analysis

- Runtime?
 - find():
 - Bound by the height of the tree
 - union():
 - Bound by the height of the tree



- What is the max height of the tree?
 - Can we modify our approach to cap its max height?

Weighted tree example

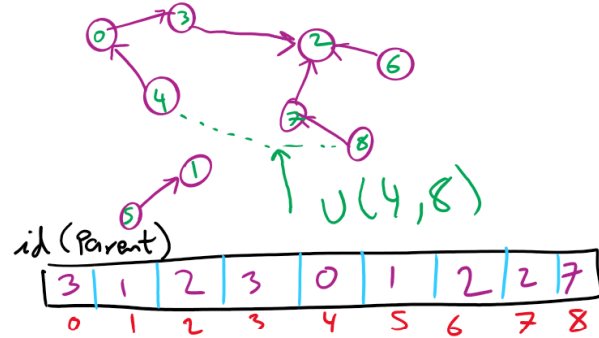


Weighted trees

```
public UF(int n) {  
    count = n;  
    id = new int[n];  
    sz = new int[n];  
    for (int i = 0; i < n; i++) { id[i] = i; sz[i] = 1; }  
}
```

```
public void union(int p, int q) {  
    int i = find(p), j = find(q);  
    if (i == j) return;  
    if (sz[i] < sz[j]) { id[i] = j; sz[j] += sz[i]; }  
    else                { id[j] = i; sz[i] += sz[j]; }  
    count--;  
}
```

Weighted Trees Union-Find Example



find(p) $\left[\begin{array}{l} \text{while}(p \neq \text{id}[p]) \\ \quad p = \text{id}[p] \\ \text{return } p \end{array} \right.$

sz

| | | | | | | | | |
|---|---|---|---|---|---|---|---|---|
| 2 | 2 | 4 | 3 | 1 | 1 | 1 | 2 | 1 |
| 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |

union(i, j)

$r_i = \text{find}(i)$
 $r_j = \text{find}(j)$
 if ($\text{sz}[r_i] > \text{sz}[r_j]$)
 $\text{id}[r_j] = r_i$
 $\text{sz}[r_i] += \text{sz}[r_j]$
 else
 $\text{id}[r_i] = r_j$
 $\text{sz}[r_j] += \text{sz}[r_i]$

