

## Chongqing Univer<u>c</u>ity Standard Code Library Extended By st.Krwlng

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```
#define square(x) (x)*(x)
#define getvec(x, y) ((y)-(x))
#define xmu1(x1, y1, x2, y2) ((x1)*(y2)-(x2)*(y1))
struct line { double a, b, c; }; //ax + by = c
struct pnt {
      double x, y;
      pnt operator+(const pnt &p) const{
             pnt ret; ret.x = x + p.x;
             ret.y = y + p.y; return ret;
      pnt operator-(const pnt &p) const{
             pnt ret; ret.x = x - p.x;
             ret.y = y - p.y; return ret;
      pnt operator*(const double &c) const{
             pnt ret; ret.x = x * c;
             ret.y = y * c; return ret;
      pnt operator/(const double &c) const{
             pnt ret; ret.x = x / c;
             ret.y = y / c; return ret;
      bool operator<(const pnt &p) const{</pre>
             if (fabs(x-p.x) < eps) return y-p.y < -eps;
             return x-p, x < -eps:
};
typedef pnt vec;
const double eps = 1e-6; //1e-8, 1e-10, 1e-16
const double pi = acos(-1.0);
double nummu1(const vec &v1, const vec &v2) { return v1.x * v2.x + v1.y * v2.y; }
double \ submul(const \ vec \ \&v1, \ const \ vec \ \&v2) \ \{ \ return \ xmul(v1. \ x, \ v1. \ y, \ v2. \ x, \ v2. \ y) \ ; \ \}
int lessver(const pnt &p1, const pnt &p2) {
      if (fabs(p1.y-p2.y) <eps) return p1.x-p2.x < -eps;
      return p1. y-p2. y < -eps;
}
double getdis(const pnt &p1, const pnt &p2) { return sqrt(square(p2.x-p1.x) + square(p2.y-p1.y)); }
double getdis(const line &1, const pnt &p) { return fabs(1.a * p.x + 1.b * p.y - 1.c) / sqrt(square(1.a) + square(1.b)); }
double oridis(const line &1, const pnt &p) { return 1.a * p.x + 1.b * p.y - 1.c; }
pnt getcrs(const line &11, const line &12) {
      pnt ret; ret.x = xmul(11.b, 11.c, 12.b, 12.c) / xmul(11.b, 11.a, 12.b, 12.a);
      ret.y = xmul(11.a, 11.c, 12.a, 12.c) / xmul(11.a, 11.b, 12.a, 12.b); return ret;
pnt rotate(const pnt &p, double ang) {
      vec v = { sin(ang), cos(ang) };
      pnt ret = { submul(p, v), nummul(p, v) }; return ret;
pnt stretch_rotate(const pnt &p, const vec &v1, const vec &v2) {
      \mbox{vec } v = \{ \mbox{ submul}(v1, \mbox{ } v2) / \mbox{nummul}(v1, \mbox{ } v1), \mbox{ nummul}(v1, \mbox{ } v2) / \mbox{nummul}(v1, \mbox{ } v1) \};
      pnt ret = { submul(p, v), nummul(p, v) }; return ret;
} // v1 \rightarrow v2, len(v1) \rightarrow len(v2).
vec uvec (const vec &v) {
      double len = sqrt(nummul(v, v));
       vec ret; ret.x = v.x / len; ret.y = v.y / len; return ret;
line getline(const pnt &p1, const pnt &p2) {
      vec v = getvec(p1, p2); v = uvec(v);
      line ret; ret.a = -v.y; ret.b = v.x;
      ret.c = ret.a * p1.x + ret.b * p1.y; return ret;
} // error if p1 == p2 \,
```

```
line getline(const pnt &p, const vec &dv) {
    vec v = uvec(v);
    line ret; ret.a = -v.y; ret.b = v.x;
    ret.c = ret.a * p.x + ret.b * p.y; return ret;
}
line getmidver(const pnt &pl, const pnt &p2) {
    vec v = getvec(pl, p2); v = v/2; pnt mid = pl + v; v = uvec(v);
    line ret; ret.a = v.x; ret.b = v.y; ret.c = ret.a*mid.x + ret.b*mid.y; return ret;
}

//it should be garanteed that p[i] < q[i]
bool checkers(pnt* p, pnt* q, line* l, int i, int j, pnt &crs) {
    if (q[i] < p[j] || q[j] < p[i]) return false;
    if (submul(getvec(p[i], q[i]), getvec(p[i], p[j])) * submul(getvec(p[i], q[i]), getvec(p[i], q[i])) > eps) return false;
    if (submul(getvec(p[j], q[j]), getvec(p[j], q[j])) * submul(getvec(p[j], q[j]), getvec(p[j], q[i])) > eps) return false;
    if (fabs(submul(getvec(p[i], q[i]), getvec(p[j], q[j]))) < eps) if (q[i] < q[j]) crs = q[i]; else crs = q[j];
    else crs = getcrs(l[i], l[j]); return true;
}</pre>
```

```
//get the line which is the reflection of 11 through 10
//Caution: the line of 10 and 11 must be unified in construction, or it'll have some precision problem!
