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Overview

- Z3 is a Satisfiability Modulo Theories (SMT) solver.
- Z3 integrates several decision procedures.
- Z3 is used in several program analysis, verification, testcase generation projects at Microsoft.
- Z3 1.2 is freely available for academic research:
 - http://research.microsoft.com/projects/z3



Satisfiability Modulo Theories (SMT)

$$x+2=y \Rightarrow f(read(write(a,x,3),y-2)) = f(y-x+1)$$

Arithmetic

Array Theory

Uninterpreted Functions

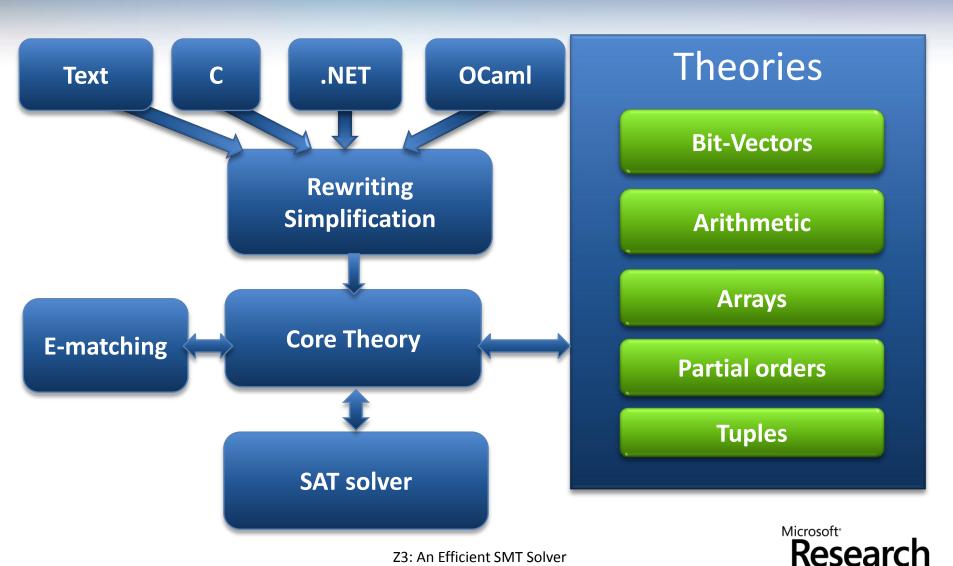


Main features

- Linear real and integer arithmetic.
- Fixed-size bit-vectors
- Uninterpreted functions
- Extensional arrays
- Quantifiers
- Model generation
- Several input formats (Simplify, SMT-LIB, Z3, Dimacs)
- Extensive API (C/C++, .Net, OCaml)

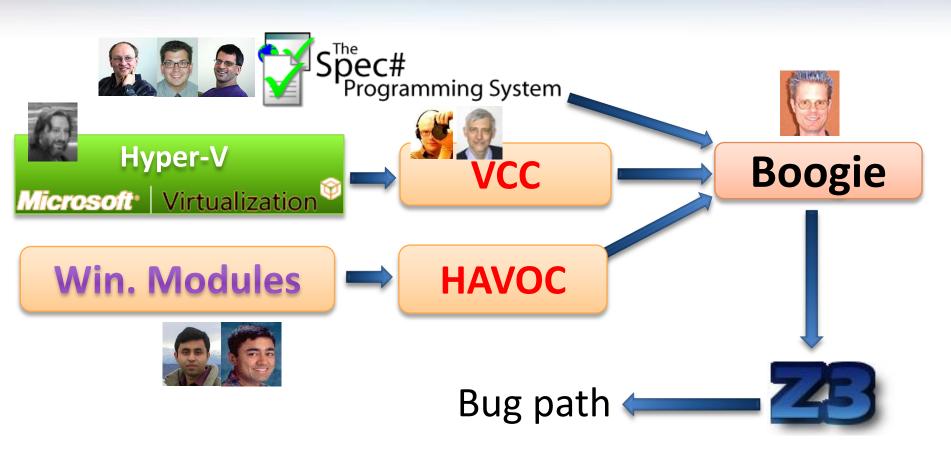


Z3: Core System Components



Z3: An Efficient SMT Solver

Clients: Program Verification



Rustan Leino, Mike Barnet, Michal Moskal, Shaz Qadeer, Shuvendu Lahiri, Herman Venter, Peter Muller, Wolfram Schulte, Ernie Cohen

