# Term-Indexing

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# References

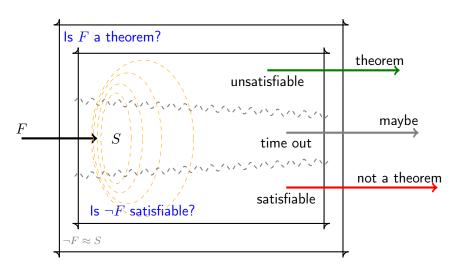


R. Sekar, I. V. Ramakrishnan, and Andrei Voronkov, *Term indexing*, Handbook of Automated Reasoning (Alan Robinson and Andrei Voronkov, eds.), Elsevier Science Publishers B. V., Amsterdam, The Netherlands, 2001, pp. 1853–1964.

# Outline

- Motivation
- 2 Term Structure
- 3 Path Indexing
- 4 Discrimination Trees
- 5 Substitution Trees
- 6 Experiment

# Refutation



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#### Clausal form

$$\left\{ \begin{array}{l} \mathsf{P}(\mathsf{f}(x)) \vee \mathsf{f}(x) \not\approx \mathsf{a}, \ \mathsf{g}(x,y) \approx \mathsf{a} \vee \neg \mathsf{Q}(x,y), \ \mathcal{C}_3 \end{array} \right\} \\ \equiv \\ \forall x \left( \mathsf{P}(\mathsf{f}(x)) \vee \mathsf{f}(x) \not\approx \mathsf{a} \right) \\ \wedge \\ \forall xy \left( \mathsf{g}(x,y) \approx \mathsf{a} \vee \neg \mathsf{Q}(x,y) \right) \\ \wedge \\ \forall \mathcal{V}\mathsf{ar}(\mathcal{C}_3) \left( \mathcal{C}_3 \right) \end{aligned}$$

#### Goal

A sound and refutation complete calculus.

### Resolution (without equality)

Resolve and factor all clauses and literals in an unsatisfiable set

$$\frac{A \vee \mathcal{C} \quad \neg B \vee \mathcal{D}}{(\mathcal{C} \vee \mathcal{D})\sigma} \ (\sigma) \ \text{resolution} \qquad \frac{A \vee B \vee \mathcal{C}}{(A \vee \mathcal{C})\sigma} \ (\sigma) \ \text{factoring}$$

$$\sigma = \mathrm{mgu}(A, B)$$

and the empty clause will be derived eventually.

#### Observation

Usually the set grows too fast to obtain a result.

#### Goal

A sound, refutation complete, and *effective* calculus.

- 1 Reduce search space
  - Ordered Resolution, Strategies, . . .
  - ... with selection functions for clauses and literals
- 2 Reduce redundancy
  - e.g. discard clauses that are subsumed by other clauses
  - ...depending on the calculus

# Example (forward subsumption)

$$S = \{^{^{1:}}\mathsf{P}(x,y),^{^{2:}}\neg\mathsf{P}(\mathsf{a},z)\} \cup \{^{^{3:}}\!\mathsf{P}(\mathsf{a},z')\}$$

$$t_1$$
 subsumes  $t_3$ 

$$\frac{\mathsf{P}(x,y) \quad \neg \mathsf{P}(\mathsf{a},z)}{\Box} \ \{x \mapsto \mathsf{a}, y \mapsto z\}$$

Resolution

$$S \perp = \{ \mathsf{P}(\perp, \perp), \neg \mathsf{P}(\mathsf{a}, \perp), \mathsf{P}(\mathsf{a}, \perp) \}$$

InstGen / SMT

#### Goal

A sound, refutation complete, and effective calculus.

- 3 Quickly find
  - variants
  - instances
  - generalizations
  - unifiable terms

of a query term in a given set of terms.

variant removal backward subsumption forward subsumption resolution, demodulation

# Observation

Deduction rate drops quickly with sequential search.

# Term Indexing

Data structures and algorithms for fast retrieval of matching terms.

