

## 10. Unconstrained minimization

- terminology and assumptions
- gradient descent method
- steepest descent method
- Newton's method
- self-concordant functions
- implementation

# Unconstrained minimization

$$\text{minimize } f(x)$$

- $f$  convex, twice continuously differentiable (hence  $\text{dom } f$  open)
- we assume optimal value  $p^* = \inf_x f(x)$  is attained (and finite)

## unconstrained minimization methods

- produce sequence of points  $x^{(k)} \in \text{dom } f$ ,  $k = 0, 1, \dots$  with

$$f(x^{(k)}) \rightarrow p^*$$

- can be interpreted as iterative methods for solving optimality condition

$$\nabla f(x^*) = 0$$

## Initial point and sublevel set

algorithms in this chapter require a starting point  $x^{(0)}$  such that

- $x^{(0)} \in \mathbf{dom} f$
- sublevel set  $S = \{x \mid f(x) \leq f(x^{(0)})\}$  is closed

2nd condition is hard to verify, except when *all* sublevel sets are closed:

- equivalent to condition that  $\mathbf{epi} f$  is closed
- true if  $\mathbf{dom} f = \mathbf{R}^n$
- true if  $f(x) \rightarrow \infty$  as  $x \rightarrow \mathbf{bd} \mathbf{dom} f$

examples of differentiable functions with closed sublevel sets:

$$f(x) = \log\left(\sum_{i=1}^m \exp(a_i^T x + b_i)\right), \quad f(x) = -\sum_{i=1}^m \log(b_i - a_i^T x)$$

# Strong convexity and implications

$f$  is strongly convex on  $S$  if there exists an  $m > 0$  such that

$$\nabla^2 f(x) \succeq mI \quad \text{for all } x \in S$$

## implications

- for  $x, y \in S$ ,

$$f(y) \geq f(x) + \nabla f(x)^T (y - x) + \frac{m}{2} \|x - y\|_2^2$$

hence,  $S$  is bounded

$$y = x - \frac{1}{m} \nabla f(x)$$

- $p^* > -\infty$ , and for  $x \in S$ ,

$$f(x) - p^* \leq \frac{1}{2m} \|\nabla f(x)\|_2^2$$

useful as stopping criterion (if you know  $m$ )

# Descent methods

$$x^{(k+1)} = x^{(k)} + t^{(k)} \Delta x^{(k)} \quad \text{with } f(x^{(k+1)}) < f(x^{(k)})$$

- other notations:  $x^+ = x + t\Delta x$ ,  $x := x + t\Delta x$
- $\Delta x$  is the *step*, or *search direction*;  $t$  is the *step size*, or *step length*
- from convexity,  $f(x^+) < f(x)$  implies  $\nabla f(x)^T \Delta x < 0$   
(*i.e.*,  $\Delta x$  is a *descent direction*)

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*General descent method.*

**given** a starting point  $x \in \text{dom } f$ .

**repeat**

1. Determine a descent direction  $\Delta x$ .
2. *Line search.* Choose a step size  $t > 0$ .
3. *Update.*  $x := x + t\Delta x$ .

**until** stopping criterion is satisfied.

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# Line search types

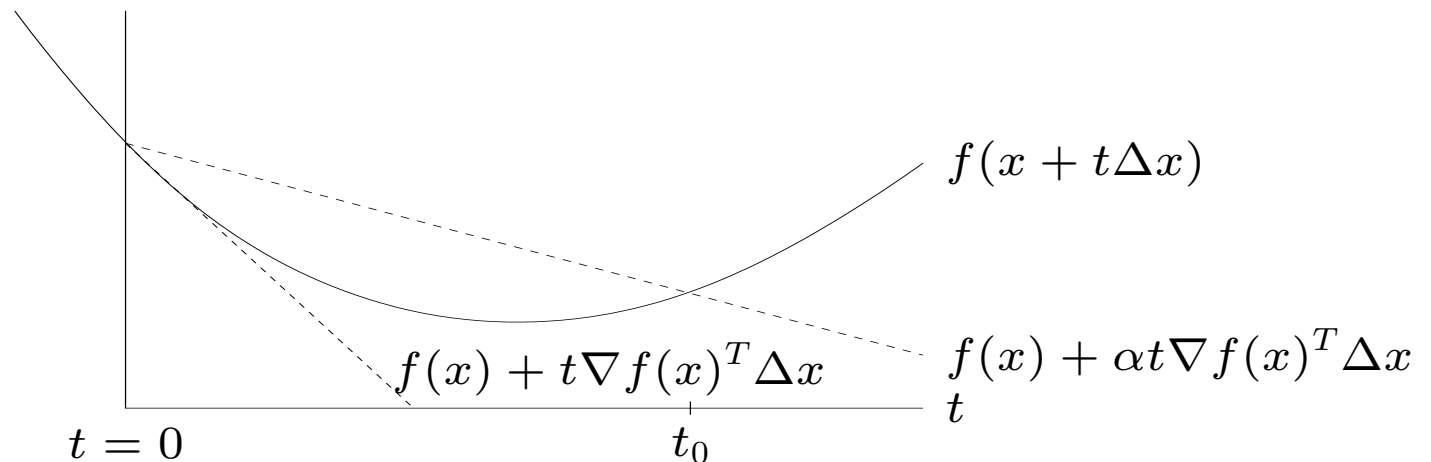
**exact line search:**  $t = \operatorname{argmin}_{t>0} f(x + t\Delta x)$

**backtracking line search** (with parameters  $\alpha \in (0, 1/2)$ ,  $\beta \in (0, 1)$ )

- starting at  $t = 1$ , repeat  $t := \beta t$  until

$$f(x + t\Delta x) < f(x) + \alpha t \nabla f(x)^T \Delta x$$

- graphical interpretation: backtrack until  $t \leq t_0$



# Gradient descent method

general descent method with  $\Delta x = -\nabla f(x)$

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**given** a starting point  $x \in \text{dom } f$ .

**repeat**

1.  $\Delta x := -\nabla f(x)$ .
2. *Line search*. Choose step size  $t$  via exact or backtracking line search.
3. *Update*.  $x := x + t\Delta x$ .

