

COURSE PROJECT

Algorithmic Methods for Mathematical Models

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Problem statement & ILP model

Problem Statement

Given

- A set of street crossings $i = 1, \dots, N$.
- Camera models $k = 1, \dots, K$ with:
 - price P_k , range R_k , autonomy A_k , daily cost E_k .
- Coverage matrix M_{ij} .
- Weekly schedule (7 days).

Goal: Help Batman **minimize the total cost** while covering all crossings during the seven days of the week.



Example (simplified)

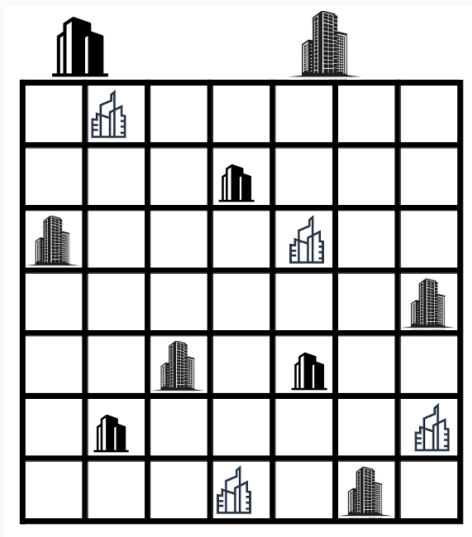


Figure 1: Example with $K = 5$, $N = 64$, blank

Example (simplified)

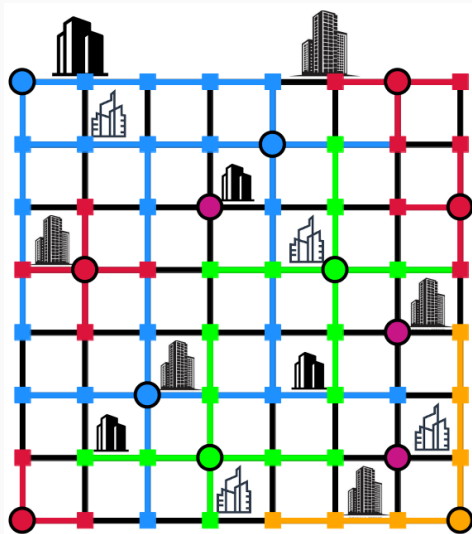


Figure 2: Example with $K = 5$, $N = 64$, complete

Decision variables

- Installation variable:

$$x_{ik} = \begin{cases} 1, & \text{if we install a camera } k \text{ at } i, \\ 0, & \text{otherwise.} \end{cases}$$

- Operation variable:

$$y_{ikd} = \begin{cases} 1, & \text{if camera } (i, k) \text{ is ON on day } d, \\ 0, & \text{otherwise.} \end{cases}$$

Objective function

$$z = \min \left(\sum_{i=1}^N \sum_{k=1}^K P_k x_{ik} + \sum_{i=1}^N \sum_{k=1}^K \sum_{d=1}^7 E_k y_{ikd} \right)$$

Constraints

- Placement

$$\sum_{k=1}^K x_{ik} \leq 1 \quad \forall i$$

- Coverage

$$\sum_{\substack{i=1..N \\ k:M_{ij} \leq R_k}} y_{ikd} \geq 1 \quad \forall j, d$$

- Consistency

$$y_{ikd} \leq x_{ik} \quad \forall i, k, d$$

- Autonomy

$$\sum_{h=0}^{A_k} y_{ik, ((d+h-1) \bmod 7)+1} \leq A_k$$

$$y_{ikd} = 1 \Rightarrow y_{ik, d+1} = 1$$

Heuristics

Algorithm 1 Greedy constructive heuristic

Require: (K, P, R, A, E, N, M)

```
1:  $C \leftarrow \text{INITCANDIDATES}()$ 
2:  $S \leftarrow \emptyset$ 
3:  $covered \leftarrow \emptyset$ 
4: while  $\neg \text{ISSOLUTION}(covered, N)$  do
5:    $(c^*, q^*) \leftarrow \text{EVALUATEQUALITY}(C, covered)$ 
6:   if  $c^* = \text{null}$  then break
7:   end if
8:   Add  $c^*$  to  $S$ 
9:   Update  $covered$ 
10:   $C \leftarrow \text{FEASIBILITY}(C, c^*)$ 
11: end while
12: return  $S$ 
```

