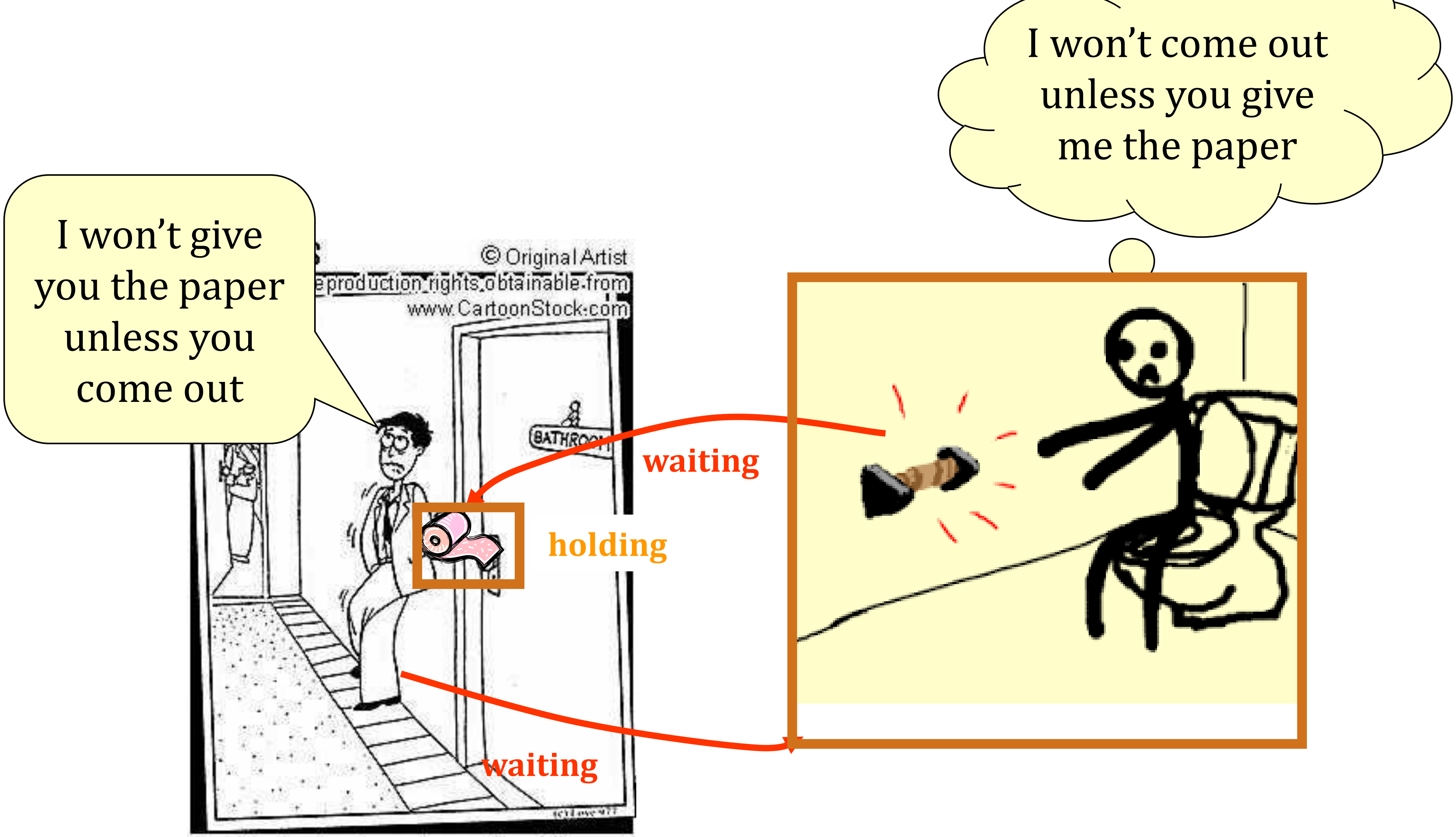
Lecture 7: Deadlock

Bo Tang @ 2021, Spring

Deadlock


I won't come out unless you give me the paper
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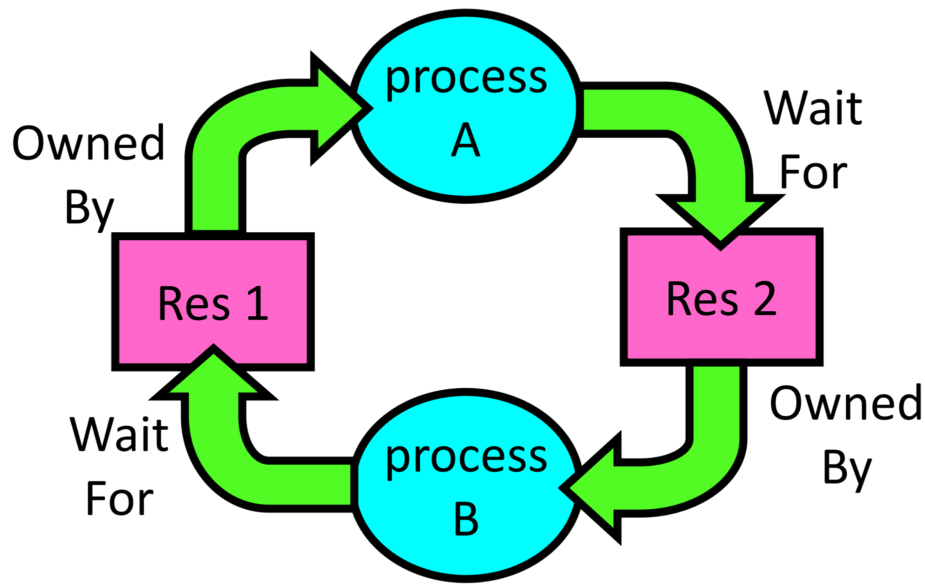
waiting
holding
waiting

Starvation vs Deadlock

- Starvation vs. Deadlock
  - Starvation: thread waits indefinitely
    - Low-priority thread waiting for resources constantly in use by high-priority threads
  - Deadlock: circular waiting for resources
    - Thread A owns Res 1 and is waiting for Res 2
      Thread B owns Res 2 and is waiting for Res 1



  - Deadlock $\Rightarrow$ Starvation but not vice versa
    - Starvation can end (but does not have to)
    - Deadlock canot end without external intervention

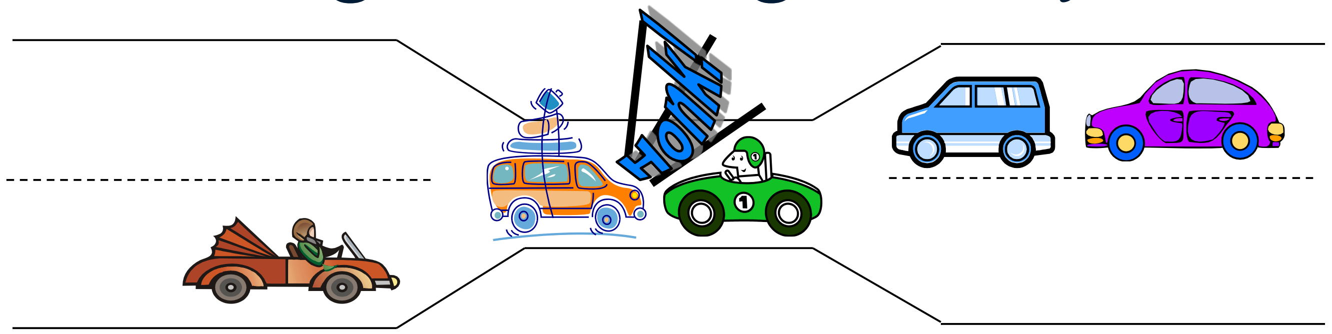
Conditions for Deadlock

- Deadlock not always deterministic

```
process A                process B
sem_wait(x);             sem_wait(y);
sem_wait(y);             sem_wait(x);
sem_post(y);             sem_wait(x);
sem_post(x);             sem_wait(y);
```

  - Deadlock will not always happen with this code
    - Have to have exactly the right timing
    - So you release a piece of software, and you tested it, and there it is, controlling a nuclear power plant…
- Deadlocks occur with multiple resources
  - Means you cannot decompose the problem
  - Cannot solve deadlock for each resource independently
  - System with 2 disk drives and two threads
    - Each process needs 2 disk drives to function
    - Each process gets one disk and waits for another one