line lineref(const line &10, const line &11)
{
      vec v; v.x = 10.b; v.y = -10.a; double len = sqrt(square(10.a) + square(10.b));
      double sa = v.y / len, ca = v.x / len;
      double a = 11.a * square(ca) - 11.a * square(sa) + 2 * 11.b * sa * ca;
      double b = 11.b * square(sa) - 11.b * square(ca) + <math>2 * 11.a * sa * ca;
      line ret; ret. a = a; ret. b = b; ret. c = 11.c + (a - 11.a) * p. x + (b - 11.b) * p. y;
//{
m get} the number of intersection of a circle and a line, as well as the point of intersection
int linexcircle(const line &1, const pnt &cen, double rad, pnt *p)
      double ver, hor: vec v1, v2:
      v1 = (vec) \{ 1.a, 1.b \}; v1 = uvec(v1); v2 = (vec) \{ -v1.y, v1.x \};
      ver = oridis(1, cen); if (fabs(ver)-rad > eps) return 0;
      if (fabs(fabs(ver)-rad) \leq p[0] = cen - v1 * ver; return 1; }
     hor = sqrt(square(rad)-square(ver));
      p[0] = cen - v1*ver + v2*hor; p[1] = cen - v1*ver - v2*hor; return 2;
//get the number of intersection of a circle and a segment, as well as the point of intersection
//If there exists two such points, the first one has the less value of dmul(p1-\ranglep2, p1-\ranglecrs).
void segxcircle(const pnt &p1, const pnt &p2, const pnt &cen, double rad, int &top, pnt *p) {
      vec v1, v2; pnt tp; double ver, hor;
      v1 = uvec(p2-p1); v2 = (vec) \{ -v1. y, v1. x \};
      \label{eq:ver} \mbox{ver = nummul(cen, v2) - nummul(p1, v2); if (fabs(ver)-rad > eps) return;}
     hor = sqrt(square(rad)-square(ver));
      tp = cen - v2 * ver - v1 * hor; if (nummul(tp-p1, tp-p2) < eps) p[top++] = tp;
      //check if the circle is totaly in the given convex.
bool circleincvx(int n, pnt *p, const pnt &cen, double rad) {
      for (int i = 0; i < n; ++i) if (submul(p[i]-cen, p[i+1]-cen)<-eps) return false;
       for \ (int \ i = {\color{red}0}; \ i \ < \ n; \ +\! +\! i) \ if \ (fabs(oridis(getline(p[i], \ p[i+1]), \ cen)) - rad < -eps) \ return \ false; 
      return true:
//check if the point is in the circle or on the border of the circle
bool pntincircle(const pnt &p, const pnt &cen, double rad) { return getdis(p, cen)-rad < eps; }
//get the cross points of two circle.
int circlecrs(const pnt &p1, const pnt &p2, double r1, double r2, pnt *p) {
      double len = getdis(p1, p2); vec v = uvec(p2-p1);
      if (len + r1 - r2 < -eps | | len + r2 - r1 < -eps | | r1 + r2 - len < -eps | return 0;
      if (fabs(len+r1-r2) < eps) { p[0] = p1 - v * r1; return 1; }
      if (fabs(len+r2-r1) < eps | | fabs(r1+r2-len) < eps) { p[0] = p1 + v * r1; return 1; }
      double \ ang1 = atan2(v.y, \ v.x), \ ang2 = acos((square(len) + square(r1) - square(r2)) / (2*len*r1));
      v = (vec) \{ cos(ang1+ang2), sin(ang1+ang2) \}; p[0] = p1 + v * r1;
      v = (vec) \{ cos(ang1-ang2), sin(ang1-ang2) \}; p[1] = p1 + v * r1;
      return 2;
```

 $/\!/ in this section, the operator + (pnt), -(pnt), *(double), /(double) should be overloaded in definition$ 

```
//\mathrm{get} the intersection of a circle and a convex, return the number of the points.
//description: the array evt reflects the condition of the same postion at q. If evt[i] equals 1, it means the part of ret from q[i]
to q[i+1] is an arc, otherwise it's a segment.
//caution: the length of arc exists in the ret must be larger than 0 (means the condition that the ret is a entire circle can be expressed
valid), but the degeneration should take extra O(n) time at last four lines of code.
int circlexcvx(int n, pnt *p, const pnt &cen, double rad, pnt *q, bool *evt) {
      int i, j, k, l, ret; pnt tp[4];
      if (circleincvx(n, p, cen, rad)) { q[0] = q[1] = (pnt) \{ cen. x + rad, cen. y \}; evt[0] = 1; return 1; }
      for (ret = i = 0; i < n; ++i) {
             tp[0] = p[i]; 1 = 1; segxcircle(p[i], p[i+1], cen, rad, 1, tp);
             for (k = j = 1; j < 1; ++j) if (tp[j] != p[i] \&\& tp[j] != p[i+1]) tp[k++] = tp[j];
              for \ (j = 0; \ j < k; \ ++j) \ if \ (incircle(tp[j], \ cen, \ rad)) \ \{ \ evt[ret] = 0; \ q[ret++] = tp[j]; \ \} 
             if (!incircle(p[i+1], cen, rad) \&\& incircle(tp[k-1], cen, rad)) evt[ret-1] = 1;
      } q[ret] = q[0]; evt[ret] = evt[0]; 1 = ret; if (!ret) return 0;
      for (ret = i = 0; i < 1; ++i) {
             if (!evt[i] \mid | q[i] != q[i+1]) \{ q[ret] = q[i]; evt[ret] = evt[i]; ret++; \}
      } if (!ret) { ret = 1; evt[0] = 0; }
      q[ret] = q[0]; evt[ret] = evt[0]; return ret;
//get the judgement whether ang 3 is between the arc, all the ang has the constrains of [-pi, pi)
bool arcbetween(double ang1, double ang2, bool ati, double ang3)
      if (fabs(ang1-ang3) \langle eps || fabs(ang2-ang3) \langle eps) return true;
      return (ati ^ ((ang1-ang3) * (ang2 - ang3) < -eps));</pre>
}
//get the enclosing circle of the triangle by the given 3 points of the vertices of triangle. return the radius of the circle and
the center of the circle will be assigned to the last paremeter.