Trade-off between **gain** (newly covered crossings) and **cost**.

$$q(c) = \frac{|covers(c) \setminus covered|}{P_{k(c)} + E_{k(c)} \cdot |days(c)|}$$

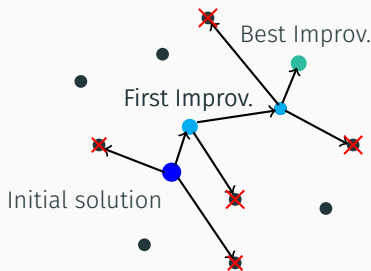
Tie-breaking rule: first candidate in the list.

Local Search (LS)

Algorithm 2 LOCALSEARCH (simplified)

Require: S_{initial} from Greedy and policy FI/BI

```
1:  $S \leftarrow S_{\text{initial}}$ 
2:  $S_{\text{best}} \leftarrow S$ 
3: while there exist neighbors to evaluate do
4:   Evaluate new neighbor  $S_{\text{new}}$ 
5:   if  $S_{\text{new}}$  is better than  $S$  then
6:     if policy = FI then
7:        $S \leftarrow S_{\text{new}}$ 
8:       break
9:     else ▷ Best Improvement
10:      if  $S_{\text{new}}$  is better than  $S_{\text{best}}$  then
11:         $S_{\text{best}} \leftarrow S_{\text{new}}$ 
12:      end if
13:    end if
14:  end if
15: end while
16: if policy = BI then
17:    $S \leftarrow S_{\text{best}}$ 
18: end if
19: return  $S$  ▷ Local optimum
```



Algorithm 3 GRASP constructive phase (simplified)

Require: $(K, P, R, A, E, N, M, \alpha)$

```
1:  $C \leftarrow \text{INITCANDIDATES}()$ 
2:  $S \leftarrow \emptyset$ ,  $\text{covered} \leftarrow \emptyset$ ,
3: while  $\neg \text{ISOLUTION}(\text{covered}, N)$  do
4:    $(q_{\max}, q_{\min}) \leftarrow \text{EVALUATEQUALITY}(C, \text{covered})$ 
5:    $\text{RCL} \leftarrow \text{BUILDRCL}(C, \text{covered}, P, E, q_{\max}, q_{\min}, \alpha)$ 
6:    $c^{\text{rand}} \leftarrow \text{RANDOMCHOICE}(\text{RCL})$ 
7:    $S \leftarrow S \cup \{c^{\text{rand}}\}$ 
8:   Update  $\text{covered}$ 
9:    $C \leftarrow \text{FEASIBILITY}(C, c^{\text{rand}})$ 
10: end while
11: return  $S$ 
```

Restricted Candidate List construction (**RCL**):

$$\text{RCL} = \{c \in C \mid q(c) \geq q_{\max} - \alpha(q_{\max} - q_{\min})\}.$$

Tuning and instances

Instance	N	K	$N \cdot K$
small	4	2	8
medium	10	5	50
large	20	10	200

Table 1: Set of generated instances used for performance evaluation.

GRASP tuning: Maximum Iterations (I)

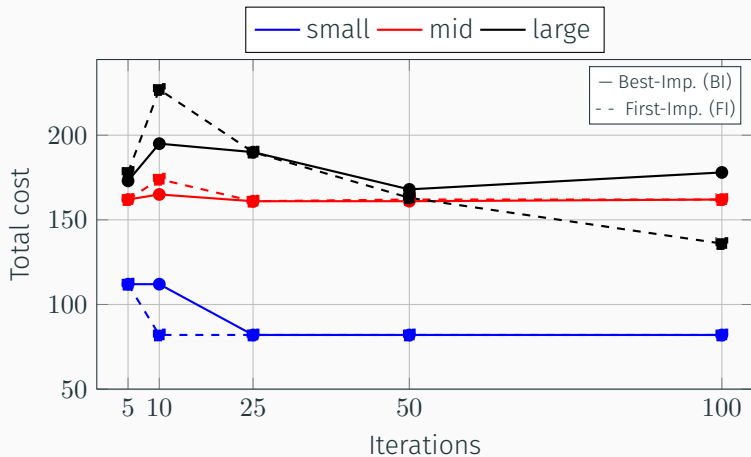


Figure 3: Total cost as a function of the maximum number of iterations ($\alpha = 0.5$).

