

bits and pieces

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References



Christoph Stickse, *Efficient equational reasoning for the Inst-Gen framework*, Ph.D. thesis, School of Computer Science, University of Manchester, 2011.

*Hofstadter's Law: It always takes longer than you expect,
even when you take into account Hofstadter's Law.*

— Douglas Hofstadter, Gödel, Escher, Bach: An Eternal Golden Braid

Definition (Ordered Resolution)

$$\frac{L \vee C \quad \neg L' \vee D}{(C \vee D)\sigma}$$

where

$L\sigma$ strictly maximal in $C\sigma$, $\neg L'\sigma$ maximal in $D\sigma$, $\sigma = \text{mgu}(L, L')$.

Definition (Inst-Gen)

$$\frac{L \vee C \quad \neg L' \vee D}{(L \vee C)\sigma \quad (\neg L' \vee D)\sigma}$$

where

$$\text{sel}(L \vee C) = L \quad \text{sel}(\neg L' \vee D) = \neg L' \quad \sigma = \text{mgu}(L, L')$$

Example (Unsatisfiable set of clauses)

$$S = \{P(x) \vee \neg P(y), \neg P(a), P(b)\}$$

Example (Resolution)

$$\frac{\frac{P(x) \vee \neg P(y) \quad \neg P(a)}{\neg P(y) \quad P(b)} \quad x \mapsto a}{\square} \quad y \mapsto b$$

Example (Inst-Gen)

$$S_0 \perp = \{P(\perp) \vee \neg P(\perp), \neg P(a), P(b)\} \quad (\text{satisfiable})$$

$$\frac{P(x) \vee \neg P(y) \quad \neg P(a)}{P(a) \vee \neg P(y)} \quad x \mapsto a$$

$$S_1 \perp \supsetneq \{\neg P(a), P(b), P(a) \vee \neg P(\perp)\} \quad (\text{satisfiable})$$

$$\frac{P(b) \quad P(a) \vee \neg P(y)}{P(a) \vee P(b)} \quad y \mapsto b$$

$$S_2 \perp \supsetneq \{\neg P(a), P(b), P(a) \vee \neg P(b)\} \quad (\text{unsatisfiable})$$

$$x = a \vee x \neq a$$

$$f(a) \neq f(b)$$

$R = \{x = a\}$ is ground complete

$$\sigma = \{x \mapsto b\} \quad (x = a)\sigma = a \rightarrow b \text{ with } a > b \quad f(a) \neq b$$

$$\perp = \perp \vee \perp \neq \perp$$

$$f(\perp) \neq f(a)$$

$$P(a), \neg P(f(a, b)), f(x, b) = x$$

$$P(a), \neg P(f(a, b)), f(\perp, b) = \perp$$

$\{f(x, b) = x\}$ is ground complete and with $\{x \mapsto a\}$ we get $\neg P(a)$