double outcircle(const pnt &p1, const pnt &p2, const pnt &p3, pnt &cen)
{
      line 11 = getmidver(p1, p2), 12 = getmidver(p2, p3);
      cen = getcrs(11, 12); return getdis(cen, p1);
//get the internal tangent circle of the triangle. return the radius of the circle and the center of the circle will be assigned
to the last paremeter.
double incircle(const pnt &p1, const pnt &p2, const pnt &p3, pnt &cen)
      vec v1, v2: line 11, 12:
      v1 = uvec(p2-p1); v2 = uvec(p3-p1); 11 = getline(p1, v1+v2);
      v1 = uvec(p3-p2); v2 = uvec(p1-p2); 12 = getline(p2, v1+v2);
      cen = getcrs(11, 12); return fabs(submul(v2, cen-p2));
//get the common point of three verticle line of a triangle.
pnt orthocenter(const pnt &p1, const pnt &p2, const pnt &p3)
      vec v1, v2, ret; line 11, 12;
      v1 = uvec(p3-p2); 11 = (1ine) \{ v1.x, v1.y, nummu1(v1, p1) \};
      v2 = uvec(p3-p1); 12 = (1ine) \{ v2. x, v2. y, nummu1(v2, p2) \};
      ret = getcrs(11, 12): return ret:
//get the barycenter of a polygon given in clockwise
pnt barycenter(int n, pnt *p)
      int i; double tot, ta; pnt ret;
      p[n] = p[0]: ret = (pnt) { 0, 0}:
       for (tot = i = 0; i < n; ++i) {
             ta = submul(p[i], p[i+1]);
             tot += ta; ret. x += (p[i].x+p[i+1].x)*ta/3; ret. y += (p[i].y+p[i+1].y)*ta/3;
      } ret.x /= tot; ret.y /= tot; return ret;
```

```
//{\rm get} the fermat point which has the least total distance to the given n points.
\#define NUM 2 //C = (NUM+1)^2
#define BD 2 //C = (BD*2+1)^2
const double ATT = 0.9, EPS = 1e-6; //when ADd = 0.9, C = log10(range/EPS) * 10
double valuate(int n, pnt *p, const pnt &tp) \{
       double ret; int i;
       for (ret = i = 0; i < n; ++i) ret += getdis(tp, p[i]);
       return ret;
pnt fermatpoint(int n, pnt* p)
       double 1ft, rit, up, dn, stx, sty, dx, dy, 1st, val, tv, ttv;
       pnt ret, vp, tp, ttp; int i, j, k, 1;
       lft = dn = inf: rit = up = -inf:
       for (i = 0; i < n; ++i) {
              1ft = getmin(1ft, p[i].x); rit = getmax(rit, p[i].x);
               dn = getmin(dn, p[i].y); up = getmax(up, p[i].y);
       } stx = rit-lft; sty = up-dn; lst = inf;
       for (i = 0; i \le NUM; ++i) for (j = 0; j \le NUM; ++j) {
               vp = (pnt) { lft+stx/NUM*i, dn+sty/NUM*j }; val = valuate(n, p, vp);
               for (dx = stx/NUM/BD, dy = sty/NUM/BD; dx > EPS || dy > EPS; dx *= ATT, dy *= ATT) {
                       for (tp = vp, tv = val, k = -BD; k \le BD; ++k) for (1 = -BD; 1 \le BD; ++1) {
                              ttp = (pnt) \{vp. x+dx*k, vp. y+dy*1\}; ttv = valuate(n, p, ttp);
                              if (ttv < tv) { tv = ttv; tp = ttp; }</pre>
                      } vp = tp; val = tv;
              } if (val < lst) { lst = val; ret = vp; }
       } return vp;
//check whether the given point is in the polygon, no matter the order of points about polygon is clockwise or conter-clockwise
bool pntinpoly(int n, pnt *p, const pnt &p1) {
       int i, cnt; pnt p2 = {-inf, p1.y}; p[n] = p[0];
       for (cnt = i = 0; i < n; ++i) {
                if \ (fabs(submul(p[i+1]-p1,\ p[i]-p1)) < eps \ \&\& \ nummul(p[i+1]-p1,\ p[i]-p1) < eps) \ return \ true; \\
               if (fabs(p[i+1], v-p[i], v)(eps) continue:
                if \ (fabs (p[i+1]. \ y-p1. \ y) \\ < eps \ \&\& \ p[i+1]. \ x-p1. \ x \\ < eps) \ \{ \ if \ (p[i+1]. \ y-p[i]. \ y \\ < eps) \ cnt \\ ++; \ continue; \ \} 
                \text{if } (fabs \, (p[i]. \, y-p1. \, y) \, < \, eps \, \&\& \, \, p[i]. \, x-p1. \, x \, < \, eps) \, \, \{ \, \, \text{if } \, (p[i]. \, y-p[i+1]. \, y < eps) \, \, cnt++; \, \, continue; \, \} \\
                \label{eq:condition}  \mbox{if } (\mbox{submul} (p2-p1, \ p[i]-p1) *\mbox{submul} (p2-p1, \ p[i+1]-p1) < \mbox{eps \&\& } 
                       submul\,(p[i+\!\!1]-p[i],\ p1-p[i])*submul\,(p[i+\!\!1]-p[i],\ p2-p[i]) < eps)\ cnt++;
       } return cnt % 2;
//get the distance of nearest pair of points.
