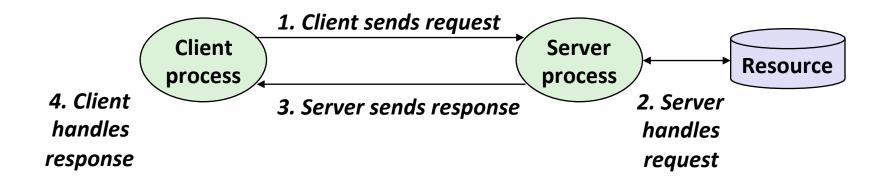
# Internetworking

15-213/18-213: Introduction to Computer Systems 20<sup>th</sup> Lecture, April. 1st, 2014

#### **Instructors:**

Anthony Rowe, Seth Goldstein, and Gregory Kesden

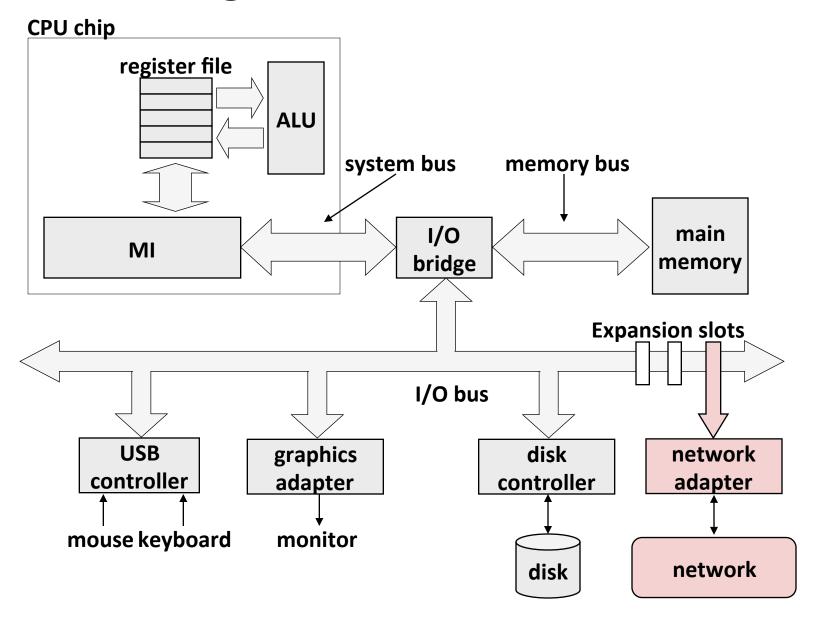
## **A Client-Server Transaction**



Note: clients and servers are processes running on hosts (can be the same or different hosts)

- Most network applications are based on the client-server model:
  - A server process and one or more client processes
  - Server manages some resource
  - Server provides service by manipulating resource for clients
  - Server activated by request from client (vending machine analogy)

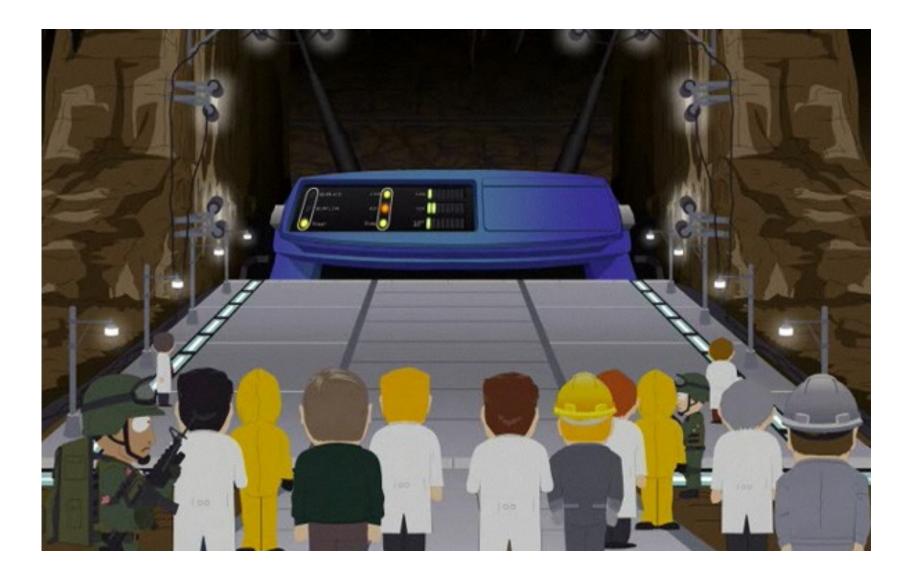
# **Hardware Organization of a Network Host**



