flea

system description

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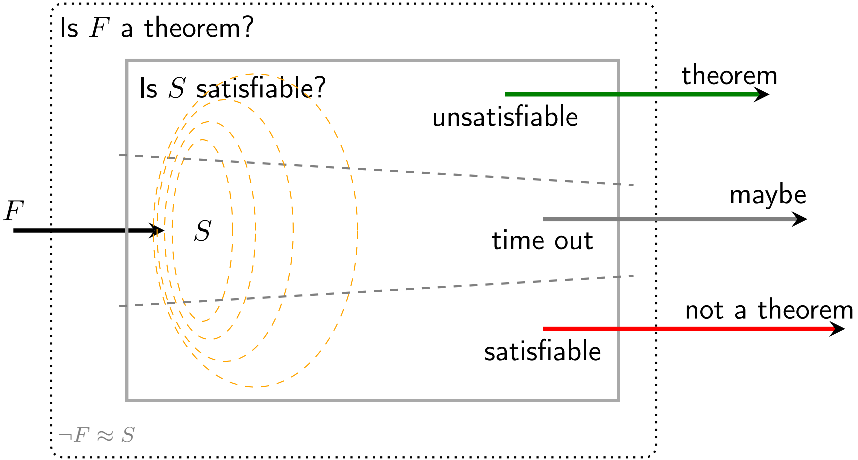
Refutation



Is F a theorem?
Is S satisfiable?
F
S
unsatisfiable
theorem
time out
maybe
satisfiable
not a theorem
$\neg F \approx S$

Goals

fisrst order logic with equality theorem attester

- maintainable
- simple
- tested
- fast
- extendable
- flexible

1. start with
    - empty list of processed clauses with selected literals
    - list of unprocessed clauses
    - empty (satisfiable) yices context
2. return satisfiable if the list of unprocessed clauses is empty
3. select a unprocessed file and process it
    - assert a ground instance of the clause in the yices context
    - return unsatisfiable if the context is unsatisfiable
    - retrieve a model from the context
    - select a literal of the clause that holds in the model
    - search for contradictions with the selected literals of the processed clauses and retrieve non-redundant instances from the contributing clauses and add them to the unprocessed clauses
    - remove the clause from the list of unprocessed clauses
    - add the clause and its selected literal to the processed clauses
4. continue with step 2.

$$S_0\bot = \{\mathsf{P}(\bot) \vee \neg\mathsf{P}(\bot), \neg\mathsf{P}(\mathsf{a}), \mathsf{P}(\mathsf{b})\}$$ satisfiable

$$\frac{\mathsf{P}(x) \vee \neg\mathsf{P}(y)) \quad \neg\mathsf{P}(\mathsf{a})}{\mathsf{P}(\mathsf{a}) \vee \neg\mathsf{P}(y)}\; x \mapsto \mathsf{a}$$

$$S_1\bot \supsetneq \{\neg\mathsf{P}(\mathsf{a}), \mathsf{P}(\mathsf{b}), \mathsf{P}(\mathsf{a}) \vee \neg\mathsf{P}(\bot)\}$$ satisfiable

$$\frac{\mathsf{P}(\mathsf{b}) \quad \mathsf{P}(\mathsf{a}) \vee \neg\mathsf{P}(y)}{\mathsf{P}(\mathsf{a}) \vee \mathsf{P}(\mathsf{b})}\; y \mapsto \mathsf{b}$$

$$S_2\bot \supsetneq \{\neg\mathsf{P}(\mathsf{a}), \mathsf{P}(\mathsf{b}), \mathsf{P}(\mathsf{a}) \vee \neg\mathsf{P}(\mathsf{b})\}$$ unsatisfiable

$$S = \{\mathsf{P}(x) \vee \neg\mathsf{P}(y), \neg\mathsf{P}(\mathsf{a}), \mathsf{P}(\mathsf{b})\}$$

$$\dfrac{\dfrac{\mathsf{P}(x) \vee \neg\mathsf{P}(y) \qquad \neg\mathsf{P}(\mathsf{a})}{\neg\mathsf{P}(y)}\; x \mapsto \mathsf{a} \qquad \mathsf{P}(\mathsf{b})}{\square}\; y \mapsto \mathsf{b}$$

$$\frac{L \vee C \quad \neg L' \vee D}{(L \vee C)\sigma \quad (\neg L' \vee D)\sigma}$$

where

$$\mathrm{sel}(L \vee C) = L \qquad \mathrm{sel}(\neg L' \vee D) = \neg L' \qquad \sigma = \mathrm{mgu}(L, L')$$

$$\frac{L \vee C \quad \neg L' \vee D}{(C \vee D)\sigma}$$

where

$L\sigma$ strictly maximal in $C\sigma$, $\neg L'\sigma$ maximal in $D\sigma$, $\sigma = \mathrm{mgu}(L, L')$.

Unit Superposition Inference Rules

$$\frac{s \approx t \quad L[s']}{(L[t]) \cdot \sigma} \text{ unit paramodulation}$$

where $\sigma = \mathrm{mgu}(s, s')$, $s' \notin \mathcal{V}$, $t\sigma \not\succeq s\sigma$

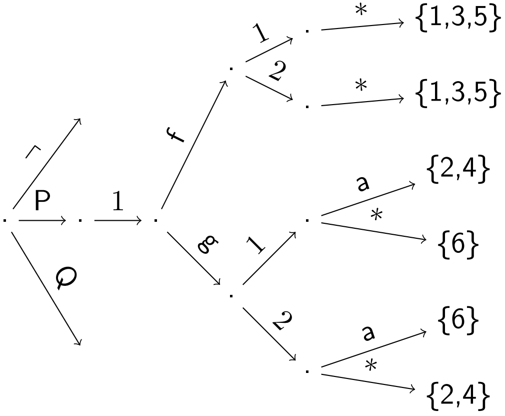
$$\frac{s \approx t \quad u[s'] \not\approx v}{(u[t] \not\approx v) \cdot \sigma} \text{ unit superposition } \frac{s \approx t \quad u[s'] \approx v}{(u[t] \approx v) \cdot \sigma}$$

where $\sigma = \mathrm{mgu}(s, s')$, $s' \notin \mathcal{V}$, $t\sigma \not\succeq s\sigma$, $v\sigma \not\succeq u[s']\sigma$

$$\frac{s \not\approx t}{\square} \text{ unit equality resolution} \qquad \frac{A \quad \neg B}{\square} \text{ unit resolution}$$

where s and t (A and B respectively) are unifiable

$$\{{}^{1:}\mathsf{P}(\mathsf{f}(x,x)), {}^{2:}\mathsf{P}(\mathsf{g}(\mathsf{a},x)), {}^{3:}\mathsf{P}(\mathsf{f}(y,z)), {}^{4:}\mathsf{P}(\mathsf{g}(\mathsf{a},y)), {}^{5:}\mathsf{P}(\mathsf{f}(y,x)), {}^{6:}\mathsf{P}(\mathsf{g}(y,a))\}$$



$$\neg\mathsf{P}(\mathsf{g}(\mathsf{b},z)) \mapsto \{\mathsf{P}.1.\mathsf{g}.1.\mathsf{b}, \mathsf{P}.1.\mathsf{g}.2.*\} \mapsto \{6\} \cap \{2,4,6\}$$