





POLITECNICO
MILANO 1863

Exercise session 1 - Structural Properties Advanced and Multivariable Control

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- ▶ Anyone is invited to raise questions during the exercise sessions.
- ▶ If you wish to further discuss the topic, drop an e-mail asking for an appointment.
- ▶ No online classroom, registrations of last year available on WeBeep.
- ▶ **Attending exercise sessions and laboratories is highly recommended!**

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- ▶ Exercise sessions are on Tuesday, exceptions will be notified in advance. One class only.
- ▶ Six laboratories of 3 hours each, held on Tuesdays, with me and Matteo Luigi De Pascali (S.0.1), One class only.
- ▶ A schedule of lessons, exercise Sessions and laboratories will be shared on Beep.

Concerning the laboratories, the following software will be necessary:

- ▶ MATLAB (my version is R2022a)
- ▶ pplane (see WeBeep)
- ▶ Control System Toolbox
- ▶ CasADI Toolbox (web.casadi.org/get)

Download the .zip and add it to your Matlab path

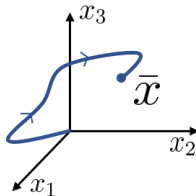
Definition - Reachability (continuous-time systems)

Given the continuous-time system $\dot{x}(t) = Ax(t) + Bu(t)$, a state \bar{x} is said to be **reachable** if there exists an arbitrary finite time \bar{t} and an input realization $\bar{u}(\tau)$, $\tau \in [0, \bar{t}]$, such that starting from the origin (i.e. $x(0) = 0$), $x(\bar{t}) = \bar{x}$.

arbitrary state \bar{x}

independently from traj taken to reach state \bar{x}

In other words, a state is reachable if it is possible to design an input sequence to **bring the state $x(t)$ from the origin to the desired value.**

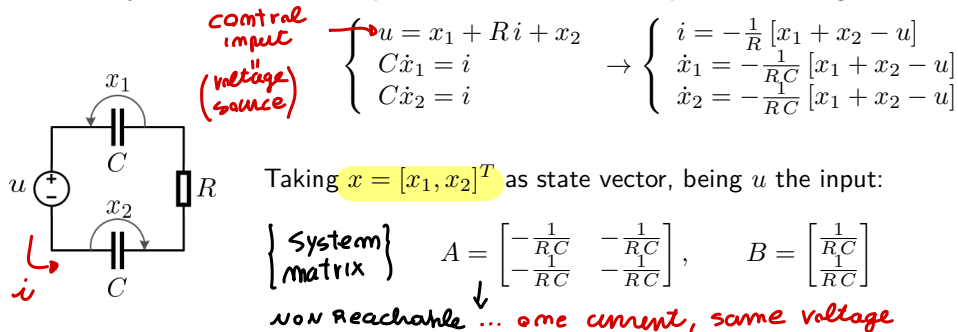


Definition

A system is said to be **fully reachable** if all its states are reachable.

two states = voltage V on capacitors

Consider a system made of two capacitors of same size, at equal initial charge.



Remark: The system is not fully reachable, since the two states x_1 and x_2 must be equal (capacitors with same size and initial states, fed by the same current i). The "target" states $\bar{x} = (\bar{x}_1, \bar{x}_2)$ are reachable if and only if $\bar{x}_1 = \bar{x}_2$.

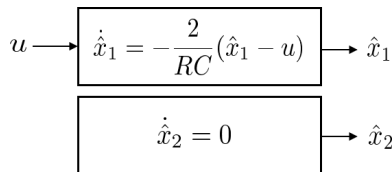
$\bar{x} = [x_1 \neq x_2]$ non reachable

↳ changing the variables

To highlight this, let's make a **change of variables**: $\hat{x}_1 = x_1 + x_2$, $\hat{x}_2 = x_1 - x_2$. Then:

$$\begin{cases} \dot{\hat{x}}_1 = \dot{x}_1 + \dot{x}_2 = -\frac{2}{RC} [\hat{x}_1 - u] \\ \dot{\hat{x}}_2 = \dot{x}_1 - \dot{x}_2 = 0 \end{cases} \rightarrow \hat{A} = \begin{bmatrix} -\frac{2}{RC} & 0 \\ 0 & 0 \end{bmatrix}, \quad \hat{B} = \begin{bmatrix} \frac{2}{RC} \\ 0 \end{bmatrix}$$

↓ In this way we can represent system: (different set of variables matrix)



By means of this change of variables, the system is decomposed into:

- ▶ A reachable part, \hat{x}_1 , function of the input u .
- ▶ An unreachable part, \hat{x}_2 , onto which no control variable is acting.

u acts only on \hat{x}_1
while \hat{x}_2 is not
reached by $u \rightarrow x_2$ is unreachable!