**until** stopping criterion is satisfied.

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- stopping criterion usually of the form  $\|\nabla f(x)\|_2 \leq \epsilon$

- convergence result: for strongly convex  $f$ ,  
$$f(x^{(k)}) - p^* \leq c^k (f(x^{(0)}) - p^*)$$
  
$$m \leq \nabla^2 f \leq M$$
  
$$c = 1 - \frac{m}{M}$$

$c \in (0, 1)$  depends on  $m$ ,  $x^{(0)}$ , line search type

- very simple, but often very slow; rarely used in practice

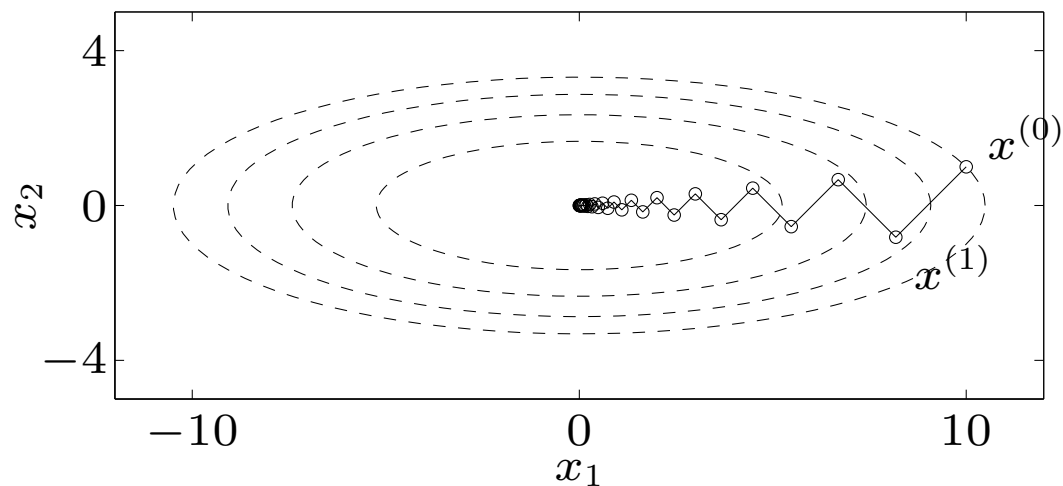
## quadratic problem in $\mathbf{R}^2$

$$f(x) = (1/2)(x_1^2 + \gamma x_2^2) \quad (\gamma > 0)$$

with exact line search, starting at  $x^{(0)} = (\gamma, 1)$ :

$$x_1^{(k)} = \gamma \left( \frac{\gamma - 1}{\gamma + 1} \right)^k, \quad x_2^{(k)} = \left( -\frac{\gamma - 1}{\gamma + 1} \right)^k$$

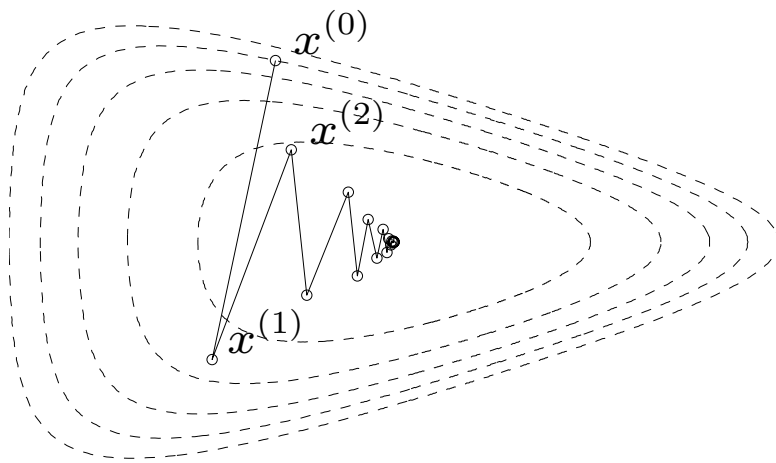
- very slow if  $\gamma \gg 1$  or  $\gamma \ll 1$
- example for  $\gamma = 10$ :