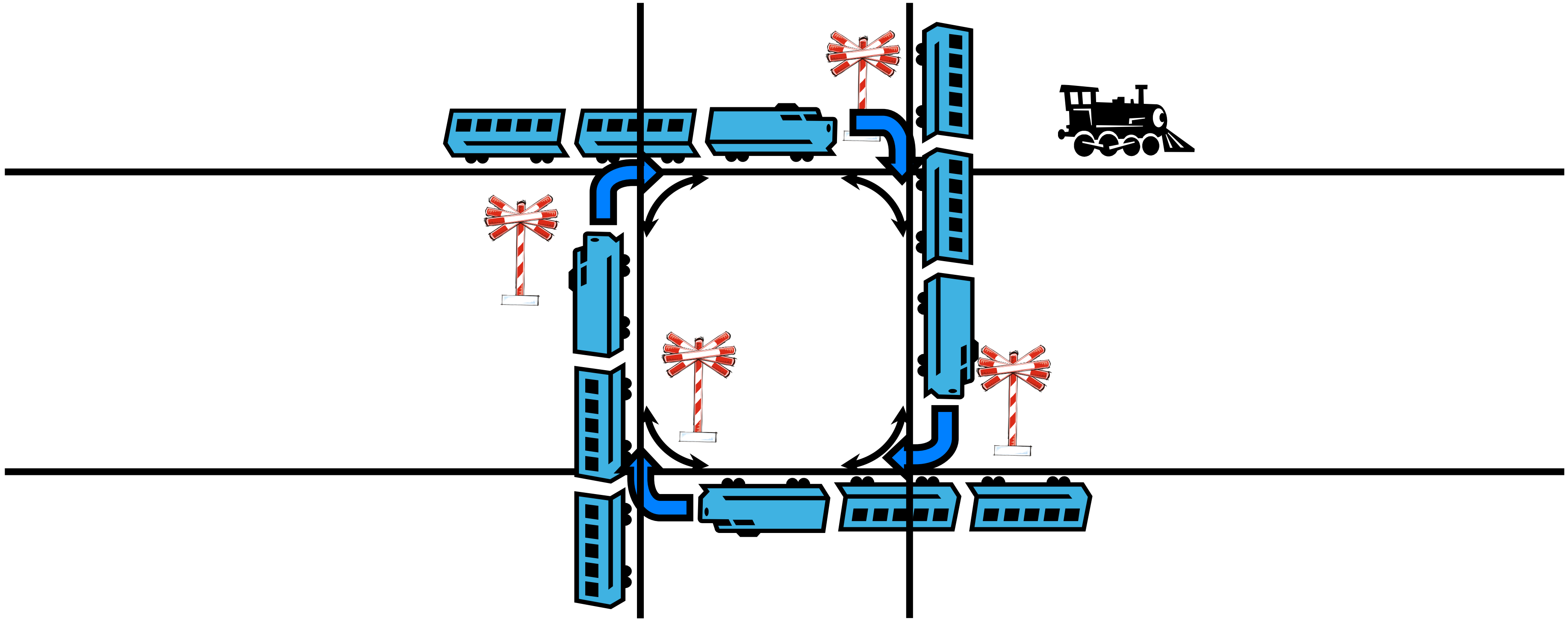
Bridge Crossing Example



- Each segment of road can be viewed as a resource
  - Car must own the segment under them
  - Must acquire segment that they are moving into
- For bridge: must acquire both halves
  - Traffic only in one direction at a time
  - Problem occurs when two cars in opposite directions on bridge: each acquires one segment and needs next
- If a deadlock occurs, it can be resolved if one car backs up (preempt resources and rollback)
  - Several cars may have to be backed up
- Starvation is possible
  - East-going traffic really fast $\Rightarrow$ no one goes west