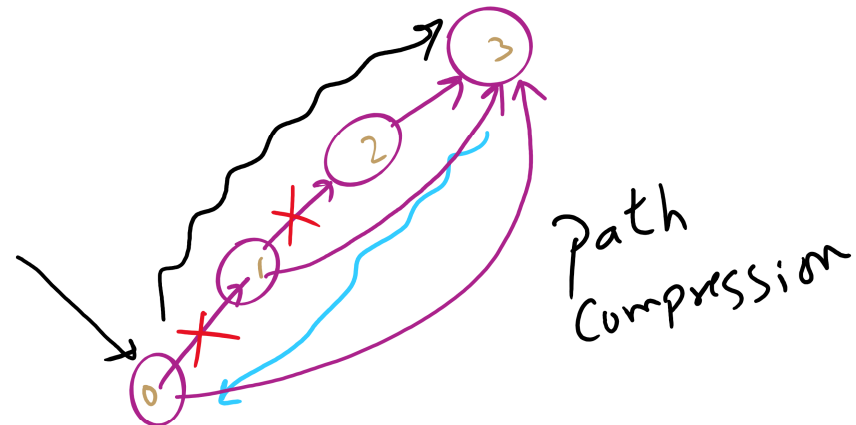
Count--

Weighted tree approach analysis

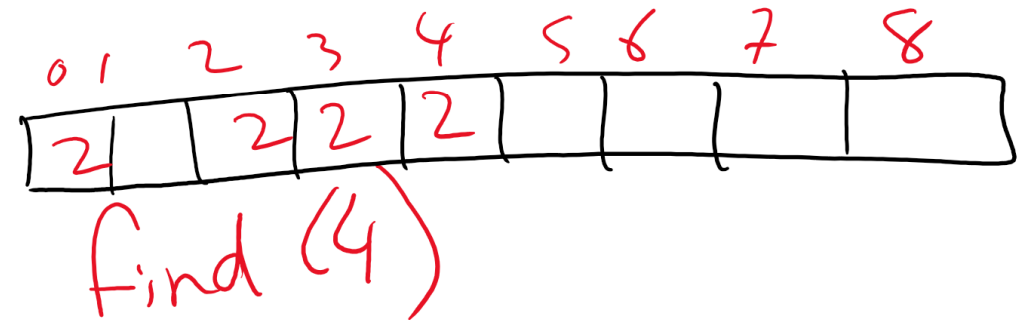
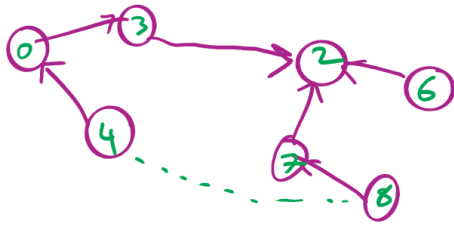
- Runtime?
 - find()?
 - union()?

find
union
 $\Theta(\log n)$

- Can we do any better?



Path Compression Example



Kruskal's algorithm, once again

- What is the runtime of Kruskal's algorithm??

e iterations

removal in $\log e \Leftarrow \text{bottleneck}$
2 find $\downarrow \log v$
1 union $\downarrow \log v$