# **Computer Networks**

- A network is a hierarchical system of boxes and wires (or not) organized by geographical proximity
  - SAN (System Area Network) spans cluster or machine room
    - Switched Ethernet, Quadrics QSW, ...
  - LAN (Local Area Network) spans a building or campus
    - Ethernet is most prominent example
  - WAN (Wide Area Network) spans country or world
    - Typically high-speed point-to-point phone lines
- An internetwork (internet) is an interconnected set of networks
  - The Global IP Internet (uppercase "I") is the most famous example of an internet (lowercase "i")
- Let's see how an internet is built from the ground up

## What is the Internet?



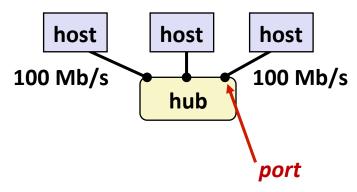
# The



# "Internet"

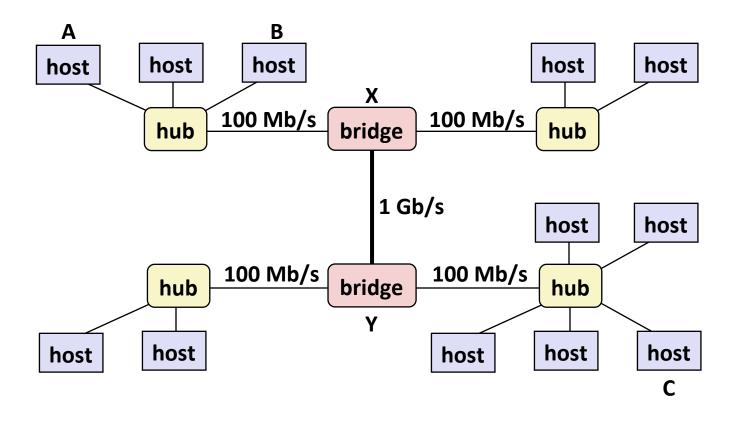
"Rough consensus and Working code"
- Dave Clarke, MIT

# **Lowest Level: Ethernet Segment**



- Ethernet segment consists of a collection of hosts connected by wires (twisted pairs) to a hub
- Spans room or floor in a building
- Operation
  - Each Ethernet adapter has a unique 48-bit address (MAC address)
    - E.g., 00:16:ea:e3:54:e6
  - Hosts send bits to any other host in chunks called frames
  - Hub slavishly copies each bit from each port to every other port
    - Every host sees every bit
    - Note: Hubs are on their way out. Bridges (switches, routers) became cheap enough to replace them (means no more broadcasting)

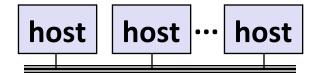
# **Next Level: Bridged Ethernet Segment**



- Spans building or campus
- Bridges cleverly learn which hosts are reachable from which ports and then selectively copy frames from port to port

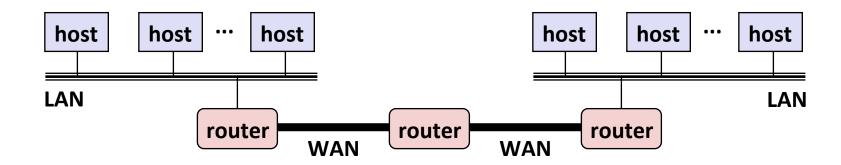
# **Conceptual View of LANs**

For simplicity, hubs, bridges, and wires are often shown as a collection of hosts attached to a single wire:



