

# Term-Indexing for Instantiation-Based First Order Theorem Proving

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# References



Robert Nieuwenhuis, Thomas Hillenbrand, Alexandre Riazanov, and Andrei Voronkov, *On the evaluation of indexing techniques for theorem proving*, Automated Reasoning (Rajeev Goré, Alexander Leitsch, and Tobias Nipkow, eds.), Lecture Notes in Computer Science, vol. 2083, Springer Berlin Heidelberg, 2001, pp. 257–271.



R. Sekar, I. V. Ramakrishnan, and Andrei Voronkov, *Term indexing*, Handbook of Automated Reasoning (Alan Robinson and Andrei Voronkov, eds.), Elsevier Science Publishers B. V., Amsterdam, The Netherlands, 2001, pp. 1853–1964.

# Outline

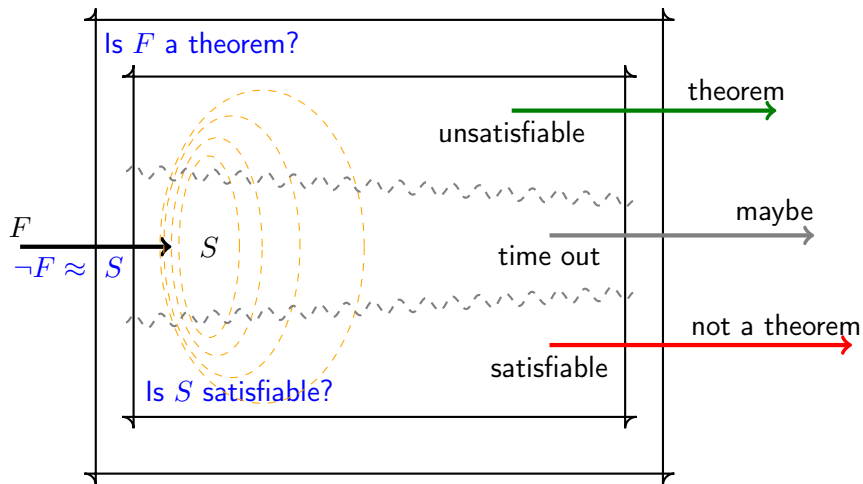
- 1 Motivation
- 2 Position
- 3 Path Indexing
- 4 Discrimination Tree
- 5 discrimination trees

## Notation (First Order Logic)

- $\Sigma = (\mathcal{V}, \mathcal{F}, \mathcal{P})$  signature
- $\mathcal{T} = \mathcal{V} \cup \{ f(t_1, \dots, t_n) \mid f \in \mathcal{F}_{(n)}, t_i \in \mathcal{T} \}$  terms
- $\mathcal{A} = \{ P(t_1, \dots, t_n) \mid P \in \mathcal{P}_{(n)}, t_i \in \mathcal{T} \} \cup \{ s \approx t \mid s, t \in \mathcal{T} \}$  atoms
- $\mathcal{L} = \mathcal{A} \cup \{ \neg A \mid A \in \mathcal{A} \}$  literals
- $\mathcal{C} = 2^{\mathcal{L}}$ , e.g.  $f(x) \approx a \vee a \not\approx b$  clauses
- $\mathcal{S} = 2^{\mathcal{C}}$ , e.g.  $\{ f(x) \approx a \vee a \not\approx b, f(x) \approx b \}$  clause sets

$$\{C_1, \dots, C_n\} \equiv \forall \text{Var}(C_1) C_1 \wedge \dots \wedge \forall \text{Var}(C_n) C_n$$

# Resolution



## Goal

A sound, refutational complete, and *effective* calculus.

