

Problem:

Write a method **double det(double matrix[][])** that takes as its argument a 2D array representing a square matrix and returns its determinant recursively computed as follows:

- Determinant of a 1-by-1 matrix is its only element;
- Determinant of a larger matrix **A** of size **int n** **det(A)** is the sum over all rows of the expressions **A[k][0] * det(B) * (-1)^k**, where **B** is a smaller matrix called a minor and made up of all elements of **A** not belonging to the row **k** and column **0** (**k = 0, 1, 2, ..., n-1**). For example, if **n = 5** and **k = 3**, then **B** consists of the shaded elements

	0	1	2	3	4
0					
1					
2					
3					
4					

Hint:

In other words, expresses the determinant as a weighted sum, where each element of the chosen row is multiplied by the determinant of a smaller matrix (*minor*) formed by deleting the row and column corresponding to that element, plus an additional cofactor which alternates between -1 and 1.

For example:

$$\begin{vmatrix} 1 & 2 & 3 \\ 4 & 5 & 6 \\ 7 & 8 & 9 \end{vmatrix} = 1 \cdot \begin{vmatrix} 5 & 6 \\ 8 & 9 \end{vmatrix} - 4 \cdot \begin{vmatrix} 2 & 3 \\ 7 & 8 \end{vmatrix} + 7 \cdot \begin{vmatrix} 2 & 3 \\ 4 & 5 \end{vmatrix}$$

$$= 1 \cdot (5 \cdot 9 - 8 \cdot 6) - 4 \cdot (2 \cdot 9 - 8 \cdot 3) + 7 \cdot (2 \cdot 6 - 5 \cdot 3) = -3 + 24 - 21 = 0$$

The base case is just that the determinant of a single matrix is the value of its single element. The method above implements the additional case when $n=2$. This is just a special case of the above formula and is probably intended as an optimization.

A *minor* of an element in a matrix is the determinant of the matrix that is formed by removing the elements in the same row and column of the element. The minor of a 4 x 4 matrix is the determinant of a 3 x 3 matrix.

P.S.

I would not recommend using this method to evaluate determinants unless the matrices are all very small because it's extremely slow for large matrices - given an $n \times n$ matrix, it takes $O(n!)$ time. This comes from the fact that each method call must evaluate n determinants of size $(n-1)$, so the total number of recursive calls is $n \cdot (n-1) \cdot (n-2) \cdot \dots \cdot 3 \cdot 2 \cdot 1$.