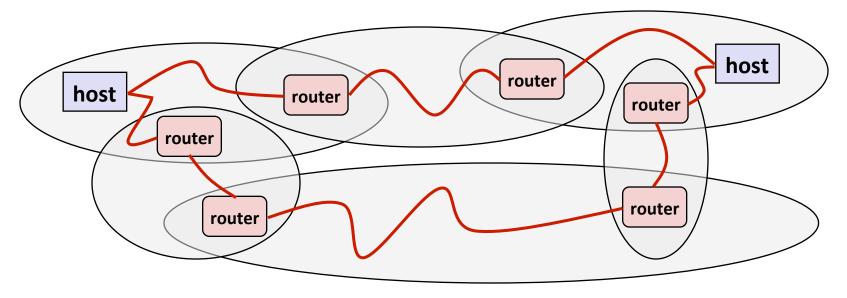
## **Next Level: internets**

- Multiple incompatible LANs can be physically connected by specialized computers called *routers*
- The connected networks are called an internet



LAN 1 and LAN 2 might be completely different, totally incompatible (e.g., Ethernet and Wifi, 802.11\*, T1-links, DSL, ...)

# Logical Structure of an internet



- Ad hoc interconnection of networks
  - No particular topology
  - Vastly different router & link capacities
- Send packets from source to destination by hopping through networks
  - Router forms bridge from one network to another
  - Different packets may take different routes

## The Notion of an internet Protocol

How is it possible to send bits across incompatible LANs and WANs?

#### Solution:

- protocol software running on each host and router
- smooths out the differences between the different networks

#### Implements an internet protocol (i.e., set of rules)

- governs how hosts and routers should cooperate when they transfer data from network to network
- TCP/IP is the protocol for the global IP Internet

## What Does an internet Protocol Do?

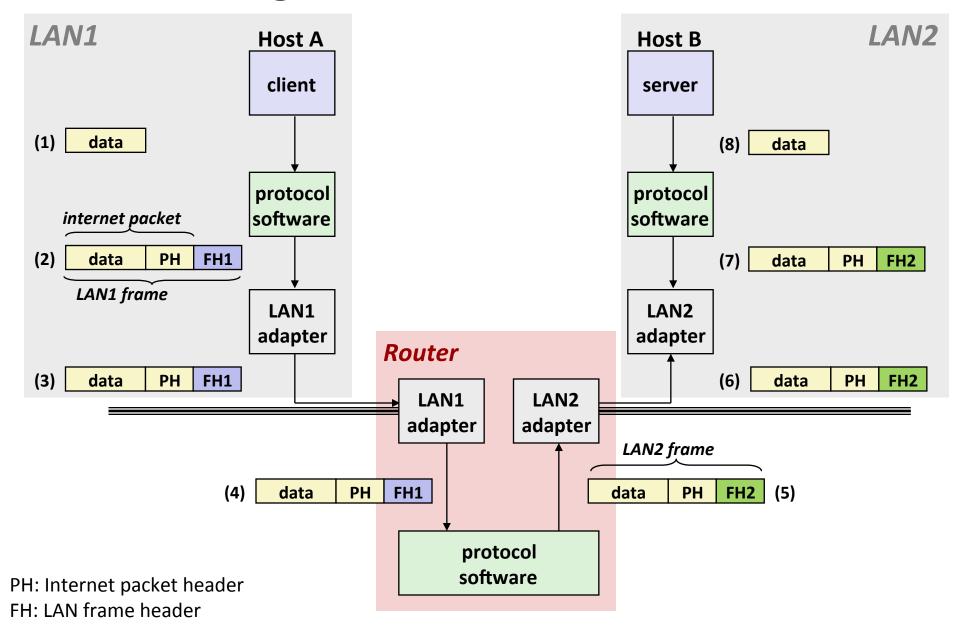
#### Provides a naming scheme

- An internet protocol defines a uniform format for host addresses
- Each host (and router) is assigned at least one of these internet addresses that uniquely identifies it

### Provides a delivery mechanism

- An internet protocol defines a standard transfer unit (packet)
- Packet consists of header and payload
  - Header: contains info such as packet size, source and destination addresses
  - Payload: contains data bits sent from source host