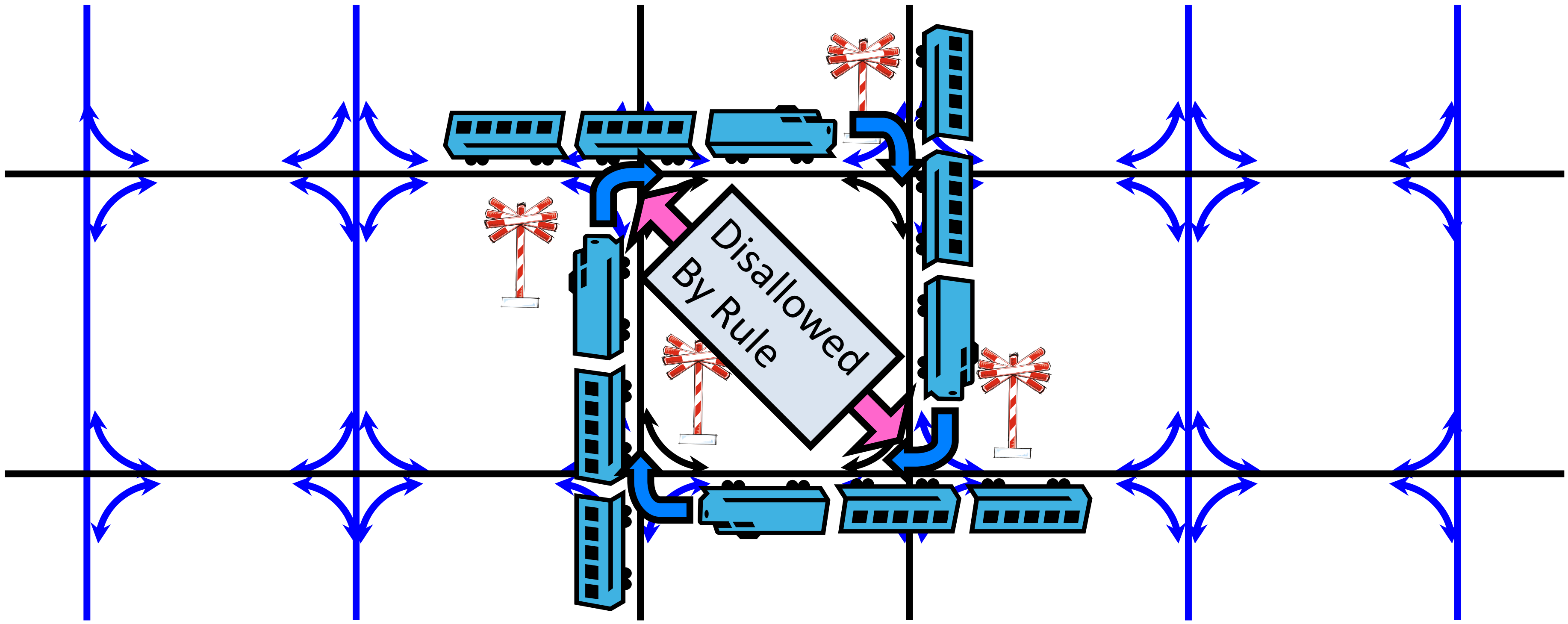
Train Example

- Circular dependency (Deadlock!)
  - Each train wants to turn right
  - Blocked by other trains
  - Similar problem to multiprocessor networks



Train Example

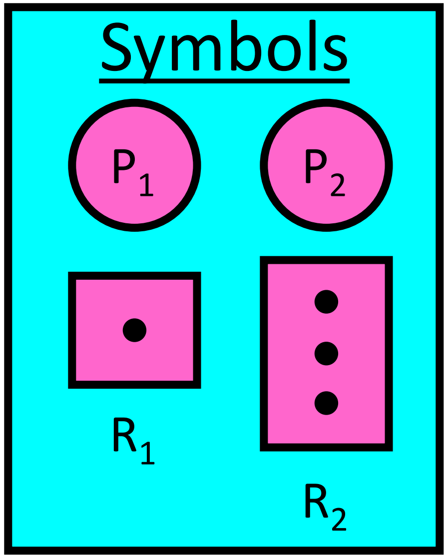
- Fix? Imagine grid extends in all four directions
  - Force ordering of channels (tracks)
    - Protocol: Always go east-west first, then north-south
  - Called “dimension ordering” (X then Y)



Four requirements for Deadlock

- Mutual exclusion
  - Only one thread at a time can use a resource.
- Hold and wait
  - Thread holding at least one resource is waiting to acquire additional resources held by other threads
- No preemption
  - Resources are released only voluntarily by the thread holding the resource, after thread is finished with it
- Circular wait
  - There exists a set $\{P_1, \ldots, P_n\}$ of waiting processes
    - P_1 is waiting for a resource that is held by P_2
    - P_2 is waiting for a resource that is held by P_3
    - …
    - P_n is waiting for a resource that is held by P_1