↳ in discrete
time \Rightarrow

Definition - Reachability (discrete-time systems)

Given the system $x(k+1) = Ax(k) + Bu(k)$, with $x(0) = 0$, a state \bar{x} is said to be **reachable in \bar{k} steps** if there exists an input sequence $u(0), \dots, u(\bar{k})$ such that $x(\bar{k}) = \bar{x}$.

Definition

A system whose states are reachable in \bar{k} steps is **fully reachable in \bar{k} steps**.

How to assess the reachability? We start from **discrete-time system** and then extend the results to continuous-time ones. **Starting from $x(0) = 0$** , one has that:

$$x(1) = \cancel{Ax(0)} + Bu(0) = Bu(0)$$

$$x(2) = Ax(1) + Bu(1) = ABu(0) + Bu(1)$$

$$x(3) = Ax(2) + Bu(2) = A^2Bu(0) + ABu(1) + Bu(2)$$

\vdots

$$x(n) = A^{n-1}Bu(0) + A^{n-2}Bu(1) + \dots + ABu(n-2) + Bu(n-1)$$

↙ group it into a matrix M_R (in reversed order)

Definition - Reachability Matrix

$$M_R = [B \quad AB \quad \dots A^{n-1}B] \quad (1)$$

Remark: $x(n) = M_R \cdot \begin{bmatrix} u(n-1) \\ \vdots \\ u(0) \end{bmatrix}$

redundant to use more than $n-1$ (pointing to $u(n-1)$)

Importance... state x in n steps we can reach $x(n)$ with this input sequence of $u(k)$ (pointing to the input vector)

⇒ analyze its properties...

→ full rank: any input sequence leads us to a state vector $x \neq 0$

Theorem 1 - Necessary and sufficient condition for reachability of linear systems

A linear system is **fully reachable** iff $\text{rank}(\mathcal{M}_R) = n$, where n is the system's order.

use until term $(m-1)$...

Remark: Why is \mathcal{M}_R constructed using $x(n)$? Suppose to consider $x(n+1)$. Then \mathcal{M}_R will contain $A^n B$ as well. But, in light of Cayley-Hamilton theorem, A^n can be written as a linear combination of I, A^1, \dots, A^{n-1} , and thus this extra term does not affect the rank of \mathcal{M}_R .

other column will be linearly depend if using more than $m-1$ terms

Remark: In discrete-time linear systems, states are reachable **at most** in n steps.

Remark - Reachability of continuous-time systems

For **continuous-time** linear systems, \mathcal{M}_R is computed as for discrete-time ones, i.e. by (1), and the same condition as Theorem 1 holds.

Reachability - Example 1 (cont'd)

(circuit with 2 capacitors)

Considering the previous example, let's compute \mathcal{M}_R (in the original coordinates). Being $n = 2$:

from A, B
of syst.. $\rightarrow \mathcal{M}_R = [B \quad AB] = \begin{bmatrix} \frac{1}{RC} & -\frac{2}{(RC)^2} \\ \frac{1}{RC} & -\frac{2}{(RC)^2} \end{bmatrix}$ \leftarrow same row...
linear dependent
 \hookrightarrow not fully reach
rank=1

Since $\text{rank}(\mathcal{M}_R) = 1 < n$, the system is not fully reachable (as previously discussed).

We have also shown that, even if the system is not fully reachable, **by means of a suitable change of variables it is possible to decompose the system into its reachable and unreachable part.**

This leads to the following theorem, which is stated for continuous-time systems, but holds for discrete-time ones as well.

Seen physically that NOT reachable...

\hookrightarrow easy to check it also analytically

*\hookrightarrow always possible
to split the
system \Rightarrow*

Theorem - Reachability decomposition

Given the system $\dot{x} = Ax + Bu$, not fully reachable, there exists a non-unique change of variables $\hat{x} = T_R x$, where $\hat{x} = [\hat{x}_r^T, \hat{x}_{nr}^T]^T$, which allows to write the system as:

$$\begin{cases} \dot{\hat{x}}_r = \hat{A}_r \hat{x}_r + \hat{A}_x \hat{x}_{nr} + \hat{B}_r u \\ \dot{\hat{x}}_{nr} = \hat{A}_{nr} \hat{x}_{nr} \end{cases} \Leftrightarrow \dot{\hat{x}} = \hat{A} \hat{x} + \hat{B} u$$

where $\hat{A} = \begin{bmatrix} \hat{A}_r & \hat{A}_x \\ \mathbf{0} & \hat{A}_{nr} \end{bmatrix}$ and $\hat{B} = \begin{bmatrix} \hat{B}_r \\ \mathbf{0} \end{bmatrix}$.

\hat{x}_r is called **reachable part** of the system, while \hat{x}_{nr} is called **unreachable part**.