//the ret can be saved in the global variable, in order to record the index of two points.
int seq[N], que[N];
int cmpx(const int &i, const int &j) {
       \label{eq:continuous} \mbox{if } (fabs\,(p[i].\,x-p[j].\,x)\ens) \mbox{ } \mbox{return } p[i].\,y-p[j].\,y \ensuremath{\,{}^{\checkmark}} -eps;
       return p[i].x-p[j].x < -eps;</pre>
int cmpv(const int &i, const int &i) {
       if (fabs(p[i].y-p[j].y) < eps) return p[i].x-p[j].x < -eps;
       return p[i].y-p[j].y < -eps;</pre>
double lstpair(pnt *p, int lft, int rit) {
       int i, j, cnt, mid; double ret;
       if (1ft >= rit) return inf:
        \label{eq:mid} \mbox{mid} = (\mbox{lft+rit}) \ / \ \mbox{2}; \ \mbox{ret} = \mbox{getmin}(\mbox{lstpair}(\mbox{p}, \mbox{ lft}, \mbox{mid}), \mbox{ lstpair}(\mbox{p}, \mbox{mid+1}, \mbox{ rit})); 
        for \ (cnt = 0, \ i = lft; \ i <= rit; \ ++i) \ if \ (fabs(p[seq[i]]. \ x-p[seq[mid]]. \ x)-ret < eps) \ que[cnt++] = seq[i]; 
       sort(que, que+cnt, cmpy);
       for (i = 0; i < cnt; ++i) for (j = i+1; j < cnt && j < i+8; ++j) ret = getmin(ret, getdis(p[que[i]], p[que[j]]));
       return ret;
double nrstpair(int n, pnt *p) {
       for (int i = 0; i < n; ++i) seq[i] = i; sort(seq, seq+n, cmpx);
       return lstpair(p, 0, n-1);
```

```
//calculate the convex of the combine of several half plane, output the number of distinct point at the border, as well as each point
and line in counter-clockwise order.
//the answer of lines is in the seq array, the endpoint of line seq[i] is p[i] and p[i+1].
//{\it Caution}: the line here must be unified before calculation.
 \\ 1 \text{ line } bdr1 = \{ \text{ 1.0, 0.0, -inf } \}, \ bdr2 = \{ \text{ -1.0, 0.0, -inf } \}, \ bdr3 = \{ \text{ 0.0, 1.0, -inf } \}, \ bdr4 = \{ \text{ 0.0, -1.0, -inf } \}; \\ \\ \\ \end{aligned} 
int hfpcmp(const int &x, const int &y) {
          if (fabs(ang[x]-ang[y]) <eps) return 1[x].c - 1[y].c > eps;
           return ang[x] - ang[y] < -eps;</pre>
int hfplane(int n, line *1, pnt *p, int *seq, double *ang)
{
           int i, j, top, bot, ret;
           1[n++] = bdr1; 1[n++] = bdr2; 1[n++] = bdr3; 1[n++] = bdr4;
           for (i = 0; i < n; ++i) seq[i] = i; std::sort(seq, seq+n, hfpcmp);
           for \ (i = j = 1; \ i < n; \ ++i) \ if \ (fabs(ang[seq[i]]-ang[seq[j-1]]) > eps) \ seq[j++] = seq[i]; \ n = j;
           for (i = 1, top = bot = 0; i < n; ++i) {
                      while (top - bot && oridis(1[seq[i]], p[top]) < -eps) --top;
                      p[top+1] = getcrs(1[seq[i]], \ 1[seq[top]]); \ seq[++top] = seq[i];
                      if (ang[seq[i]] > -eps) break;
           } p[bot] = (pnt) { -inf*2, -inf*2 };
           for (++i; i < n; ++i) {
                      if (oridis(1[seq[i]], p[bot]) > -eps) continue;
                      while (top - bot && oridis(1[seq[i]], p[top]) < -eps) --top;
                      if (top == bot) break; p[top+1] = getcrs(1[seq[i]], 1[seq[top]]); seq[++top] = seq[i];
                       \label{eq:while original} \mbox{while (original leaf top]], p[bot+1]) < -eps) ++bot; p[bot] = getcrs(1[seq[top]], 1[seq[bot]]); 
           } if (i < n) return 0;
            for (ret = 0, i = bot; i < top; ++i) if (fabs(p[i].x-p[i+1].x) > ps || fabs(p[i].y-p[i+1].y) > ps) \\  \{ (ret = 0, i = bot; i < top; ++i) || fabs(p[i].y-p[i+1].y) > ps || fabs(p[i].y-p
                      seq[ret] = seq[i]; p[ret] = p[i]; ret++;
           } if (!ret || fabs(p[i].x-p[0].x)>eps || fabs(p[i].y-p[0].y)>eps) {
                      seq[ret] = seq[i]; p[ret] = p[i]; ret++;
          } p[ret] = p[0]; return ret;
//graham method to get convex
int tmp[N];
int vercmp(const int &x, const int &y) { return lessver(p[x], p[y]); }
int graham(int n, pnt* p, int *seq) {
           int i, top, bot;
           for (i = 0; i < n; ++i) tmp[i] = i; sort(tmp, tmp+n, vercmp);
           for (seq[top = bot = 0] = tmp[0], i = 1; i < n; ++i) {
