



# CONTENT

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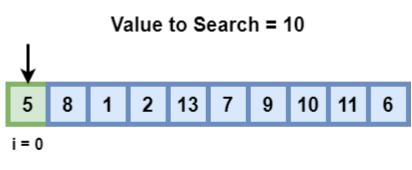
## LINEAR SEARCH

The traditional search method (Brute force)

Time complexity is O(N)

Does NOT need an array to be sorted

If we are given an array of integers A without any further information and have to decide if an element x is in A, we just have to search through it, element by element







### LINEAR SEARCH

```
// Linked list example
public boolean hasItem(T item) {
   MyNode<T> current = head;
   while (current != null) {
      if (current.data.equals(item))
          return true;
   }
   current = current.next;
}

return false;
}

// Array based example
public boolean hasItem(T item) {
   int current = 0;
   while (current < size) {
      if (array[current].equals(item))
          return true;
   }
   return false;
}</pre>
```





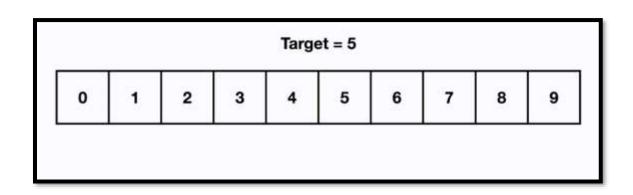
### BINARY SEARCH

The bisection search method

Time complexity is O(logN)

List **must be sorted** to give correct answer

We start by examining the middle element of the array



If it smaller than x, then x must be in the upper half of the array (if it is there at all); if is greater than x then it must be in the lower half

Now we continue by **restricting** our attention to either the upper or lower half, again finding the middle element and proceeding as before



### BINARY SEARCH

```
public boolean hasItem(T item) {
   return binarySearch(item, start: 0, end: size-1);
// Binary Search example (Recursive)
private boolean binarySearch(T item, int start, int end)
   if (start > end) return false;
   int mid = (start + end) / 2;
   int cmp = array[mid].compareTo(item);
   if (cmp == 0) { // if we found the item
       return true;
   } else if (cmp > 0) { // if middle element is more
       return binarySearch(item, start, end: mid - 1);
    } else {
       return binarySearch(item, start: mid + 1, end);
```

```
/ Binary Search example (Iterative)
private boolean binarySearch(T item) {
   int start = 0, end = size - 1;
    while (start <= end) {
        int mid = (start + end) / 2;
        int cmp = array[mid].compareTo(item);
        if (cmp == 0) {
            return true;
        } else if (cmp > 0) {
            end = mid - 1;
        } else {
            start = mid + 1;
   return false;
```





### BINARY SEARCH

Binary search gives better performance

Is it good to sort before we search to apply binary search instead of linear?

- only if it satisfies the inequality: Sort + O(logN) < O(N)
- Sort < O(N) O(log N)
- no such sorting (never true)

#### Multiple Searches Case (search k times)

- $Sort + kO(logN) < kO(N) \Rightarrow Sort < k[O(N) O(logN)]$
- for large k, Sort time becomes irrelevant





### HASHING

The HashTable is another good solution for searching

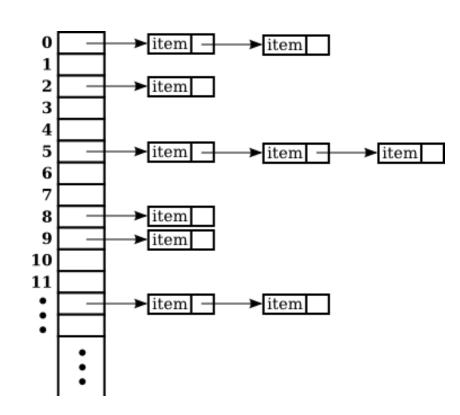
Searching Time complexity:

Average case: O(1)

Worst case: O(N)

Needs well-designed **hashing** method for making **less** collisions (see Lecture 5)

 Hash tables become quite inefficient when there are many collisions (Time complexity tends to O(N) as the number of collisions increases)







### HASHING

```
public boolean has(K key) {
   int index = hash(key);
    HashNode<K, V> temp = chainArray[index];
    while (temp != null) {
                                                This can be accepted as O(1) for
         if (temp.key.equals(key)) {
             return true;
                                                small number of collisions (it occurs
         temp = temp.next;
                                                in average case)
    return false;
                                                     The number of buckets (chains)
private int hash(K key) {
    return (key.hashCode() & 0x7fffffff) % M;
```



### FIND A PAIR WITH THE GIVEN SUM

Straightforward solution  $O(N^2)$ :

```
for (int <u>i</u> = 0; <u>i</u> < arr.length - 1; <u>i</u>++)
    for (int j = <u>i</u> + 1; j < arr.length; j++)
        if (arr[<u>i</u>] + arr[<u>j</u>] == sum)
            return true;
return false;
```

Using Binary search < O(NlogN)

(only for sorted list):

```
for (int i = 0; i < arr.length; i++)
   if (binarySearch(item: sum - arr[i], start: i + 1, end: arr.length - 1))
     return true;
return false;</pre>
```

Using HashTable O(N) (we use HashSet<K>, since we are not dealing with key-value pairs):

```
MyHashSet<Integer> previousItems = new MyHashSet<>();
for (int i = 0; i < arr.length; i++) {
   if (previousItems.has(key: sum - arr[i])) // O(1)
      return true;
   previousItems.put(arr[i]); // O(1)
}
return false;</pre>
```



## LITERATURE

Algorithms, 4th Edition, by Robert Sedgewick and Kevin Wayne, Addison-Wesley

Chapter 3

Grokking Algorithms, by Aditya Y. Bhargava, Manning

Chapter 5