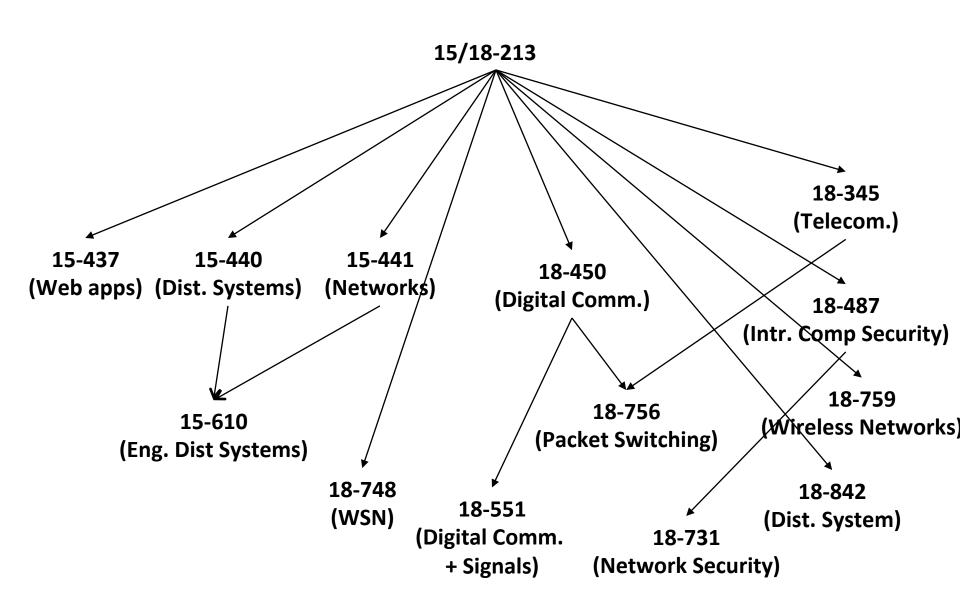
# **Transferring Data Over an internet**



## Other Issues

- We are glossing over a number of important questions:
  - What if different networks have different maximum frame sizes? (segmentation)
  - How do routers know where to forward frames?
  - How are routers informed when the network topology changes?
  - What if packets get lost?
- These (and other) questions are addressed by the area of systems known as computer networking

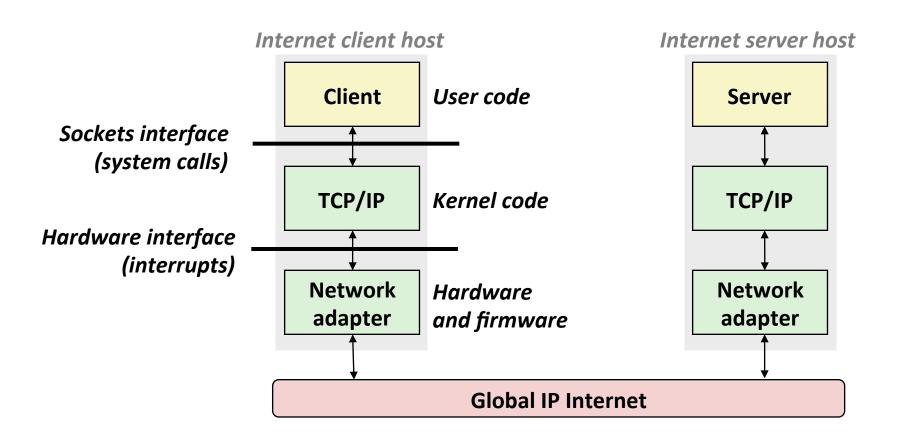
# If you are interested in networking...