In particular, if $\text{rank}(\mathcal{M}_R) = n_r$, then \hat{A}_r is $n_r \times n_r$, \hat{B}_r is $n_r \times m$, and

$$\text{rank} \left(\hat{\mathcal{M}}_R = \begin{bmatrix} \hat{B}_r & \hat{A}_r \hat{B}_r & \dots & \hat{A}_r^{n_r-1} \hat{B}_r \end{bmatrix} \right) = n_r$$

↑ not trivial to find... it is possible to find which states are reach/ unreachable...

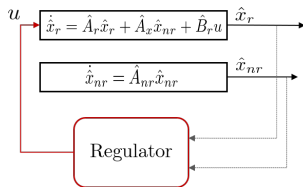
Stabilizability



more general concept than Reachability

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Usefull to know during design of Regulator
to know what we can control..



Remark: The regulator can act **only** on the reachable part.

- ▶ If the unreachable part is asymptotically stable, \hat{x}_{nr} goes to zero and its effect on \hat{x}_r vanishes.
- ▶ If the unreachable part is unstable, **nothing can be done to stabilize the system.**

Definition - Stabilizability

A system whose unreachable part is asymptotically stable is said to be **stabilizable**.

Stabilizability is a milder condition than reachability. If a system is not fully reachable, we must check that the dynamics of the unreachable part are asymptotically stable, i.e.:

- ↳
- ▶ All the eigenvalues of \hat{A}_{nr} have negative real part (continuous-time systems)
 - ▶ All the eigenvalues of \hat{A}_{nr} lie within the unit circle (discrete-time systems)

(to check stabilizability)

the unreachable part MUST be asymp stable or it will
tends to diverge and we cannot do anything.

Checking stabilizability requires to find a change of variables that allows to decompose the system into its reachable and unreachable part, which is not straightforward.

PBH reachability test

The system is fully reachable if and only if

$$P_R(s) = [sI - A \quad B]$$

has rank n for any complex value s .

only dependency on s ($sI - A$) is
how we find eig(A)

Note that the only values of s that could decrease the rank of $P_R(s)$ are the eigenvalues of A , since for those values it holds that $\det(sI - A) = 0$. only $s = \text{eig}(A)$ can reduce rank of P_R

↓ CONSEQUENCES

Remark

To assess the reachability of a system, it is enough to check that for any s_* eigenvalue of A :

$$\text{rank}(P_R(s_*)) = n$$

(see if parts of unreachable parts)

The PBH reachability test allows to derive a condition for the stabilizability of the system.

PBH stabilizability condition

If the rank of $P_R(s)$ is decreased only in correspondence of asymptotically stable eigenvalues, the system is stabilizable.

Example: Consider the system $\dot{x} = \overbrace{\begin{bmatrix} -1 & 2 \\ 2 & -1 \end{bmatrix}}^{[A]} x + \overbrace{\begin{bmatrix} 1 \\ 1 \end{bmatrix}}^{[B]} u$. Its eigenvalues are $s_1 = -3, s_2 = 1$.

If you compute \mathcal{M}_R , $\text{rank}(\mathcal{M}_R) < n$. The system is not fully reachable. Is it at least stabilizable?

$$P_R(s) = [sI - A \quad B] = \begin{bmatrix} s+1 & -2 & 1 \\ -2 & s+1 & 1 \end{bmatrix}$$

← {PBH TEST}

check the rank in correspondence of eig values: to determine which belongs to reach/unreach

$\blacktriangleright \text{rank}(P_R(s_1)) = \text{rank} \left(\begin{bmatrix} -2 & -2 & 1 \\ -2 & -2 & 1 \end{bmatrix} \right) = 1 < n$ (same row)
 $\blacktriangleright \text{rank}(P_R(s_2)) = \text{rank} \left(\begin{bmatrix} 2 & -2 & 1 \\ -2 & 2 & 1 \end{bmatrix} \right) = 2 = n$ (OK, REACH PART)

\Rightarrow

The system is stabilizable, because the eigenvalue causing the loss of rank is asymptotically stable ($s_1 = -3$). belongs to NON reach part

↪ opposite from Reach... starting from \bar{x} can we reach the origin?
 $\bar{x} \rightarrow 0$

- ▶ Reachability: bring the state from the origin to any \bar{x} .
- ▶ **Controllability**: bring the states from any initial \bar{x}_0 to the origin.