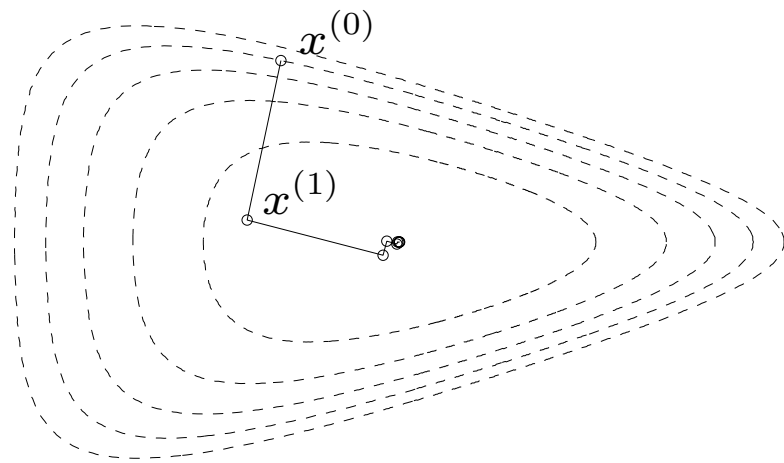


## nonquadratic example

$$f(x_1, x_2) = e^{x_1+3x_2-0.1} + e^{x_1-3x_2-0.1} + e^{-x_1-0.1}$$



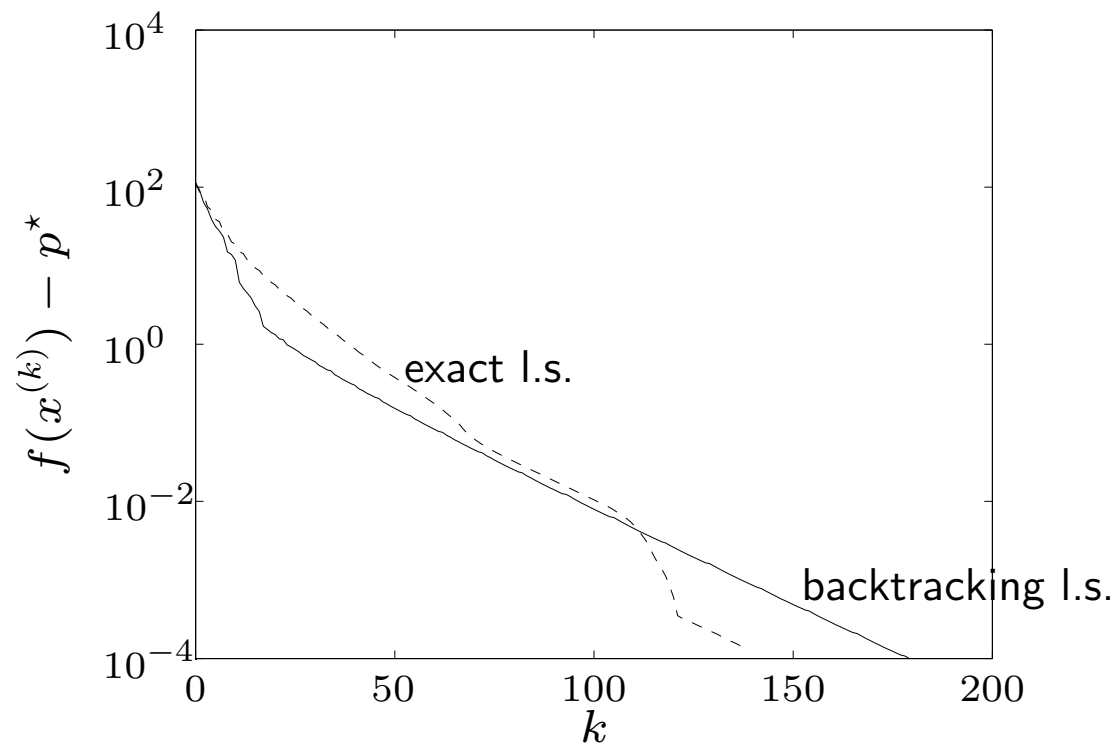
backtracking line search



exact line search

a problem in  $\mathbf{R}^{100}$

$$f(x) = c^T x - \sum_{i=1}^{500} \log(b_i - a_i^T x)$$



‘linear’ convergence, *i.e.*, a straight line on a semilog plot

# Steepest descent method

**normalized steepest descent direction** (at  $x$ , for norm  $\|\cdot\|$ ):

$$\Delta x_{\text{nsd}} = \operatorname{argmin}\{\nabla f(x)^T v \mid \|v\| = 1\}$$

interpretation: for small  $v$ ,  $f(x+v) \approx f(x) + \nabla f(x)^T v$ ;

direction  $\Delta x_{\text{nsd}}$  is unit-norm step with most negative directional derivative

**(unnormalized) steepest descent direction**

$$\Delta x_{\text{sd}} = \|\nabla f(x)\|_* \Delta x_{\text{nsd}}$$

satisfies  $\nabla f(x)^T \Delta x_{\text{sd}} = -\|\nabla f(x)\|_*^2$

$$\Delta x_{\text{sd}} = \operatorname{argmin}_v \left( \nabla f(x)^T v + \frac{1}{2} \|v\|^2 \right)$$

**steepest descent method**

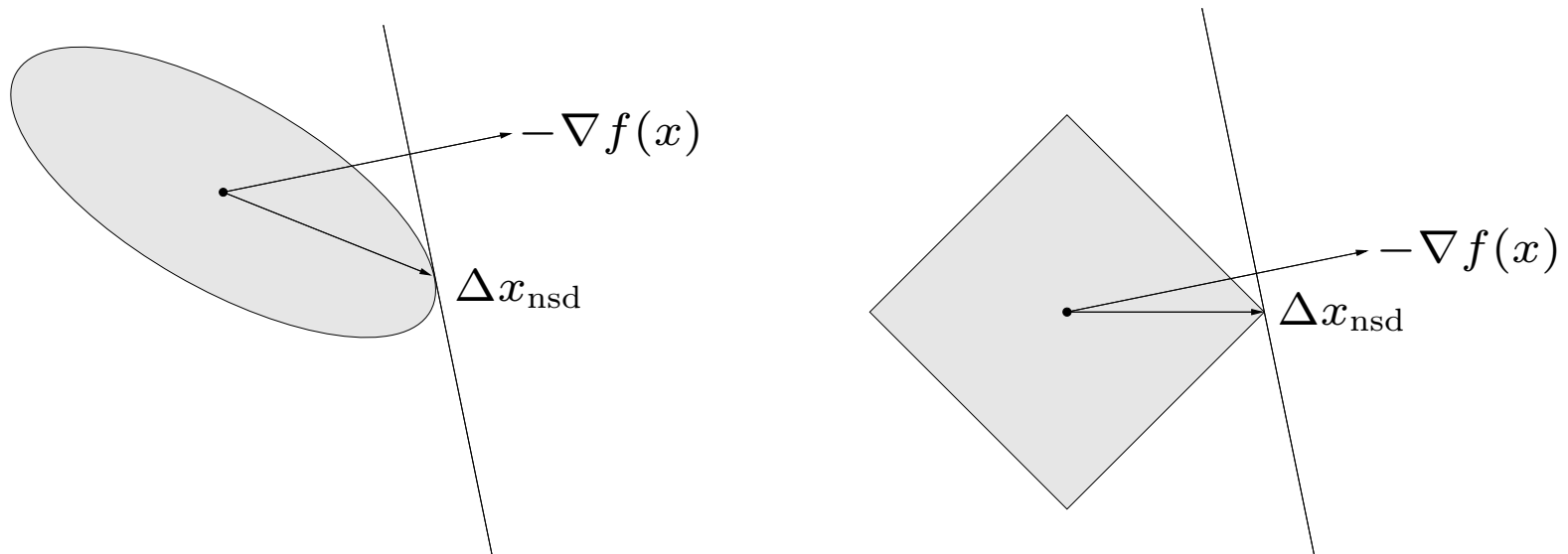
$$\|\Delta x_{\text{sd}}\| = \|\nabla f(x)\|_*$$