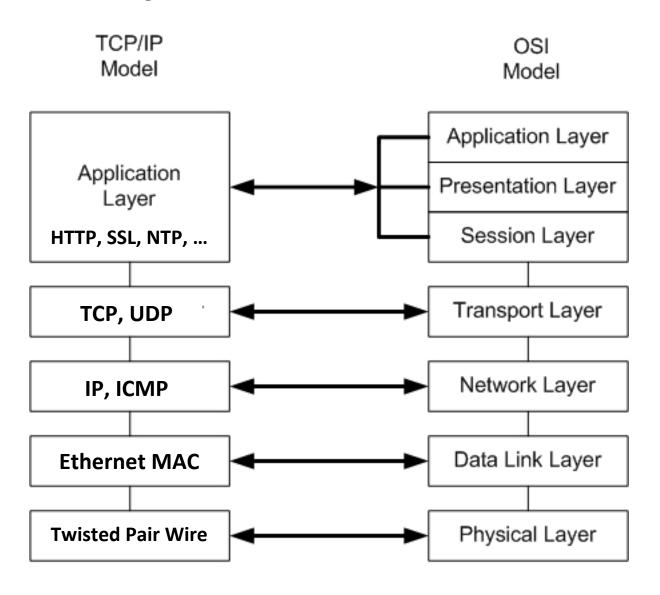
## **Global IP Internet**

- Most famous example of an internet
- Based on the TCP/IP protocol family
  - IP (Internet protocol) :
    - Provides basic naming scheme and unreliable delivery capability of packets (datagrams) from host-to-host
  - UDP (Unreliable Datagram Protocol)
    - Uses IP to provide unreliable datagram delivery from process-to-process
  - TCP (Transmission Control Protocol)
    - Uses IP to provide *reliable* byte streams from process-to-process over connections
- Accessed via a mix of Unix file I/O and functions from the sockets interface

# Hardware and Software Organization of an Internet Application



# **Protocol Layers**



# **Basic Internet Components**

#### Internet backbone:

 collection of routers (nationwide or worldwide) connected by high-speed point-to-point networks

#### Network Access Point (NAP):

router that connects multiple backbones (often referred to as peers)

#### Regional networks:

 smaller backbones that cover smaller geographical areas (e.g., cities or states)

#### Point of presence (POP):

machine that is connected to the Internet

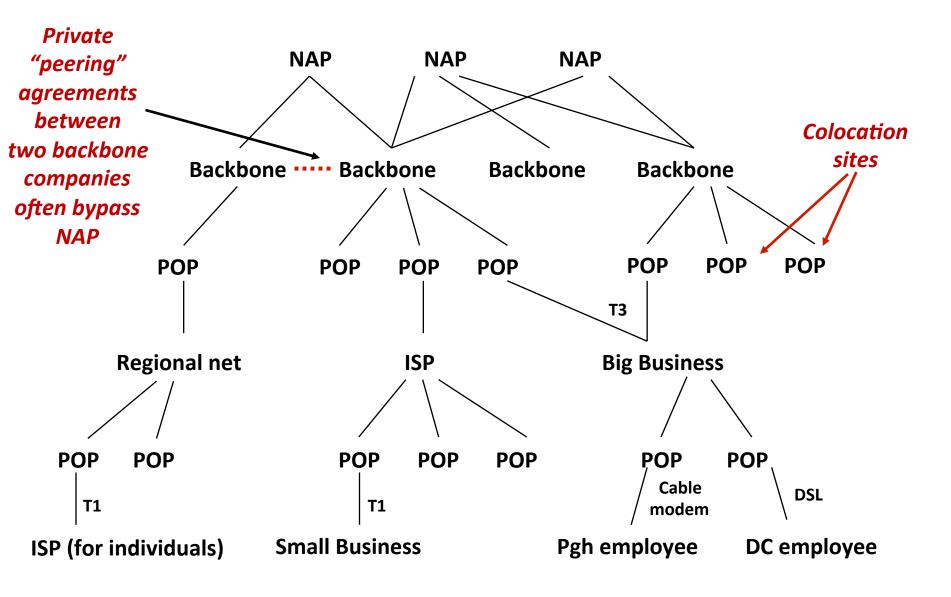
#### Internet Service Providers (ISPs):

provide dial-up or direct access to POPs

## **NAP-Based Internet Architecture**

- NAPs link together commercial backbones provided by companies such as AT&T and Worldcom
- Currently in the US there are about 50 commercial backbones connected by ~12 NAPs (peering points)
- Similar architecture worldwide connects national networks to the Internet
- Now called Internet Exchange Points (IXP)
  - Link together ISPs