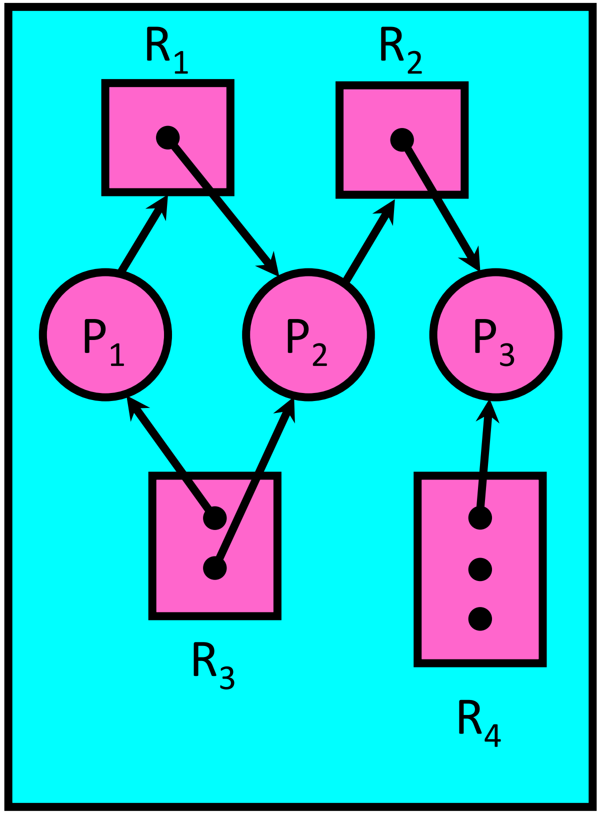
Resource-Allocation Graph



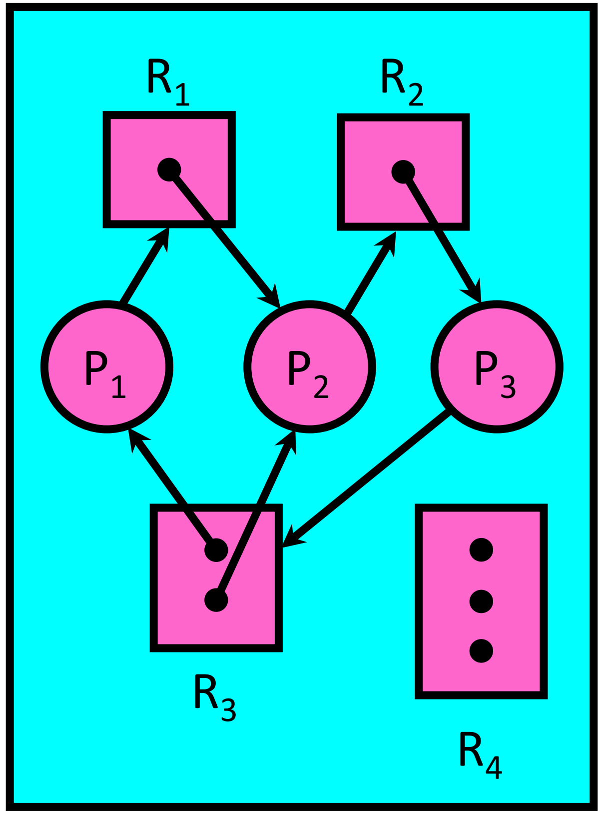
- System Model
  - A set of Threads $P_1, P_2, \ldots, P_n$
  - Resource types $R_1, R_2, \ldots, R_m$
    CPU cycles, memory space, I/O devices
  - Each resource type R_i has W_i instances
  - Each thread utilizes a resource as follows:
    - `Request() / Use() / Release()`
- Resource-Allocation Graph:
  - V is partitioned into two types:
    - $P = \{P_1, P_2, \ldots, P_n\}$, the set processes in the system.
    - $R = \{R_1, R_2, \ldots, R_m\}$, the set of resource types in system
  - request edge – directed edge $P_1 \rightarrow R_j$
  - assignment edge – directed edge $R_j \rightarrow P_i$