- general descent method with  $\Delta x = \Delta x_{\text{sd}}$
- convergence properties similar to gradient descent

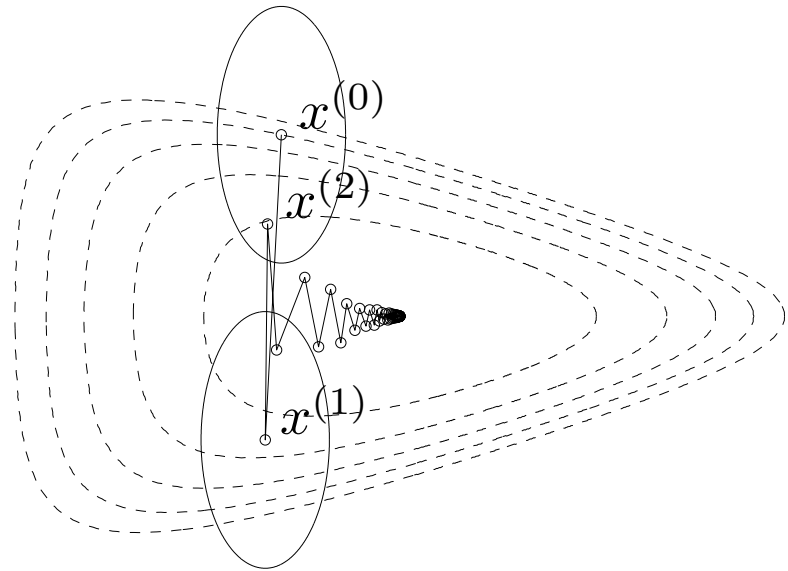
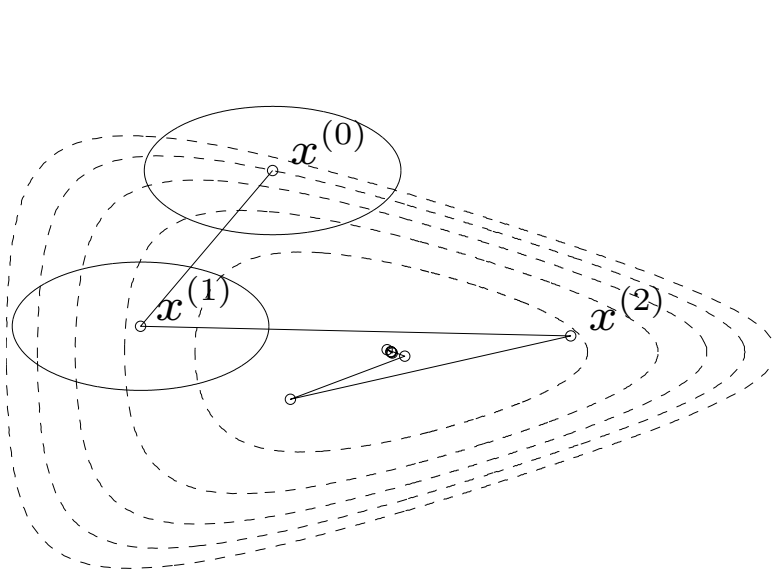
## examples

- Euclidean norm:  $\Delta x_{\text{sd}} = -\nabla f(x)$
- quadratic norm  $\|x\|_P = (x^T P x)^{1/2}$  ( $P \in \mathbf{S}_{++}^n$ ):  $\Delta x_{\text{sd}} = -P^{-1} \nabla f(x)$
- $\ell_1$ -norm:  $\Delta x_{\text{sd}} = -(\partial f(x)/\partial x_i) e_i$ , where  $|\partial f(x)/\partial x_i| = \|\nabla f(x)\|_\infty$

unit balls and normalized steepest descent directions for a quadratic norm and the  $\ell_1$ -norm:



## choice of norm for steepest descent



- steepest descent with backtracking line search for two quadratic norms
- ellipses show  $\{x \mid \|x - x^{(k)}\|_P = 1\}$
- equivalent interpretation of steepest descent with quadratic norm  $\|\cdot\|_P$ : gradient descent after change of variables  $\bar{x} = P^{1/2}x$

