

# Term-Indexing for Instantiation-Based First Order Theorem Proving

Alexander Maringele

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## References



Alexandre Riazanov and Andrei Voronkov, *Efficient instance retrieval with standard and relational path indexing*, Automated Deduction – CADE-19 (Franz Baader, ed.), Lecture Notes in Computer Science, vol. 2741, Springer Berlin Heidelberg, 2003, pp. 380–396 (English).



R. Sekar, I. V. Ramakrishnan, and Andrei Voronkov, *Term indexing*, Handbook of Automated Reasoning (Alan Robinson and Andrei Voronkov, eds.), Elsevier Science Publishers B. V., Amsterdam, The Netherlands, The Netherlands, 2001, pp. 1853–1964.

# Outline

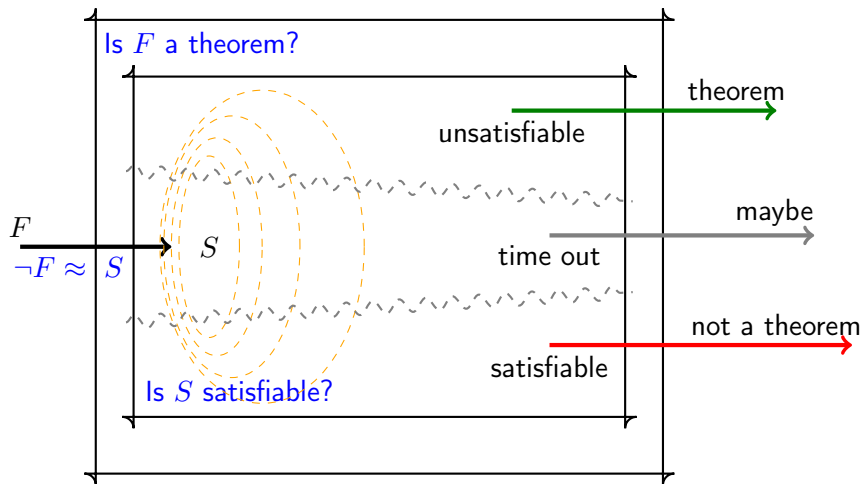
- 1 Motivation
- 2 path indexing
- 3 discrimination trees

## Notation (First Order Logic)

- $\Sigma = (\mathcal{V}, \mathcal{F}, \mathcal{P})$  signature
- $\mathcal{T} = \mathcal{V} \cup \{ f(t_1, \dots, t_n) \mid f \in \mathcal{F}_{(n)}, t_i \in \mathcal{T} \}$  terms
- $\mathcal{A} = \{ P(t_1, \dots, t_n) \mid P \in \mathcal{P}_{(n)}, t_i \in \mathcal{T} \} \cup \{ s \approx t \mid s, t \in \mathcal{T} \}$  atoms
- $\mathcal{L} = \mathcal{A} \cup \{ \neg A \mid A \in \mathcal{A} \}$  literals
- $\mathcal{C} = 2^{\mathcal{L}}$ , e.g.  $f(x) \approx a \vee a \not\approx b$  clauses
- $\mathcal{S} = 2^{\mathcal{C}}$ , e.g.  $\{ f(x) \approx a \vee a \not\approx b, f(x) \approx b \}$  clause sets

$$\{C_1, \dots, C_n\} \equiv \forall \text{Var}(C_1) C_1 \wedge \dots \wedge \forall \text{Var}(C_n) C_n$$

# Resolution



## Goal

A sound, refutational complete, and *effective* calculus.

- *Reduce* search space
  - Ordered Resolution
  - ...selection functions for clauses and literals
- *Reduce* redundancy
  - Ignore clause  $C$ , if  $C$  subsumes  $D$ , i.e.  $C\tau \subseteq D$ ?
  - ...depends on the calculus

## Example

$$S = \{^{t_1}P(x, y), ^{t_2}\neg P(a, z), ^{t_3}P(a, z')\}$$

$t_1$  subsumes  $t_3$

$$S\perp = \{P(\perp, \perp), \neg P(a, \perp), P(a, \perp)\}$$

InstGen/SMT

$$\frac{P(x, y) \quad \neg P(a, z)}{\square} \quad \{x \mapsto a, y \mapsto z\}$$

Resolution

## unit superposition inference rules

$$\frac{s \approx t \quad L[s']}{(L[t]) \cdot \sigma} \quad \begin{array}{l} \text{unit} \\ \text{paramodulation} \end{array}$$

where  $\sigma = \text{mgu}(s, s')$ ,  $s' \notin \mathcal{V}$ ,  $t\sigma \neq s\sigma$

$$\frac{s \approx t \quad u[s'] \not\approx v}{(u[t] \not\approx v) \cdot \sigma} \quad \begin{array}{l} \text{unit} \\ \text{superposition} \end{array} \quad \frac{s \approx t \quad u[s'] \approx v}{(u[t] \approx v) \cdot \sigma}$$

where  $\sigma = \text{mgu}(s, s')$ ,  $s' \notin \mathcal{V}$ ,  $t\sigma \neq s\sigma$ ,  $v\sigma \neq u[s']\sigma$