$$e * (\log e + 2/\log v)$$

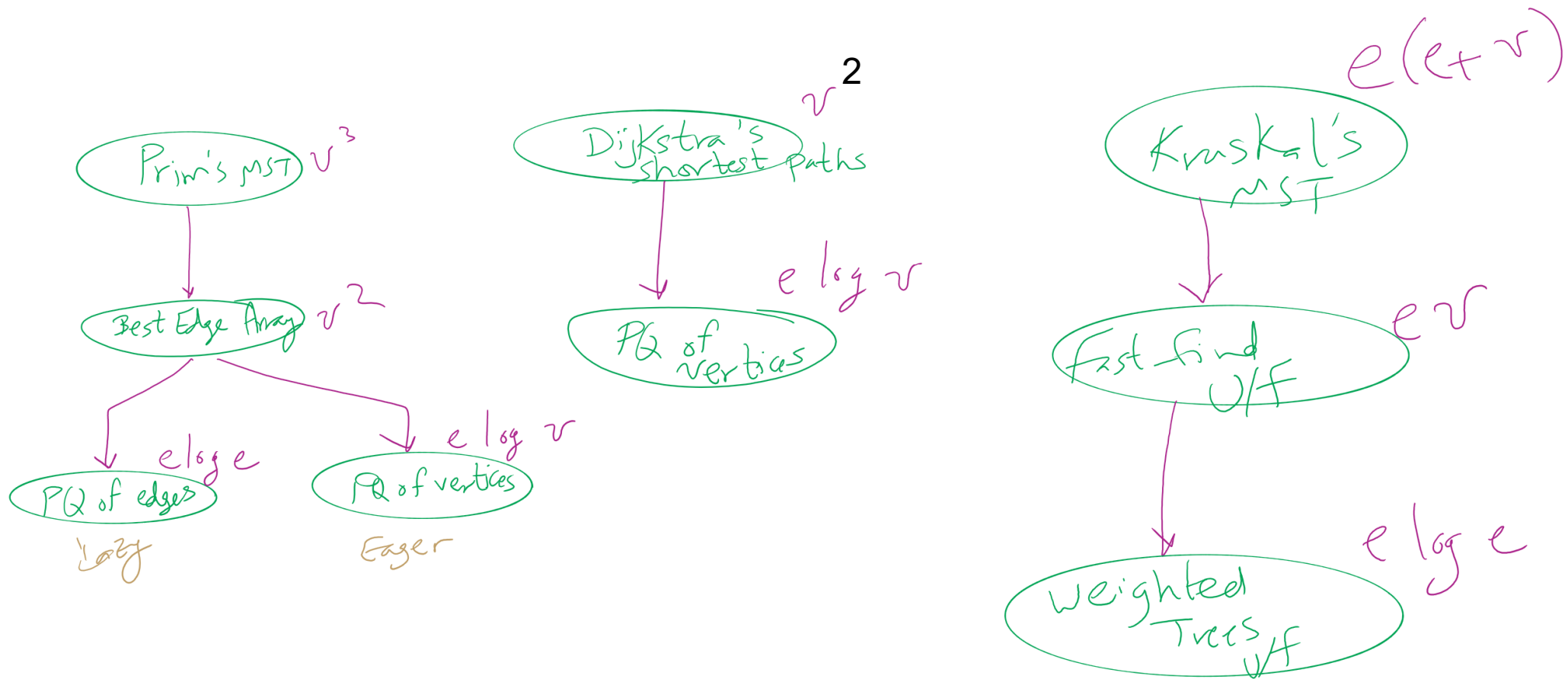
$$\Theta(e \log e) = \Theta(e \log v)$$

$$e = v^2$$
$$\log e = 2 \log v = \Theta(\log v)$$

Union-Find Implementations

| | Find | Union |
|------------------|------------------|------------------|
| Fast - Find | $\Theta(1)$ | $\Theta(n)$ |
| Forest of Trees | $\Theta(n)$ | $\Theta(n)$ |
| Weighted Trees | $\Theta(\log n)$ | $\Theta(\log n)$ |
| Path Compression | $\sim \Theta(1)$ | $\sim \Theta(1)$ |

Algorithm Optimizations



Problem of the Day 1: Weighted Shortest Path

- Input:
 - A road network
 - Road segments and intersections
 - Road segments are labeled by travel time
 - From length and maximum speed
 - How do we get max speed?
 - Starting address and destination address
- Output:
 - A shortest path from source to destination