↳ Definition - Controllability (continuous-time systems)

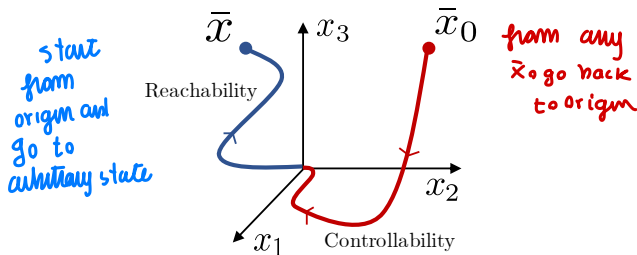
Given a continuous-time system $\dot{x}(t) = Ax(t) + Bu(t)$, a state $x(0) = \bar{x}_0$ is said **controllable** if there exists a finite arbitrary $\bar{t} > 0$, and an input profile $u(\tau)$, $\tau \in [0, \bar{t})$, such that $x(\bar{t}) = 0$.

↳ Definition - Controllability (discrete-time systems)

Given a discrete-time system $x(k+1) = Ax(k) + Bu(k)$, a state $x(0) = \bar{x}_0$ is said **controllable** in \bar{k} steps if there exists an input sequence $u(0), \dots, u(\bar{k})$ such that $x(\bar{k}) = 0$.

Definition - Full controllability

A system is said to be **fully controllable** if all its states are controllable.



Remark

- For continuous systems, the set of reachable states matches the set of controllable states.
- For discrete systems, if a state \bar{x} is reachable it is also controllable, but not viceversa!



Example: Given the system $x(k+1) = 0$:

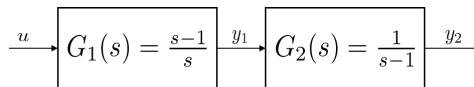
- ▶ All the states are controllable, since they are brought to 0 at $\bar{k} = 1$;
- ▶ Only the origin is reachable. → no input sequence to reach a state \bar{x}

Where does the unreachability of the system come from?

- From a problem of the model (e.g. the two-capacitors circuit)
- From a zero-pole cancellation

Example

a system sequence



{ state space }

$$G_1 : \begin{cases} \dot{x}_1 = -u \\ y_1 = x_1 + u \end{cases} \quad G_2 : \begin{cases} \dot{x}_2 = x_2 + y_1 = x_1 + x_2 + u \\ y_2 = x_2 \end{cases}$$

The state-space equation of the system is hence

$$\begin{bmatrix} \dot{x}_1 \\ \dot{x}_2 \end{bmatrix} = \begin{bmatrix} 0 & 0 \\ 1 & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} + \begin{bmatrix} -1 \\ 1 \end{bmatrix} u$$

$$A = \begin{bmatrix} 0 & 0 \\ 1 & 1 \end{bmatrix} \quad B = \begin{bmatrix} -1 \\ 1 \end{bmatrix} \quad C = [0 \quad 1]$$

Let's now check the reachability of the system: \hookrightarrow by reachability matrix

$$\mathcal{M}_R = [B \quad AB] = \begin{bmatrix} -1 & 0 \\ 1 & 0 \end{bmatrix} \Rightarrow \text{rank}(\mathcal{M}_R) = 1 < n$$

NOT FULLY Reachable

Due to the zero/pole cancellation, an unreachable part is created. Being such unreachable part associated to the unstable pole ($s = 1$), the system is not stabilizable!

*critical, because
on $s=1$ ($\text{Re}(s) > 0$)*

\hookrightarrow

Comment

Reachability is a **fundamental property** describing systems' structure and the possibility to regulate them. However, **it does not "describe" how** the system reaches the target state \bar{x} , i.e. if it exhibits overshoots, oscillations, etc.

↑ only guarantee we can reach that state

Comment

Typically, systems are characterized by the saturation of control variables. In these cases, all the previous results are **not valid**. In these cases, one shall resort to the concept of *constrained reachability*.

reach in n steps only in theory, BUT if saturation it does not holds

Definition - Observability (continuous-time system)

Consider a continuous-time autonomous linear system:

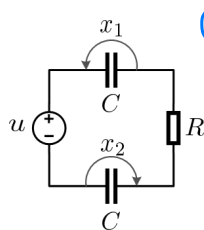
$$\begin{cases} \dot{x}(t) = Ax(t) \\ y(t) = Cx(t) \end{cases}$$

A non-null state $x(0) = \bar{x}_0$ is **non-observable** if, $\forall \bar{t} > 0$ finite, the corresponding free movement due to \bar{x}_0 , denoted by $\bar{y}(t)$, is constantly zero, i.e. $\bar{y}(\tau) = 0 \quad \forall \tau \in [0, \bar{t}]$.

A system **without non observable** states is said to be **fully observable**.

- For observable systems, there do not exist non-null states \bar{x} causing null outputs.
- If the system is not observable, there exist states \bar{x} not "showing up" on the output.
Measuring the outputs is not sufficient to know what's going on with the system.

Consider the previous system made of two capacitors of same size, at equal initial charge.