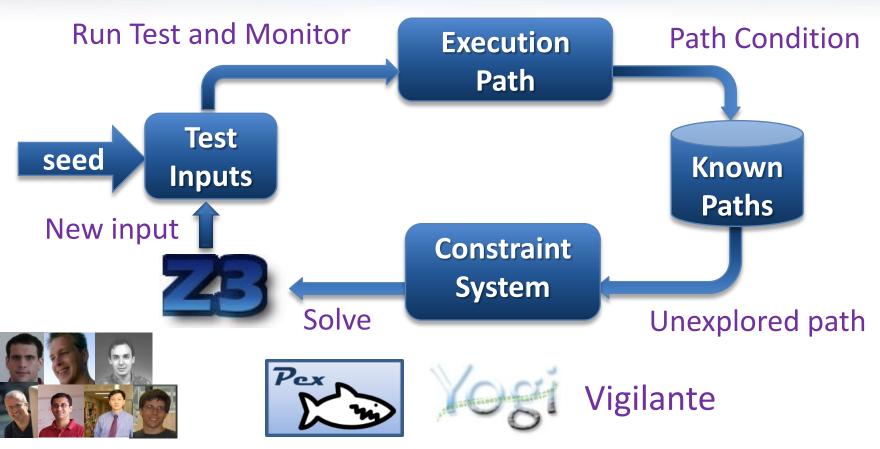
Research

Z3 & Program Verification

- Quantifiers, quantifiers, quantifiers, ...
 - Modeling the runtime
 - Frame axioms ("what didn't change")
 - Users provided assertions (e.g., the array is sorted)
 - Prototyping decision procedures (e.g., reachability, heaps, ...)
- Solver must be fast in satisfiable instances.
- Trade-off between precision and performance.
- Candidate (Potential) Models



Clients: Test case generation



Nikolai Tillmann, Peli de Halleux, Patrice Godefroid Aditya Nori, Jean Philippe Martin, Miguel Castro, Manuel Costa, Lintao Zhang

Research

Z3 & Test case generation

- Formulas may be a big conjunction
 - Pre-processing step
 - Eliminate variables and simplify input format
- Incremental: solve several similar formulas
 - New constraints are asserted.
 - push and pop: (user) backtracking
 - Lemma reuse
- "Small Models"
 - Given a formula F, find a model M, that minimizes the value of the variables $x_0 \dots x_n$



Client: Static Driver Verifier

- Z3 is part of SDV 2.0 (Windows 7)
- It is used for:
 - Predicate abstraction (c2bp)
 - Counterexample refinement (newton)







Ella Bounimova, Vlad Levin, Jakob Lichtenberg, Tom Ball, Sriram Rajamani, Byron Cook



Z3 & Static Driver Verifier

- All-SAT
 - Fast Predicate Abstraction
- Unsatisfiable cores
 - Why the abstract path is not feasible?



More Microsoft clients

Bounded model-checking of model programs



- Termination
- 9
- Security protocols



Business application modeling



Cryptography



- Model Based Testing (SQL-Server)
- Your killer-application here



Some Technical goodies

- Model-based Theory Combination
 - How to efficiently combine theory solvers?
 - Use models to control Theory Combination.
- E-matching abstract machine
 - Term indexing data-structures for incremental matching modulo equalities.
- Relevancy propagation
 - Use Tableau advantages with DPLL engine



Example: CAPI

```
Given arrays:
for (n = 2; n \le 5; n++) {
   printf("n = %d\n", n);
   ctx = Z3 mk context(cfg);
                                                                     bool a1[bool];
   bool type = Z3 mk bool type(ctx);
                                                                     bool a2[bool];
   array type = Z3 mk array type(ctx, bool type, bool type);
                                                                     bool a3[bool];
   /* create arrays */
                                                                     bool a4[bool];
   for (i = 0; i < n; i++) {
        Z3 symbol s = Z3 mk int symbol(ctx, i);
        a[i]
              = Z3 mk const(ctx, s, array type);
                                                                     All can be distinct.
   /* assert distinct(a[0], ..., a[n]) */
    d = Z3 mk distinct(ctx, n, a);
                                                                     Add:
    printf("%s\n", Z3 ast to string(ctx, d));
    Z3 assert cnstr(ctx, d);
                                                                     bool a5[bool];
   /* context is satisfiable if n < 5 */
   if (Z3 check(ctx) == 1 false)
        printf("unsatisfiable, n: %d\n", n);
                                                                     Two of a1,..,a5 must
    Z3 del context(ctx);
                                                                     be equal.
```

Future/Current Work

- Coming soon (Z3 2.0):
 - Proofs & Unsat cores
 - Superposition Calculus
 - Decidable Fragments
 - Machine Learning
 - Non linear arithmetic (Gröbner Bases)
 - Inductive Datatypes
 - Improved Array & Bit-vector theories
- Several performance improvements
- More "customers" & Applications



Conclusions

- Z3 is a new SMT solver from Microsoft Research.
- Z3 is used in several projects.
- Z3 is freely available for academic research:
 - http://research.microsoft.com/projects/z3