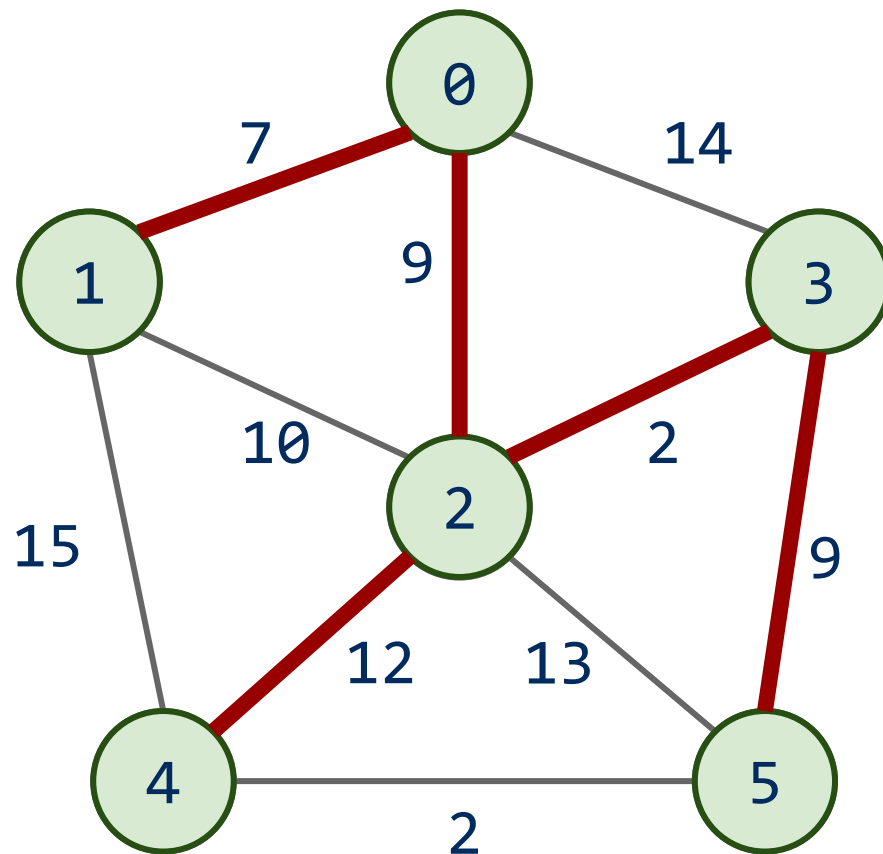
Weighted shortest path

- Dijkstra's algorithm:
 - Set a distance value of MAX_INT for all vertices but start
 - Set cur = start
 - While destination is not visited:
 - For each unvisited neighbor of cur:
 - Compute tentative distance from start to the unvisited neighbor through cur
 - Update any vertices for which a lesser distance is computed
 - Mark cur as visited
 - Let cur be the unvisited vertex with the smallest tentative distance from start

Tentative Distance

$\text{distance}[\text{cur}] + \text{edge weight}$
between
cur & neighbor

Dijkstra's example

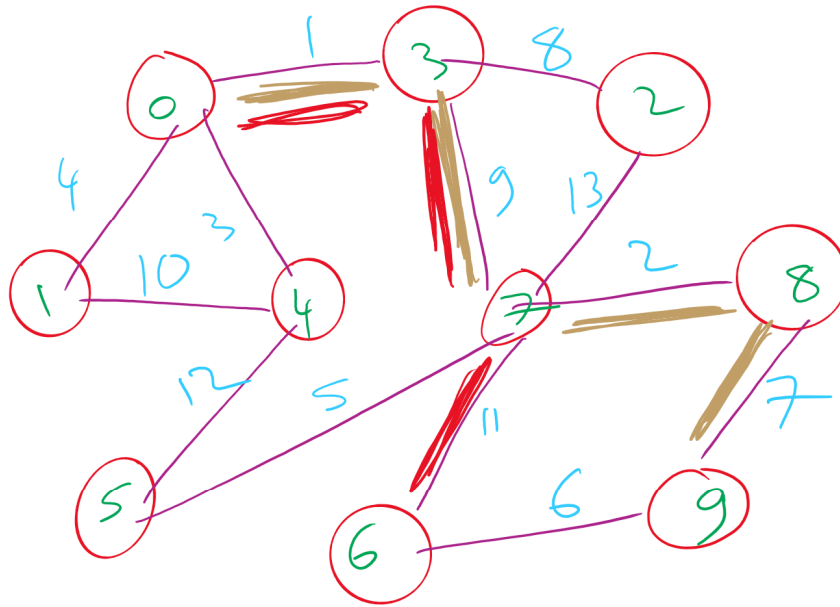


| | Distance | Parent |
|---|----------|--------|
| 0 | 0 | -- |
| 1 | 7 | 0 |
| 2 | 9 | 0 |
| 3 | 11 | 2 |
| 4 | 21 | 2 |
| 5 | 20 | 3 |