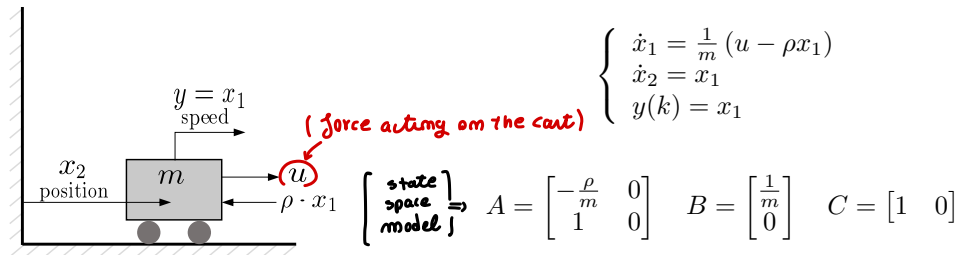
($x_1 = x_2$ bounded equals)

$$\begin{cases} i = -\frac{1}{R} [x_1 + x_2 - u] \\ \dot{x}_1 = -\frac{1}{RC} [x_1 + x_2 - u] \\ \dot{x}_2 = -\frac{1}{RC} [x_1 + x_2 - u] \end{cases}$$

It is reminded that $x_1(t) = x_2(t)$.

- Assume to measure $y_1 = x_1 + x_2$: *($\forall t$ we know our state) $\forall x_1, x_2$ combination \rightarrow output mom Q*
The **states are observable** (when x_1 and x_2 are non-null, y_1 is not-null)

- Assume instead to measure $y_2 = x_1 - x_2$: *(measuring 0 voltage)*
The states are **non-observable** (y_2 is null even when x_1 and x_2 are non-null)
 \uparrow NO information on the syst dym



Since we measure the speed, we don't have any information on the position!

- If we initialize the system in $\bar{x}_0 = [0, 1]^T$, it stays still, since $y(t) = x_1(t) = 0$.
The position does not affect the output. *so it is non obs... meas output we cannot reconstruct position*
- Measuring the output (the speed) we cannot reconstruct the position, unless we know the exact initial position).

The system is not fully observable.

Definition - Observability (discrete-time system)

Consider a discrete-time autonomous linear system:

$$\begin{cases} x(k+1) = Ax(k) \\ y(k) = Cx(k) \end{cases}$$

A non-null state $x(0) = \bar{x}_0$ is **non-observable** if, $\forall \bar{k} \in \mathbb{N}$ finite, the corresponding free movement due to \bar{x}_0 , denoted by $\bar{y}(k)$, is **constantly zero**, i.e. $\bar{y}(k) = 0 \quad \forall k \in \{0, \bar{k}\}$.

A system **without non observable states** is said to be **fully observable**.

How to assess the observability? We start from discrete-time systems, and then extend the results to continuous-time ones. Considering an initial state $x(0) = \bar{x}_0 \neq 0$, and no input:

Similarly \hookrightarrow
to reachability,
from output sequence
of evolution...

$$\bar{y}(0) = Cx(0) = C\bar{x}_0$$

$$\bar{y}(1) = C(Ax(0)) = CA\bar{x}_0$$

$$\vdots$$

$$\bar{y}(n-1) = CA^{n-1}\bar{x}_0$$

collecting the coefficients

Definition - Observability Matrix

$$\mathcal{M}_O = \begin{bmatrix} C \\ CA \\ \vdots \\ CA^{n-1} \end{bmatrix}$$

used to understand
what happens on
its kernel

(2)

Remark: $\bar{y} = \mathcal{M}_O \cdot \bar{x}_0$ \longleftrightarrow if full rank, no vector multiplied by \mathcal{M}_O give 0 \rightarrow

Since $\bar{x}_0 \neq 0$, the only way in which $\bar{y}(k)$ can be always zero, is that $\bar{x}_0 \in \text{Ker}(\mathcal{M}_O)$.

Theorem 2 - Necessary and sufficient condition for observability of linear systems

A linear system is fully observable iff $\text{rank}(\mathcal{M}_O) = n$, where n is system's order.

↗ FULL RANK

Remark - Observability of continuous-time systems

For continuous-time linear systems, \mathcal{M}_O is computed as for discrete-time ones, i.e. by (2), and the same condition as Theorem 2 holds.