shows choice of  $P$  has strong effect on speed of convergence

# Newton step

$$\Delta x_{\text{nt}} = -\nabla^2 f(x)^{-1} \nabla f(x)$$

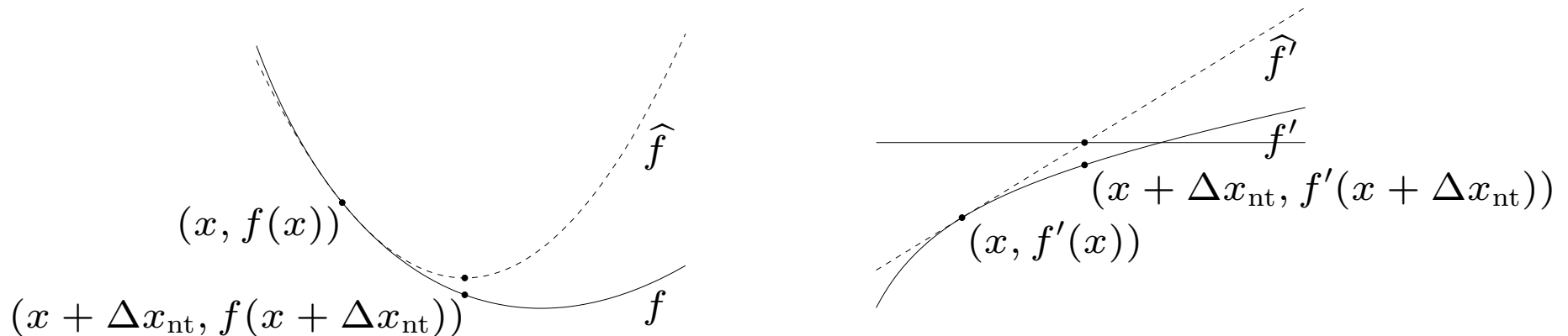
## interpretations

- $x + \Delta x_{\text{nt}}$  minimizes second order approximation

$$\hat{f}(x + v) = f(x) + \nabla f(x)^T v + \frac{1}{2} v^T \nabla^2 f(x) v$$

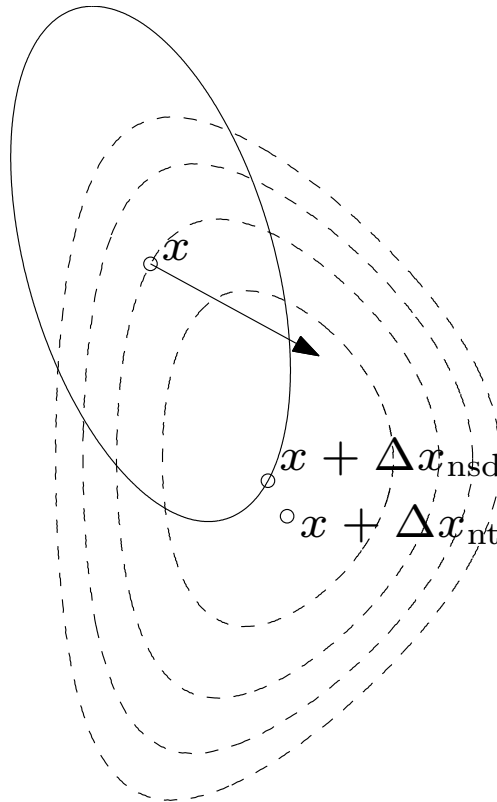
- $x + \Delta x_{\text{nt}}$  solves linearized optimality condition

$$\nabla f(x + v) \approx \nabla \hat{f}(x + v) = \nabla f(x) + \nabla^2 f(x) v = 0$$



- $\Delta x_{\text{nt}}$  is steepest descent direction at  $x$  in local Hessian norm

$$\|u\|_{\nabla^2 f(x)} = (u^T \nabla^2 f(x) u)^{1/2}$$



dashed lines are contour lines of  $f$ ; ellipse is  $\{x + v \mid v^T \nabla^2 f(x) v = 1\}$

arrow shows  $-\nabla f(x)$

# Newton decrement

$$\lambda(x) = \left( \nabla f(x)^T \nabla^2 f(x)^{-1} \nabla f(x) \right)^{1/2}$$

a measure of the proximity of  $x$  to  $x^*$

## properties

- gives an estimate of  $f(x) - p^*$ , using quadratic approximation  $\hat{f}$ :

$$f(x) - \inf_y \hat{f}(y) = \frac{1}{2} \lambda(x)^2$$

- equal to the norm of the Newton step in the quadratic Hessian norm

$$\lambda(x) = \left( \Delta x_{\text{nt}}^T \nabla^2 f(x) \Delta x_{\text{nt}} \right)^{1/2}$$

- directional derivative in the Newton direction:  $\nabla f(x)^T \Delta x_{\text{nt}} = -\lambda(x)^2$
- affine invariant (unlike  $\|\nabla f(x)\|_2$ )



# Newton's method

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**given** a starting point  $x \in \text{dom } f$ , tolerance  $\epsilon > 0$ .

**repeat**

1. *Compute the Newton step and decrement.*

$$\Delta x_{\text{nt}} := -\nabla^2 f(x)^{-1} \nabla f(x); \quad \lambda^2 := \nabla f(x)^T \nabla^2 f(x)^{-1} \nabla f(x).$$

2. *Stopping criterion.* **quit** if  $\lambda^2/2 \leq \epsilon$ .

3. *Line search.* Choose step size  $t$  by backtracking line search.

4. *Update.*  $x := x + t\Delta x_{\text{nt}}$ .

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affine invariant, *i.e.*, independent of linear changes of coordinates:

Newton iterates for  $\tilde{f}(y) = f(Ty)$  with starting point  $y^{(0)} = T^{-1}x^{(0)}$  are

$$y^{(k)} = T^{-1}x^{(k)}$$

# Classical convergence analysis

## assumptions

- $f$  strongly convex on  $S$  with constant  $m$
- $\nabla^2 f$  is Lipschitz continuous on  $S$ , with constant  $L > 0$ :

$$\|\nabla^2 f(x) - \nabla^2 f(y)\|_2 \leq L\|x - y\|_2$$

( $L$  measures how well  $f$  can be approximated by a quadratic function)

**outline:** there exist constants  $\eta \in (0, m^2/L)$ ,  $\gamma > 0$  such that

- if  $\|\nabla f(x)\|_2 \geq \eta$ , then  $f(x^{(k+1)}) - f(x^{(k)}) \leq -\gamma$
- if  $\|\nabla f(x)\|_2 < \eta$ , then

$$\frac{L}{2m^2} \|\nabla f(x^{(k+1)})\|_2 \leq \left( \frac{L}{2m^2} \|\nabla f(x^{(k)})\|_2 \right)^2$$

### **damped Newton phase** ( $\|\nabla f(x)\|_2 \geq \eta$ )

- most iterations require backtracking steps
- function value decreases by at least  $\gamma$
- if  $p^* > -\infty$ , this phase ends after at most  $(f(x^{(0)}) - p^*)/\gamma$  iterations

### **quadratically convergent phase** ( $\|\nabla f(x)\|_2 < \eta$ )

- all iterations use step size  $t = 1$
- $\|\nabla f(x)\|_2$  converges to zero quadratically: if  $\|\nabla f(x^{(k)})\|_2 < \eta$ , then

$$\frac{L}{2m^2} \|\nabla f(x^l)\|_2 \leq \left( \frac{L}{2m^2} \|\nabla f(x^k)\|_2 \right)^{2^{l-k}} \leq \left( \frac{1}{2} \right)^{2^{l-k}}, \quad l \geq k$$