                      seq[++top] = tmp[i];
           } bot = top;
           for (i = n-2; i \ge 0; --i) {
                      seq[++top] = tmp[i];
          } return top;
```

```
//melkman method to get convex
int tmp[N];
int melkman(int n, pnt *p, int *seq)
                    int i, j, bot, top;
                   for (i = top = 0; i < n; ++i) if (p[i] < p[top]) top = i;
                   for (j = top, i = 0; i < n; ++i, j = (j+1)%n) tmp[i] = j;
                   seq[n] = tmp[0]; seq[n-1] = tmp[1]; seq[n+1] = tmp[1];
                    for \ (i = 2; \ i < n; \ ++i) \ if \ (fabs(submul(getvec(p[seq[n]], \ p[seq[n-1]]), \ getvec(p[seq[n-1]], \ p[tmp[i]]))) < eps) \ \{ (i = 2; \ i < n; \ ++i) \ if \ (fabs(submul(getvec(p[seq[n-1]], \ p[seq[n-1]]), \ getvec(p[seq[n-1]], \ p[tmp[i]]))) < eps) \ \{ (i = 2; \ i < n; \ ++i) \ if \ (fabs(submul(getvec(p[seq[n-1]], \ p[tmp[i]]))) < eps) \ \{ (i = 2; \ i < n; \ ++i) \ if \ (fabs(submul(getvec(p[seq[n-1]], \ p[tmp[i]]))) < eps) \ \{ (i = 2; \ i < n; \ ++i) \ if \ (fabs(submul(getvec(p[seq[n-1]], \ p[tmp[i]]))) < eps) \ \{ (i = 2; \ i < n; \ ++i) \ if \ (fabs(submul(getvec(p[seq[n-1]], \ p[tmp[i]]))) < eps) \ \{ (i = 2; \ i < n; \ ++i) \ if \ (fabs(submul(getvec(p[seq[n-1]], \ p[tmp[i]]))) < eps) \ \{ (i = 2; \ i < n; \ ++i) \ if \ (fabs(submul(getvec(p[seq[n-1]], \ p[tmp[i]]))) < eps) \ \{ (i = 2; \ i < n; \ ++i) \ if \ (fabs(submul(getvec(p[seq[n-1]], \ p[tmp[i]]))) < eps) \ \{ (i = 2; \ i < n; \ ++i) \ if \ (fabs(submul(getvec(p[seq[n-1]], \ p[tmp[i]]))) < eps) \ \{ (i = 2; \ i < n; \ ++i) \ if \ (fabs(submul(getvec(p[seq[n-1]], \ p[tmp[i]]))) < eps) \ \{ (i = 2; \ i < n; \ ++i) \ if \ (fabs(submul(getvec(p[seq[n-1]], \ p[tmp[i]]))) < eps) \ \{ (i = 2; \ i < n; \ ++i) \ if \ (fabs(submul(getvec(p[seq[n-1]], \ p[tmp[i]]))) < eps) \ \{ (i = 2; \ i < n; \ ++i) \ if \ (fabs(submul(getvec(p[seq[n-1]], \ p[tmp[i]]))) < eps) \ \{ (i = 2; \ i < n; \ ++i) \ if \ (fabs(submul(getvec(p[seq[n-1]], \ p[tmp[i]]))) < eps) \ \{ (i = 2; \ i < n; \ ++i) \ if \ (fabs(submul(getvec(p[seq[n-1]], \ p[tmp[i]]))) < eps) \ \{ (i = 2; \ i < n; \ ++i) \ if \ (fabs(submul(getvec(p[seq[n-1]], \ p[tmp[i]]))) < eps) \ \{ (i = 2; \ i < n; \ ++i) \ if \ (fabs(submul(getvec(p[seq[n-1]], \ p[tmp[i]]))) < eps) \ \{ (i = 2; \ i < n; \ ++i) \ if \ (fabs(submul(getvec(p[seq[n-1]], \ p[tmp[i]]))) < eps) \ \{ (i = 2; \ i < n; \ ++i) \ if \ (fabs(submul(getvec(p[seq[n-1], \ p[seq[n-1]]))) < eps) \ \{ (i = 2; \ i < n; \ ++i) \ if \ (fabs(submu(getvec(p[seq[n-1], \ p[seq[n-1]]))) < eps) \ \{ (i = 2; \ i < n; \ ++i) \ if \ (fabs(submu(getvec(p[seq[n-1], \ p[seq[n-
                                     } else break; top = n+1; bot = n-1;
                   for (; i < n; ++i) {
                                      if \ (submul(getvec(p[seq[top-1]], \ p[seq[top]]), \ getvec(p[seq[top]], \ p[tmp[i]])) \ > -eps \ \&\& \ + (between the content of the conten
                                                           submul\left(getvec\left(p[seq[bot+1]],\ p[seq[bot]]\right),\ getvec\left(p[seq[bot]],\ p[tmp[i]]\right)\right) \ < \ eps\right) \ continue;
                                       while \ (submul(getvec(p[seq[top-1]], \ p[seq[top]]), \ getvec(p[seq[top]], \ p[tmp[i]])) \ < \ eps) \ --top;
                                       seq[++top] = seq[--bot] = tmp[i];
                  } for (i = bot; i <= top; ++i) seq[i-bot] = seq[i]; return top - bot;
//Randomize method to get smallest enclosing circle for N given points
int tmp[N], seq[N];
double enclosng_circle(int n, pnt *p, pnt &cen) {
                   int i, j, k, 1; double rad;
                   for (i = 0; i < n; ++i) seq[i] = i; cen = p[0]; rad = 0.0;
                   for (i = 2; i < n; ++i) {
                                       j = rand()\%(n-i) + i; k = seq[j]; seq[j] = seq[i]; seq[i] = k;
                                       if (getdis(cen, p[seq[i]])-rad < eps) continue;</pre>
