

TRIE

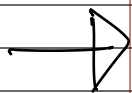
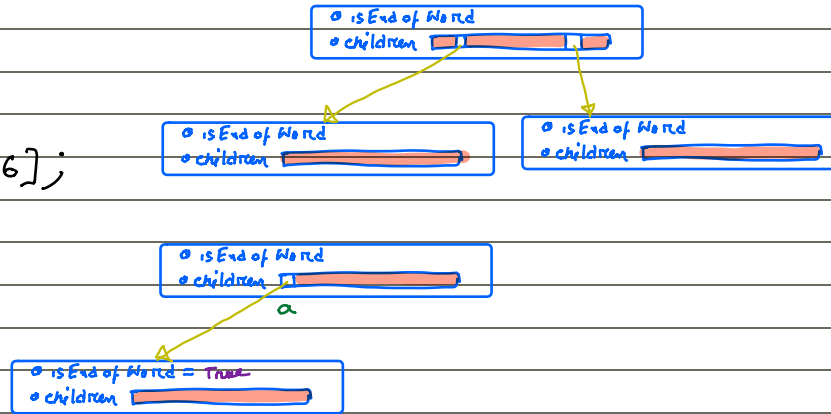


Prefix Tree

insert(word) search(word) prefixSearch(prefix)
 insert("apple") search("apple") prefixSearch("app")
 True True

```
struct TrieNode {
    bool isEndOfWord;
    TrieNode* children[26];
};
```

Ex:- insert("a")



Longest Common Suffix

3093. Longest Common Suffix Queries

Hard Topics Companies Hint

You are given two arrays of strings `wordsContainer` and `wordsQuery`.

For each `wordsQuery[i]`, you need to find a string from `wordsContainer` that has the **longest common suffix** with `wordsQuery[i]`. If there are two or more strings in `wordsContainer` that share the longest common suffix, find the string that is the **smallest** in length. If there are two or more such strings that have the **same** smallest length, find the one that occurred **earlier** in `wordsContainer`.

Return an array of integers `ans`, where `ans[i]` is the index of the string in `wordsContainer` that has the **longest common suffix** with `wordsQuery[i]`.

Input: `wordsContainer = ["abcd","bcd","xbcd"]`, `wordsQuery = ["cd","bcd","xyz"]`

Output: `[1,1,1]`

Explanation:

Let's look at each `wordsQuery[i]` separately:

- For `wordsQuery[0] = "cd"`, strings from `wordsContainer` that share the longest common suffix `"cd"` are at indices 0, 1, and 2. Among these, the answer is the string at index 1 because it has the shortest length of 3.
- For `wordsQuery[1] = "bcd"`, strings from `wordsContainer` that share the longest common suffix `"bcd"` are at indices 0, 1, and 2. Among these, the answer is the string at index 1 because it has the shortest length of 3.
- For `wordsQuery[2] = "xyz"`, there is no string from `wordsContainer` that shares a common suffix. Hence the longest common suffix is `""`, that is shared with strings at index 0, 1, and 2. Among these, the answer is the string at index 1 because it has the shortest length of 3.

Example 2:

Input: `wordsContainer = ["abcdefgh","poiuygh","ghghgh"]`, `wordsQuery = ["gh","acbfgh","acbfegh"]`

Output: `[2,0,2]`

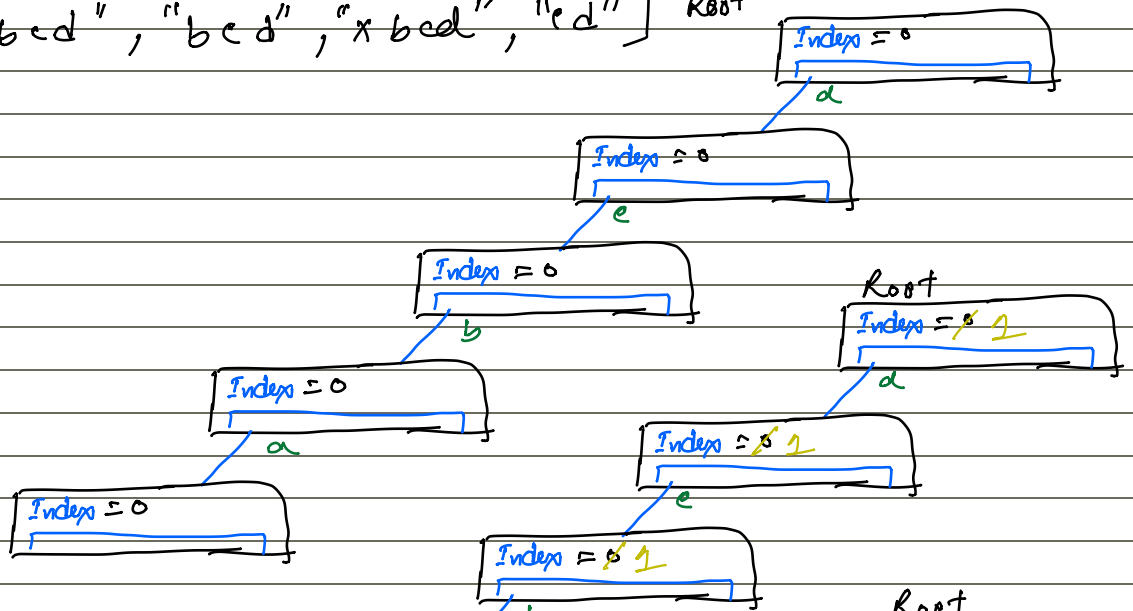
Explanation:

Let's look at each `wordsQuery[i]` separately:

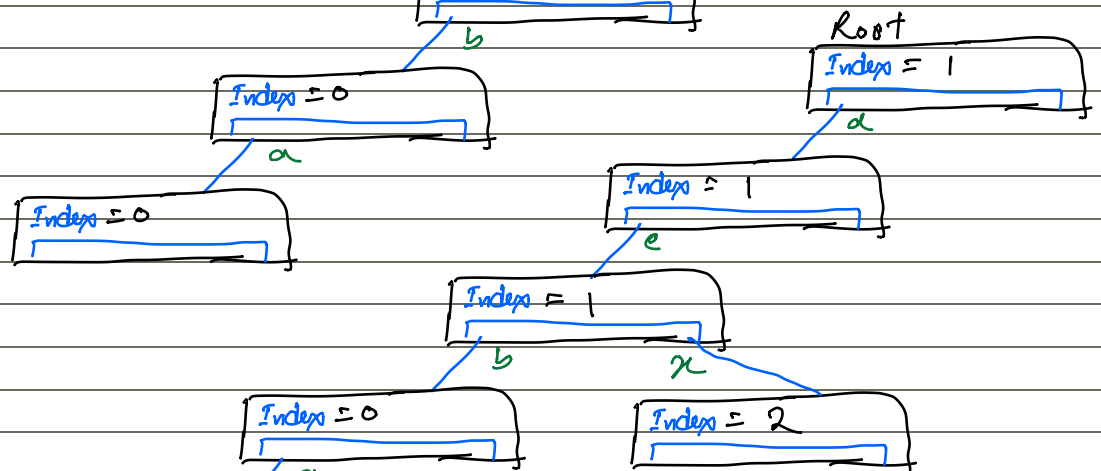
- For `wordsQuery[0] = "gh"`, strings from `wordsContainer` that share the longest common suffix `"gh"` are at indices 0, 1, and 2. Among these, the answer is the string at index 2 because it has the shortest length of 6.
- For `wordsQuery[1] = "acbfgh"`, only the string at index 0 shares the longest common suffix `"fgh"`. Hence it is the answer, even though the string at index 2 is shorter.
- For `wordsQuery[2] = "acbfegh"`, strings from `wordsContainer` that share the longest common suffix `"gh"` are at indices 0, 1, and 2. Among these, the answer is the string at index 2 because it has the shortest length of 6.