# Definition (Position Strings)

$$\mathcal{P} \mathsf{os}^\Sigma(t) = \begin{cases} \{\langle \epsilon, x \rangle\} & \text{if } t = x \in \mathcal{V} \\ \{\langle \epsilon, f \rangle\} \cup \{\langle ip, s \rangle \mid (p, s) \in \mathcal{P} \mathsf{os}^\Sigma(t_i)\} & \text{if } t = f(t_1, \dots, t_n) \end{cases}$$

#### Term traversals

$$\mathcal{P} os^{\Sigma}(\mathsf{h}(\mathsf{f}(\mathsf{a},\mathsf{y}))) = \{\langle \epsilon,\mathsf{h}\rangle, \langle 1,\mathsf{f}\rangle, \langle 11,\mathsf{a}\rangle, \langle 12,y\rangle\}$$
 
$$\langle \epsilon,\mathsf{h}\rangle$$
 
$$\langle 1,\mathsf{f}\rangle$$
 
$$\langle 1,\mathsf{f}\rangle$$
 
$$\langle \epsilon,\mathsf{h}\rangle\langle 1,\mathsf{f}\rangle\langle 12,y\rangle$$
 path from root to leaf 
$$\langle \epsilon,\mathsf{h}\rangle\langle 1,\mathsf{f}\rangle\langle 11,\mathsf{a}\rangle\langle 12,y\rangle$$
 pre-order traversal 
$$\langle 11,\mathsf{a}\rangle$$
 
$$\langle 12,\mathsf{y}\rangle$$

#### **Variables**

Variants of terms generate the same position strings

• if variable names are ignored

$$f(y,z) \Rightarrow \langle \epsilon, f \rangle \langle 1, * \rangle \langle 2, * \rangle$$

or normalized

$$f(y,z) \Rightarrow \langle \epsilon, \mathsf{f} \rangle \langle 1, x_1 \rangle \langle 2, x_2 \rangle$$
  
$$f(y,y) \Rightarrow \langle \epsilon, \mathsf{f} \rangle \langle 1, x_1 \rangle \langle 2, x_1 \rangle$$

In the first case even non-variants of terms generate the same strings.

#### Notation

We abbreviate

• path strings  $\langle \epsilon, \mathsf{h} \rangle \langle 1, \mathsf{f} \rangle \langle 12, * \rangle$ 

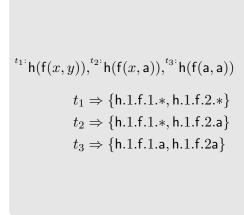
h.1.f.2.\*

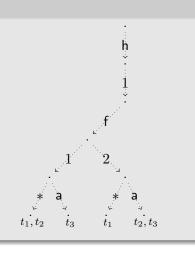
• and traversal strings  $\langle \epsilon, \mathsf{h} \rangle \langle 1, \mathsf{f} \rangle \langle 11, * \rangle \langle 12, * \rangle$  when traversal order and arities of symbols are fixed.

h.f.a.\*

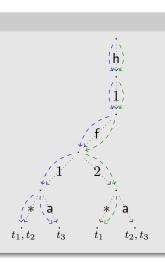
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$$\begin{split} ^{t_1:} \mathsf{h}(\mathsf{f}(x,y)), ^{t_2:} \mathsf{h}(\mathsf{f}(x,\mathsf{a})), ^{t_3:} \mathsf{h}(\mathsf{f}(\mathsf{a},\mathsf{a})) \\ & \mathsf{h}(\mathsf{f}(z,\mathsf{b}))) \Rightarrow \{\mathsf{h}.\mathsf{f}.*,\mathsf{h}.\mathsf{f}.\mathsf{b}\} \\ & u: \mathsf{h}(\mathsf{f}(z,\mathsf{b})) \mapsto \{t_1,t_2,t_3\} \cap \{t_1\} \\ & i: \mathsf{h}(\mathsf{f}(z,\mathsf{b})) \mapsto \{t_1,t_2,t_3\} \cap \{\} \\ & g: \mathsf{h}(\mathsf{f}(z,\mathsf{b})) \mapsto \{t_1,t_2\} \cap \{t_1\} \\ & v: \mathsf{h}(\mathsf{f}(z,\mathsf{b})) \mapsto \{t_1,t_2\} \cap \{\} \\ & v: \mathsf{h}(\mathsf{f}(z,z)) \mapsto \{t_1,t_2\} \cap \{t_1\} \end{split}$$