                                       tmp[0] = seq[i]; tmp[1] = seq[1]; seq[i] = seq[0];
                                       {\tt cen = (p[tmp[0]]+p[tmp[1]]) \ / \ 2.0; \ rad = getdis(cen, \ p[tmp[0]]);}
                                       for (j = 2; j \le i; ++j) {
                                                           if (getdis(cen, p[seq[j]])-rad < eps) { tmp[j] = seq[j]; continue; }
                                                           tmp[j] = tmp[1]; tmp[1] = seq[j];
                                                           {\tt cen = (p[tmp[0]]+p[tmp[1]]) \ / \ 2.0; \ rad = getdis(cen, \ p[tmp[0]]);}
                                                           for (k = 2; k \le j; ++k) if (getdis(cen, p[tmp[k]])-rad > eps) {
                                                                             1 = tmp[k]; tmp[k] = tmp[2]; tmp[2] = 1;
                                                                             rad = outcircle(p[tmp[0]], p[tmp[1]], p[tmp[2]], cen);
                                      } for (j = 0; j \le i; ++j) seq[j] = tmp[j];
                  } return rad;
```

```
//get the farthest point for each given point in convex, store in array ati(i's farthest point is at index ati[i])
//Caution: 1. The order of convex must be counter-clockwise
            2. The length of array(N) should be 3 times larger than the data region. (The same for array p)
            3. The line construct function here has no need to unify.
int pre[N], pro[N], pos[N];
bool vst[N]; line l1[N];
void farestpair(int n, pnt *p, int *ati) {
      int i, j, k, 1, st, ed;
      for (i = 0; i < n; ++i) \{ p[i].x \neq 2; p[i].y \neq 2; \}
       for \ (i = 0; \ i < n; \ +\!+i) \ \{ \ pre[i] = i-\!\!\!-1; \ pro[i] = i+\!\!\!\!+1; \ \} \ pre[0] = n-\!\!\!\!-1; \ pro[n-\!\!\!\!-1] = 0; 
      for (i = 0; i < n; ++i) p[i+n] = p[i+2*n] = p[i];
      for (1 = st = i = 0; 1 \le n; i = pro[i], ++1) {
            11[i] = getline(p[i], p[pro[i]]);
             while ((i!=st | |1==n) \& submul (p[pro[i]]-p[i]), p[pos[pre[i]]]-p[i]) >= 0 \& soridis (11[i], p[pos[pre[i]]]) + (1!=n-1?0:1) > 0 ) 
                   if (1 == n \&\& i == st) { st = pro[i]; 1 = n-1; }
                   pro[pre[i]] = pro[i]; pre[pro[i]] = pre[i]; i = pre[i];
                   11[i] = getline(p[i], p[pro[i]]);
            } if (i == st && 1 != n) ed = pro[i]; else ed = pos[pre[i]];
            } memset(vst, 0, sizeof(vst));
      for (i = st; !vst[i]; i = pro[i]) {
            j = pos[i]; k = pos[pro[i]]; vst[i] = 1; while (k < j) k += n;
            for (1 = j; 1 < k; ++1) ati[1%n] = pro[i];</pre>
/*Hi-light-expression-1: (A?B:C)
                                      1. If the input is float, change the B and C to the 2*eps if its value equals 1 below.
                                      2. For the same farthest, A is diffierent according to the requirement to choose the index:
                                            1).the smallest index: A <- 1==n-1
                                            2).the largest index : A <- 1!=n-1
                                            3).arbitrary: A <- 1==n-1, 1!=n-1, 1, 0 are all ok.
```

## \*Trick Note:

1.after erase of mid-point, the terminal of new seg's previous seg may before the endpoint pro[i], even on the left side of judge line. In this condition, the erase operation should be stopped. And when calculate terminal, the points before pro[i] should be ignored. The farthest points to them will not be determined by both ends of the current seg. using submul to judge such kind of point.

2.originally I set the l=n to l=n-1 after the erase operation. But the state l become n-1 since the st point has erased. There's a trick data in requirement of 2) if this small error is ignored.

\*Optimize Note: The operation of calculate the terminal can be optimized using dichonomy(from ed->i\*k[k is least positive integer to make ed<i\*k]).