# **Internet Connection Hierarchy**



# Naming and Communicating on the Internet

#### Original Idea

- Every node on Internet would have unique IP address
  - Everyone would be able to talk directly to everyone
- No secrecy or authentication
  - Messages visible to routers and hosts on same LAN
  - Possible to forge source field in packet header

#### Shortcomings

- There aren't enough IP addresses available
- Don't want everyone to have access or knowledge of all other hosts
- Security issues mandate secrecy & authentication

# **Evolution of Internet: Naming**

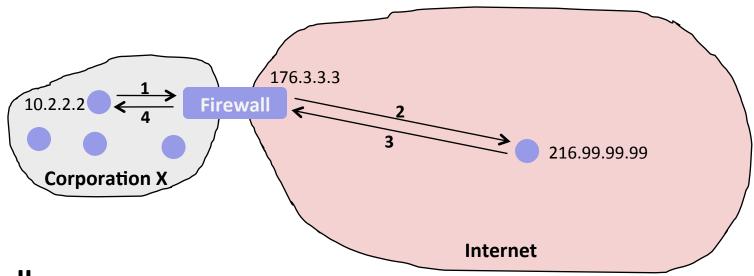
#### Dynamic address assignment

- Most hosts don't need to have known address
  - Only those functioning as servers
- DHCP (Dynamic Host Configuration Protocol)
  - Local ISP assigns address for temporary use

#### Example:

- My laptop at CMU (wired connection)
  - IP address 128.237.207.3 (corvina.ece.cmu.edu)
  - Assigned statically
- My laptop at home
  - IP address 192.168.1.5
  - Only valid within home network

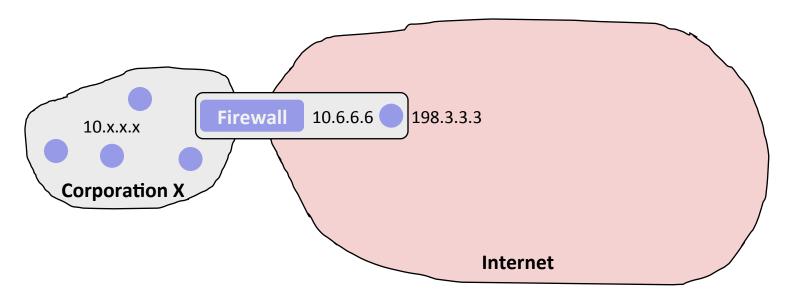
## **Evolution of Internet: Firewalls**



#### Firewalls

- Hides organizations nodes from rest of Internet
- Use local IP addresses within organization
- For external service, provides proxy service
  - 1. Client request: src=10.2.2.2, dest=216.99.99.99
  - 2. Firewall forwards: src=176.3.3.3, dest=216.99.99.99
  - 3. Server responds: src=216.99.99.99, dest=176.3.3.3
  - 4. Firewall forwards response: src=216.99.99.99, dest=10.2.2.2

## **Virtual Private Networks**



#### Supporting road warrior

- Employee working remotely with assigned IP address 198.3.3.3
- Wants to appear to rest of corporation as if working internally
  - From address 10.6.6.6
  - Gives access to internal services (e.g., ability to send mail)

### Virtual Private Network (VPN)

Overlays private network on top of regular Internet

# A Programmer's View of the Internet

- Hosts are mapped to a set of 32-bit IP addresses
  - **128.2.203.179**
- The set of IP addresses is mapped to a set of identifiers called Internet *domain names* 
  - 128.2.203.179 is mapped to www.cs.cmu.edu
- A process on one Internet host can communicate with a process on another Internet host over a connection

## **IP Addresses**

#### ■ 32-bit IP addresses are stored in an IP address struct

- IP addresses are always stored in memory in network byte order (big-endian byte order)
- True in general for any integer transferred in a packet header from one machine to another.
  - E.g., the port number used to identify an Internet connection.

```
/* Internet address structure */
struct in_addr {
   unsigned int s_addr; /* network byte order (big-endian) */
};
```

