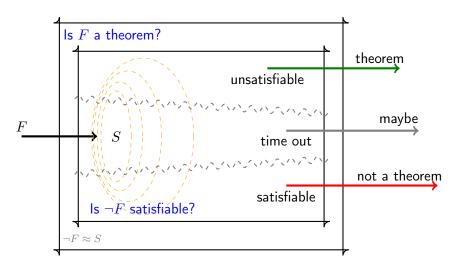
References

Outline

- Motivation
- 2 Term Structure
- 3 Path Indexing
- 4 Discrimination Trees
- 5 Substitution Trees
- 6 Experiment

Refutation



Alexander Maringele

Clausal form

$$\{ \begin{array}{l} \mathsf{P}(\mathsf{f}(x)) \vee \mathsf{f}(x) \not\approx \mathsf{a}, \ \mathsf{g}(x,y) \approx \mathsf{a} \vee \neg \mathsf{Q}(x,y), \ \mathcal{C}_3 \end{array} \} \\ & \equiv \\ \forall x \left(\mathsf{P}(\mathsf{f}(x)) \vee \mathsf{f}(x) \not\approx \mathsf{a} \right) \\ & \wedge \\ \forall xy \left(\mathsf{g}(x,y) \approx \mathsf{a} \vee \neg \mathsf{Q}(x,y) \right) \\ & \wedge \\ \forall \mathcal{V}\mathsf{ar}(\mathcal{C}_3) \left(\mathcal{C}_3 \right) \end{array}$$

Goal

A sound and refutation complete calculus.

Resolution (without equality)

Resolve and factor all clauses and literals in an unsatisfiable set

$$\frac{A \vee \mathcal{C} \quad \neg B \vee \mathcal{D}}{(\mathcal{C} \vee \mathcal{D})\sigma} \ (\sigma) \ \text{resolution} \qquad \frac{A \vee B \vee \mathcal{C}}{(A \vee \mathcal{C})\sigma} \ (\sigma) \ \text{factoring}$$

$$\sigma = \mathrm{mgu}(A, B)$$

and the empty clause will be derived eventually.

Observation

Usually the set grows too fast to obtain a result.

Goal

A sound, refutation complete, and *effective* calculus.

- 1 Reduce search space
 - Ordered Resolution, Strategies, . . .
 - ... with selection functions for clauses and literals
- 2 Reduce redundancy
 - e.g. discard clauses that are subsumed by other clauses
 - ...depending on the calculus

Example (forward subsumption)

 t_1 subsumes t_3

$$\frac{\mathsf{P}(x,y) \quad \neg \mathsf{P}(\mathsf{a},z)}{\Box} \ \{x \mapsto \mathsf{a}, y \mapsto z\}$$

Resolution

$$S \perp = \{ \mathsf{P}(\perp, \perp), \neg \mathsf{P}(\mathsf{a}, \perp), \mathsf{P}(\mathsf{a}, \perp) \}$$

InstGen / SMT

Goal

A sound, refutation complete, and effective calculus.

- 3 Quickly find
 - variants
 - instances
 - generalizations
 - unifiable terms

of a query term in a given set of terms.

variant removal backward subsumption forward subsumption resolution, demodulation

Observation

Deduction rate drops quickly with sequential search.

Term Indexing

Data structures and algorithms for fast retrieval of matching terms.

Position Strings

Definition

$$\mathcal{P} \mathsf{os}^\Sigma(t) = \begin{cases} \{\langle \epsilon, x \rangle\} & \text{if } t = x \in \mathcal{V} \\ \{\langle \epsilon, f \rangle\} \cup \{\langle ip, s \rangle \mid \langle p, s \rangle \in \mathcal{P} \mathsf{os}^\Sigma(t_i)\} & \text{if } t = f(t_1, \dots, t_n) \end{cases}$$

Term traversals

Variables

Variants of terms generate the same position strings

• if variable names are ignored

$$f(y,z) \Rightarrow \langle \epsilon, f \rangle \langle 1, * \rangle \langle 2, * \rangle$$

or normalized

$$f(y,z) \Rightarrow \langle \epsilon, \mathsf{f} \rangle \langle 1, x_1 \rangle \langle 2, x_2 \rangle$$

$$f(y,y) \Rightarrow \langle \epsilon, \mathsf{f} \rangle \langle 1, x_1 \rangle \langle 2, x_1 \rangle$$

In the first case even non-variants of terms generate the same strings.