\*/

VIII

 $\begin{tabular}{ll} while $(submul(v, getvec(p[i], p[j+1]))-submul(v, getvec(p[i], p[j])) > -eps) $j = (j+1) \% $n$; \\ while $(nummul(v, getvec(p[i], p[k+1]))-nummul(v, getvec(p[i], p[k])) > -eps) $k = (k+1) \% $n$; \\ while $(nummul(v, getvec(p[i], p[1+1]))-nummul(v, getvec(p[i], p[1])) < eps) $l = (l+1) \% $n$; \\ \end{tabular}$ 

wi = nummul(v, getvec(p[1], p[k])); hi = submul(v, getvec(p[i], p[j]));

area = getmin(area, wi\*hi); pra = getmin(pra, (wi+hi)\*2);

}

```
#define square(x) (x)*(x)
#define xmu1(x1, y1, x2, y2) ((x1)*(y2)-(x2)*(y1))
#define dmul(x1, y1, x2, y2) ((x1)*(x2)+(y1)*(y2))
const double pi = acos(-1.0), eps = 1e-8;
struct pnt {
      double x, y, z;
      pnt operator+(const pnt &p) const {
            pnt ret; ret.x = x + p.x; ret.y = y + p.y;
            ret. z = z + p. z; return ret;
      } pnt operator-(const pnt &p) const {
            pnt ret; ret.x = x - p.x; ret.y = y - p.y;
            ret.z = z - p.z; return ret;
      } pnt operator*(const double c) const {
            pnt ret; ret.x = x * c; ret.y = y * c;
            ret.z = z * c; return ret;
      } pnt operator/(const double c) const {
            pnt ret; ret.x = x / c; ret.y = y / c;
            ret.z = z / c; return ret;
}:
typedef pnt vec;
struct plane { vec ori; double val; };
vec submul(const vec &v1, const vec &v2) {
      vec ret; ret.x = xmul(v1.y, v1.z, v2.y, v2.z); ret.y = xmul(v1.z, v1.x, v2.z, v2.x);
      ret.z = xmul(v1.x, v1.y, v2.x, v2.y); return ret;
vec uvec (vec v) {
      double len = sqrt(square(v.x)+square(v.y)+square(v.z));
      vec ret; ret.x = v.x / len; ret.y = v.y / len; ret.z = v.z / len; return ret;
plane getplane(const pnt &p1, const pnt &p2, const pnt &p3) {
      vec v1 = p2 - p1, v2 = p3 - p1;
      plane ret; ret.ori = uvec(submul(v1, v2));
      ret.val = nummul(ret.ori, p1); return ret;
plane uplane(const plane &s) {
      double len = sqrt(square(s.ori.x)+square(s.ori.y)+square(s.ori.z));
      plane ret; ret.ori = s.ori / len; ret.val = s.val / len; return ret;
//give the rotate angle(conter-clockwise) and the normal vector for rotate, output the position of given point after the rotate.
#define OFF(i, j) ((i) == (j) ? 1 : 0)
pnt axis_rotate(const pnt &p, const pnt &ori, double ang) {
      vec v = uvec(ori);
      \label{eq:conditional_condition} \mbox{double pv[3] = { p. x, p. y, p. z}, val[3] = { v. x, v. y, v. z }, \mbox{ off[3] = { 0, -1, 1 };} \\
      double sa = sin(ang), ca = cos(ang), rv[3];
      for (int i = 0; i < 3; ++i) for (int j = rv[i] = 0; j < 3; ++j)
           rv[i] += pv[j] * ((1 - ca) * val[i] * val[j] + sa * val[(6-i-j)%3] * off[(3-i+j)%3] + ca * OFF(i, j));
      pnt ret = { rv[0], rv[1], rv[2] }; return ret;
/* Note: the principle shows below -->
                                                             | 0, -vz, vy |
            vx*vx, vx*vy, vx*vz
      A1 = | vy*vx, vy*vy, vy*vz |
                                                        A2 = | vz, 0, -vx |
                                                              |-vy, vx, 0 |
            vz*vz, vz*vy, vz*vz
      p' = ((1-cosa)*A1+sina*A2+cosa*I) * p;
      It's a simplified expression for axis_rotate2, but it doesn't always run faster than the latter.
```

```
//{\rm the} same as the axis_rotate. Sometimes it runs better but the code is longer.
pnt axis_rotate2(const pnt &p, const pnt &ori, double ang) {
      double ca[5], sa[5], tv1, tv2, xyv; vec v = uvec(ori); int i, j, k;
      xyv = sqrt(square(v.x)+square(v.y));
      if (xyv > eps) \{ ca[0] = v.x / xyv; sa[0] = -v.y / xyv; \}
      else { ca[0] = 1.0; sa[0] = 0.0; }
      ca[1] = v.z; sa[1] = -xyv; ca[2] = cos(ang); sa[2] = sin(ang);
      ca[3] = ca[1]; sa[3] = -sa[1]; ca[4] = ca[0]; sa[4] = -sa[0];
      double rv[3] = \{ p. x, p. y, p. z \};
      for (i = 0; i < 5; ++i) {
             j = (3-i\%2) \% 3; k = (j+1)\%3;
             tv1 = xmul(rv[j], rv[k], sa[i], ca[i]); tv2 = dmul(rv[j], rv[k], sa[i], ca[i]);
             rv[i] = tv1: rv[k] = tv2:
      } pnt ret = { rv[0], rv[1], rv[2] }; return ret;
//Transformation with matrix
//The advantage of it is to simplify the multi transformation operation. A matrix can contain the combination of the operation. If
there're n points and m operation, using matrix only takes 0\,(\text{n+m}) time.
//Caution1: right-hand coordinate system and right-hand rule in rotation only.