$$\frac{s \not\approx t}{\square} \quad \begin{array}{l} \text{unit equality} \\ \text{resolution} \end{array}$$

$$\frac{A \quad \neg B}{\square} \quad \begin{array}{l} \text{unit} \\ \text{resolution} \end{array}$$

where  $s$  and  $t$  ( $A$  and  $B$  respectively) are unifiable

# Term retrieval problems

## Definition

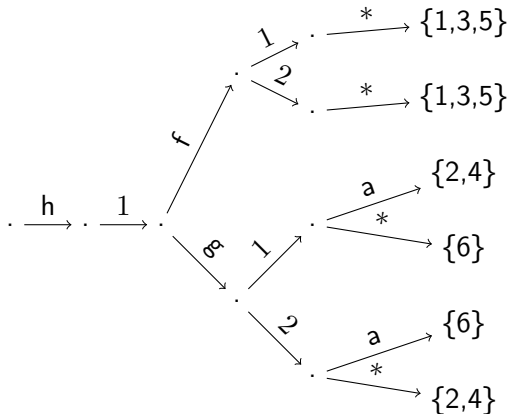
In a given set of terms

- find terms that are variants of a given term.  
 $\text{variant}(\ell, t) \Leftrightarrow \exists \sigma \ell \sigma = t$  and  $\sigma$  is renaming.
- find terms that are unifiable with a given term.  
 $\text{unifiable}(\ell, t) \Leftrightarrow \exists \sigma \ell \sigma = t \sigma$
- find terms that are instances of a given term.  
 $\text{instance}(\ell, t) \Leftrightarrow \exists \sigma \ell = t \sigma$
- find terms that are generalizations of a given term.  
 $\text{generalization}(\ell, t) \Leftrightarrow \exists \sigma \ell \sigma = t$



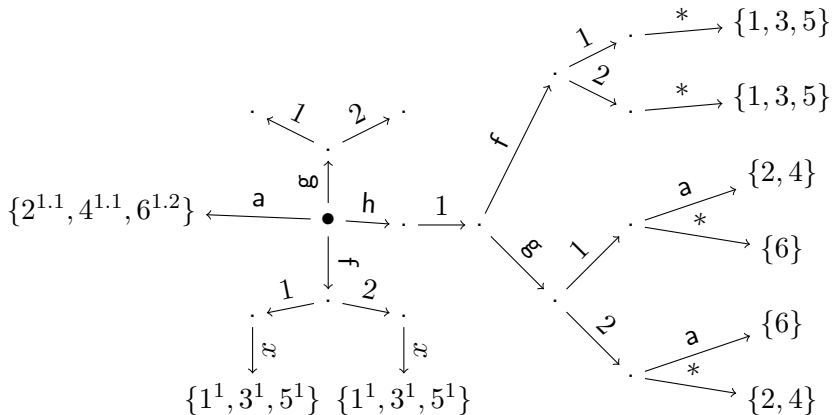
## Definition

$$\{^1: h(f(x, x)), ^2: h(g(a, x)), ^3: h(f(y, z)) ^4: h(g(a, y)), ^5: h(f(y, x)), ^6: h(g(y, a))\}$$

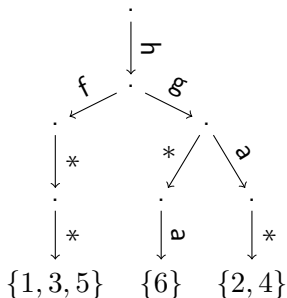


$$h(g(y, x)) \mapsto \{ h.1.g.1.*, h.1.g.2.* \}$$

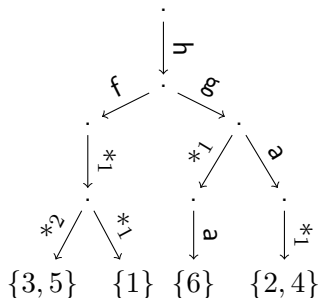
$$\{^1\text{h}(\text{f}(\text{x}, \text{x})), ^2\text{h}(\text{g}(\text{a}, \text{x})), ^3\text{h}(\text{f}(\text{y}, \text{z})) ^4\text{h}(\text{g}(\text{a}, \text{y})), ^5\text{h}(\text{f}(\text{y}, \text{x})), ^6\text{h}(\text{g}(\text{y}, \text{a}))\}$$



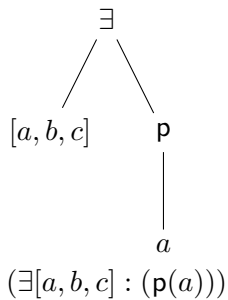
$$\left\{ \begin{array}{lll} 1: h(f(x, x)), & 2: h(g(a, x)), & 3: h(f(y, z)), \\ 4: h(g(a, y)), & 5: h(f(y, x)), & 6: h(g(y, a)) \end{array} \right\}$$



$$h(f(y, x)) \mapsto h.f.*.*$$



$$h(f(y, x)) \mapsto' h.f.*_1.*_2$$



$f(a, x)$  $f$  $a$  $x$  $f(a, x) \approx f(x, a)$  $\approx$  $f$  $a$  $x$  $f$  $x$  $a$