0 1 2 3
 ["abcd", "bcd", "xbed", "ed"] Root
 i

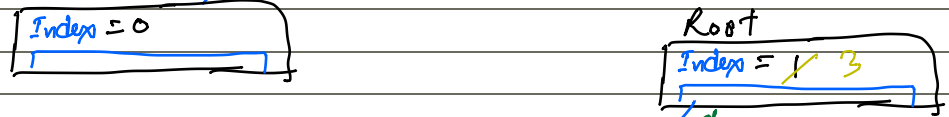
①



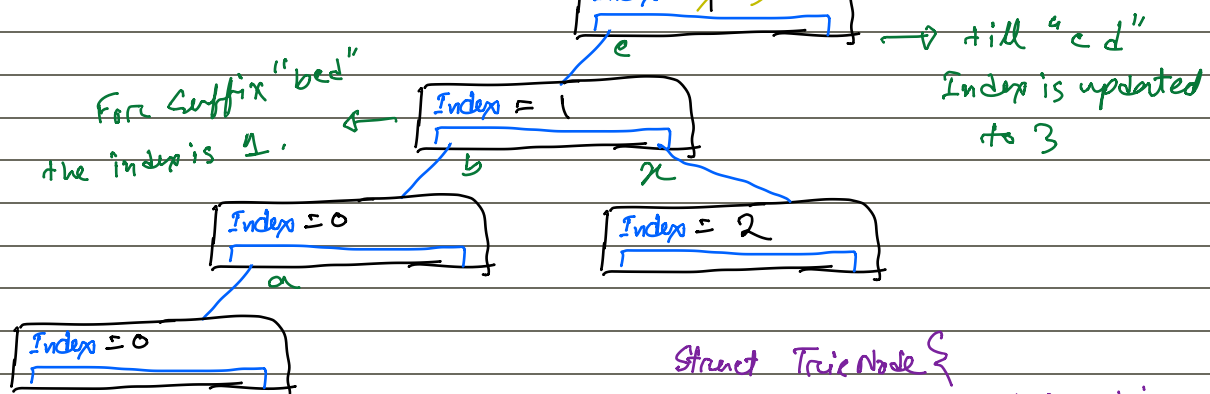
②



③



④



struct TrieNode {

int idx of smallest till now;

TrieNode* children[26];

* at each node update the idx of smallest till now for the prefix till that character.

✓ multiple things taken care of (just by inserting in this data structure)

i.e. smallest index; when abc, xbc

length = 3 & 3 but so node 'c' was there idx of smallest till now as 0 not 1

✓ if we have abc, xbc, bc the node 'c' has value 2

Maximum XOR of 2 numbers in an Array

arr [3, 5, 2, 7]

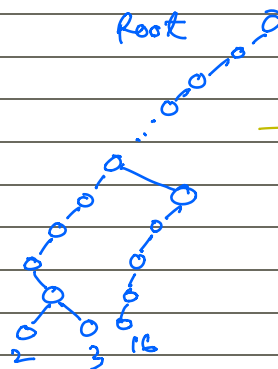
$$\begin{array}{r} 3 - 0 \ 1 \ 1 \\ 5 - 1 \ 0 \ 1 \\ 2 - 0 \ 1 \ 0 \\ \hline 7 - 1 \ 1 \ 1 \end{array}$$

	0	1	2
0	0	1	1
1	1	0	1

need to store whole 32 bit otherwise
problem of not counting might
occur.

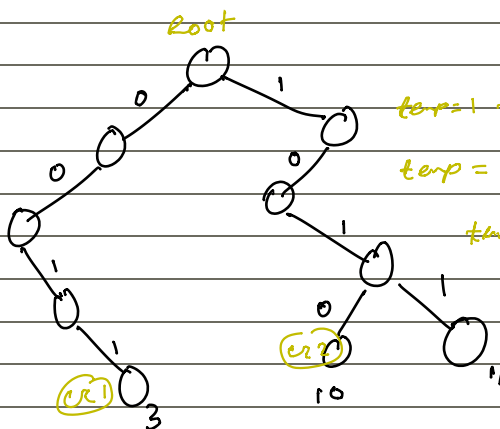
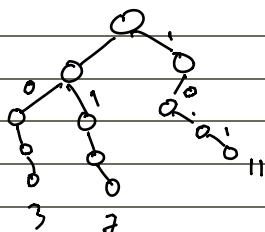
& also in post word order

$3 = 0000 \dots 0001$
 $2 = 0000 \dots 0010$
 $16 = 0000 \dots 10000$



} this part is not too much redundant

3, 11, 7, 10


$$i=3$$
$$3 = \begin{matrix} & 3 & 2 & 1 & 0 \\ & 0 & 0 & 1 & 1 \\ & \uparrow & & & \end{matrix}$$
$$t_{\text{er}} = 1 \rightarrow 1$$

temp = 10

$$k_{mp} = 100$$

$q_{exp} = 1001$

$$\text{temp} = 0$$

3 3 0 0 1 1

$$(0 \rightarrow) \quad 1 \quad 0 \quad 1 \quad 0$$
$$x02 \quad \overline{1001}$$

→ Replace Words:-

word container: ["cat", "bat"]

sentence: "a cattle was battling"

O/p: "A cat was bat"

✓ Approach: Using Hash (unordered Set).

✓ Approach: Using Trie

Why Trie is better;

map: {"cat", "bat"} words to be checked
⇒ a, cattle, was, battle.

cattle
i → i } check map.find("c")
 map.find("ca")
 map.find("cat") ✓

was
i → } check map.find("w") → # In Trie we can cut
 map.find("wa") the time to traverse
 map.find("was") x whole "was", could
conclude that any word
with 'w' isn't present
in the word container.

→ Palindrome Pairs :-

!! Difficult AF !!
o o o o