Notation

We abbreviate

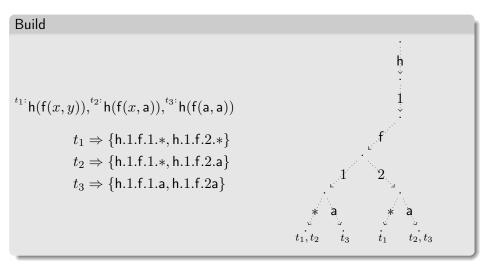
• path strings $\langle \epsilon, \mathsf{h} \rangle \langle 1, \mathsf{f} \rangle \langle 12, * \rangle$

h.1.f.2.*

• and pre-order traversal strings $\langle \epsilon, \mathsf{h} \rangle \langle 1, \mathsf{f} \rangle \langle 11, * \rangle \langle 12, * \rangle$ when the arities of function symbols are fixed.

h.f.a.*

Path Indexing



Retrieve

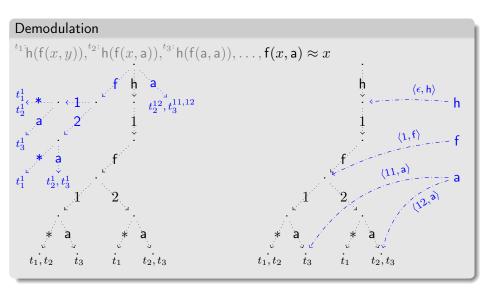
$$\frac{s \approx t \quad L[s']}{(L[t]) \cdot \sigma} \quad \underset{\text{paramodulation}}{\text{unit}}$$

where $\sigma = \text{mgu}(s, s')$, $s' \notin \mathcal{V}$, $t\sigma \not\geq s\sigma$

$$\frac{s \approx t \quad u[s'] \not\approx v}{(u[t] \not\approx v) \cdot \sigma} \quad \underset{\text{superposition}}{\text{unit}} \quad \frac{s \approx t \quad u[s'] \approx v}{(u[t] \approx v) \cdot \sigma}$$

where $\sigma = \text{mgu}(s, s')$, $s' \notin \mathcal{V}$, $t\sigma \not\succeq s\sigma$, $v\sigma \not\succeq u[s']\sigma$

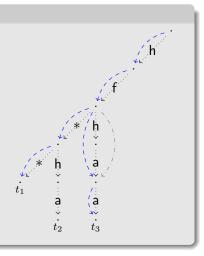
where s and t (A and B respectively) are unifiable



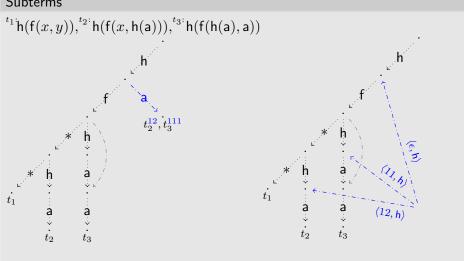
Discrimination Trees

Insert t_1 : $h(f(x,y)), {}^{t_2}$: $h(f(x,h(a))), {}^{t_3}$:h(f(h(a),a)) $t_1 \Rightarrow \mathsf{h.f.}*.*$ $t_2 \Rightarrow \text{h.f.*.h.a}$ $t_3 \Rightarrow h.f.h.a.a$

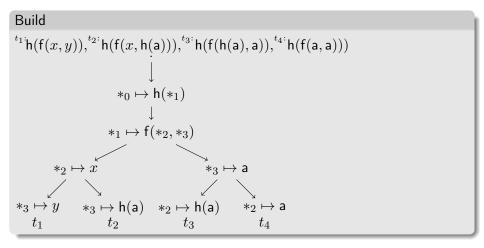
$$\begin{split} ^{t_1:}\mathsf{h}(\mathsf{f}(x,y)),^{t_2:}\mathsf{h}(\mathsf{f}(x,\mathsf{h}(\mathsf{a}))),^{t_3:}\mathsf{h}(\mathsf{f}(\mathsf{h}(\mathsf{a}),\mathsf{a})) \\ & \mathsf{h}(\mathsf{f}(x',\mathsf{a})) \Rightarrow \mathsf{h.f.*.a} \\ \\ u:\mathsf{h}(\mathsf{f}(x',\mathsf{a})) \mapsto \{t_1,t_3\} \\ \\ i:\mathsf{h}(\mathsf{f}(x',\mathsf{a})) \mapsto \{t_3\} \\ \\ g:\mathsf{h}(\mathsf{f}(x',\mathsf{a})) \mapsto \{t_1\} \\ \\ v:\mathsf{h}(\mathsf{f}(x',\mathsf{a})) \mapsto \{\ \} \end{split}$$







Substitution Trees



TPTP/Problems/HWV/HWV134-1.p 2 332 428 formulae, 6 570 884 literals

checking afterwards	g 1000 new lite (ℓ_1,ℓ_2)	rals $A, \neg B$	sequential search	path index	speed up
1 000	500 000	761	726ms	70ms	10
2 000	1 500 000	812	2s	69ms	29
4 000	3 500 000	723	4s	75ms	53
8 000	7 500 000	433	9s	125ms	72
16 000	15 500 000	742	21s	221ms	95
32 000	31 500 000	592	40s	489ms	81
64 000	63 500 000	1 167	80s	697ms	114
128 000	127 500 000	1 479	160s	13s	12
256 000	255 500 000	1 097	320s	440s	<1
512 000	511 500 000	1 440	640s	348s	<2
1 024 000	1023 500 000	1 534	1280s	330s	<4