

## Zces 0.53.2

This document is in the Stable state. Assume anything could still change, but limited change should be expected. For more information see: <https://riscv.org/spec-state>

Zces is the set of sequenced or more complex instructions for code-size reduction.

This extension reuses encodings from the D-extension. Therefore it is *incompatible* with D. It is fully compatible with F and also with Zdinx.

jt and jalt require [JVT CSR](#), [table jump base vector](#) and [control register](#).

RV32	RV64	Mnemonic	Instruction
✓	✓	c.push {reg_list_1_4}, {areg_list_1_3}, -stack_adj	c.pusha: Create stack frame: push registers, move A to S registers, allocate additional stack space.
✓	✓	c.push {reg_list_5_13}, {areg_list_4}, -stack_adj	c.pusha.l: Create stack frame: push registers, move A to S registers, allocate additional stack space.
✓	✓	c.push {reg_list_1_4}, -stack_adj	c.push: Create stack frame: push registers, allocate additional stack space.
✓	✓	c.push {reg_list_5_13}, -stack_adj	c.push.l: Create stack frame: push registers, allocate additional stack space.
✓	✓	c.pop {reg_list_1_4}, stack_adj	c.pop: Destroy stack frame: pop registers, deallocate stack frame.
✓	✓	c.popret {reg_list_1_4}, stack_adj	c.popret: Destroy stack frame: pop registers, deallocate stack frame, return.
✓	✓	c.popret {reg_list_5_13}, stack_adj	c.popret.l: Destroy stack frame: pop registers, deallocate stack frame, return.
✓	✓	c.popretz {reg_list_1_4}, stack_adj	c.popretz: Destroy stack frame: pop registers, deallocate stack frame, return zero.
✓	✓	c.popretz {reg_list_5_13}, stack_adj	c.popretz.l: Destroy stack frame: pop registers, deallocate stack frame, return zero.
✓	✓	c.jt #index	c.jt: jump via table without link, 16-bit encoding
✓	✓	c.jalt #index	c.jalt: jump via table and link to ra, 16-bit encoding
✓	✓	c.mva01s sreg1, sreg2	c.mva01s: move two s0-s7 registers into a0-a1, 16-bit encoding

## c.pusha

### Synopsis

Create stack frame: store ra and 1 to 3 saved registers to the stack frame, move arguments into saved registers, optionally allocate additional stack space.

### Mnemonic

c.push {*reg\_list\_2\_4*}, {*areg\_list\_1\_3*}, -*stack\_adj*

### Encoding (RV32, RV64)

15					10	9	8	7	6		4	3	2	1	0
1	0	1	0	1	0	0	spimm[5:4]		0	1	1	rlist2		1	0
FUNCT6							reg_list_2_4							OP=C2	

#### NOTE

*rlist2=0* is reserved

### Syntax

The variables used in the syntax are defined below.

<reg\_list\_2\_4> ::= <ra> ", " [<s0> | <s0-sN>] (where N is 1,2)

```
if (<reg_list_2_4>=="ra, s0")      <xreg_list_2_4>="x1, x8"
if (<reg_list_2_4>=="ra, s0-s1")   <xreg_list_2_4>="x1, x8-x9"
if (<reg_list_2_4>=="ra, s0-s2")   <xreg_list_2_4>="x1, x8-x9, x18"
```

```
if (<reg_list_2_4>=="ra, s0")      <areg_list_1_3>="a0"
if (<reg_list_2_4>=="ra, s0-s1")   <areg_list_1_3>="a0-a1"
if (<reg_list_2_4>=="ra, s0-s2")   <areg_list_1_3>="a0-a2"
```

```
if (<reg_list_2_4>=="ra, s0")      <xareg_list_1_3>="x10"
if (<reg_list_2_4>=="ra, s0-s1")   <xareg_list_1_3>="x10-x11"
if (<reg_list_2_4>=="ra, s0-s2")   <xareg_list_1_3>="x10-x12"
```

RV32:

<stack\_adj>=[16|32|48|64]

RV64:

```
if (<reg_list_2_4>=="ra, s0")      <stack_adj>=[16|32|48|64|96]
if (<reg_list_2_4>=="ra, s0-s1")   <stack_adj>=[32|48|64|96|112]
if (<reg_list_2_4>=="ra, s0-s2")   <stack_adj>=[32|48|64|96|112]
```

---

**Description**

This instruction pushes (stores) the registers in *reg\_list\_2\_4* to stack memory, moves *areg\_list\_1\_3* into correspondingly numbered *s* registers, and then adjusts the stack pointer by *-stack\_adj*.

**NOTE**

All ABI register mapping are for the UABI. An EABI version is planned once the EABI is frozen.

For further information see [PUSH/POP Register Instructions](#).

Field decoding

The mapping from the *rlist2* field in the encoding are as shown below.

Table 1. *rlist2* decoding

rlist2	reg_list_2_4	stack_adj_base	
		RV32	RV64
0	reserved encoding		
1	ra, s0	16	16
2	ra, s0-s1	16	32
3	ra, s0-s2	16	32

*stack\_adj\_base* is the minimum number of bytes, in multiples of 16-byte blocks, required to store the registers in *reg\_list\_2\_4*.

*spimm* is the number of additional 16-byte blocks allocated for the stack frame.

The total stack adjustment is calculated as shown.

```
stack_adj = stack_adj_base+spimm[5:4]*16
```

Prerequisites

The C-extension must also be configured.

32-bit equivalent

No direct equivalent encoding exists

## Operation

The first section of pseudo-code may be executed multiple times before the instruction successfully completes.

```
//This is not SAIL, it's pseudo-code. The SAIL hasn't been written yet.

if (misa.MXL==1) bytes=4; else bytes=8;

addr=sp-bytes;
for(i=31;i>0;i--) {
    //if register i is in xreg_list
    if (xreg_list[i]) {
        switch(bytes) {
            4:  asm("sw x[i], 0(addr)");
            8:  asm("sd x[i], 0(addr)");
        }
        addr-=bytes;
    }
}
```

The final section of pseudo-code executes atomically, and only executes if the section above completes without any exceptions or interrupts.

```
//This is not SAIL, it's pseudo-code. The SAIL hasn't been written yet.

if (areg_list[a0]) asm("mv s0, a0");
if (areg_list[a1]) asm("mv s1, a1");
if (areg_list[a2]) asm("mv s2, a2");

sp+=stack_adj;
```

RV32 Assembly examples

```
c.push {ra, s0-s2}, {a0-a2}, -64
```

Encoding: *rlist2*=3, *spimm*=3

Equivalent sequence:

```
sw s2, -4(sp);
sw s1, -8(sp);
sw s0, -12(sp);
sw ra, -16(sp);
mv s0, a0
mv s1, a1
mv s2, a2
addi sp, sp, -64;
```

```
c.push {ra, s0-s1}, {a0-a1}, -32
```

Encoding: *rlist2*=2, *spimm*=1

Equivalent sequence:

```
sw s1, -4(sp);
sw s0, -8(sp);
sw ra, -12(sp);
mv s0, a0
mv s1, a1
addi sp, sp, -32;
```

Included in

Extension	Minimum version	Lifecycle state
Zces ( <a href="#">Zces 0.53.2</a> )	0.53.2	Stable

## c.pusha.l

### Synopsis

Create stack frame: store ra and 4 to 12 saved registers, move argument registers into saved registers, optionally allocate additional stack space.

### Mnemonic

c.push {*reg\_list\_5\_13*}, {*areg\_list\_4*}, -*stack\_adj*

### Encoding (RV32, RV64)

15					10	9		7	6	5	4		2	1	0
1	0	1	0	1	0	spimm[6:4]			1	0	rlist3			1	0
FUNCT6						reg_list_5_13						OP=C2			

### Syntax

```
c.push {<reg_list_5_13> | <xreg_list_5_13>}, {<areg_list_4> | <xarg_list_4>},
-<stack_adj>
```

The variables used in the syntax are defined below.

`<reg_list_5_13> ::= <ra> ", " <s0-sN> (where N is 3,4,5,6,7,8,9,11)`

if (`<reg_list_5_13>=="ra, s0-sN"`) `<xreg_list_5_13>="x1, x8-x9, x18-xM"`  
(where  $M=N+16$  and N is 3,4,5,6,7,8,9,11)

RV32:

if (`<reg_list_5_13>=="ra, s0-sN"`) `<stack_adj>=[32|48|..|528]` (where N is 3-6)  
 if (`<reg_list_5_13>=="ra, s0-sN"`) `<stack_adj>=[48|64|..|544]` (where N is 7-9)  
 if (`<reg_list_5_13>=="ra, s0-s11"`) `<stack_adj>=[64|80|..|560]`

RV64:

if (`<reg_list_5_13>=="ra, s0-sN"`) `<stack_adj>=[ 48| 64|..|544]` (where N is 3,4)  
 if (`<reg_list_5_13>=="ra, s0-sN"`) `<stack_adj>=[ 64| 80|..|560]` (where N is 5,6)  
 if (`<reg_list_5_13>=="ra, s0-sN"`) `<stack_adj>=[ 80| 96|..|576]` (where N is 7,8)  
 if (`<reg_list_5_13>=="ra, s0-s9"`) `<stack_adj>=[ 96|112|..|576]`  
 if (`<reg_list_5_13>=="ra, s0-s11"`) `<stack_adj>=[112|128|..|592]`

RV32/RV64:

`<areg_list_4>="a0-a3"`  
`<xareg_list_4>="x10-x13"`

Description

This instruction pushes (stores) the registers in *reg\_list\_5\_13* to stack memory, moves *areg\_list\_4* into correspondingly numbered *s* registers, and then adjusts the stack pointer by *-stack\_adj*.

NOTE

All ABI register mapping are for the UABI. An EABI version is planned once the EABI is frozen.