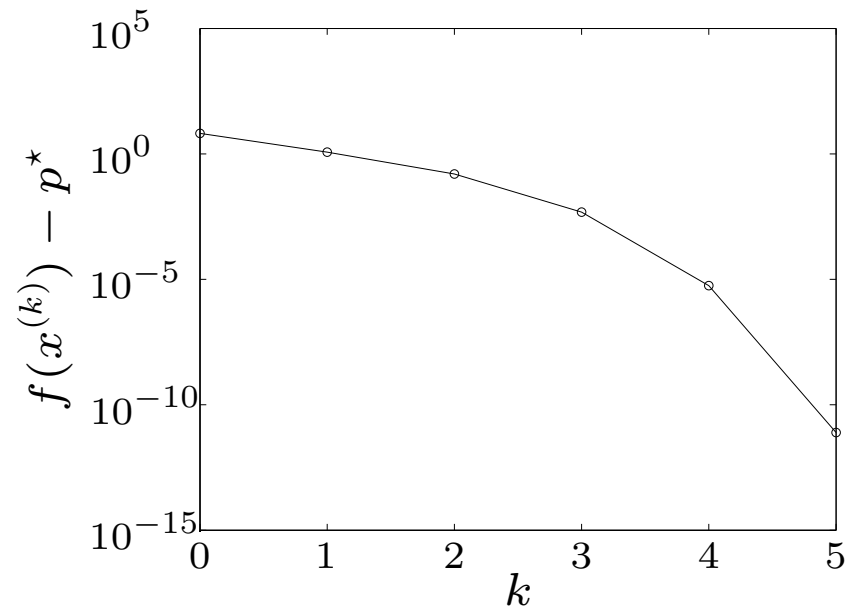
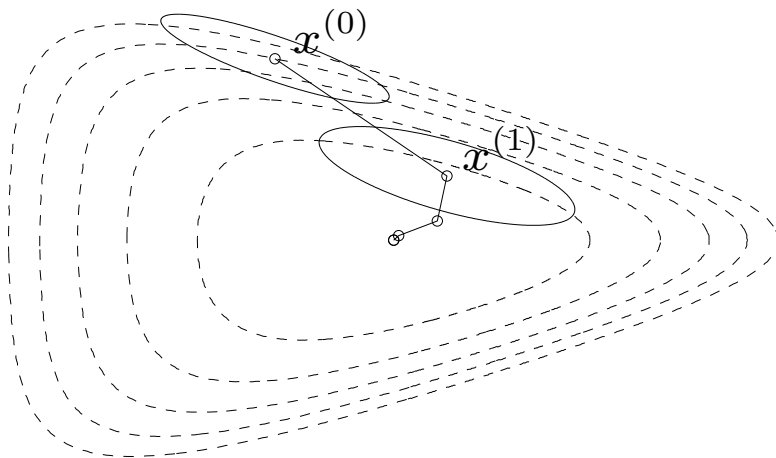
**conclusion:** number of iterations until  $f(x) - p^* \leq \epsilon$  is bounded above by

$$\frac{f(x^{(0)}) - p^*}{\gamma} + \log_2 \log_2(\epsilon_0/\epsilon)$$

- $\gamma, \epsilon_0$  are constants that depend on  $m, L, x^{(0)}$
- second term is small (of the order of 6) and almost constant for practical purposes
- in practice, constants  $m, L$  (hence  $\gamma, \epsilon_0$ ) are usually unknown
- provides qualitative insight in convergence properties (*i.e.*, explains two algorithm phases)

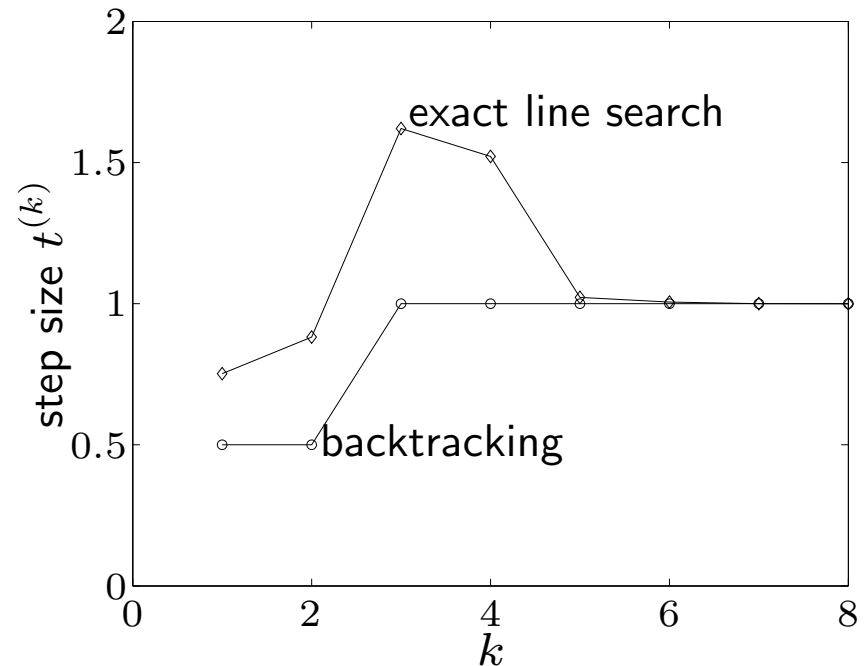
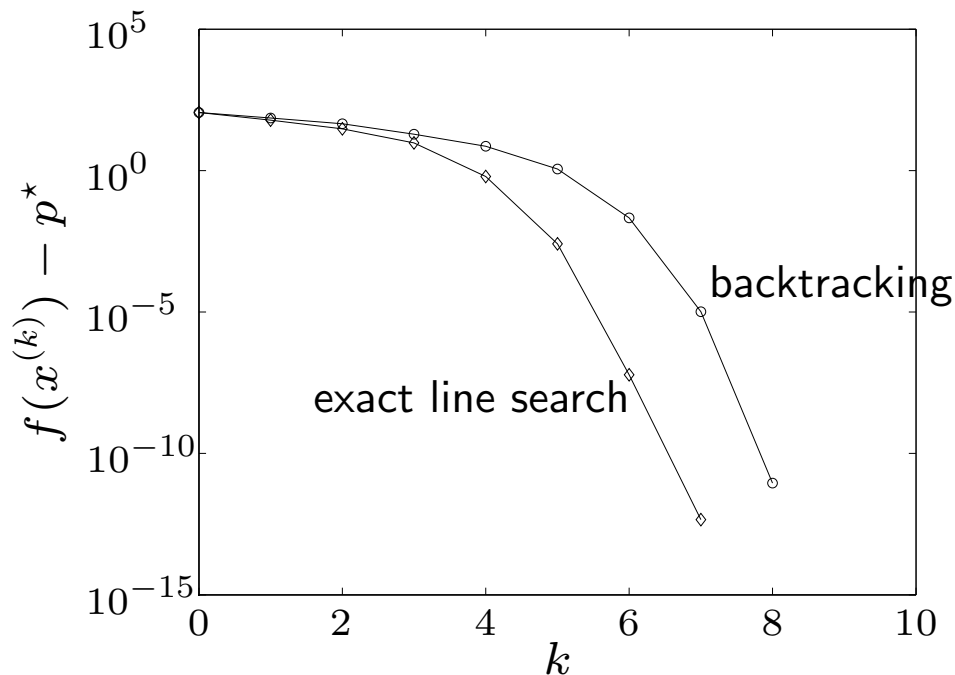
# Examples