GRASP tuning: Maximum Iterations (II)

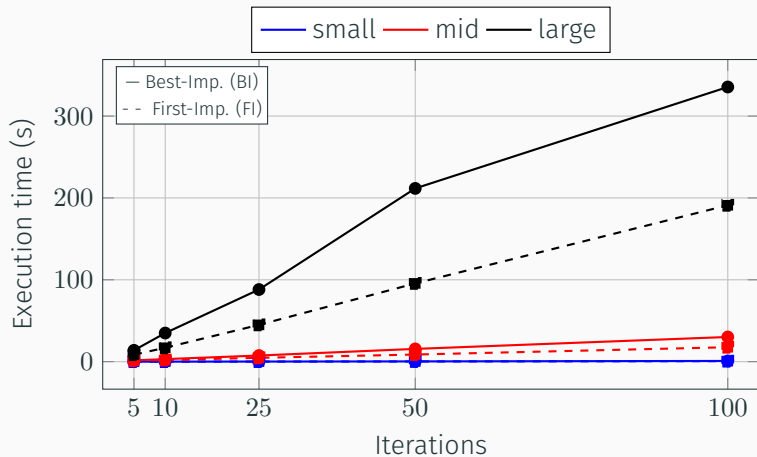


Figure 4: Execution time as a function of the maximum number of iterations ($\alpha = 0.5$).

GRASP tuning: Alpha (α)

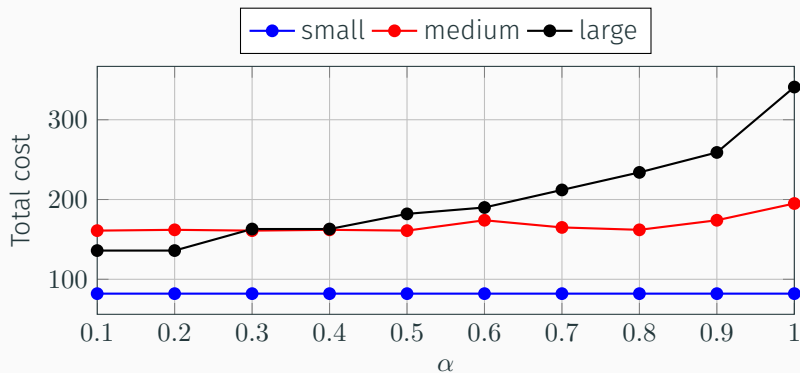


Figure 5: Effect of α on total cost (50 iterations, First Improvement policy).

Instance Generation (list)

Instance	N	K	$N \cdot K$
1	5	5	25
2	7	7	49
3	10	10	100
4	20	20	400
5	30	30	900
6	50	50	2500
7	75	75	5625
8	2	20	40
10	10	20	200
11	40	2	80
12	40	20	800

Table 2: Set of generated instances used for performance evaluation.

Experimental results

- ILP (CPLEX)
- Greedy
 - Greedy constructive (simple)
 - Greedy constructive + LS with First improvement policy
 - Greedy constructive + LS with Best improvement policy
- GRASP
 - GRASP constructive (simple)
 - GRASP constructive + LS with First improvement policy
 - GRASP constructive + LS with Best improvement policy

Procedure : find the best setting for Greedy / GRASP, and compare them with ILP.

Results: Greedy



Figure 6: Objective value as a function of the instance size for Greedy.