Example: In the cart example, the observability matrix is not full rank, therefore the system is not fully observable:

$$\mathcal{M}_O = \begin{bmatrix} C \\ CA \end{bmatrix} = \begin{bmatrix} 1 & 0 \\ -\frac{\rho}{m} & 0 \end{bmatrix} \Rightarrow \text{rank}(\mathcal{M}_O) = 1 < n$$

↓
not fully obs

as on Reachability... obs allows to change var. and split the system into obs/non obs part

Theorem - observability decomposition

If a system is not fully observable, there exists a non-unique change of variables, $\hat{x} = T_o x$, where $\hat{x} = [\hat{x}_o^T, \hat{x}_{no}^T]^T$, which allows to write the system as:

$$\begin{cases} \dot{\hat{x}}_o = \hat{A}_o \hat{x}_o + \hat{B}_o u \\ \dot{\hat{x}}_{no} = \hat{A}_x \hat{x}_o + \hat{A}_{no} \hat{x}_{no} + \hat{B}_{no} u \\ \hat{y} = \hat{C}_o \hat{x}_o \end{cases} \Leftrightarrow \begin{cases} \dot{\hat{x}} = \hat{A} \hat{x} + \hat{B} u \\ \hat{y} = \hat{C} \hat{x} \end{cases}$$

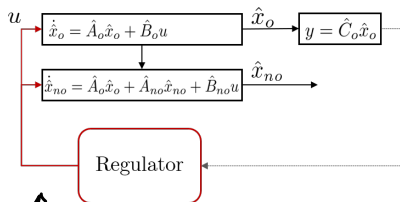
where $\hat{A} = \begin{bmatrix} \hat{A}_o & \mathbf{0} \\ \hat{A}_x & \hat{A}_{no} \end{bmatrix}$, $\hat{B} = \begin{bmatrix} \hat{B}_o \\ \hat{B}_{no} \end{bmatrix}$, and $\hat{C} = [\hat{C}_o \quad \mathbf{0}]$.

\hat{x}_o is called **observable part** of the system, while \hat{x}_{no} is called **non-observable part**.

Being $\text{rank}(\mathcal{M}_O) = \text{rank}(\hat{\mathcal{M}}_O) = n_o$, then \hat{A}_o is $n_o \times n_o$, \hat{B}_o is $n_o \times m$, \hat{C}_o is $p \times n_o$.

$$\begin{cases} \text{rank}(\mathcal{M}_O) = \text{dimension of obs part} \\ \text{while } n - \text{rank}(\mathcal{M}_O) = \text{non obs part} \end{cases}$$

Obs: to choose measurements of
syst \rightarrow we must deal with non obs part
so we define.. Detectability



↑ split the system into OBS/NON OBS
if non obs is asymp stable \rightarrow we can assume it NOT diverge

Remark: When we close the control loop, only the observable states are "accounted" by the regulator

- ▶ If the non-observable part is asymptotically stable, \hat{x}_{no} vanishes.
- ▶ If the non-observable part is unstable, **diverging states cannot be detected from outputs**. The system cannot be controlled.

Definition - Detectability

A system whose non-observable part is asymptotically stable is said to be **detectable**.

Detectability is a milder condition than observability. If a system is not fully observable, we must check that the dynamics of the non-observable part are asymptotically stable.

to check observability
detectability of syst..

PBH observability test

The system is fully observable if and only if

$$P_O(s) = \begin{bmatrix} sI - A \\ C \end{bmatrix}$$

has rank n for any complex value s .

checked $\forall s$ in
theory... but the only s that
could decrease P_O rank are $\text{eig}(A)$

As for reachability, the only values of s that *could* decrease the rank of $P_O(s)$ are the eigenvalues of A .

↓ we check rank only for $s = \text{eig}(A)$

Remark

To assess the observability of a system, it is enough to check that for any s_* , eigenvalue of A :

$$\text{rank}(P_O(s_*)) = n$$

full rank \Rightarrow fully obs

$\Leftarrow n$: can be detectable in
correspondence of asymp. st eig val

PBH detectability condition

If the rank of $P_O(s)$ is decreased only in correspondence of asymptotically stable eigenvalues, the system is detectable.

Example: Consider the system $\begin{cases} \dot{x} = \overbrace{\begin{bmatrix} -1 & -2 \\ -4 & 1 \end{bmatrix}}^A x + \overbrace{\begin{bmatrix} 1 \\ 1 \end{bmatrix}}^B u, \\ y = \underbrace{\begin{bmatrix} 1 & -1 \end{bmatrix}}_C x \end{cases}$ with eigenvalues $s_1 = -3$ (asymptotically stable), $s_2 = 3$ (unstable).

If you compute \mathcal{M}_O , $\text{rank}(\mathcal{M}_O) < n$. The system is not fully observable. Is it at least detectable?

Check for $s = s_1, s_2$ for which it can lose rank

$\text{rank}(P_O(s_1)) = \text{rank} \left(\begin{bmatrix} -2 & 2 \\ 4 & -4 \\ 1 & -1 \end{bmatrix} \right) = 1 < n$ (column m lin dep.)