example in  $\mathbf{R}^2$  (page 10–9)



- backtracking parameters  $\alpha = 0.1$ ,  $\beta = 0.7$
- converges in only 5 steps
- quadratic local convergence

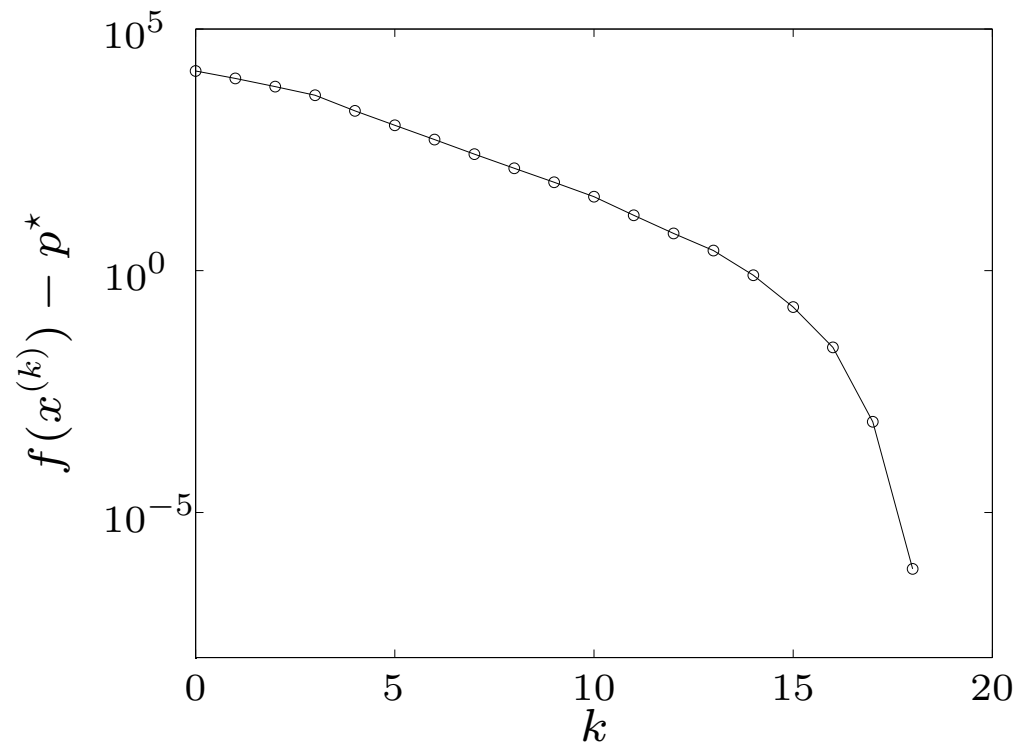
## example in $\mathbf{R}^{100}$ (page 10–10)



- backtracking parameters  $\alpha = 0.01$ ,  $\beta = 0.5$
- backtracking line search almost as fast as exact l.s. (and much simpler)
- clearly shows two phases in algorithm

example in  $\mathbf{R}^{10000}$  (with sparse  $a_i$ )

$$f(x) = - \sum_{i=1}^{10000} \log(1 - x_i^2) - \sum_{i=1}^{100000} \log(b_i - a_i^T x)$$



- backtracking parameters  $\alpha = 0.01$ ,  $\beta = 0.5$ .
- performance similar as for small examples

# Self-concordance

## shortcomings of classical convergence analysis

- depends on unknown constants  $(m, L, \dots)$
- bound is not affinely invariant, although Newton's method is

## convergence analysis via self-concordance (Nesterov and Nemirovski)

- does not depend on any unknown constants
- gives affine-invariant bound
- applies to special class of convex functions ('self-concordant' functions)
- developed to analyze polynomial-time interior-point methods for convex optimization



# Self-concordant functions

## definition

- convex  $f : \mathbf{R} \rightarrow \mathbf{R}$  is self-concordant if  $|f'''(x)| \leq 2f''(x)^{3/2}$  for all  $x \in \text{dom } f$
- $f : \mathbf{R}^n \rightarrow \mathbf{R}$  is self-concordant if  $g(t) = f(x + tv)$  is self-concordant for all  $x \in \text{dom } f$ ,  $v \in \mathbf{R}^n$

## examples on $\mathbf{R}$

- linear and quadratic functions
- negative logarithm  $f(x) = -\log x$
- negative entropy plus negative logarithm:  $f(x) = x \log x - \log x$