Resource Allocation Graph Examples

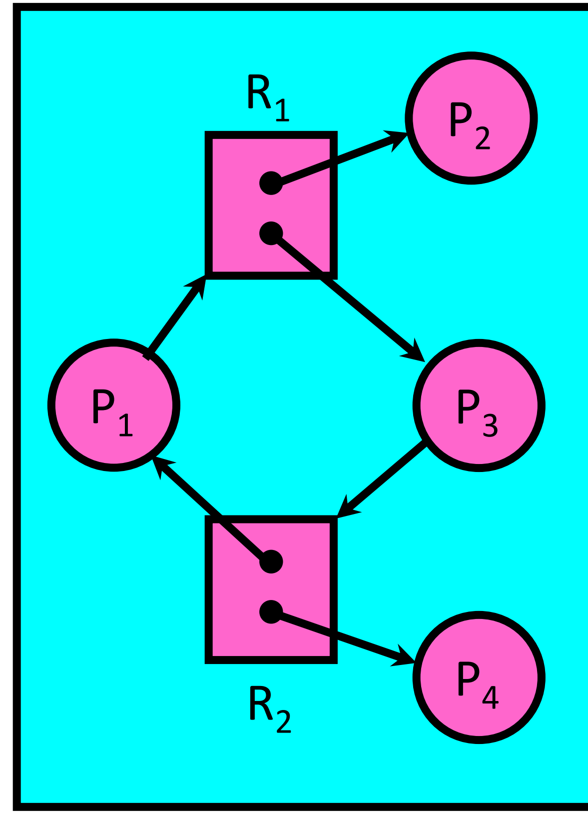
- Recall:
  - request edge – directed edge $P_1 \rightarrow R_j$
  - assignment edge – directed edge $R_j \rightarrow P_i$



Simple Resource Allocation Graph



Allocation Graph With Deadlock



Allocation Graph with Cycle, but No Deadlock

Methods for Handling Deadlocks

- Allow system to enter deadlock and then recover
  - Requires deadlock detection algorithm
  - Some technique for forcibly preempting resources and/or terminating tasks

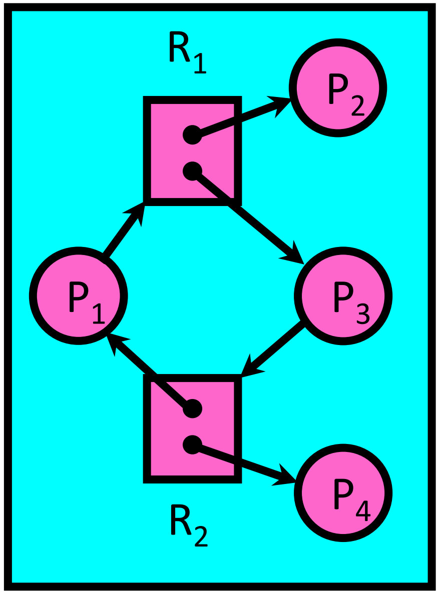
- Ensure that system will *never* enter a deadlock
  - Need to monitor all lock acquisitions
  - Selectively deny those that *might* lead to deadlock

- Ignore the problem and pretend that deadlocks never occur in the system
  - Used by most operating systems, including UNIX

Deadlock Detection Algorithm

- Only one of each type of resource ⇒ look for loops
- More general deadlock detection algorithm
  - Let [X] represent an m-ary vector of non-negative integers (quantities of resources of each type):
    - `[FreeResources]`: current free resources each type
    - `[Request`$_X$`]`: current requests from process X
    - `[Alloc`$_X$`]`: current resources held by process X
  - See if tasks can eventually terminate on their own

```
[Avail] = [FreeResources]
Add all nodes to UNFINISHED
do {
    done = true
    Foreach node in UNFINISHED {
        if ([Request_node] <= [Avail]) {
            remove node from UNFINISHED
            [Avail] = [Avail] + [Alloc_node]
            done = false
        }
    }
} until(done)
```



  - Nodes left in `UNFINISHED` ⇒ deadlocked

What to do when detect deadlock?

- Terminate thread, force it to give up resources
  - In Bridge example, Godzilla picks up a car, throw it into the river. Deadlock solved!
  - Shoot a dining philosopher
  - But, not always possible – killing a process holding a mutex leaves world inconsistent
- Preempt resources without killing off process
  - Take away resources from process temporarily
  - Does not always fit with semantics of computation
- Roll back actions of deadlocked threads
  - For bridge example, make one car roll backwards (may require others behind him)
  - Common technique in databases (transactions)
  - Of course, if you restart in exactly the same way, may reenter deadlock once again
- Many operating systems use other options