### Unit Superposition Inference Rules

$$\frac{s \approx t \quad L[s']}{(L[t]) \cdot \sigma} \quad \underset{\text{paramodulation}}{\text{unit}}$$

where  $\sigma = \text{mgu}(s, s'), s' \notin \mathcal{V}, t\sigma \not\succeq s\sigma$ 

$$\frac{s \approx t \quad u[s'] \not\approx v}{(u[t] \not\approx v) \cdot \sigma} \text{ } \underset{\text{superposition}}{\text{unit}} \quad \frac{s \approx t \quad u[s'] \approx v}{(u[t] \approx v) \cdot \sigma}$$

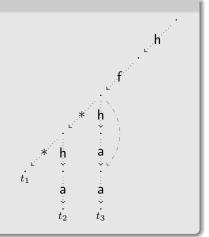
where  $\sigma = \text{mgu}(s, s'), s' \notin \mathcal{V}, t\sigma \not\succeq s\sigma, v\sigma \not\succeq u[s']\sigma$ 

where s and t (A and B respectively) are unifiable

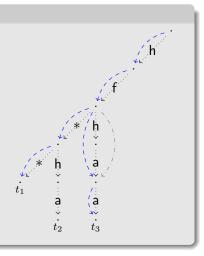
# Insert

 ${}^{t_1}$ : $h(f(x,y)), {}^{t_2}$ : $h(f(x,h(a))), {}^{t_3}$ :h(f(h(a),a)) $t_1 \Rightarrow \text{h.f.}*.*$  $t_2 \Rightarrow h.f.*.h.a$ 

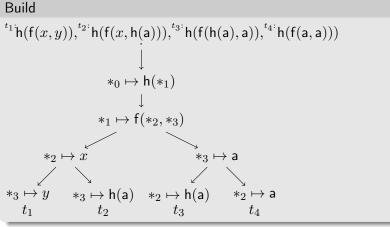
 $t_3 \Rightarrow h.f.h.a.a$ 



$$\begin{split} ^{t_1:}\mathsf{h}(\mathsf{f}(x,y)),^{t_2:}\mathsf{h}(\mathsf{f}(x,\mathsf{h}(\mathsf{a}))),^{t_3:}\mathsf{h}(\mathsf{f}(\mathsf{h}(\mathsf{a}),\mathsf{a})) \\ & \mathsf{h}(\mathsf{f}(x',\mathsf{a})) \Rightarrow \mathsf{h.f.*.a} \\ \\ u:\mathsf{h}(\mathsf{f}(x',\mathsf{a})) \mapsto \{t_1,t_3\} \\ \\ i:\mathsf{h}(\mathsf{f}(x',\mathsf{a})) \mapsto \{t_3\} \\ \\ g:\mathsf{h}(\mathsf{f}(x',\mathsf{a})) \mapsto \{t_1\} \\ \\ v:\mathsf{h}(\mathsf{f}(x',\mathsf{a})) \mapsto \{\ \} \end{split}$$



# Subterms $^{t_1}\dot{\mathsf{h}}(\mathsf{f}(x,y)),^{t_2}\dot{\mathsf{h}}(\mathsf{f}(x,\mathsf{h}(\mathsf{a}))),^{t_3}\dot{\mathsf{h}}(\mathsf{f}(\mathsf{h}(\mathsf{a}),\mathsf{a}))$ $t_2^{12}, t_3^{111}$ $\langle 12, h \rangle$



TPTP/Problems/HWV/HWV134-1.p 2 332 428 formulae, 6 570 884 literals

checkin afterwards	ng 1000 new lite $(\ell_1,\ell_2)$	rals $A, \neg B$	sequential search	path index	speed up
1 000	500 000	761	726ms	70ms	10
2 000	1 500 000	812	2s	69ms	29
4 000	3 500 000	723	4s	75ms	53
8 000	7 500 000	433	9s	125ms	72
16 000	15 500 000	742	21s	221ms	95
32 000	31 500 000	592	40s	489ms	82
64 000	63 500 000	1 167	80s	697ms	115
128 000	127 500 000	1 479	160s	13s	12
256 000	255 500 000	1 097	320s	440s	1
512 000	511 500 000	1 440	640s	348s	2
1 024 000	1023 500 000	1 534	1280s	330s	4