//Caution2: the whole transformation matrix should initially set to I.
void matmul(double a[][4], double b[][4]) {
      int i, j, k; double c[4][4];
      for (i = 0; i < 4; ++i) for (j = 0; j < 4; ++j) { c[i][j] = a[i][j]; a[i][j] = 0; }
       \text{for } (i = 0; \ i < 4; \ ++i) \ \text{for } (j = 0; \ j < 4; \ ++j) \ \text{for } (k = 0; \ k < 4; \ ++k) \ a[i][j] \ += \ c[i][k] \ * \ b[k][j]; 
pnt pnttran(const pnt &p, double a[][4]) {
      double val[4] = { p. x, p. y, p. z, 1 }, rv[4] = \{ 0, 0, 0, 0 \};
       \text{for (int i = 0; i < 4; ++i) for (int j = 0; j < 4; ++j) rv[i] += val[j] * a[j][i]; } \\
      pnt ret = { rv[0], rv[1], rv[2] }; return ret;
void mattranslate(double a[][4], double dx, double dy, double dz) {
      int i; double b[4][4];
      memset(b, 0, sizeof(b)); for (i = 0; i < 4; ++i) b[i][i] = 1;
      b[3][0] = dx; b[3][1] = dy; b[3][2] = dz; matmul(a, b);
void matrotate(double a[][4], const vec &axis, double dlt) {
      int i, j; double b[4][4]; vec v = uvec(axis);
      double sind = sin(dlt), cosd = cos(dlt), val[3] = { v. x, v. y, v. z }, off[3] = { 0, 1, -1 };
      memset(b, 0, sizeof(b)); for (i = 0; i < 4; ++i) b[i][i] = 1;
      for (i = 0; i < 3; ++i) for (j = 0; j < 3; ++j)
            b[i][j] = (1-\cos d) * val[i] * val[j] + \sin d * off[(3-i+j)\%3] * val[(6-i-j)\%3] + \cos d * OFF(i, j);
      matmul(a, b);
void matscale(double a[][4], double cx, double cy, double cz) {
      int i; double b[4][4];
      memset(b, 0, sizeof(b)); for (i = 0; i < 4; ++i) b[i][i] = 1;
      b[0][0] = cx; b[1][1] = cy; b[2][2] = cz; matmul(a, b);
}
                           1, 0, 0, 0
                                                                     cx, 0, 0, 0
TRANSLATE(dx, dy, dz) =
                            \mid \ 0, \ 1, \ 0, \ 0 \mid \ \text{SCALE}\left(\text{cx}, \text{cy}, \text{cz}\right) = \ \mid \ 0, \ \text{cy}, \ 0, \ 0 \mid \ \text{ROTATE}\left(\text{vx}, \text{vy}, \text{vz}, \text{delta}\right) \text{ is decribed above}. 
                           0, 0, 1, 0
                                                                     0, 0, cz, 0
                           | dx, dy, dz, 1 |
                                                                     0, 0, 0, 1
```

```
//2-Sat
//The meaning of the edge (u->v) means u must be chosen before v is chosen.
struct node { int des, next; };
struct edge { int st, ed; };
#define N 10010
#define M 1000010
node way[M];
edge e[M];
int rd[N], seq[N], col[N], pre[N], cpn[N];
bool vst[N], use[M], chs[N];
int n, cnt, cc, rc;
void dfs(int x)
{
      int i:
       for (i = rd[x]; i; i = way[i].next) if (!vst[way[i].des]) dfs(way[i].des);
      seq[++cnt] = x;
void dfs2(int x, int c)
{
       int i;
      vst[x] = true; col[x] = c;
      for (i = rd[cpn[x]]; i; i = way[i].next) if (!vst[cpn[way[i].des]]) dfs2(cpn[way[i].des], c);
void release(int x)
{
       int i;
      vst[x] = true; vst[cpn[x]] = true; chs[x] = true; chs[cpn[x]] = false; cc += 2;
      for (i = rd[x]; i; i = way[i].next) pre[way[i].des]--;
void twosat()
{
      //iudge section
      memset(vst, false, sizeof(vst)); cnt = 0;
       for (i = 1; i \le n; ++i) if (!vst[i]) dfs(i);
      memset(vst, false, sizeof(vst)); cnt = 0;
       for (i = n; i \geq 0; --i) if (!vst[seq[i]]) dfs2(seq[i], ++cnt);
       for (i = 1; i \leq n; ++i) if (col[i] == col[cpn[i]]) { printf("NO\n"); return; } printf("YES\n");
       //output section
        for \ (i = 1; \ i <= n; \ +\!+i) \ seq[col[i]] = col[cpn[i]]; \ for \ (i = 1; \ i <= cnt; \ +\!+i) \ cpn[i] = seq[i]; 
      memset(use, true, sizeof(use));
       for (i = 1; i \le rc; ++i) if (col[e[i].st] == col[e[i].ed]) use[i] = false;
      else { e[i].st = col[e[i].st]; e[i].ed = col[e[i].ed]; }
      \texttt{memset}(\texttt{rd}, \ \ 0, \ \texttt{sizeof}(\texttt{rd})) \, ; \ \texttt{memset}(\texttt{pre}, \ \ 0, \ \texttt{sizeof}(\texttt{pre})) \, ;
       for (i = 1; i <= rc; ++i) if (use[i]) { way[i].des = e[i].ed; way[i].next = rd[e[i].st]; rd[e[i].st] = i; pre[e[i].ed]++; }
      memset(vst, false, sizeof(vst)); cc = 0;
      while (cc < cnt) for (i = 1; i <= cnt; ++i) if (!vst[i] \&\& !pre[i]) release(i);
       for (i = 1; i \le n; ++i) if (chs[col[i]]) printf("%d ", i); printf("\n");
```