For further information see [PUSH/POP Register Instructions](#).

Field decoding

The mapping from the *rlist3* and *spimm* fields in the encoding are as shown below.

Table 2. *rlist3* decoding

rlist3	reg_list_5_13	stack_adj_base	
		RV32	RV64
0	ra, s0-s3	32	48
1	ra, s0-s4	32	48
2	ra, s0-s5	32	64
3	ra, s0-s6	32	64
4	ra, s0-s7	48	80
5	ra, s0-s8	48	80
6	ra, s0-s9	48	96
7	ra, s0-s11	64	112

*stack\_adj\_base* is the minimum number of bytes, in multiples of 16-byte blocks, required to store the registers in *reg\_list\_5\_13*.

*spimm* is the number of additional 16-byte blocks allocated for the stack frame.

The total stack adjustment is calculated as shown.

$$\text{stack\_adj} = \text{stack\_adj\_base} + \text{spimm}[6:4] * 16$$

Prerequisites

The C-extension must also be configured.

32-bit equivalent

No direct equivalent encoding exists.



## Operation

The first section of pseudo-code may be executed multiple times before the instruction successfully completes.

```
//This is not SAIL, it's pseudo-code. The SAIL hasn't been written yet.

if (misa.MXL==1) bytes=4; else bytes=8;

addr=sp-bytes;
for(i=31;i>0;i--) {
    //if register i is in xreg_list
    if (xreg_list[i]) {
        switch(bytes) {
            4:  asm("sw x[i], 0(addr)");
            8:  asm("sd x[i], 0(addr)");
        }
        addr-=bytes;
    }
}
```

The final section of pseudo-code executes atomically, and only executes if the section above completes without any exceptions or interrupts.

```
//This is not SAIL, it's pseudo-code. The SAIL hasn't been written yet.

asm("mv s0, a0");
asm("mv s1, a1");
asm("mv s2, a2");
asm("mv s3, a3");

sp+=stack_adj;
```

**RV32 Assembly examples**

```
c.push {ra, s0-s3}, {a0-a3}, -64
```

Encoding: *rlist3*=0, *spimm*=2

Equivalent sequence:

```
sw s3, -4(sp);  
sw s2, -8(sp);  
sw s1, -12(sp);  
sw s0, -16(sp);  
sw ra, -20(sp);  
mv s0, a0  
mv s1, a1  
mv s2, a2  
mv s3, a3  
addi sp, sp, -64;
```

```
c.push {ra, s0-s11}, {a0-a3}, -128
```

Encoding: *rlist3*=7, *spimm*=4

Equivalent sequence:

```
sw s11, -4(sp);  
sw s10, -8(sp);  
sw s9, -12(sp);  
sw s8, -16(sp);  
sw s7, -20(sp);  
sw s6, -24(sp);  
sw s5, -28(sp);  
sw s4, -32(sp);  
sw s3, -36(sp);  
sw s2, -40(sp);  
sw s1, -44(sp);  
sw s0, -48(sp);  
sw ra, -52(sp);  
mv s0, a0  
mv s1, a1  
mv s2, a2  
mv s3, a3  
addi sp, sp, -128;
```

**Included in**

Extension	Minimum version	Lifecycle state
Zces ( <a href="#">Zces 0.53.2</a> )	0.53.2	Stable

# c.push

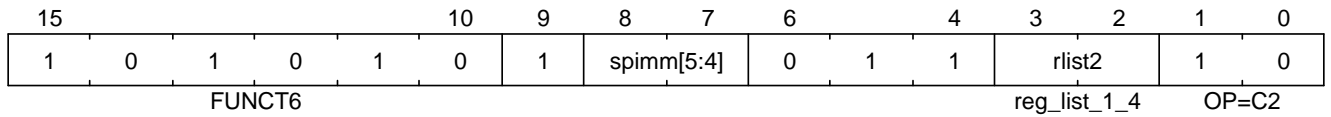
## Synopsis

Create stack frame: store ra and 0 to 3 saved registers to the stack frame, optionally allocate additional stack space.

## Mnemonic

```
c.push {reg_list_1_4}, -stack_adj
```

## Encoding (RV32, RV64)



## Syntax

```
c.push {<reg_list_1_4> | <xreg_list_1_4>}, -<stack_adj>
```

The variables used in the syntax are defined below.

```
<reg_list_1_4> ::= <ra> ["," <s0> | <s0-sN> ] (where N is 1,2)

if (<reg_list_1_4>=="ra")           <xreg_list_1_4>="x1"
if (<reg_list_1_4>=="ra, s0")       <xreg_list_1_4>="x1, x8"
if (<reg_list_1_4>=="ra, s0-s1")    <xreg_list_1_4>="x1, x8-x9"
if (<reg_list_1_4>=="ra, s0-s2")    <xreg_list_1_4>="x1, x8-x9, x18"

RV32:
<stack_adj>=[16|32|48|64]

RV64:
if (<reg_list_1_4>=="ra")           <stack_adj>=[16|32|48|64|96]
if (<reg_list_1_4>=="ra, s0")       <stack_adj>=[16|32|48|64|96]
if (<reg_list_1_4>=="ra, s0-s1")    <stack_adj>=[32|48|64|96|112]
if (<reg_list_1_4>=="ra, s0-s2")    <stack_adj>=[32|48|64|96|112]
```

---

**Description**

This instruction pushes (stores) the registers in *reg\_list\_1\_4* to stack memory, and then adjusts the stack pointer by *-stack\_adj*.

**NOTE**

All ABI register mapping are for the UABI. An EABI version is planned once the EABI is frozen.

For further information see [PUSH/POP Register Instructions](#).

Field decoding

The mapping from the *rlist2* field in the encoding are as shown below.

Table 3. *rlist2* decoding

rlist2	reg_list_1_4	stack_adj_base	
		RV32	RV64
0	ra	16	16
1	ra, s0	16	16
2	ra, s0-s1	16	32
3	ra, s0-s2	16	32

*stack\_adj\_base* is the minimum number of bytes, in multiples of 16-byte blocks, required to store the registers in *reg\_list\_1\_4*.

*spimm* is the number of additional 16-byte blocks allocated for the stack frame.

The total stack adjustment is calculated as shown.

```
stack_adj = stack_adj_base+spimm[5:4]*16
```

Prerequisites

The C-extension must also be configured.

32-bit equivalent

No direct equivalent encoding exists

## Operation

The first section of pseudo-code may be executed multiple times before the instruction successfully completes.

```
//This is not SAIL, it's pseudo-code. The SAIL hasn't been written yet.

if (misa.MXL==1) bytes=4; else bytes=8;

addr=sp-bytes;
for(i=31;i>0;i--) {
    //if register i is in xreg_list
    if (xreg_list[i]) {
        switch(bytes) {
            4:  asm("sw x[i], 0(addr)");
            8:  asm("sd x[i], 0(addr)");
        }
        addr-=bytes;
    }
}
```

The final section of pseudo-code executes atomically, and only executes if the section above completes without any exceptions or interrupts.

```
//This is not SAIL, it's pseudo-code. The SAIL hasn't been written yet.

sp+=stack_adj;
```

RV32 Assembly examples

```
c.push {ra, s0-s2}, -64
```

Encoding: *rlist2*=3, *spimm*=3

Equivalent sequence:

```
sw s2, -4(sp);
sw s1, -8(sp);
sw s0, -12(sp);
sw ra, -16(sp);
addi sp, sp, -64;
```

```
c.push {ra, s0-s1}, -32
```

Encoding: *rlist2*=2, *spimm*=1

Equivalent sequence:

```
sw s1, -4(sp);
sw s0, -8(sp);
sw ra, -12(sp);
addi sp, sp, -32;
```

Included in

Extension	Minimum version	Lifecycle state
Zces ( <a href="#">Zces 0.53.2</a> )	0.53.2	Stable



## c.push.l

### Synopsis

Create stack frame: store ra and 4 to 12 saved registers, optionally allocate additional stack space.

### Mnemonic

c.push {*reg\_list\_5\_13*}, -*stack\_adj*

### Encoding (RV32, RV64)

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	0	1	0	spimm[6:4]		1	1	rlist3			1	0	
FUNCT6						reg_list_5_13						OP=C2			

### Syntax

```
c.push {<reg_list_5_13> | <xreg_list_5_13>}, -<stack_adj>
```

The variables used in the syntax are defined below.

**<reg\_list\_5\_13> ::= <ra> ", " <s0-sN> (where N is 3,4,5,6,7,8,9,11)**

if (<reg\_list\_5\_13>=="ra, s0-sN") <xreg\_list\_5\_13>="x1, x8-x9, x18-xM"  
(where M=N+16 and N is 3,4,5,6,7,8,9,11)

**RV32:**

if (<reg\_list\_5\_13>=="ra, s0-sN") <stack\_adj>=[32|48|..|528] (where N is 3-6)  
if (<reg\_list\_5\_13>=="ra, s0-sN") <stack\_adj>=[48|64|..|544] (where N is 7-9)  
if (<reg\_list\_5\_13>=="ra, s0-s11") <stack\_adj>=[64|80|..|560]

**RV64:**

if (<reg\_list\_5\_13>=="ra, s0-sN") <stack\_adj>=[ 48| 64|..|544] (where N is 3,4)  
if (<reg\_list\_5\_13>=="ra, s0-sN") <stack\_adj>=[ 64| 80|..|560] (where N is 5,6)  
if (<reg\_list\_5\_13>=="ra, s0-sN") <stack\_adj>=[ 80| 96|..|576] (where N is 7,8)  
if (<reg\_list\_5\_13>=="ra, s0-s9") <stack\_adj>=[ 96|112|..|576]  
if (<reg\_list\_5\_13>=="ra, s0-s11") <stack\_adj>=[112|128|..|592]

Description

This instruction pushes (stores) the registers in *reg\_list\_5\_13* to stack memory, and then adjusts the stack pointer by *-stack\_adj*.