- *Reduce* search space
  - e.g. Ordered Resolution
  - ...selection functions for clauses and literals
- *Reduce* redundancy
  - Ignore clause  $D$ , if  $C$  subsumes  $D$ , i.e.  $C\tau \subseteq D$ ?
  - ...depends on the calculus

## Example

$$S = \{^1P(x, y), ^2\neg P(a, z)\} \cup \{^3P(a, z')\}$$

$t_1$  subsumes  $t_3$

$$S\perp = \{P(\perp, \perp), \neg P(a, \perp), P(a, \perp)\}$$

InstGen/SMT

$$\frac{P(x, y) \quad \neg P(a, z)}{\square} \quad \{x \mapsto a, y \mapsto z\}$$

Resolution

## Unit Superposition Inference Rules

$$\frac{s \approx t \quad L[s']}{(L[t]) \cdot \sigma} \quad \begin{array}{l} \text{unit} \\ \text{paramodulation} \end{array}$$

where  $\sigma = \text{mgu}(s, s')$ ,  $s' \notin \mathcal{V}$ ,  $t\sigma \neq s\sigma$

$$\frac{s \approx t \quad u[s'] \not\approx v}{(u[t] \not\approx v) \cdot \sigma} \quad \begin{array}{l} \text{unit} \\ \text{superposition} \end{array} \quad \frac{s \approx t \quad u[s'] \approx v}{(u[t] \approx v) \cdot \sigma}$$

where  $\sigma = \text{mgu}(s, s')$ ,  $s' \notin \mathcal{V}$ ,  $t\sigma \neq s\sigma$ ,  $v\sigma \neq u[s']\sigma$

$$\frac{s \not\approx t}{\square} \quad \begin{array}{l} \text{unit equality} \\ \text{resolution} \end{array}$$

$$\frac{A \quad \neg B}{\square} \quad \begin{array}{l} \text{unit} \\ \text{resolution} \end{array}$$

where  $s$  and  $t$  ( $A$  and  $B$  respectively) are unifiable

# Term retrieval problems

## Definition

In a given set of terms

- find terms that are variants of a given term  $t$   
 $\text{variant}(s, t) \Leftrightarrow \exists \sigma \ s \sigma = t$  and  $\sigma$  is renaming
- find terms that are instances of a given term  
 $\text{instance}(s, t) \Leftrightarrow \exists \sigma \ s = t \sigma$
- find terms that are generalizations of a given term  
 $\text{generalization}(s, t) \Leftrightarrow \exists \sigma \ s \sigma = t$
- find terms that are unifiable with a given term  
 $\text{unifiable}(s, t) \Leftrightarrow \exists \sigma \ s \sigma = t \sigma$

backward subsumption

forward subsumption

resolution, etc.

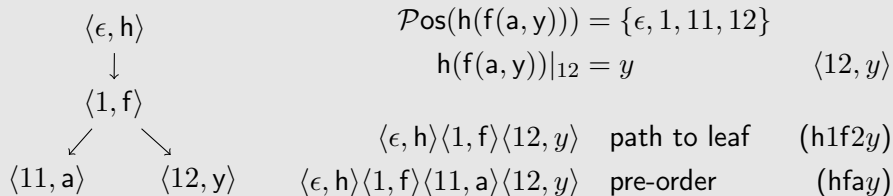


# Position-Strings

## Positions in a term

$$\mathcal{P}\text{os}(t) = \begin{cases} \{\epsilon\} & \text{if } t = x \in \mathcal{V} \\ \{\epsilon\} \cup \{ip \mid 1 \leq i \leq n \wedge p \in \mathcal{P}\text{os}(t_i)\} & \text{if } t = f(t_1, \dots, t_n) \end{cases}$$

$h(f(a, y))$



# Normalization of variables

## Amnesia

$$\text{path} : f(x, y) \mapsto \{f1*, f2*\}$$

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$$\text{pre-order} : f(x, y) \mapsto f**$$