$\text{rank}(P_O(s_2)) = \text{rank} \left(\begin{bmatrix} 4 & 2 \\ 4 & 2 \\ 1 & -1 \end{bmatrix} \right) = 2 = n$ (belongs to OBS PART)

$m=2 \leftarrow s_1 \text{ belongs to non obs part}$

\Rightarrow The system is detectable, because the eigenvalue causing the loss of rank is asymptotically stable ($s_1 = -3$).

not fully obs but detectable

PBH test for detectability

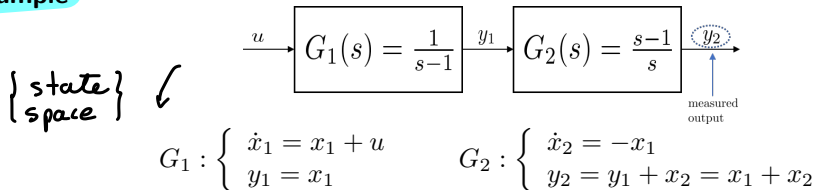
Cancellations and non-observability

↓ what causes loss of obs? → measurement issues... from ZERO/POLE cancell.

Where does the non-observability of the system come from?

- ▶ From a problem of the model (e.g. the cart with speed measurement)
- ▶ From a zero-pole cancellation

Example



The state-space equation of the system is hence

$$\begin{bmatrix} \dot{x}_1 \\ \dot{x}_2 \end{bmatrix} = \begin{bmatrix} 1 & 0 \\ -1 & 0 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} + \begin{bmatrix} 1 \\ 0 \end{bmatrix} u$$
$$y_2 = \begin{bmatrix} 1 & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix}$$

$$A = \begin{bmatrix} 1 & 0 \\ -1 & 0 \end{bmatrix} \quad B = \begin{bmatrix} 1 \\ 0 \end{bmatrix} \quad C = [1 \quad 1]$$

Let's now check both the reachability and the observability of the system:

$$\mathcal{M}_R = [B \quad AB] = \begin{bmatrix} 1 & 1 \\ 0 & -1 \end{bmatrix} \Rightarrow \text{rank}(\mathcal{M}_R) = 2 = n \quad \text{fully Reachable}$$

$$\mathcal{M}_O = \begin{bmatrix} C \\ CA \end{bmatrix} = \begin{bmatrix} 1 & 1 \\ 0 & 0 \end{bmatrix} \Rightarrow \text{rank}(\mathcal{M}_O) = 1 < n \quad \text{not fully observable}$$

Due to the zero/pole cancellation, a non-observable part is created. Being such non-observable part associated to the unstable pole ($s = 1$), the system is not detectable!

↑
cancellation on unstable part!

PROPERTIES

- Reachability \rightarrow decomposable in Reach/Unreach

- stabilizability

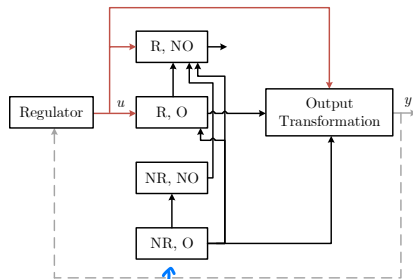
- controllability

- observability

- detectability



multiple possible decomposition...



during Regulator design, consider that properties (critical cancellation) -> unstable parts hidden

Observations:

- ▶ Any transfer function represents the reachable and observable (R, O) part only.
- ▶ If the other parts are asymptotically stable, this is not a problem. Otherwise, it is not possible to regulate the system.
- ▶ Remember that the cancellations of unstable poles and zeros are forbidden, because they create unstable unreachable/non-observable parts. \uparrow to avoid

Kalman Canonical Decomposition

For any linear system there exists a change of variable that allows to decompose the system into four parts:

- ▶ (R, O): Reachable and observable part
- ▶ (NR, NO): Unreachable and non-observable part
- ▶ (R, NO): Reachable and non-observable part
- ▶ (NR, O): Unreachable and observable part



Realization (SISO)

↪ process that allows us to go from T.F to S.S form

POLITECNICO
MILANO 1863

Given the transfer function of a SISO linear system, we want to find the underlying state-space model in its **minimal form**. (minimal ss)

↪ ss model using least number of space to represent our system

Consider a generic transfer function:

$$G(s) = \frac{\beta_n s^n + \beta_{n-1} s^{n-1} + \dots + \beta_1 s + \beta_0}{s^n + \alpha_{n-1} s^{n-1} + \dots + \alpha_1 s + \alpha_0} = \hat{\beta}_n + \underbrace{\frac{\hat{\beta}_{n-1} s^{n-1} + \dots + \hat{\beta}_1 s + \hat{\beta}_0}{s^n + \alpha_{n-1} s^{n-1} + \dots + \alpha_1 s + \alpha_0}}_{\text{Strictly proper}}$$

where $\hat{\beta}_n = \beta_n$ and $\hat{\beta}_i = \beta_i - \alpha_i \beta_n$, for $i = 0, \dots, n-1$.