NOTE

All ABI register mapping are for the UABI. An EABI version is planned once the EABI is frozen.

For further information see [PUSH/POP Register Instructions](#).

Field decoding

The mapping from the *rlist3* and *spimm* fields in the encoding are as shown below.

Table 4. *rlist3* decoding

rlist3	reg_list_5_13	stack_adj_base	
		RV32	RV64
0	ra, s0-s3	32	48
1	ra, s0-s4	32	48
2	ra, s0-s5	32	64
3	ra, s0-s6	32	64
4	ra, s0-s7	48	80
5	ra, s0-s8	48	80
6	ra, s0-s9	48	96
7	ra, s0-s11	64	112

*stack\_adj\_base* is the minimum number of bytes, in multiples of 16-byte blocks, required to store the registers in *reg\_list\_5\_13*.

*spimm* is the number of additional 16-byte blocks allocated for the stack frame.

The total stack adjustment is calculated as shown.

$$\text{stack\_adj} = \text{stack\_adj\_base} + \text{spimm}[6:4] * 16$$

Prerequisites

The C-extension must also be configured.

32-bit equivalent

No direct equivalent encoding exists.

## Operation

The first section of pseudo-code may be executed multiple times before the instruction successfully completes.

```
//This is not SAIL, it's pseudo-code. The SAIL hasn't been written yet.

if (misa.MXL==1) bytes=4; else bytes=8;

addr=sp-bytes;
for(i=31;i>0;i--) {
    //if register i is in xreg_list
    if (xreg_list[i]) {
        switch(bytes) {
            4:  asm("sw x[i], 0(addr)");
            8:  asm("sd x[i], 0(addr)");
        }
        addr-=bytes;
    }
}
```

The final section of pseudo-code executes atomically, and only executes if the section above completes without any exceptions or interrupts.

```
//This is not SAIL, it's pseudo-code. The SAIL hasn't been written yet.

sp+=stack_adj;
```

RV32 Assembly examples

```
c.push {ra, s0-s3}, -64
```

Encoding: *rlist3*=0, *spimm*=2

Equivalent sequence:

```
sw s3, -4(sp);
sw s2, -8(sp);
sw s1, -12(sp);
sw s0, -16(sp);
sw ra, -20(sp);
addi sp, sp, -64;
```

```
c.push {ra, s0-s11}, -128
```

Encoding: *rlist3*=7, *spimm*=4

Equivalent sequence:

```
sw s11, -4(sp);
sw s10, -8(sp);
sw s9, -12(sp);
sw s8, -16(sp);
sw s7, -20(sp);
sw s6, -24(sp);
sw s5, -28(sp);
sw s4, -32(sp);
sw s3, -36(sp);
sw s2, -40(sp);
sw s1, -44(sp);
sw s0, -48(sp);
sw ra, -52(sp);
addi sp, sp, -128;
```

Included in

Extension	Minimum version	Lifecycle state
Zces ( <a href="#">Zces 0.53.2</a> )	0.53.2	Stable

## c.pop

### Synopsis

Destroy stack frame: load ra and 0 to 3 saved registers from the stack frame, deallocate the stack frame.

### Mnemonic

c.pop {*reg\_list\_1\_4*}, *stack\_adj*

### Encoding (RV32, RV64)

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	0	1	1	0	spimm[5:4]		0	0	1		rlist2	1	0
FUNCT6						reg_list_1_4						OP=C2			

### Syntax

```
c.pop {<reg_list_1_4> | <xreg_list_1_4>}, <stack_adj>
```

The variables used in the syntax are defined below.

<reg\_list\_1\_4> ::= <ra> [", " <s0> | <s0-sN> ] (where N is 1,2)

```
if (<reg_list_1_4>=="ra")           <xreg_list_1_4>="x1"
if (<reg_list_1_4>=="ra, s0")       <xreg_list_1_4>="x1, x8"
if (<reg_list_1_4>=="ra, s0-s1")    <xreg_list_1_4>="x1, x8-x9"
if (<reg_list_1_4>=="ra, s0-s2")    <xreg_list_1_4>="x1, x8-x9, x18"
```

RV32:

<stack\_adj>=[16|32|48|64]

RV64:

```
if (<reg_list_1_4>=="ra")           <stack_adj>=[16|32|48|64|96]
if (<reg_list_1_4>=="ra, s0")       <stack_adj>=[16|32|48|64|96]
if (<reg_list_1_4>=="ra, s0-s1")    <stack_adj>=[32|48|64|96|112]
if (<reg_list_1_4>=="ra, s0-s2")    <stack_adj>=[32|48|64|96|112]
```

### Description

This instruction pop (loads) the registers in *reg\_list\_1\_4* from stack memory, and then adjusts the stack pointer by *stack\_adj*.

#### NOTE

All ABI register mapping are for the UABI. An EABI version is planned once the EABI is frozen.

For further information see [PUSH/POP Register Instructions](#).

Field decoding

The mapping from the *rlist2* field in the encoding are as shown below.

Table 5. *rlist2* decoding

rlist2	reg_list_1_4	stack_adj_base	
		RV32	RV64
0	ra	16	16
1	ra, s0	16	16
2	ra, s0-s1	16	32
3	ra, s0-s2	16	32

*stack\_adj\_base* is the minimum number of bytes, in multiples of 16-byte blocks, required to store the registers in *reg\_list\_1\_4*.

*spimm* is the number of additional 16-byte blocks allocated for the stack frame.

The total stack adjustment is calculated as shown.

```
stack_adj = stack_adj_base+spimm[5:4]*16
```

Prerequisites

The C-extension must also be configured.

32-bit equivalent

No direct equivalent encoding exists

## Operation

The first section of pseudo-code may be executed multiple times before the instruction successfully completes.

```
//This is not SAIL, it's pseudo-code. The SAIL hasn't been written yet.

if (misa.MXL==1) bytes=4; else bytes=8;

addr=sp+stack_adj-bytes;
for(i=31;i>0;i--) {
    //if register i is in xreg_list
    if (xreg_list[i]) {
        switch(bytes) {
            4:  asm("lw x[i], 0(addr)");
            8:  asm("ld x[i], 0(addr)");
        }
        addr-=bytes;
    }
}
```

The final section of pseudo-code executes atomically, and only executes if the section above completes without any exceptions or interrupts.

```
//This is not SAIL, it's pseudo-code. The SAIL hasn't been written yet.

sp+=stack_adj;
```

RV32 Assembly examples

```
c.pop    {ra, s0-s2}, 48
```

Encoding: *rlist2*=3, *spimm*=2

Equivalent sequence:

```
lw    s2, 44(sp);
lw    s1, 40(sp);
lw    s0, 36(sp);
lw    ra, 32(sp);
addi  sp, sp, 48;
```

```
c.pop    {ra}, 16
```

Encoding: *rlist2*=0, *spimm*=0

Equivalent sequence:

```
lw    ra, 12(sp);
addi  sp, sp, 16;
```

Included in

Extension	Minimum version	Lifecycle state
Zces ( <a href="#">Zces 0.53.2</a> )	0.53.2	Stable



## c.popret

### Synopsis

Destroy stack frame: load ra and 0 to 3 saved registers from the stack frame, deallocate the stack frame, return to ra.

### Mnemonic

c.popret {*reg\_list\_1\_4*}, *stack\_adj*

### Encoding (RV32, RV64)

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	1	0	1	1	0	spimm[5:4]		0	1	1	rlist2		1	0
FUNCT6						reg_list_1_4						OP=C2			

### Syntax

```
c.popret {<reg_list_1_4> | <xreg_list_1_4>}, <stack_adj>
```

The variables used in the syntax are defined below.

```
<reg_list_1_4> ::= <ra> [", " <s0> | <s0-sN> ] (where N is 1,2)
```

```
if (<reg_list_1_4>=="ra")           <xreg_list_1_4>="x1"
if (<reg_list_1_4>=="ra, s0")       <xreg_list_1_4>="x1, x8"
if (<reg_list_1_4>=="ra, s0-s1")    <xreg_list_1_4>="x1, x8-x9"
if (<reg_list_1_4>=="ra, s0-s2")    <xreg_list_1_4>="x1, x8-x9, x18"
```

RV32:

```
<stack_adj>=[16|32|48|64]
```

RV64:

```
if (<reg_list_1_4>=="ra")           <stack_adj>=[16|32|48|64|96]
if (<reg_list_1_4>=="ra, s0")       <stack_adj>=[16|32|48|64|96]
if (<reg_list_1_4>=="ra, s0-s1")    <stack_adj>=[32|48|64|96|112]
if (<reg_list_1_4>=="ra, s0-s2")    <stack_adj>=[32|48|64|96|112]
```

### Description

This instruction pop (loads) the registers in *reg\_list\_1\_4* from stack memory, adjusts the stack pointer by *stack\_adj* and then returns.

#### NOTE

All ABI register mapping are for the UABI. An EABI version is planned once the EABI is frozen.

For further information see [PUSH/POP Register Instructions](#).

Field decoding

The mapping from the *rlist2* field in the encoding are as shown below.

Table 6. *rlist2* decoding

rlist2	reg_list_1_4	stack_adj_base	
		RV32	RV64
0	ra	16	16
1	ra, s0	16	16
2	ra, s0-s1	16	32
3	ra, s0-s2	16	32

*stack\_adj\_base* is the minimum number of bytes, in multiples of 16-byte blocks, required to store the registers in *reg\_list\_1\_4*.

*spimm* is the number of additional 16-byte blocks allocated for the stack frame.

The total stack adjustment is calculated as shown.

```
stack_adj = stack_adj_base+spimm[5:4]*16
```

Prerequisites

The C-extension must also be configured.