Techniques for Preventing Deadlock

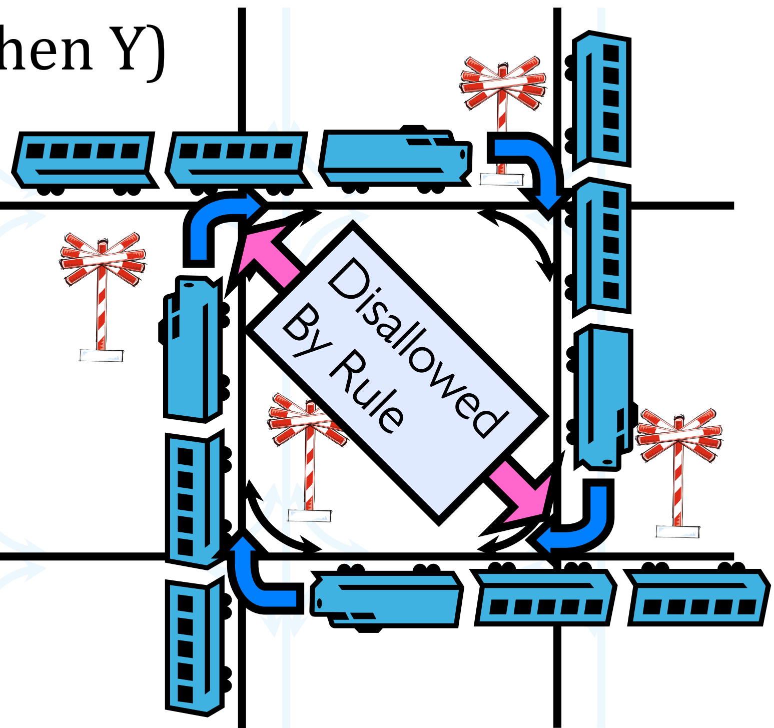
- Infinite resources
  - Include enough resources so that no one ever runs out of resources. Examples:
    - Bay bridge with 12,000 lanes. Never wait!
    - Infinite disk space (not realistic yet?)
- No sharing of resources (totally independent threads)
  - Not very realistic
- Do not allow waiting
  - Technique used in Ethernet/some multiprocessor nets
    - Everyone speaks at once. On collision, back off and retry
  - Inefficient, since have to keep retrying
    - Consider: driving to SUSTech; when hit traffic jam, suddenly you are transported back home and told to retry!

Techniques for Preventing Deadlock

- Make all processes request everything they will need at the beginning.
  - Problem: Predicting future is hard, tend to over-estimate resources. Example:
    - If need 2 chopsticks, request both at same time
    - Do not leave home until we know no one is using any intersection between home and SUSTech; only one car on the Bay Bridge at a time
- Force all processes to request resources in a particular order preventing any cyclic use of resources
  - Thus, preventing deadlock
  - Example (x.P, y.P, z.P,…)
    - Make tasks request disk, then memory, then…
    - Keep from deadlock on freeways around SF by requiring everyone to go clockwise

Review: Train Example

- Circular dependency (Deadlock!)
  - Each train wants to turn right, blocked by other trains
  - Similar problem to multiprocessor networks
- Fix? Imagine grid extends in all four directions
  - Force ordering of channels (tracks)
    - Protocol: Always go east-west first, then north-south
  - Called "dimension ordering" (X then Y)



Banker's Algorithm for Preventing Deadlock

- Toward right idea:
  - State maximum (max) resource needs in advance
  - Allow particular process to proceed if:
    - (available resources - #requested) ≥ max

    Remaining that might be needed by any process

- Banker's algorithm:
  - Allocate resources dynamically
    - Evaluate each request and grant if some ordering of processes is still deadlock free afterward
    - Technique: pretend each request is granted, then run deadlock detection algorithm, substituting
      ([Maxnode]-[Allocnode] ≤ [Avail]) for ([Requestnode] ≤ [Avail])
      Grant request if result is deadlock free (conservative!)

(a) [8 pts] Consider the following snapshot of a system

	Allocation				Max				Available			
	A	B	C	D	A	B	C	D	A	B	C	D
P0	0	0	1	2	0	0	1	2	1	5	2	0
P1	1	0	0	0	1	7	5	0				
P2	1	3	5	4	2	3	5	6				
P3	0	0	1	4	0	6	5	6				

Answer the following questions using the banker's algorithm:

1. What is the content of the matrix **Need?**
2. Is the system in a safe state? If yes, please give a safe sequence.
3. If a request from process P1 arrives for (0,4,2,0), can the request be granted immediately? Please explain your answer.