Analysis of Dijkstra's algorithm

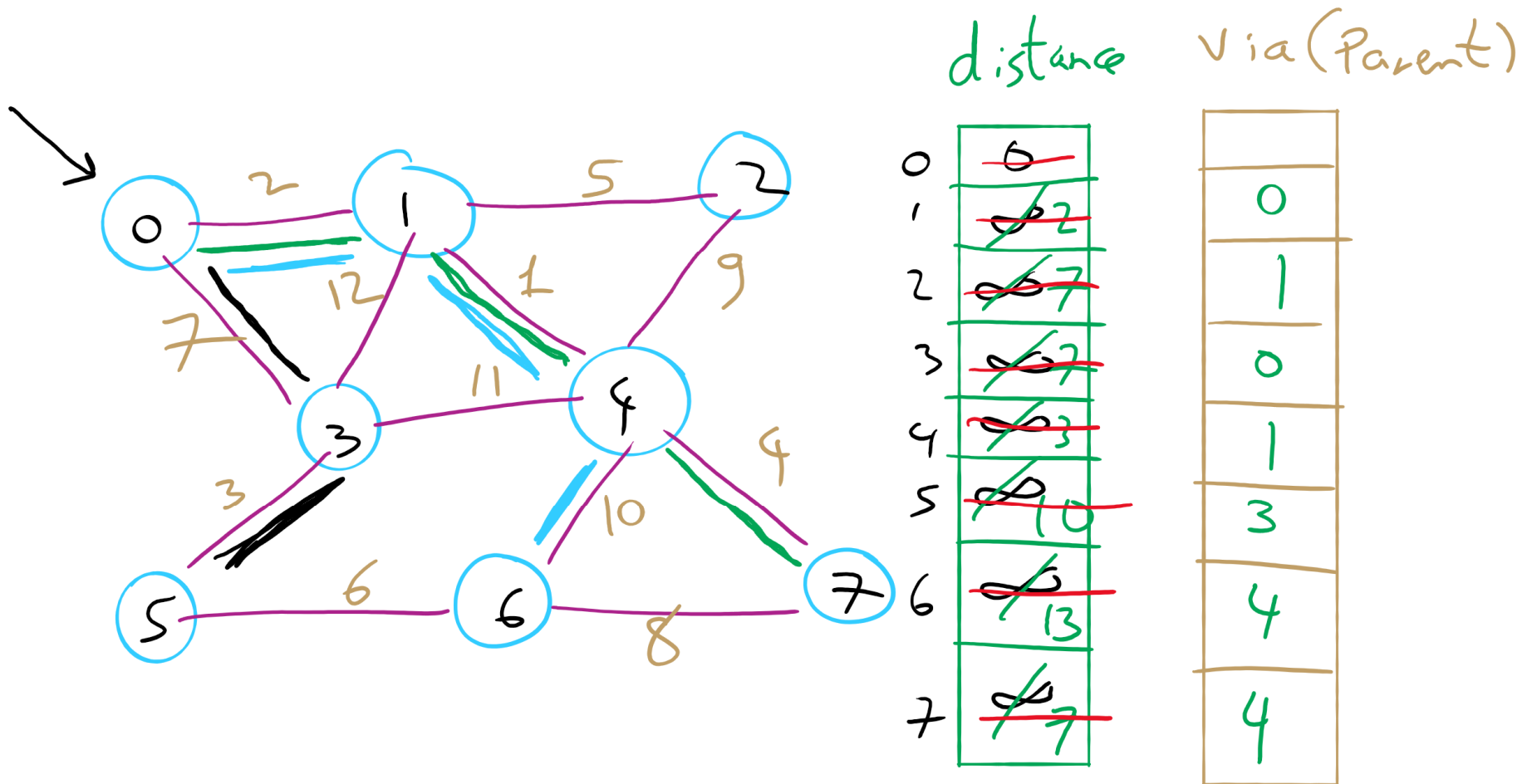
- How to implement?
 - Best path/parent array?
 - Runtime?
 - PQ?
 - Turns out to be very similar to Eager Prim's
 - Storing paths instead of edges
 - Runtime?

Dijkstra's Shortest Paths Example 1



| | distance | via (Parent) |
|---|---------------|--------------|
| 0 | 0 | |
| 1 | 4 | 0 |
| 2 | 9 | 3 |
| 3 | 1 | 0 |
| 4 | 3 | 0 |
| 5 | 15 | 4 |
| 6 | 21 | 7 |
| 7 | 10 | 3 |
| 8 | 12 | 7 |
| 9 | 19 | 8 |

Dijkstra's Shortest Paths Example 2



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8/29/2022