#### Useful network byte-order conversion functions ("I" = 32 bits, "s" = 16 bits)

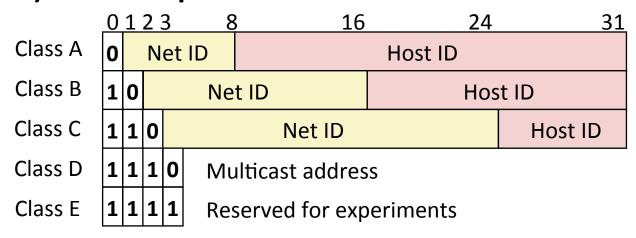
```
htonl: convert uint32_t from host to network byte order
htons: convert uint16_t from host to network byte order
ntohl: convert uint32_t from network to host byte order
ntohs: convert uint16_t from network to host byte order
```

## **Dotted Decimal Notation**

- By convention, each byte in a 32-bit IP address is represented by its decimal value and separated by a period
  - IP address: 0x8002C2F2 = 128.2.194.242
- Functions for converting between binary IP addresses and dotted decimal strings:
  - inet\_aton: dotted decimal string → IP address in network byte order
  - inet\_ntoa: IP address in network byte order → dotted decimal string
  - "n" denotes network representation
  - "a" denotes application representation

## **IP Address Structure**

IP (V4) Address space divided into classes:



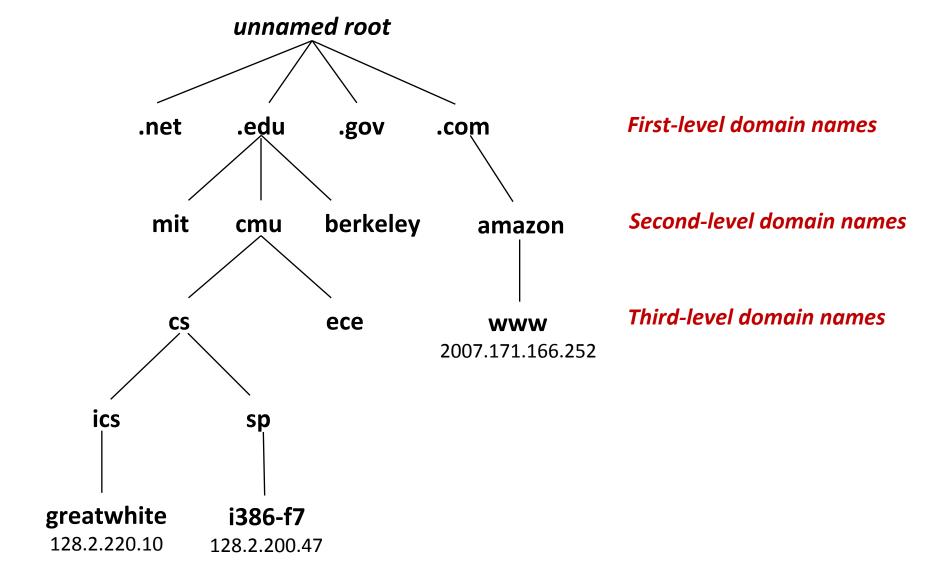
### Network ID Written in form w.x.y.z/n

- n = number of bits in host address
- E.g., CMU written as 128.2.0.0/16
  - Class B address

#### Unrouted (private) IP addresses:

10.0.0.0/8 172.16.0.0/12 192.168.0.0/16

## **Internet Domain Names**



# **Domain Naming System (DNS)**

- The Internet maintains a mapping between IP addresses and domain names in a huge worldwide distributed database called DNS
  - Conceptually, programmers can view the DNS database as a collection of millions of host entry structures:

- Functions for retrieving host entries from DNS:
  - **gethostbyname:** query key is a DNS domain name.
  - gethostbyaddr: query key is an IP address.