$$f(x, x) \mapsto f**$$

## Renaming

$$\text{path} : f(x, y) \mapsto \{f1*_1, f2*_2\}$$

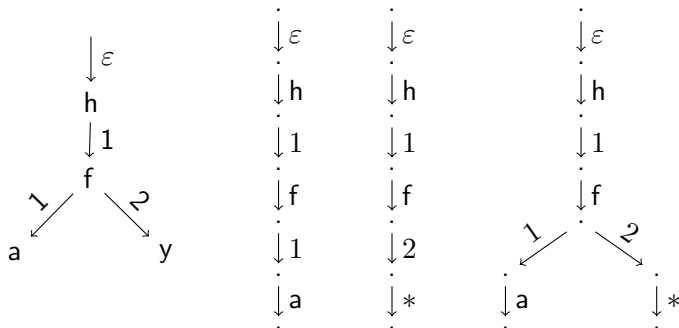
$$f(x, x) \mapsto \{f1*_1, f2*_1\}$$

$$\text{pre-order} : f(x, y) \mapsto f*_1*_2$$

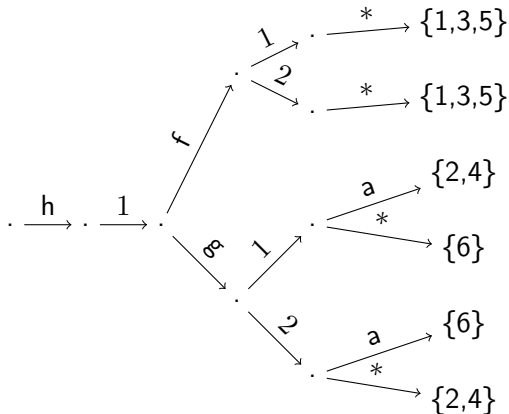
$$f(x, x) \mapsto f*_1*_1$$

# Path-Strings

The path-strings of  $t = h(f(a, y))$  are  $\epsilon.h.1.f.1.a$  and  $\epsilon.h.1.f.1.2.*$ :

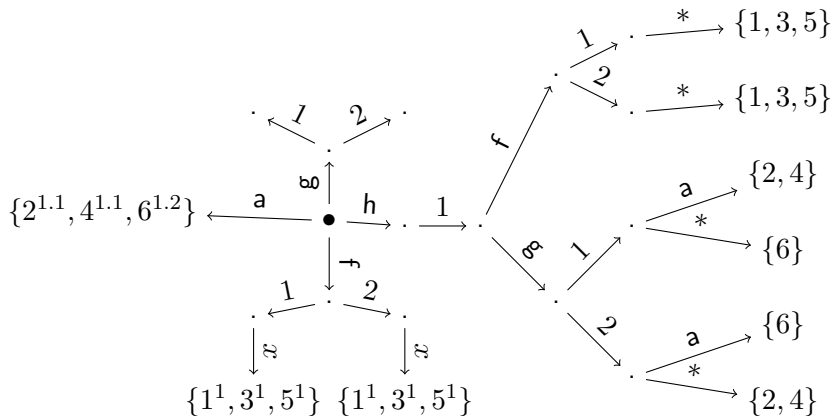


$$\{^1: h(f(x, x)), ^2: h(g(a, x)), ^3: h(f(y, z)) ^4: h(g(a, y)), ^5: h(f(y, x)), ^6: h(g(y, a))\}$$

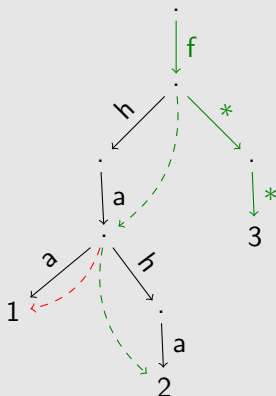


$$h(g(y, x)) \mapsto \{h.1.g.1.*, h.1.g.2.*\}$$

$$\{^1\text{h}(\text{f}(\text{x}, \text{x})), ^2\text{h}(\text{g}(\text{a}, \text{x})), ^3\text{h}(\text{f}(\text{y}, \text{z})) ^4\text{h}(\text{g}(\text{a}, \text{y})), ^5\text{h}(\text{f}(\text{y}, \text{x})), ^6\text{h}(\text{g}(\text{y}, \text{a}))\}$$



# Non-linear terms



$$\{^1: f(h(a), a), ^2: f(h(a), h(a)), ^3: f(x, y)\}$$

The terms  $f(x, y)$  and  $f(z, z)$   
share the preorder-term  $f**$ .

But  $f(h(a), a)$  is not an instance of  $f(z, z)$ .

$f(a, x)$  $f$  $a$  $x$  $f(a, x) \approx f(x, a)$  $\approx$  $f$  $a$  $x$  $f$  $x$  $a$