On SISO: mim → can form to represent T.F $G(s)$

↪ from one canonical rep... we apply long division

↪

starting from can form... S.S representation by this MATRIX

Definition - reachability canonical form

$$A = \begin{bmatrix} 0 & 1 & 0 & \dots & 0 \\ 0 & 0 & 1 & \dots & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 0 & \dots & 1 \\ -\alpha_0 & -\alpha_1 & -\alpha_2 & \dots & -\alpha_{n-1} \end{bmatrix} \quad B = \begin{bmatrix} 0 \\ 0 \\ \vdots \\ 0 \\ 1 \end{bmatrix} \quad C = [\hat{\beta}_0 \quad \hat{\beta}_1 \quad \dots \quad \hat{\beta}_{n-1}] \quad D = \hat{\beta}_n$$

Definition - observability canonical form

$$A = \begin{bmatrix} 0 & 0 & \dots & 0 & -\alpha_0 \\ 1 & 0 & \dots & 0 & -\alpha_1 \\ 0 & 1 & \dots & 0 & -\alpha_2 \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & \dots & 1 & -\alpha_{n-1} \end{bmatrix} \quad B = \begin{bmatrix} \hat{\beta}_0 \\ \hat{\beta}_1 \\ \vdots \\ \hat{\beta}_{n-1} \end{bmatrix} \quad C = [0 \quad 0 \quad \dots \quad 0 \quad 1] \quad D = \hat{\beta}_n$$

The realization problem is significantly more complex for MIMO systems.

Given $G(s) = \begin{bmatrix} G_{11}(s) & \dots & G_{1m}(s) \\ \vdots & \ddots & \vdots \\ G_{p1}(s) & \dots & G_{pm}(s) \end{bmatrix}$ we want to find the state-space model, such that

$$G(s) = C(sI - A)^{-1}B + D$$

↙ no minimal
form easy to
find...

In case of MIMO system, an extensive use of canonical forms is required, which is out of scope.

Consider the system $\begin{bmatrix} y_1(s) \\ y_2(s) \end{bmatrix} = \begin{bmatrix} G_{11}(s) & G_{12}(s) \\ G_{21}(s) & G_{22}(s) \end{bmatrix} \begin{bmatrix} u_1(s) \\ u_2(s) \end{bmatrix}$, which can be re-written as:

$$\begin{cases} y_1(s) = \tilde{y}_{11} + \tilde{y}_{12} = G_{11}(s) u_1(s) + G_{12}(s) u_2(s) \\ y_2(s) = \tilde{y}_{21} + \tilde{y}_{22} = G_{21}(s) u_1(s) + G_{22}(s) u_2(s) \end{cases}$$

To find a **non-minimal** realization of the system, we can find the state-space realization of each SISO element of the transfer matrix $G_{ij}(s)$ separately:

↳
apply
separately
∀ SISO system

$$\begin{cases} \dot{\tilde{x}}_{ij} = A_{ij} \tilde{x}_{ij} + B_{ij} u_j \\ \tilde{y}_i = C_{ij} \tilde{x}_{ij} + D_{ij} u_j \end{cases}$$

The state-space model of the entire system is therefore:

↪ non minimal
but usable in MIMO syst.

$$\begin{bmatrix} \dot{\tilde{x}}_{11} \\ \dot{\tilde{x}}_{12} \\ \dot{\tilde{x}}_{21} \\ \dot{\tilde{x}}_{22} \end{bmatrix} = \begin{bmatrix} A_{11} & & & \\ & A_{12} & & \\ & & A_{21} & \\ & & & A_{22} \end{bmatrix} \begin{bmatrix} \tilde{x}_{11} \\ \tilde{x}_{12} \\ \tilde{x}_{21} \\ \tilde{x}_{22} \end{bmatrix} + \begin{bmatrix} B_{11} & B_{12} \\ B_{21} & B_{22} \end{bmatrix} \begin{bmatrix} u_1 \\ u_2 \end{bmatrix}$$

$$\begin{bmatrix} y_1 \\ y_2 \end{bmatrix} = \begin{bmatrix} C_{11} & C_{12} & & \\ & & C_{21} & C_{22} \end{bmatrix} \begin{bmatrix} \tilde{x}_{11} \\ \tilde{x}_{12} \\ \tilde{x}_{21} \\ \tilde{x}_{22} \end{bmatrix} + \begin{bmatrix} D_{11} & D_{12} \\ D_{21} & D_{22} \end{bmatrix} \begin{bmatrix} u_1 \\ u_2 \end{bmatrix}$$

Remark

In general, this is a non-minimal realization.