32-bit equivalent

No direct equivalent encoding exists

## Operation

The first section of pseudo-code may be executed multiple times before the instruction successfully completes.

```
//This is not SAIL, it's pseudo-code. The SAIL hasn't been written yet.

if (misa.MXL==1) bytes=4; else bytes=8;

addr=sp+stack_adj-bytes;
for(i=31;i>0;i--) {
    //if register i is in xreg_list
    if (xreg_list[i]) {
        switch(bytes) {
            4:  asm("lw x[i], 0(addr)");
            8:  asm("ld x[i], 0(addr)");
        }
        addr-=bytes;
    }
}
```

The final section of pseudo-code executes atomically, and only executes if the section above completes without any exceptions or interrupts.

```
//This is not SAIL, it's pseudo-code. The SAIL hasn't been written yet.

sp+=stack_adj;
asm("ret");
```

RV32 Assembly examples

```
c.popret {ra, s0-s2}, 48
```

Encoding: *rlist2*=3, *spimm*=2

Equivalent sequence:

```
lw    s2, 44(sp);
lw    s1, 40(sp);
lw    s0, 36(sp);
lw    ra, 32(sp);
addi  sp, sp, 48;
ret;
```

```
c.popret {ra}, 16
```

Encoding: *rlist2*=0, *spimm*=0

Equivalent sequence:

```
lw    ra, 12(sp);
addi  sp, sp, 16;
ret;
```

Included in

Extension	Minimum version	Lifecycle state
Zces ( <a href="#">Zces 0.53.2</a> )	0.53.2	Stable

## c.popret.l

### Synopsis

Destroy stack frame: load ra and 4 to 12 saved registers from the stack frame, deallocate the stack frame, return to ra.

### Mnemonic

c.popret {*reg\_list\_5\_13*}, *stack\_adj*

### Encoding (RV32, RV64)

15	10	9	7	6	5	4	2	1	0
1	0	1	0	1	1	spimm[6:4]		1	0
FUNCT6						reg_list_5_13			OP=C2

### Syntax

```
c.popret {<reg_list_5_13> | <xreg_list_5_13>}, -<stack_adj>
```

The variables used in the syntax are defined below.

**<reg\_list\_5\_13> ::= <ra> ", " <s0-sN> (where N is 3,4,5,6,7,8,9,11)**

if (<reg\_list\_5\_13>=="ra, s0-sN") <xreg\_list\_5\_13>="x1, x8-x9, x18-xM"  
(where M=N+16 and N is 3,4,5,6,7,8,9,11)

**RV32:**

if (<reg\_list\_5\_13>=="ra, s0-sN") <stack\_adj>=[32|48|..|528] (where N is 3-6)  
if (<reg\_list\_5\_13>=="ra, s0-sN") <stack\_adj>=[48|64|..|544] (where N is 7-9)  
if (<reg\_list\_5\_13>=="ra, s0-s11") <stack\_adj>=[64|80|..|560]

**RV64:**

if (<reg\_list\_5\_13>=="ra, s0-sN") <stack\_adj>=[ 48| 64|..|544] (where N is 3,4)  
if (<reg\_list\_5\_13>=="ra, s0-sN") <stack\_adj>=[ 64| 80|..|560] (where N is 5,6)  
if (<reg\_list\_5\_13>=="ra, s0-sN") <stack\_adj>=[ 80| 96|..|576] (where N is 7,8)  
if (<reg\_list\_5\_13>=="ra, s0-s9") <stack\_adj>=[ 96|112|..|576]  
if (<reg\_list\_5\_13>=="ra, s0-s11") <stack\_adj>=[112|128|..|592]

Description

This instruction pop (loads) the registers in *reg\_list\_5\_13* from stack memory, adjusts the stack pointer by *stack\_adj* and then returns.

NOTE

All ABI register mapping are for the UABI. An EABI version is planned once the EABI is frozen.

For further information see [PUSH/POP Register Instructions](#).

Field decoding

The mapping from the *rlist3* and *spimm* fields in the encoding are as shown below.

Table 7. *rlist3* decoding

rlist3	reg_list_5_13	stack_adj_base	
		RV32	RV64
0	ra, s0-s3	32	48
1	ra, s0-s4	32	48
2	ra, s0-s5	32	64
3	ra, s0-s6	32	64
4	ra, s0-s7	48	80
5	ra, s0-s8	48	80
6	ra, s0-s9	48	96
7	ra, s0-s11	64	112

*stack\_adj\_base* is the minimum number of bytes, in multiples of 16-byte blocks, required to store the registers in *reg\_list\_5\_13*.

*spimm* is the number of additional 16-byte blocks allocated for the stack frame.

The total stack adjustment is calculated as shown.

$$\text{stack\_adj} = \text{stack\_adj\_base} + \text{spimm}[6:4] * 16$$

Prerequisites

The C-extension must also be configured.

32-bit equivalent

No direct equivalent encoding exists.

## Operation

The first section of pseudo-code may be executed multiple times before the instruction successfully completes.

```
//This is not SAIL, it's pseudo-code. The SAIL hasn't been written yet.

if (misa.MXL==1) bytes=4; else bytes=8;

addr=sp+stack_adj-bytes;
for(i=31;i>0;i--) {
    //if register i is in xreg_list
    if (xreg_list[i]) {
        switch(bytes) {
            4:  asm("lw x[i], 0(addr)");
            8:  asm("ld x[i], 0(addr)");
        }
        addr-=bytes;
    }
}
```

The final section of pseudo-code executes atomically, and only executes if the section above completes without any exceptions or interrupts.

```
//This is not SAIL, it's pseudo-code. The SAIL hasn't been written yet.

sp+=stack_adj;
asm("ret");
```

RV32 Assembly examples

```
c.popret {ra, s0-s3}, 48
```

Encoding: *rlist2*=0, *spimm*=1

Equivalent sequence:

```
lw    s3, 44(sp);
lw    s2, 40(sp);
lw    s1, 36(sp);
lw    s0, 32(sp);
lw    ra, 28(sp);
addi  sp, sp, 48;
ret;
```

```
c.popret {ra, s0-s4}, 64
```

Encoding: *rlist2*=1, *spimm*=2

Equivalent sequence:

```
lw    s4, 60(sp);
lw    s3, 56(sp);
lw    s2, 52(sp);
lw    s1, 48(sp);
lw    s0, 44(sp);
lw    ra, 40(sp);
addi  sp, sp, 64;
ret;
```

Included in

Extension	Minimum version	Lifecycle state
Zces ( <a href="#">Zces 0.53.2</a> )	0.53.2	Stable



## c.popretz

### Synopsis

Destroy stack frame: load ra and 0 to 3 saved registers from the stack frame, deallocate the stack frame, move zero into a0, return to ra.

### Mnemonic

c.popretz {*reg\_list\_1\_4*}, *stack\_adj*

### Encoding (RV32, RV64)

15					10	9	8	7	6		4	3	2	1	0		
1	0	1	0	1	1	1	spimm[5:4]		0	1	1	rlist2		1	0		
FUNCT6								reg_list 1 4								OP=C2	

### Syntax

```
c.popretz {<reg_list_1_4> | <xreg_list_1_4>}, <stack_adj>
```

The variables used in the syntax are defined below.

```
<reg_list_1_4> ::= <ra> [", " <s0> | <s0-sN> ] (where N is 1,2)
```

```
if (<reg_list_1_4>=="ra")           <xreg_list_1_4>="x1"
if (<reg_list_1_4>=="ra, s0")        <xreg_list_1_4>="x1, x8"
if (<reg_list_1_4>=="ra, s0-s1")     <xreg_list_1_4>="x1, x8-x9"
if (<reg_list_1_4>=="ra, s0-s2")     <xreg_list_1_4>="x1, x8-x9, x18"
```

RV32:

```
<stack_adj>=[16|32|48|64]
```

RV64:

```
if (<reg_list_1_4>=="ra")           <stack_adj>=[16|32|48|64|96]
if (<reg_list_1_4>=="ra, s0")        <stack_adj>=[16|32|48|64|96]
if (<reg_list_1_4>=="ra, s0-s1")     <stack_adj>=[32|48|64|96|112]
if (<reg_list_1_4>=="ra, s0-s2")     <stack_adj>=[32|48|64|96|112]
```

### Description

This instruction pop (loads) the registers in *reg\_list\_1\_4* from stack memory, adjusts the stack pointer by *stack\_adj*, moves zero into a0 and then returns.

#### NOTE

All ABI register mapping are for the UABI. An EABI version is planned once the EABI is frozen.

For further information see [PUSH/POP Register Instructions](#).

Field decoding

The mapping from the *rlist2* field in the encoding are as shown below.

Table 8. *rlist2* decoding

rlist2	reg_list_1_4	stack_adj_base	
		RV32	RV64
0	ra	16	16
1	ra, s0	16	16
2	ra, s0-s1	16	32
3	ra, s0-s2	16	32

*stack\_adj\_base* is the minimum number of bytes, in multiples of 16-byte blocks, required to store the registers in *reg\_list\_1\_4*.

*spimm* is the number of additional 16-byte blocks allocated for the stack frame.

The total stack adjustment is calculated as shown.

```
stack_adj = stack_adj_base+spimm[5:4]*16
```

Prerequisites

The C-extension must also be configured.

32-bit equivalent

No direct equivalent encoding exists

## Operation

The first section of pseudo-code may be executed multiple times before the instruction successfully completes.

```
//This is not SAIL, it's pseudo-code. The SAIL hasn't been written yet.

if (misa.MXL==1) bytes=4; else bytes=8;

addr=sp+stack_adj-bytes;
for(i=31;i>0;i--) {
    //if register i is in xreg_list
    if (xreg_list[i]) {
        switch(bytes) {
            4:  asm("lw x[i], 0(addr)");
            8:  asm("ld x[i], 0(addr)");
        }
        addr-=bytes;
    }
}
```

The final section of pseudo-code executes atomically, and only executes if the section above completes without any exceptions or interrupts.