Results: Greedy

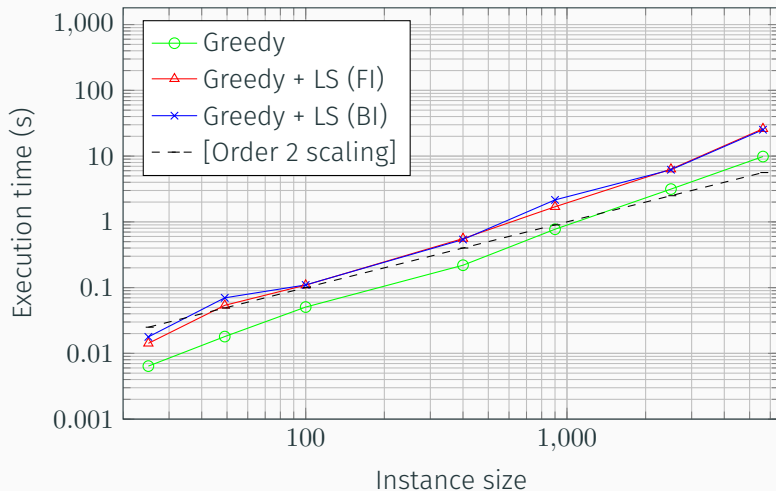


Figure 7: Execution time as a function of the instance size for Greedy.

Results: Greedy

- **Objective function:** FI = BI for all instances; simple Greedy stays close (especially for large instances).
- **Execution time :** simple Greedy is significantly faster; FI little faster than BI.

Choice: Computation is fast in all cases. Accuracy should be prioritised. Between FI and BI, FI is faster with similar performance.

⇒ **Greedy + LS (FI).**

Results: GRASP

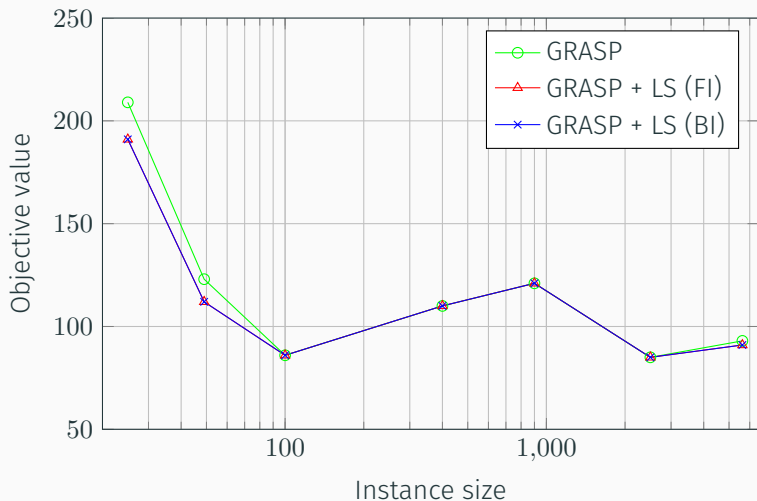


Figure 8: Objective value as a function of the instance size for GRASP.

Results: GRASP

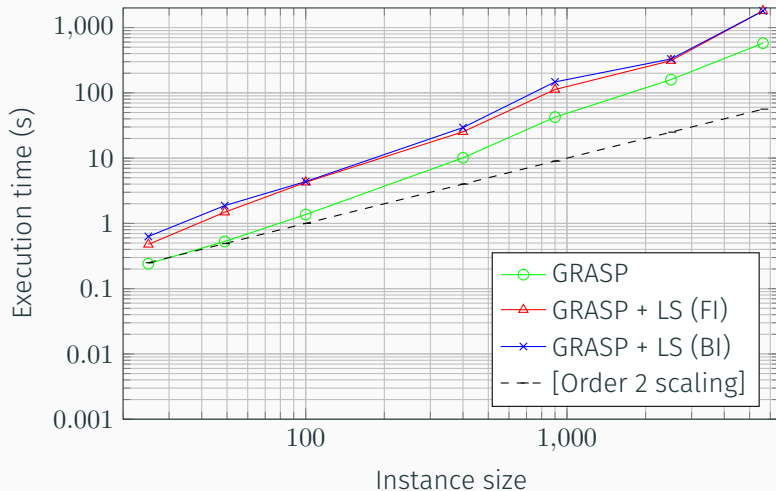


Figure 9: Execution time as a function of the instance size for GRASP.

