

This is a 50 minute exam, commencing at 10:10am and finishing at 11:00am. You may not use any reference material during the exam, including books or notes. You may not use any electronics equipment (cell phones, calculators, computers, etc.) during the exam. *Please read every question carefully before answering* and express your answers as completely as you can.

There are 4 questions on the exam, each worth 10 points; **you may choose any three to complete**. Thus the highest score that can be achieved on the exam is 30 points. Please indicate which 3 questions you wish to be graded in the check boxes next to the question numbers below.

Name: _____

Question	Grade?	Score
1		
2		
3		
4		
Total		

1. (10 points.) Define SCHEME functions with the following specifications.

(a) (2 points.) Define a SCHEME function `max2` which takes two numeric inputs (call them x and y) and returns the larger of the two. (You may not use the built-in scheme function `max` for this purpose—define your own function from scratch using a conditional.)

(b) (2 points.) Define a SCHEME function `max3` which takes three numeric inputs (call them x , y , and z) and returns the largest of the three. (You may not use the built-in scheme function `max` for this purpose, but you may use your function `max2` from the previous problem.)

(c) (2 points.) Define a SCHEME function `crazy` which takes a single input x and returns

$$\frac{(x + 10)(x + 10) + 10}{x + 10}.$$

- (d) **(2 points.)** Define a SCHEME function `monster-fact` which takes a single numeric argument x and returns $(x!)!$. (Yes, that's the factorial of the factorial of x , so `(monster-fact 4)` should return $(4!)! = 24! = 620448401733239439360000$. You must define the `factorial` function from scratch, if you intend to use it.)

- (e) **(2 points.)** Define a SCHEME function `dfact` (which stands for “double factorial”). The double factorial function is defined (for the natural numbers $\{0, 1, 2, \dots\}$) by the recursive rule:

$$\text{dfact}(n) = \begin{cases} 1 & \text{if } n = 0, \\ 1 & \text{if } n = 1, \\ n \cdot \text{dfact}(n-2) & \text{otherwise.} \end{cases}$$

2. (10 points.) In this problem, you will implement *Halley's method* for extracting square roots. Consider the following mysterious function of two variables Q and g :

$$h(Q, g) = \frac{g(g^2 + 3Q)}{3g^2 + Q}.$$

Edmond Halley (after whom the comet is named) noticed the following remarkable fact. For any (positive) number Q and any (positive) number g , $h(Q, g)$ is always closer to \sqrt{Q} than g was. In particular, for any specific fixed value of Q we can consider the sequence of numbers

$$g_s = \begin{cases} 1 & \text{if } s = 0, \\ h(Q, g_{s-1}) & \text{for } s > 0. \end{cases}$$

Then these numbers g_0, g_1, \dots converge very quickly to \sqrt{Q} .

- (a) (1 points.) Write a SCHEME function `h` which computes the function above (so, the function `h` should take two arguments Q and g).

- (b) (3 points.) Write a function `Halley-iterate` which takes two numeric arguments, Q and t , and returns g_t (defined by the sequence above). Thus, if $t = 0$, your function should return 1; if $t = 1$, your function should return $g_1 = h(Q, 1)$, etc.

- (c) (4 points.) Write a function `Halley-approx` which takes a single numeric argument Q and returns an approximation α to the square root of Q with the property that $|\alpha^2 - Q| \leq .00001$. To do this, your function should effectively compute the sequence g_0, g_1, g_2, \dots until it finds a g_t for which $|g_t^2 - Q| \leq .00001$ (at which point it can simply return the value g_t).

- (d) (2 points.) Show how to restructure your code from the previous problems so that h is defined in the scope of `Halley-approx`, and use this restructuring to make h a function of a single parameter g .

3. (10 points.) The integers 1, 2, 4, and 5 can be written as the *sum of two perfect squares*:

$$1 = 0^2 + 1^2, \quad 2 = 1^2 + 1^2, \quad 4 = 0^2 + 2^2, \quad \text{and} \quad 5 = 1^2 + 2^2.$$

On the other hand, neither 3, 6, nor 7 can be expressed this way.

In this problem, you will define a function `sum-of-squares` so that `(sum-of-squares n)` returns `#t` if the positive integer n can be written as a sum of squares of two integers and `#f` otherwise. You may assume that n is positive. If you wish, you may use the following function `is-square`, which returns `#t` if k is a perfect square, and `#f` otherwise:

```
(define (is-square k) (let ((int-root (round (sqrt k))))  
                        (= k (* int-root int-root))))
```

- (a) (2 points.) Define a SCHEME function `square-pieces` so that `(square-pieces x n)` returns `#t` if both x and $(n - x)$ are perfect squares, and `#f` otherwise.
- (b) (4 points.) Observe that n can be written as a sum of two squares exactly when there is a number $x \in \{0, \dots, n\}$ for which `(square-pieces x n)` returns `#t`. Write a SCHEME function `(test-up-to k n)` which returns `#t` if there is a number $x \in \{0, \dots, k\}$ for which `(square-pieces x n)` is true.
- (c) (4 points.) Using the above functions, define the SCHEME function `sum-of-squares`. For full credit, indicate how the definitions of your helper functions can be made private; do you need to keep passing around the parameter n ?

4. (10 points.)

(a) (2 points.) Consider the following version of `factorial`:

```
(define (factorial n)
  (let ((recursive-value (factorial (- n 1))))
    (if (= n 0)
        1
        (* n recursive-value))))
```

For which, if any, values of n does this correctly compute $n!$? Explain.

(b) (2 points.) Consider the following declaration:

```
(define (f x)
  (define (g y) (+ x y))
  (define (h x) (+ x (g x)))
  (h (+ x 10)))
```

After this, what would `(f 100)` return?

(c) (2 points.) To what does the following expression evaluate?

```
(let ((x 10)
      (y 20)
      (z 40))
  (let ((x (+ x 10))
        (y (+ x 20)))
    (+ z (- x y))))
```

- (d) (2 points.) Consider the following two implementations of multiplication (the first is called `times`; the second is called `ftimes`).

```
(define (times x y)
  (cond ((= x 0) 0)
        ((= y 0) 0)
        (else (+ x (times x (- y 1))))))
```

and

```
(define (double x) (* 2 x))
(define (half x) (floor (/ x 2)))
(define (even? x) (= 0 (modulo x 2)))
(define (ftimes x y)
  (cond ((= x 0) 0)
        ((even? x) (double (ftimes (half x) y)))
        (else (+ y (ftimes (- x 1) y)))))
```

These both correctly compute multiplication (if the arguments are non-negative integers). However, the behavior of a call to `(times 1000000 1000001)` will be quite different from that of a call to `(ftimes 1000000 1000001)`. Explain.

- (e) (2 points.) Consider the following flawed implementation of the Fibonacci sequence.

```
(define (fib n)
  (cond ((= n 0) 0)
        (else (+ (fib (- n 1))
                  (fib (- n 2))))))
```

For which values, if any, of n will this correctly compute the n th Fibonacci number? What's wrong?

SCRATCH SPACE

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