### NOTE

The *li a0, 0* **could** be executed more than once, but is included in the atomic section for convenience.

```
//This is not SAIL, it's pseudo-code. The SAIL hasn't been written yet.

asm("li a0, 0");
sp+=stack_adj;
asm("ret");
```

RV32 Assembly examples

```
c.popretz {ra, s0-s2}, 48
```

Encoding: *rlist2*=3, *spimm*=2 Equivalent sequence:

```
lw    s2, 44(sp);
lw    s1, 40(sp);
lw    s0, 36(sp);
lw    ra, 32(sp);
addi  sp, sp, 48;
li    a0, 0;
ret;
```

```
c.popretz {ra}, 16
```

Encoding: *rlist2*=0, *spimm*=0

Equivalent sequence:

```
lw    ra, 12(sp);
addi  sp, sp, 16;
li    a0, 0;
ret;
```

Included in

Extension	Minimum version	Lifecycle state
Zces ( <a href="#">Zces 0.53.2</a> )	0.53.2	Stable

## c.popretz.l

### Synopsis

Destroy stack frame: load ra and 4 to 12 saved registers from the stack frame, deallocate the stack frame, move zero into a0, return to ra.

### Mnemonic

c.popretz {*reg\_list\_5\_13*}, *stack\_adj*

### Encoding (RV32, RV64)

15					10	9		7	6	5		4		2	1	0
1	0	1	0	1	1	spimm[6:4]			1	1		rlist3			1	0
FUNCT6						reg_list_5_13						OP=C2				

### Syntax

```
c.popretz {<reg_list_5_13> | <xreg_list_5_13>}, -<stack_adj>
```

The variables used in the syntax are defined below.

<reg\_list\_5\_13> ::= <ra> ", " <s0-sN> (where N is 3,4,5,6,7,8,9,11)

if (<reg\_list\_5\_13>=="ra, s0-sN") <xreg\_list\_5\_13>="x1, x8-x9, x18-xM"  
(where M=N+16 and N is 3,4,5,6,7,8,9,11)

RV32:

if (<reg\_list\_5\_13>=="ra, s0-sN") <stack\_adj>=[32|48|..|528] (where N is 3-6)  
if (<reg\_list\_5\_13>=="ra, s0-sN") <stack\_adj>=[48|64|..|544] (where N is 7-9)  
if (<reg\_list\_5\_13>=="ra, s0-s11") <stack\_adj>=[64|80|..|560]

RV64:

if (<reg\_list\_5\_13>=="ra, s0-sN") <stack\_adj>=[ 48| 64|..|544] (where N is 3,4)  
if (<reg\_list\_5\_13>=="ra, s0-sN") <stack\_adj>=[ 64| 80|..|560] (where N is 5,6)  
if (<reg\_list\_5\_13>=="ra, s0-sN") <stack\_adj>=[ 80| 96|..|576] (where N is 7,8)  
if (<reg\_list\_5\_13>=="ra, s0-s9") <stack\_adj>=[ 96|112|..|576]  
if (<reg\_list\_5\_13>=="ra, s0-s11") <stack\_adj>=[112|128|..|592]

Description

This instruction pop (loads) the registers in *reg\_list\_5\_13* from stack memory, adjusts the stack pointer by *stack\_adj*, moves zero into a0 and then returns.

NOTE

All ABI register mapping are for the UABI. An EABI version is planned once the EABI is frozen.

For further information see [PUSH/POP Register Instructions](#).

Field decoding

The mapping from the *rlist3* and *spimm* fields in the encoding are as shown below.

Table 9. *rlist3* decoding

rlist3	reg_list_5_13	stack_adj_base	
		RV32	RV64
0	ra, s0-s3	32	48
1	ra, s0-s4	32	48
2	ra, s0-s5	32	64
3	ra, s0-s6	32	64
4	ra, s0-s7	48	80
5	ra, s0-s8	48	80
6	ra, s0-s9	48	96
7	ra, s0-s11	64	112

*stack\_adj\_base* is the minimum number of bytes, in multiples of 16-byte blocks, required to store the registers in *reg\_list\_5\_13*.

*spimm* is the number of additional 16-byte blocks allocated for the stack frame.

The total stack adjustment is calculated as shown.

$$\text{stack\_adj} = \text{stack\_adj\_base} + \text{spimm}[6:4] * 16$$

Prerequisites

The C-extension must also be configured.

32-bit equivalent

No direct equivalent encoding exists.

## Operation

The first section of pseudo-code may be executed multiple times before the instruction successfully completes.

```
//This is not SAIL, it's pseudo-code. The SAIL hasn't been written yet.

if (misa.MXL==1) bytes=4; else bytes=8;

addr=sp+stack_adj-bytes;
for(i=31;i>0;i--) {
    //if register i is in xreg_list
    if (xreg_list[i]) {
        switch(bytes) {
            4:  asm("lw x[i], 0(addr)");
            8:  asm("ld x[i], 0(addr)");
        }
        addr-=bytes;
    }
}
```

The final section of pseudo-code executes atomically, and only executes if the section above completes without any exceptions or interrupts.

### NOTE

The *li a0, 0* **could** be executed more than once, but is included in the atomic section for convenience.

```
//This is not SAIL, it's pseudo-code. The SAIL hasn't been written yet.

asm("li a0, 0");
sp+=stack_adj;
asm("ret");
```

RV32 Assembly examples

```
c.popretz {ra, s0-s3}, 48
```

Encoding: *rlist2*=0, *spimm*=1

Equivalent sequence:

```
lw    s3, 44(sp);
lw    s2, 40(sp);
lw    s1, 36(sp);
lw    s0, 32(sp);
lw    ra, 28(sp);
addi  sp, sp, 48;
li    a0, 0;
ret;
```

```
c.popretz {ra, s0-s4}, 64
```

Encoding: *rlist2*=1, *spimm*=2

Equivalent sequence:

```
lw    s4, 60(sp);
lw    s3, 56(sp);
lw    s2, 52(sp);
lw    s1, 48(sp);
lw    s0, 44(sp);
lw    ra, 40(sp);
addi  sp, sp, 64;
li    a0, 0;
ret;
```

Included in

Extension	Minimum version	Lifecycle state
Zces ( <a href="#">Zces 0.53.2</a> )	0.53.2	Stable



## c.jt

### Synopsis

jump via table without link, 16-bit encoding

### Mnemonic

c.jt *#index*

### Encoding (RV32, RV64)



### NOTE

For this encoding to decode as *c.jt*, *index8*<64, otherwise it decodes as [c.jalt: jump via table and link to ra, 16-bit encoding](#).

### Syntax

```
c.jt #index
```

### Description

This instruction is used to dereference a table of PCs, and then jumps without linking to the dereferenced PC.

For further information see [Table Jump Instructions](#).

### Prerequisites

The C-extension must also be configured.

### 32-bit equivalent

No direct equivalent encoding exists.

## Operation

```
//This is not SAIL, it's pseudo-code. The SAIL hasn't been written yet.

# target_address is temporary internal state, it doesn't represent a real register
# Mem is byte indexed
# "index8" is the field from the encoding, not "index" passed to the C.JT/C.JALT
  in the assembler
# which is formed below

if (OPCODE=="C.JALT") {
    index = index8 - 64;
} else {
    index = index8;
}

switch(XLEN) {
    32: table_address[XLEN-1:0] = JVT.base + (index<<2);
    64: table_address[XLEN-1:0] = JVT.base + (index<<3);
}

//check for debug mode entry, trigger with timing=0 and action=1, haltreq or step
if ((debug_trigger(table_address) && MCONTROL.timing==0 && MCONTROL.action==1) ||
    external_debug_haltreq() || DCSR.step==1) {
    DPC          = current_PC;
    DCSR.cause = DCSR.step==1 ? 4 : external_debug_haltreq() ? 3 : 2;
    enter_debug_mode();
}
//check for breakpoint trigger which takes an exception with timing=0
} else if ((debug_trigger(table_address) && MCONTROL.timing==0) ||
            !can_access_instruction_memory(table_address)) {
    MEPC  = current_PC;
    MTVAL = table_address;
    MCAUSE = debug_trigger(table_address) ? BREAKPOINT : INSTRUCTION_ACCESS_FAULT;
    take_exception();
} else {
    //access the jump table
    switch(XLEN) {
        32: LW target_address, InstMemory[table_address][XLEN-1:0];
        64: LD target_address, InstMemory[table_address][XLEN-1:0];
    }

    //don't use haltreq or step here, only check the addresses
    //check for table_address after reading if timing=1
    if (debug_trigger(table_address) && MCONTROL.timing==1 && MCONTROL.action==1) {
        DPC          = current_PC;
        DCSR.cause = 2;
        enter_debug_mode();
    } else if (debug_trigger(table_address) && MCONTROL.timing==1) {
```

```

    MEPC      = current_PC;
    MTVAL     = table_address;
    MCAUSE    = BREAKPOINT;
    take_exception();
} else if ((debug_trigger(target_address) && MCONTROL.timing==0 &&
MCONTROL.action==1) {
    DPC       = target_address;
    DCSR.cause = 2;
    enter_debug_mode();
} else if (((debug_trigger(target_address) && MCONTROL.timing==0) ||
            !can_access_instruction_memory(target_address)) {
    MEPC      = target_address;
    MTVAL     = target_address;
    MCAUSE    = debug_trigger(target_address) ? BREAKPOINT :
INSTRUCTION_ACCESS_FAULT;
    take_exception();
} else {
    //jump to the target address
    if (OPCODE=="C.JALT") {
        JALR ra, target_address[XLEN-1:0]&~0x1;
    } else {
        JR target_address[XLEN-1:0]&~0x1;
    }
}
}
}

```

#### Included in

Extension	Minimum version	Lifecycle state
Zces ( <a href="#">Zces 0.53.2</a> )	0.53.2	Stable

# c.jalt

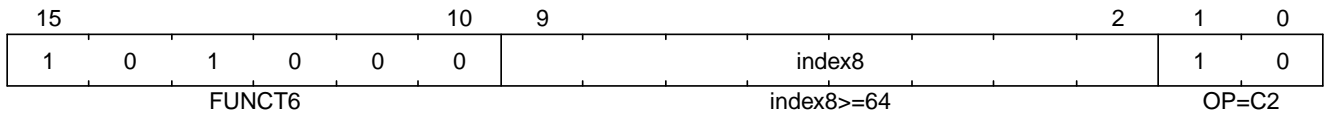
## Synopsis

jump via table and link to ra, 16-bit encoding

## Mnemonic

c.jalt *#index*

## Encoding (RV32, RV64)



### NOTE

For this encoding to decode as *c.jalt*, *index8*>=64, otherwise it decodes as *c.jt: jump via table without link, 16-bit encoding*.

