# First-Order Term-Indexing

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## References



R. Sekar, I. V. Ramakrishnan, and Andrei Voronkov, *Term indexing*, Handbook of Automated Reasoning (Alan Robinson and Andrei Voronkov, eds.), Elsevier Science Publishers B. V., Amsterdam, The Netherlands, 2001, pp. 1853–1964.

## Outline

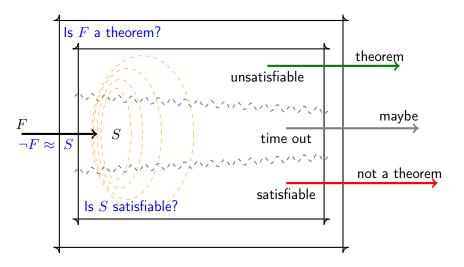
- Motivation
- 2 Term Structure
- 3 Path-Indexing
- 4 Discrimination Trees
- **5** Experiences

## Notation

### Clausal form

$$\left\{ \begin{array}{l} \mathsf{P}(\mathsf{f}(x)) \vee \mathsf{f}(x) \not\approx \mathsf{a}, \ \mathsf{g}(x,y) \approx \mathsf{a} \vee \neg \mathsf{Q}(x,y), \ \mathcal{C}_3 \end{array} \right\} \\ \equiv \\ \forall x \left( \mathsf{P}(\mathsf{f}(x)) \vee \mathsf{f}(x) \not\approx \mathsf{a} \right) \\ \wedge \\ \forall xy \left( \mathsf{g}(x,y) \approx \mathsf{a} \vee \neg \mathsf{Q}(x,y) \right) \\ \wedge \\ \forall \mathcal{V}\mathsf{ar}(\mathcal{C}_3) \left( \mathcal{C}_3 \right) \end{aligned}$$

### FOL Theorem Proving



A sound and refutation complete calculus.

Resolution (without equality)

Resolve and factor all clauses and literals in an unsatisfiable set.

$$\frac{A \vee \mathcal{C} \quad \neg B \vee \mathcal{D}}{(\mathcal{C} \vee \mathcal{D})\sigma} \ (\sigma) \ \text{resolution} \qquad \frac{A \vee B \vee \mathcal{C}}{(A \vee \mathcal{C})\sigma} \ (\sigma) \ \text{factoring}$$

$$\sigma = \mathrm{mgu}(A, B)$$

With a fair strategy the empty clause will be derived eventually.

Observation

Usually the set grows too fast to obtain a result.

## A sound, refutation complete, and effective calculus.

- 1 Reduce search space
  - e.g. Ordered Resolution
  - ... selection functions for clauses and literals
- 2 Reduce redundancy
  - e.g. ignore clause  $\mathcal{D}$ , if  $\mathcal{C}$  subsumes  $\mathcal{D}$ , i.e.  $\mathcal{C}\tau\subseteq\mathcal{D}$ .
  - ...depends on the calculus

## Example (forward subsumption)

$$S = \{^{^{1:}}\mathsf{P}(x,y),^{^{2:}}\neg\mathsf{P}(\mathsf{a},z)\} \cup \{^{^{3:}}\mathsf{P}(\mathsf{a},z')\}$$

$$t_1$$
 subsumes  $t_3$ 

$$\frac{\mathsf{P}(x,y) \quad \neg \mathsf{P}(\mathsf{a},z)}{\Box} \ \{x \mapsto \mathsf{a}, y \mapsto z\}$$

Resolution

$$S \perp = \{ \mathsf{P}(\perp, \perp), \neg \mathsf{P}(\mathsf{a}, \perp), \mathsf{P}(\mathsf{a}, \perp) \}$$

InstGen/SMT

### Goal

A sound, refutation complete, and *effective* calculus.

- 3 Quickly find
  - variants
  - instances
  - generalizations
  - unifiable terms

of a query term in a given set of terms.

variant removal backward subsumption forward subsumption resolution, etc.

### Observation

Deduction rate drops quickly with linear search.

### **Improvement**

Term-Indexing

# Position-Strings

#### Positions of a term

$$\mathcal{P}\mathsf{os}(t) = \begin{cases} \{\epsilon\} & \text{if } t = x \in \mathcal{V} \\ \{\epsilon\} \cup \{ip \mid 1 \le i \le n \land p \in \mathcal{P}\mathsf{os}(t_i)\} & \text{if } t = f(t_1, \dots, t_n) \end{cases}$$

## Normalization of variable names

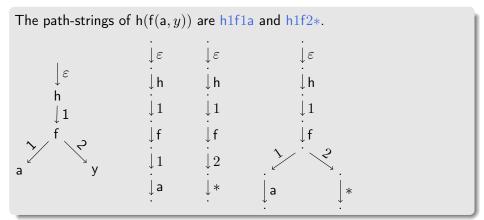
### **Forget**

$$\begin{aligned} \mathsf{path}: \mathsf{f}(x,y) \mapsto \{\mathsf{f}1*,\mathsf{f}2*\} \\ \mathsf{f}(x,x) \mapsto \{\mathsf{f}1*,\mathsf{f}2*\} \end{aligned}$$
 
$$\mathsf{pre-order}: \mathsf{f}(x,y) \mapsto \mathsf{f}** \\ \mathsf{f}(x,x) \mapsto \mathsf{f}**$$

#### Enumerate

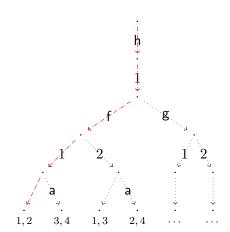
$$\begin{aligned} \mathsf{path}: \mathsf{f}(x,y) \mapsto \{\mathsf{f}1*_1,\mathsf{f}2*_2\} \\ \mathsf{f}(x,x) \mapsto \{\mathsf{f}1*_1,\mathsf{f}2*_1\} \end{aligned}$$
 
$$\mathsf{pre-order}: \mathsf{f}(x,y) \mapsto \mathsf{f}*_1*_2 \\ \mathsf{f}(x,x) \mapsto \mathsf{f}*_1*_1 \end{aligned}$$

# Path-Strings



# Pre-order Strings

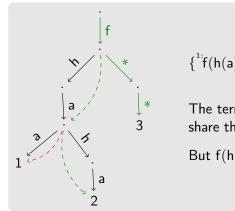
# Path-Index



## Discrimination Tree

hi!

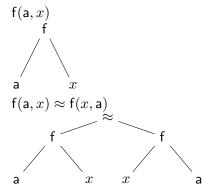
## Non-linear terms



$$\{{}^{1:}\mathsf{f}(\mathsf{h}(\mathsf{a}),\mathsf{a}),{}^{2:}\mathsf{f}(\mathsf{h}(\mathsf{a}),\mathsf{h}(\mathsf{a}))\},{}^{3:}\mathsf{f}(x,y)\}$$

The terms f(x, y) and f(z, z) share the preorder-term f\*\*.

But f(h(a), a) is not an instance of f(z, z).



# Perfect filtering

$$\begin{cases} \ ^{1:}\mathsf{h}(\mathsf{f}(x,x)), & \ ^{2:}\mathsf{h}(\mathsf{g}(\mathsf{a},x)), & \ ^{3:}\mathsf{h}(\mathsf{f}(y,z)), \\ \ ^{4:}\mathsf{h}(\mathsf{g}(\mathsf{a},y)), & \ ^{5:}\mathsf{h}(\mathsf{f}(y,x)), & \ ^{6:}\mathsf{h}(\mathsf{g}(y,a)) \end{cases}$$