**affine invariance:** if  $f : \mathbf{R} \rightarrow \mathbf{R}$  is s.c., then  $\tilde{f}(y) = f(ay + b)$  is s.c.:

$$\tilde{f}'''(y) = a^3 f'''(ay + b), \quad \tilde{f}''(y) = a^2 f''(ay + b)$$

# Self-concordant calculus

## properties

- preserved under positive scaling  $\alpha \geq 1$ , and sum
- preserved under composition with affine function
- if  $g$  is convex with  $\text{dom } g = \mathbf{R}_{++}$  and  $|g'''(x)| \leq 3g''(x)/x$  then

$$f(x) = \log(-g(x)) - \log x$$

is self-concordant

**examples:** properties can be used to show that the following are s.c.

- $f(x) = -\sum_{i=1}^m \log(b_i - a_i^T x)$  on  $\{x \mid a_i^T x < b_i, i = 1, \dots, m\}$
- $f(X) = -\log \det X$  on  $\mathbf{S}_{++}^n$
- $f(x) = -\log(y^2 - x^T x)$  on  $\{(x, y) \mid \|x\|_2 < y\}$

# Convergence analysis for self-concordant functions

**summary:** there exist constants  $\eta \in (0, 1/4]$ ,  $\gamma > 0$  such that

- if  $\lambda(x) > \eta$ , then

$$f(x^{(k+1)}) - f(x^{(k)}) \leq -\gamma$$

- if  $\lambda(x) \leq \eta$ , then

$$2\lambda(x^{(k+1)}) \leq \left(2\lambda(x^{(k)})\right)^2$$

( $\eta$  and  $\gamma$  only depend on backtracking parameters  $\alpha, \beta$ )

**complexity bound:** number of Newton iterations bounded by

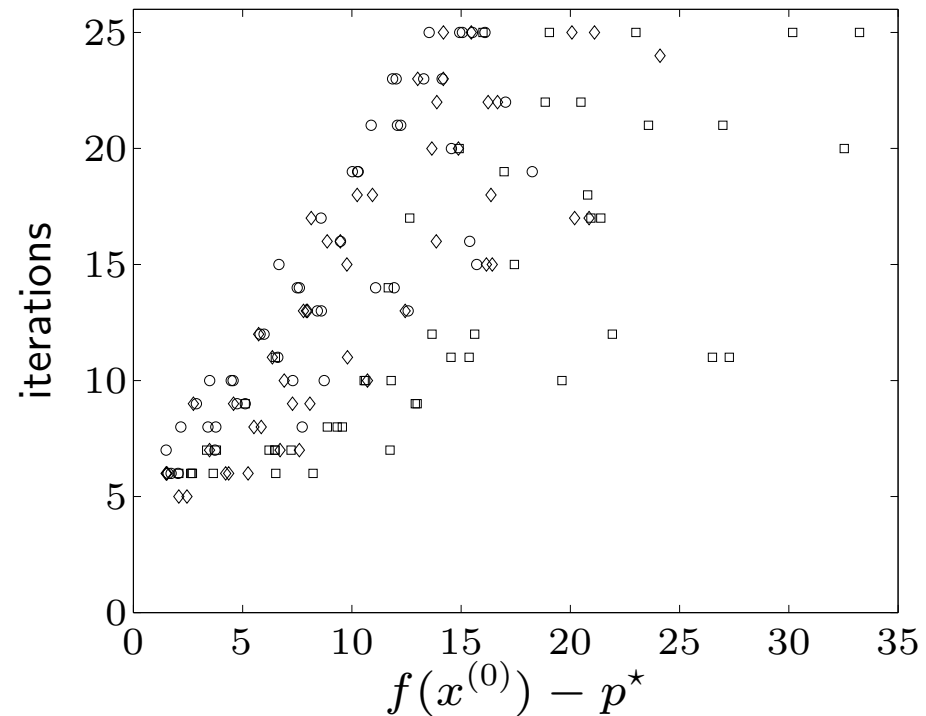
$$\frac{f(x^{(0)}) - p^*}{\gamma} + \log_2 \log_2(1/\epsilon)$$

for  $\alpha = 0.1$ ,  $\beta = 0.8$ ,  $\epsilon = 10^{-10}$ , bound evaluates to  $375(f(x^{(0)}) - p^*) + 6$

**numerical example:** 150 randomly generated instances of

$$\text{minimize } f(x) = -\sum_{i=1}^m \log(b_i - a_i^T x)$$

○:  $m = 100, n = 50$   
□:  $m = 1000, n = 500$   
◇:  $m = 1000, n = 50$



- number of iterations much smaller than  $375(f(x^{(0)}) - p^*) + 6$
- bound of the form  $c(f(x^{(0)}) - p^*) + 6$  with smaller  $c$  (empirically) valid

# Implementation

main effort in each iteration: evaluate derivatives and solve Newton system

$$H\Delta x = -g$$

where  $H = \nabla^2 f(x)$ ,  $g = \nabla f(x)$

**via Cholesky factorization**

$$H = LL^T, \quad \Delta x_{\text{nt}} = -L^{-T}L^{-1}g, \quad \lambda(x) = \|L^{-1}g\|_2$$

- cost  $(1/3)n^3$  flops for unstructured system
- cost  $\ll (1/3)n^3$  if  $H$  sparse, banded

## example of dense Newton system with structure

$$f(x) = \sum_{i=1}^n \psi_i(x_i) + \psi_0(Ax + b), \quad H = D + A^T H_0 A$$

- assume  $A \in \mathbf{R}^{p \times n}$ , dense, with  $p \ll n$
- $D$  diagonal with diagonal elements  $\psi_i''(x_i)$ ;  $H_0 = \nabla^2 \psi_0(Ax + b)$

**method 1:** form  $H$ , solve via dense Cholesky factorization: (cost  $(1/3)n^3$ )

**method 2** (page 9–15): factor  $H_0 = L_0 L_0^T$ ; write Newton system as

$$D\Delta x + A^T L_0 w = -g, \quad L_0^T A \Delta x - w = 0$$

eliminate  $\Delta x$  from first equation; compute  $w$  and  $\Delta x$  from

$$(I + L_0^T A D^{-1} A^T L_0)w = -L_0^T A D^{-1} g, \quad D\Delta x = -g - A^T L_0 w$$

cost:  $2p^2n$  (dominated by computation of  $L_0^T A D^{-1} A^T L_0$ )