## Syntax

```
c.jalt #index
```

### NOTE

*index* in the assembly syntax is valid from 0-192. *index8* in the encoding is valid from 64-255, so *index* = *index8*-64.

## Description

This instruction is used to dereference a table of PCs, and then jumps to the dereferenced PC and links to ra.

For further information see [Table Jump Instructions](#).

## Prerequisites

The C-extension must also be configured.

## 32-bit equivalent

No direct equivalent encoding exists.

## Operation

```
//This is not SAIL, it's pseudo-code. The SAIL hasn't been written yet.

# target_address is temporary internal state, it doesn't represent a real register
# Mem is byte indexed
# "index8" is the field from the encoding, not "index" passed to the C.JT/C.JALT
  in the assembler
# which is formed below

if (OPCODE=="C.JALT") {
    index = index8 - 64;
} else {
    index = index8;
}

switch(XLEN) {
    32: table_address[XLEN-1:0] = JVT.base + (index<<2);
    64: table_address[XLEN-1:0] = JVT.base + (index<<3);
}

//check for debug mode entry, trigger with timing=0 and action=1, haltreq or step
if ((debug_trigger(table_address) && MCONTROL.timing==0 && MCONTROL.action==1) ||
    external_debug_haltreq() || DCSR.step==1) {
    DPC          = current_PC;
    DCSR.cause = DCSR.step==1 ? 4 : external_debug_haltreq() ? 3 : 2;
    enter_debug_mode();
}
//check for breakpoint trigger which takes an exception with timing=0
} else if ((debug_trigger(table_address) && MCONTROL.timing==0) ||
            !can_access_instruction_memory(table_address)) {
    MEPC  = current_PC;
    MTVAL = table_address;
    MCAUSE = debug_trigger(table_address) ? BREAKPOINT : INSTRUCTION_ACCESS_FAULT;
    take_exception();
} else {
    //access the jump table
    switch(XLEN) {
        32: LW target_address, InstMemory[table_address][XLEN-1:0];
        64: LD target_address, InstMemory[table_address][XLEN-1:0];
    }

    //don't use haltreq or step here, only check the addresses
    //check for table_address after reading if timing=1
    if (debug_trigger(table_address) && MCONTROL.timing==1 && MCONTROL.action==1) {
        DPC          = current_PC;
        DCSR.cause = 2;
        enter_debug_mode();
    } else if (debug_trigger(table_address) && MCONTROL.timing==1) {
```

```
MEPC      = current_PC;
MTVAL     = table_address;
MCAUSE    = BREAKPOINT;
take_exception();
} else if ((debug_trigger(target_address) && MCONTROL.timing==0 &&
MCONTROL.action==1) {
    DPC      = target_address;
    DCSR.cause = 2;
    enter_debug_mode();
} else if (((debug_trigger(target_address) && MCONTROL.timing==0) ||
            !can_access_instruction_memory(target_address)) {
    MEPC      = target_address;
    MTVAL     = target_address;
    MCAUSE    = debug_trigger(target_address) ? BREAKPOINT :
INSTRUCTION_ACCESS_FAULT;
    take_exception();
} else {
    //jump to the target address
    if (OPCODE=="C.JALT") {
        JALR ra, target_address[XLEN-1:0]&~0x1;
    } else {
        JR target_address[XLEN-1:0]&~0x1;
    }
}
}
```

Included in

Extension	Minimum version	Lifecycle state
Zces ( <a href="#">Zces 0.53.2</a> )	0.53.2	Stable

# JVT CSR

## Synopsis

Table jump base vector and control register

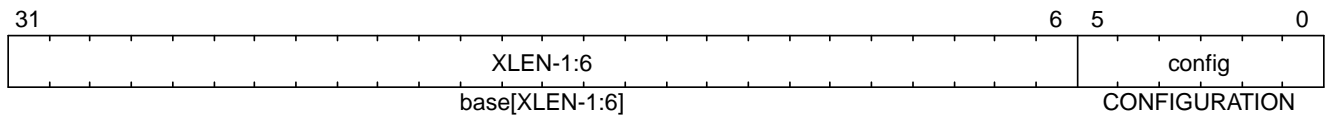
## Address

TBD

## Permissions

URW

## Format (RV32, RV64)



## Description

*JVT.base* is a virtual address, whenever virtual memory is enabled.

Using *JVT.base*[5:0] is implicitly zero, and is naturally aligned for all legal values of *XLEN*.

The memory pointed to by *JVT.base* is treated as instruction memory for the purpose of executing table jump instructions.

Table 10. *JVT.config* definition

JVT.config	Comment
000000	Jump table mode
others	<b>reserved for future standard use</b>

*JVT.config* is a WARL field, so can only be programmed to modes which are implemented. Therefore the discovery mechanism is to attempt to program different modes and read back the values to see which are available. Jump table mode *must* be implemented.

## Architectural State

JVT adds architectural state to the context, therefore must be saved/restored on context switches.

Additional architectural state requires a state enable to be allocated. Accesses when the state is disabled will throw an illegal instruction exception. The state enable is not specified in this document.

## Included in

Extension	Minimum version	Lifecycle state
Zces ( <a href="#">Zces 0.53.2</a> )	0.53.2	Stable

# c.mva01s

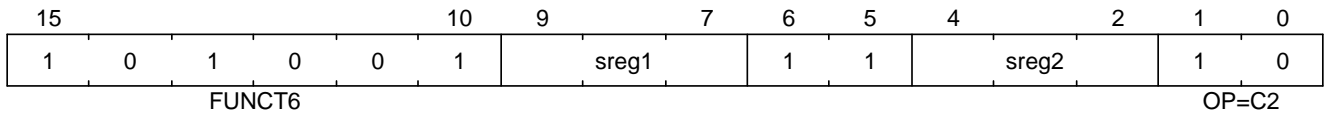
## Synopsis

Move two s0-s7 registers into a0-a1, 16-bit encoding

## Mnemonic

c.mva01s *sreg1*, *sreg2*

## Encoding (RV32, RV64)



## Syntax

```
c.mva01s sreg1, sreg2
```

## Description

This instruction moves *sreg1* into *a0* and *sreg2* into *a1*. The execution is atomic, so it is not possible to observe state where only one of *a0* or *a1* have been updated.



## Field decoding

Table 11. *sreg* decoding

<b>sreg*</b>	<b>xreg</b>
0	x8
1	x9
2	x18
3	x19
4	x20
5	x21
6	x22
7	x23

The encoding has two *sreg* number specifiers to save encoding space.

### NOTE

The *s* register mapping is taken from the UABI, and may not match the currently unratified EABI. *c.mva01s.e* may be included in the future.

## Prerequisites

The C-extension must also be configured.

## 32-bit equivalent

No direct equivalent encoding exists.

Operation

```
//This is not SAIL, it's pseudo-code. The SAIL hasn't been written yet.  
  
xreg1 = {sreg1[2:1]>0,sreg1[2:1]==0,sreg1[2:0]}  
xreg2 = {sreg2[2:1]>0,sreg2[2:1]==0,sreg2[2:0]}  
  
X[10] = X[xreg1]  
X[11] = X[xreg2]
```

Included in

Extension	Minimum version	Lifecycle state
Zces ( <a href="#">Zces 0.53.2</a> )	0.53.2	Stable

# PUSH/POP register instructions

These instructions are collectively referred to as PUSH/POP:

- *c.pusha*: Create stack frame: push registers, move A to S registers, allocate additional stack space.
- *c.pusha.l*: Create stack frame: push registers, move A to S registers, allocate additional stack space.
- *c.push*: Create stack frame: push registers, allocate additional stack space.
- *c.push.l*: Create stack frame: push registers, allocate additional stack space.
- *c.pop*: Destroy stack frame: pop registers, deallocate stack frame.
- *c.popret*: Destroy stack frame: pop registers, deallocate stack frame, return.
- *c.popret.l*: Destroy stack frame: pop registers, deallocate stack frame, return.
- *c.popretz*: Destroy stack frame: pop registers, deallocate stack frame, return zero.
- *c.popretz\_l*: Destroy stack frame: pop registers, deallocate stack frame, return zero.

The term PUSH refers to *c.push*, *c.pusha*, *c.push.l* and *c.pusha.l*. The assembly syntax for all of these uses the mnemonic *c.push*.

The term POP refers to *c.pop*.

The term POPRET refers to *c.popret*, *c.popretz*, *c.popret.l* and *c.popretz.l*. The assembly syntax for all of these uses the mnemonics *c.popret*/*c.popretz*.

Common details for these instructions are in this section.

## PUSH/POP overview

PUSH, POP, POPRET are used to reduce the size of function prologues and epilogues.