Banker's Algorithm for Preventing Deadlock

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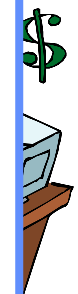
[Avail] = [FreeResources]
Add all nodes to UNFINISHED
do {
      done = true
      Foreach node in UNFINISHED {
         if ([Request_node] <= [Avail]) {
            remove node from UNFINISHED
            [Avail] = [Avail] + [Alloc_node]
            done = false
         }
      }
} until(done)
```



- Allocate resources dynamically
  - Evaluate each request and grant if some ordering of threads is still deadlock free afterward
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Banker's Algorithm for Preventing Deadlock

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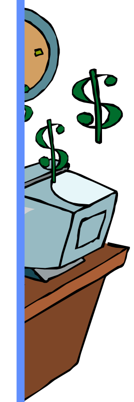

ordering of threads is still deadlock free afterward

```
[Avail] = [FreeResources]
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         }
      }
} until(done)
```

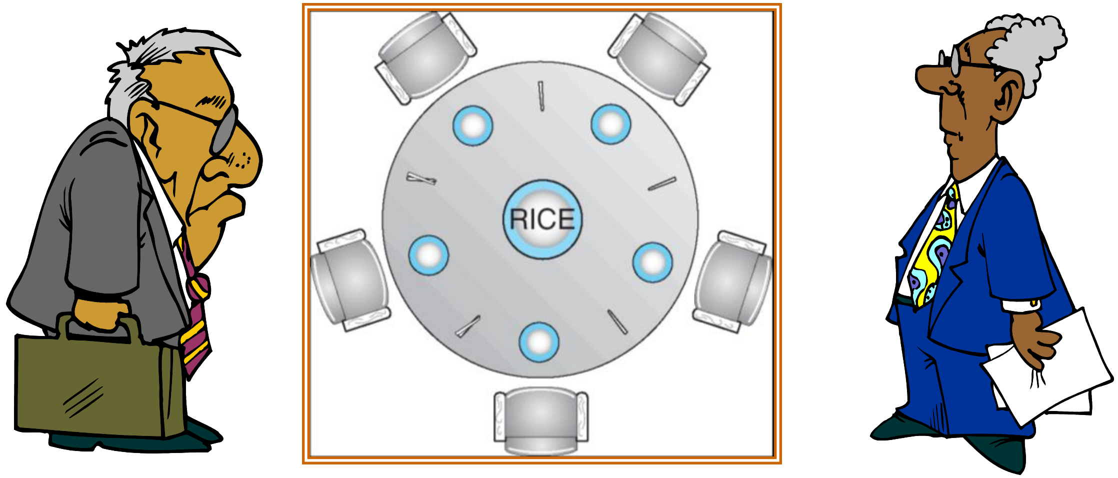
Banker's Algorithm for Preventing Deadlock

- Toward right idea:
  - State maximum resource needs in advance
  - Allow particular process to proceed if:

    (available resources - #requested) ≥ max
    remaining that might be needed by any process

- Banker's algorithm (less conservative):
  - Allocate resources dynamically
    - Evaluate each request and grant if some ordering of threads is still deadlock free afterward
    - Technique: pretend each request is granted, then run deadlock detection algorithm, substituting
      ($[\text{Max}_{\text{node}}]-[\text{Alloc}_{\text{node}}] \leq [\text{Avail}]$) for ($[\text{Request}_{\text{node}}] \leq [\text{Avail}]$)
      Grant request if result is deadlock free (conservative!)
    - Keeps system in a "SAFE" state, i.e. there exists a sequence $\{P_1, P_2, \ldots P_n\}$ with P_1 requesting all remaining resources, finishing, then P_2 requesting all remaining resources, etc..
  - Algorithm allows the sum of maximum resource needs of all current threads to be greater than total resources

Banker's Algorithm Example



- Banker's algorithm with dining philosophers
  - "Safe" (will not cause deadlock) if when try to grab chopstick either:
    - Not last chopstick
    - Is last chopstick but someone will have two afterwards
  - What if k-handed philosophers? Do not allow if:
    - It is the last one, no one would have k
    - It is 2nd to last, and no one would have k-1
    - It is 3rd to last, and no one would have k-2
    - ...

Thank You!

14. In a single-processor multi-process system, if there are several ready processes, then the wrong statement about process scheduling is ().

A. process scheduling can execute when a process terminates
B. process scheduling can execute when a process creates
C. process scheduling can execute when a process is in critical section
D. process scheduling can execute when system call finished and system returns to user space