- **Objective function** : FI = BI for all instances, simple GRASP stays close (especially for big instances)
- **Execution time** : simple GRASP is significantly faster, FI is close second above BI.

Choice: As with the Greedy case, FI outperforms BI. However, here the computation time is no longer negligible: simple GRASP is 2–3× faster, with only a small loss in accuracy.

⇒ **GRASP + LS (FI)** (simple GRASP is also a valid option).

- ILP (CPLEX)
- Greedy
 - Greedy constructive (simple)
 - Greedy constructive + Local search with first improvement policy
 - Greedy constructive + Local search with best improvement policy
- GRASP
 - GRASP constructive (simple)
 - GRASP constructive + Local search with first improvement policy
 - GRASP constructive + Local search with best improvement policy

Final comparison

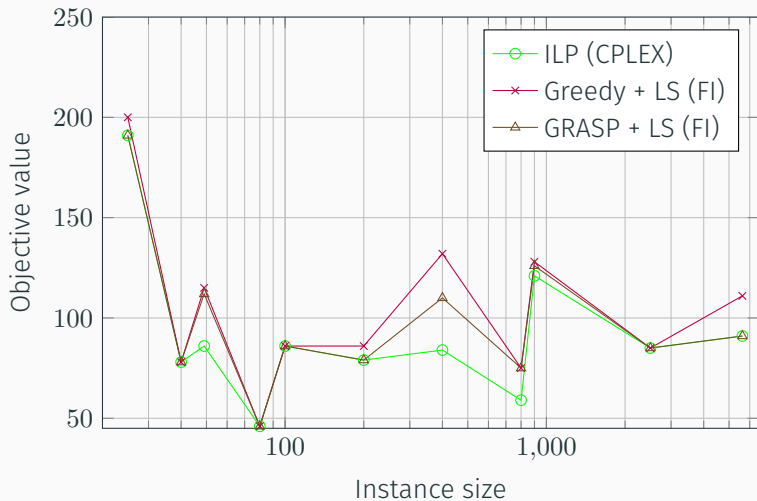


Figure 10: Objective value as a function of the instance size.

Final comparison

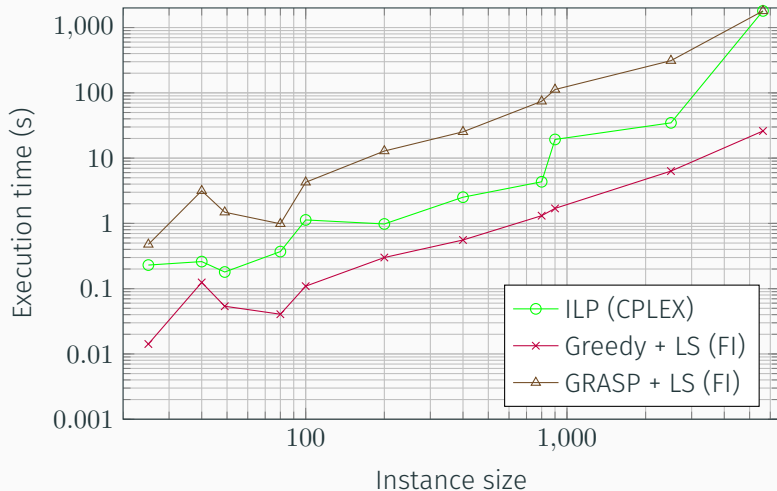


Figure 11: Execution time as a function of the instance size

Rankings :

- **Objective function** : ILP > GRASP > Greedy
- **Execution time** : Greedy > ILP > GRASP

Choice : ILP outperforms GRASP in all aspects, therefore GRASP is not viable.

Greedy is $5\text{--}10\times$ faster than ILP, scales stably, and maintains acceptable accuracy even on large instances.

ILP execution time increases sharply beyond size 50×50 and is computationally expensive (high CPU usage).

⇒ **Greedy + LS (FI)** : Suitable when computing resources are limited or when dealing with very large instances (particularly with many crossings).

⇒ **ILP** : If computing resources are not an issue and instances are moderate; best for long-term planning where operational savings justify longer solving times.

Thank you for listening !



Figure 12: Batman on his way to place all the cameras