### 1. The PUSH instructions

- push(store) the registers specified in the register list to the stack
- *c.pusha*, *c.pusha.l* also move the registers in *areg\_list* into *s* registers
  - *areg\_list* is determined automatically from the register list, it cannot be arbitrarily specified to save encoding space.
- adjust the stack pointer

### 2. The POP instruction

- pops(loads) the registers in the register list from the stack.
- adjusts the stack pointer

### 3. The POPRET instructions

- pop(load) the registers in the register list from the stack.
- *c.popretz*, *c.popretz.l* also move zero into *a0* as the return value
- adjust the stack pointer
- execute a *ret* instruction.

## Example usage

This example gives an illustration of the use of PUSH and POPRET.

```
int function(void *buf, size_t len)
{
    return function2(buf, len);
}
```

compiles with GCC10 to:

```
20405458 <function>:
20405458: 1141          addi sp,sp,-16      ;#c.pusha(1)
2040545a: c04a          sw  s2,0(sp)      ;#c.pusha(2)
...
20405464: c422          sw  s0,8(sp)      ;#c.pusha(3)
20405466: c226          sw  s1,4(sp)      ;#c.pusha(4)
20405468: c606          sw  ra,12(sp)     ;#c.pusha(5)
2040546a: 842a          mv  s0,a0          ;#c.pusha(6)
2040546c: 84ae          mv  s1,a1          ;#c.pusha(7)
<function body>
20405494: 4501          li  a0,0           ;#c.popretz(1)
20405496: 40b2          lw  ra,12(sp)     ;#c.popretz(2)
20405498: 4422          lw  s0,8(sp)      ;#c.popretz(3)
2040549a: 4492          lw  s1,4(sp)      ;#c.popretz(4)
2040549c: 4902          lw  s2,0(sp)      ;#c.popretz(5)
2040549e: 0141          addi sp,sp,16     ;#c.popretz(6)
204054a0: 8082          ret              ;#c.popretz(7)
```

with the GCC option *-msave-restore* the output is the following:

```
204089ac <function>:
204089ac: f97f72ef      jal  t0,20400942 <__riscv_save_0> ;#c.pusha(1)
...
204089b8: 842a          mv  s0,a0          ;#c.pusha(2)
204089ba: 84ae          mv  s1,a1          ;#c.pusha(3)
<function_body>
204089e2: 4501          li  a0,0           ;#c.popretz(1)
204089e4: f83f706f      j    20400966 <__riscv_restore_0> ;#c.popretz(2)
```

with PUSH/POPRET this reduces to

```

20405458 <function>:
20405458: <16-bit>          c.push    {ra,s0-s2},{a0-a2},-16
<function body>
20405496: <16-bit>          c.popretz {ra,s0-s2}, 16

```

The prologue / epilogue reduce from 28-bytes in the original code, to 14-bytes with *-msave-restore*, and to 4-bytes with PUSH and POPRET. As well as reducing the code-size PUSH and POPRET eliminate the branches from calling the millicode *save/restore* routines so also perform better.

**NOTE**

The calls to `<riscv_save_0>/<riscv_restore_0>` become 64-bit when the target functions are out of the  $\pm 1\text{MB}$  range, increasing the prologue/epilogue size to 22-bytes.

**NOTE**

*c.pusha* has an additional register move included *mv s2, a2* which wasn't in the original prologue. This is included to simplify the encoding and definition of *c.pusha* and will cost some performance.

**NOTE**

POP is used for tail-calling which is not included in this example.

## Compiler implementation

The technique used in the initial implementation in LLVM is to let the compiler generate the function prologue and epilogue, and then replace the instruction sequences with the relevant PUSH/POP instructions.

### *spimm* handling

The instructions have a restricted range of *spimm* available. If this is insufficient then a separate *c.addi16sp* can be used to increase the range.

### register list handling

*c.pop* has a restricted register list. If more registers are required then *c.pop* can be used for some of them, and the rest can be explicitly loaded.

The other instructions do not directly support  $\{ra, s0-s10\}$  to reduce the amount of encoding space required. If this register list is required then *s11* should also be included. This costs a small amount of memory and performance, but saves code-size.

### *areg\_list* handling

*c.pusha*, *c.pusha.l* include *areg\_list*. This may not match what was generated by the compiler.

## Example: *c.pusha* fits perfectly

In this real world example generated by GCC10, *c.pusha* fits perfectly.

```

00e010b8 <function>:
e010b8:      1141                addi    sp,sp,-16 ; #c.pusha
e010ba:      c422                sw      s0,8(sp) ; #c.pusha
e010bc:      c226                sw      s1,4(sp) ; #c.pusha
e010be:      c04a                sw      s2,0(sp) ; #c.pusha
e010c0:      c606                sw      ra,12(sp) ; #c.pusha
e010c2:      842a                mv      s0,a0 ; #c.pusha
e010c4:      84ae                mv      s1,a1 ; #c.pusha
e010c6:      4908                lw      a0,16(a0)
e010c8:      4d8c                lw      a1,24(a1)
e010ca:      8932                mv      s2,a2 ; #c.pusha
e010cc:      726040ef          jal     ra,e057f2 <function2>

```

this is replaced by

```

00e010b8 <function1>:
e010b8:      xxxx                c.push {ra,s0-s2}, {a0-a2}, -16
e010c6:      4908                lw      a0,16(a0)
e010c8:      4d8c                lw      a1,24(a1)
e010cc:      726040ef          jal     ra,e057f2 <function2>

```

## Example: *areg\_list* doesn't fit

In this other real world example *areg\_list* doesn't fit:

```

00e01126 <function3>:
e01126:      1101                addi    sp,sp,-32 ; #c.push.l
e01128:      ce06                sw      ra,28(sp) ; #c.push.l
e0112a:      cc22                sw      s0,24(sp) ; #c.push.l
e0112c:      ca26                sw      s1,20(sp) ; #c.push.l
e0112e:      c84a                sw      s2,16(sp) ; #c.push.l
e01130:      c64e                sw      s3,12(sp) ; #c.push.l
e01132:      c452                sw      s4,8(sp) ; #c.push.l
e01134:      c256                sw      s5,4(sp) ; #c.push.l
e01136:      c05a                sw      s6,0(sp) ; #c.push.l
e01138:      0e050363          beqz    a0,e0121e <function3+0xf8>
e0113c:      8a2a                mv      s4,a0
e0113e:      852e                mv      a0,a1
e01140:      89ae                mv      s3,a1

```

In this case, the move instructions are not part of the same basic block so *c.push.l* is selected:

```

00e01126 <function4>:
e01126:      xxxx                c.push {ra,s0-s6}, -32
e01138:      0e050363           beqz    a0,e0121e <function4+0xf8>
e0113c:      8a2a                mv     s4,a0
e0113e:      852e                mv     a0,a1
e01140:      89ae                mv     s3,a1

```

## Example: *areg\_list* needs register allocation changes

The next case is where none of the register moves match the *areg\_list* moves because the register allocator in the compiler did not allocate suitable registers:

```

00e01842 <function5>:

e01e7e:      1101                addi    sp,sp,-32
e01e80:      cc22                sw     s0,24(sp)
e01e82:      c84a                sw     s2,16(sp)
e01e84:      c64e                sw     s3,12(sp)
e01e86:      c452                sw     s4,8(sp)
e01e88:      c256                sw     s5,4(sp)
e01e8a:      ce06                sw     ra,28(sp)
e01e8c:      ca26                sw     s1,20(sp)
e01e8e:      892a                mv     s2,a0
e01e90:      89ae                mv     s3,a1
e01e92:      8a32                mv     s4,a2
e01e94:      8ab6                mv     s5,a3
e01e96:      3f41                jal     e01e26 <function6>

```

With *c.pusha.l* this becomes:

```

e01e7e <function5>:
# c.push includes moving {a0-a3} into {s0-s3}
e01e7e:      1101                c.push {ra,s0-s5}, {a0-a3}, -32
e01e8e:      892a                mv     s2,a0;# <-- switch dest to s0
e01e90:      89ae                mv     s3,a1;# <-- switch dest to s1
e01e92:      8a32                mv     s4,a2;# <-- switch dest to s2
e01e94:      8ab6                mv     s5,a3;# <-- switch dest to s3
e01e96:      3f41                jal     e01e26 <function6>

```

In this case all four moves can be deleted if the register allocation can be altered. if the register allocation *cannot* be altered, then *c.push.l* should be used instead.



## Example: *areg\_list* partially fits

In this final case, one register move can be deleted and one must be retained unless the register allocation can be changed.

```
00e02368 <function7>:
e02368:      1141                addi    sp,sp,-16
e0236a:      c226                sw      s1,4(sp)
e0236c:      03450493            addi    s1,a0,52
e02370:      c422                sw      s0,8(sp)
e02372:      842a                mv      s0,a0;# <-- delete this one
e02374:      8526                mv      a0,s1;# <-- doesn't fit areg_list
e02376:      c04a                sw      s2,0(sp)
e02378:      c606                sw      ra,12(sp)
e0237a:      892e                mv      s2,a1;# <-- switch dest to s1
e0237c:      df3fd0ef            jal     ra,e0016e <function8>
```

```
00e02368 <function7>:
e02368:      xxxx                c.push {ra,s0-s2}, {a0-a2}, -16
e0236c:      03450493            addi    s1,a0,52
e02374:      8526                mv      a0,s1;# <-- doesn't fit areg_list
e0237a:      892e                mv      s2,a1;# <-- switch dest to s1
e0237c:      df3fd0ef            jal     ra,e0016e <function8>
```

In this case one move is deleted, but one remains because unless the target register can be reallocated.

For the smallest code-size the compiler should reallocate the target registers so that the moves in *areg\_list* are not wasted.