# **Properties of DNS Host Entries**

- Each host entry is an equivalence class of domain names and IP addresses
- Each host has a locally defined domain name localhost which always maps to the loopback address 127.0.0.1
- Different kinds of mappings are possible:
  - Simple case: one-to-one mapping between domain name and IP address:
    - greatwhile.ics.cs.cmu.edu maps to 128.2.220.10
  - Multiple domain names mapped to the same IP address:
    - eecs.mit.edu and cs.mit.edu both map to 18.62.1.6
  - Multiple domain names mapped to multiple IP addresses:
    - google.com maps to multiple IP addresses
  - Some valid domain names don't map to any IP address:
    - for example: ics.cs.cmu.edu

# A Program That Queries DNS

```
int main(int argc, char **argv) { /* argv[1] is a domain name */
                                  /* or dotted decimal IP addr */
    char **pp;
    struct in addr addr;
    struct hostent *hostp;
    if (inet aton(argv[1], &addr) != 0)
       hostp = Gethostbyaddr((const char *)&addr, sizeof(addr),
                AF INET);
   else
       hostp = Gethostbyname(argv[1]);
   printf("official hostname: %s\n", hostp->h name);
    for (pp = hostp->h aliases; *pp != NULL; pp++)
       printf("alias: %s\n", *pp);
    for (pp = hostp->h addr list; *pp != NULL; pp++) {
        addr.s addr = ((struct in addr *)*pp)->s addr;
       printf("address: %s\n", inet ntoa(addr));
```

# **Using DNS Program**

```
linux> ./dns greatwhite.ics.cs.cmu.edu
official hostname: greatwhite.ics.cs.cmu.edu
address 128.2.220.10
linux> ./dns 128.2.220.11
official hostname: ANGELSHARK.ICS.CS.CMU.EDU
address: 128.2.220.11
linux> ./dns www.google.com
official hostname: www.l.google.com
alias: www.google.com
address: 72.14.204.99
address: 72.14.204.103
address: 72.14.204.104
address: 72.14.204.147
linux> dig +short -x 72.14.204.103
iad04s01-in-f103.1e100.net.
```

# **Querying DIG**

 Domain Information Groper (dig) provides a scriptable command line interface to DNS

```
linux> dig +short greatwhite.ics.cs.cmu.edu
128.2.220.10
linux> dig +short -x 128.2.220.11
ANGELSHARK.ICS.CS.CMU.EDU.
linux> dig +short google.com
72.14.204.104
72.14.204.147
72.14.204.99
72.14.204.103
linux> dig +short -x 72.14.204.103
iad04s01-in-f103.1e100.net.
```

## **More Exotic Features of DIG**

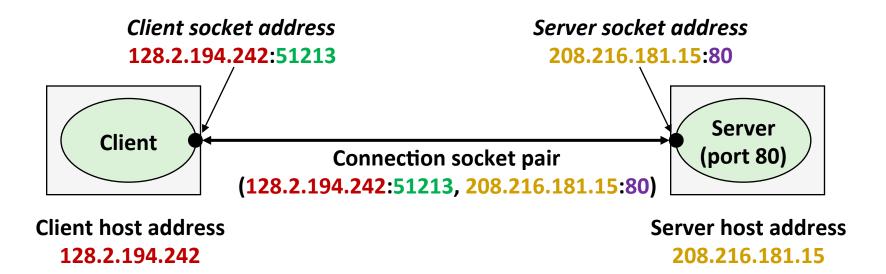
Provides more information than you would ever want about DNS

```
linux> dig www.phys.msu.ru a +trace
128.2.220.10
linux> dig www.google.com a +trace
```

## **Internet Connections**

- Clients and servers communicate by sending streams of bytes over connections:
  - Point-to-point, full-duplex (2-way communication), and reliable.
- A socket is an endpoint of a connection
  - Socket address is an IPaddress:port pair
- A *port* is a 16-bit integer that identifies a process:
  - Ephemeral port: Assigned automatically on client when client makes a connection request
  - Well-known port: Associated with some service provided by a server (e.g., port 80 is associated with Web servers)
- A connection is uniquely identified by the socket addresses of its endpoints (socket pair)
  - (cliaddr:cliport, servaddr:servport)

# Putting it all Together: Anatomy of an Internet Connection



## **Next Time**

- How to use the sockets interface to establish Internet connections between clients and servers
- How to use Unix I/O to copy data from one host to another over an Internet connection