## PUSH/POP Fault handling

The sequence required to execute the PUSH/POP instruction may be interrupted, or may not be able to start execution for several reasons.

- virtual memory page fault or PMP fault
  - these can be detected before execution, or during execution if the memory addresses cross a page/PMP boundary
  - MTVAL is set to any address which causes the fault
- watchpoint trigger
  - these can be detected before execution, or during execution depending on the trigger type (load data triggers require the sequence to have started executing, for example)
  - MTVAL is set to any address which causes the fault
- external debug halt
  - the halt can treat the whole sequence atomically, or interrupt mid sequence (implementation defined)

- debug halt caused by a trigger
  - same comment as watchpoint trigger above
- load access fault
  - these are detected while the sequence is executing
  - MTVALL is set to the fault address.
- store access fault (precise or imprecise)
  - these may be detected while the sequence is executing, or afterwards if imprecise
  - MTVALL is set to the fault address.
- interrupts
  - these may arrive at any time. An implementation can choose whether to interrupt the sequence or not.

**NOTE**

MTVALL may be hardwired to zero in an implementation. The section above assumes it is implemented.

In all case MEPC contain the PC of the PUSH/POP instruction, and MCAUSE is set as expected for the type of fault.

For debug halts DPC is set to the PC of the PUSH/POP instruction.

Because some faults can only be detected during the sequence the core implementation is able to recover from the fault and re-execute the sequence. This may involve executing some or all of the loads and stores from the sequence multiple times before the sequence completes (as multiple faults or multiple interrupts are possible).

Therefore correct execution requires that *sp* refers to idempotent memory (also see [Non-idempotent memory handling](#)).

## Software view of execution

### Software view of the PUSH sequence

From a software perspective the PUSH sequence appears as:

- A sequence of stores writing a contiguous block of memory. Any of the bytes may be written multiple times.
- An optional series of register moves
- A stack pointer adjustment

Because the memory is idempotent and the stores are non-overlapping, they may be reordered, grouped into larger accesses, split into smaller access or any combination of these.

If an implementation allows interrupts during the sequence, and the interrupt handler uses *sp* to allocate stack memory, then any stores which were executed before the interrupt may be overwritten by the handler. This is safe because the memory is idempotent and the stores will be re-executed execution resumes.

The stack pointer adjustment must only be committed once it is certain that all of the stores will complete within triggering any precise faults (stores may return imprecise bus errors which are received after the instruction has completed execution).

For example:

```
c.push {ra, s0-s5}, {a0-a3}, -64
```

Appears to software as:

```
# any bytes from sp-1 to sp-28 may be written multiple times before the
instruction completes
sw  s5, -4(sp);
sw  s4, -8(sp);
sw  s3, -12(sp);
sw  s2, -16(sp);
sw  s1, -20(sp);
sw  s0, -24(sp);
sw  ra, -28(sp);

# these must only execute once, and will only execute after all stores complete
sucessfully
mv  s0, a0
mv  s1, a1
mv  s2, a2
mv  s3, a3
addi sp, sp, -64;
```

## Software view of the POP/POPRET sequence

From a software perspective the POP/POPRET sequence appears as:

- A sequence of loads, any of which may be executed multiple times
- A stack pointer adjustment
- An optional LI zero into a0
- An optional RET

If an implementation allows interrupts during the sequence, then any loads which were executed before the interrupt may update architectural state. The loads will be re-executed once the handler completes, so the values will be overwritten. Therefore it is permitted for an implementation to update some of the destination registers before taking the interrupt or other fault.

The optional load immediate and stack pointer adjustment must only be committed once it is certain that all of the loads will complete successfully.

For POPRET once the stack pointer adjustment has been committed the RET must execute.

For example:

```
c.popretz {ra, s0-s3}, 32 ;
```

Appears to software as:

```
# any or all of these load instructions may execute multiple times
lw  s3, 28(sp);
lw  s2, 24(sp);
lw  s1, 20(sp);
lw  s0, 16(sp);
lw  ra, 12(sp);

# must only execute once, will only execute after all loads complete successfully
# all instructions must execute atomically
li a0, 0
addi sp, sp, 32;
ret;
```

## Non-idempotent memory handling

An implementation may have a requirement to issue a PUSH/POP instruction to non-idempotent memory.

### Error detection

If the core implementation does not have a requirement to support PUSH/POP to non-idempotent memories, and the core can use a PMA to detect that the memory is non-idempotent, then take a load (POP/POPRET) or store (PUSH) access fault exception.

### Non-idempotent support

It is possible to support non-idempotent memory. One reason is to re-use PUSH/POP as a restricted form of a load/store multiple instruction to a peripheral, as there is no generic load/store multiple instruction in the RISC-V ISA.

If accessing non-idempotent memory then it is *recommended* to:

1. Not allow interrupts during execution
2. Not allow external debug halt during execution
3. Detect any virtual memory page faults or PMP faults for the whole instruction before starting execution (instead of during the sequence)
4. Not split / merge / reorder the generated memory accesses

It is possible that one of the following will still occur during execution:

1. Watchpoint trigger
2. Load/store access fault

In these cases the core will jump to the debug or exception handler. If execution is required to continue afterwards (so the event is not fatal to the code execution), then the handler is required to do so in software.

By following these rules memory accesses will only ever be issued once, and in the order listed in the SAIL.

It is possible for implementations to follow these restricted rules and to safely access both types of memory. It is also possible for an implementation to use PMAs to detect the memory type and apply different rules, such as only allowing interrupts if accessing cacheable memory, for example.

**Included in**

Extension	Minimum version	Lifecycle state
Zces ( <a href="#">Zces 0.53.2</a> )	0.52	Stable

# Table Jump Instructions

These instructions are collectively referred to as table jump:

- [c.jt](#): jump via table without link, 16-bit encoding
- [c.jalt](#): jump via table and link to ra, 16-bit encoding

Common details for these instructions are in this section.

## Table Jump Overview

Table jump is a form of dictionary compression used to reduce the code size of JAL / AUIPC+JALR / JR / AUIPC+JR instructions.

Function calls and jumps to fixed labels typically take 32-bit or 64-bit instruction sequences.

Table jump allows the linker to:

- replace 32-bit *j* calls with *c.jt*
- replace 32-bit *jal ra* calls with *c.jalt*
- replace 64-bit *auipc/jalr* calls to fixed locations with *c.jt*
- replace 64-bit *auipc/jalr ra* calls to fixed locations with *c.jalt*
  - The AUIPC+JR/JALR sequence is used because the offset from the PC is out of the  $\pm 1\text{MB}$  range.

## JVT

The base of the table is in the JVT CSR (see [JVT CSR, table jump base vector and control register](#)), each table entry is XLEN bits.

The table entry number is from the *index8* field in the encoding, which controls the link register.

- *c.jt* : entries 0-63, link to *zero*
- *c.jalt* : entries 64-255, link to *ra*

Note that the LSB of every jump table entry is *ignored* which matches standard JALR behaviour.

If the same function is called with and without linking then it must have two entries in the table. This case does happen in practice but only affects a small number of entries so it does not waste much space in the table. It is typically caused by the same function being called with and without tail calling.

## Recommended algorithm for allocating entries in the jump table

Calls to each function are categorised as shown in [Table jump code size saving for each function call replacement](#).

Table 12. Table jump code size saving for each function call replacement

original sequence	Table Jump saving
J	$A*2-(XLEN/8)$ bytes
AUIPC+JR	$B*6-(XLEN/8)$ bytes
JAL ra	$C*2-(XLEN/8)$ bytes
AUIPC+JALR ra	$D*6-(XLEN/8)$ bytes

Each function is called by using one of the two link registers. The total saving per function is calculated by counting the number of calls and adding up the total saving from each replacement of the existing sequence with a Table Jump instruction, as follows:

```
saving_per_function_c_jt    = A * 2 + B * 6 - 2*(XLEN-8)
saving_per_function_c_jalt  = C * 2 + D * 6 - 2*(XLEN-8)
```

The functions are sorted so that the one with the highest saving is in table entry 0, the second highest in entry 1 etc. for that encoding.

### NOTE

This algorithm assumes that each function is only called with one link register. If the same function is called with more than one link register, then it must have two entries in the table.

This allows the core to cache the most frequent targets by caching the lowest numbered entries of each section of the jump table. Only caching a few entries will greatly improve the performance.

## Table Jump Fault handling

Table Jump involves two instruction fetches from a single instruction, and either fetch can cause a fault.

The sequence required to execute the table jump instruction may be interrupted, or may not be able to start execution for several reasons.

- virtual memory page fault or PMP fault
  - these can be detected before execution, or during execution if the memory addresses cross a page/PMP boundary
  - MTVAL is set to any address which causes the fault
- watchpoint trigger
  - these can be detected before execution, or during execution depending on the trigger type (load data triggers require the sequence to have started executing, for example)
  - MTVAL is set to any address which causes the fault
- external debug halt
  - the halt can treat the whole sequence atomically, or interrupt mid sequence (implementation defined)
- debug halt caused by a trigger
  - same comment as watchpoint trigger above
- load access fault
  - these are detected while the sequence is executing
  - MTVAL is set to the fault address.
- store access fault (precise or imprecise)
  - these may be detected while the sequence is executing, or afterwards if imprecise
  - MTVAL is set to the fault address.
- interrupts
  - these may arrive at any time. An implementation can choose whether to interrupt the sequence or not.

In all case MEPC contain the PC of the table jump instruction, and MCAUSE is set as expected for the type of fault.

For debug halts DPC is set to the PC of the table jump instruction.

This section gives an overview of the behaviour, the exact operation is documented in the SAIL code for each instruction

- [c.jalt SAIL code](#)
- [c.jt SAIL code](#)

### Included in

Extension
Minimum version
Lifecycle state



Zces ( <a href="#">Zces 0.53.2</a> )
0.53.2
Stable