

# Complexity of Neural Networks

Vasquez Alessandro

Université Nice-Sophia-Antipolis

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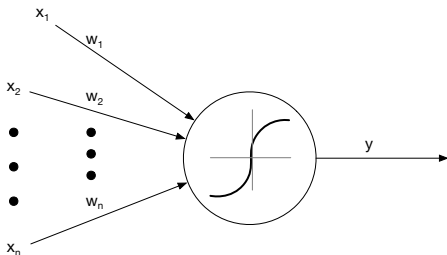
Neuron's representation :

$$y = \sigma(\sum_{j=1}^n w_j x_j - w_0)$$

Transfer functions :

- $\text{sign}(t) = \begin{cases} 0 & \text{if } t \leq 0 \\ 1 & \text{if } t > 0 \end{cases}$

- $\sigma(t) = (1 + e^{-t})^{-1}$



Neural nets can be

- cyclic or acyclic
- symmetric or asymmetric w.r.t. interconnections
- continuous or discrete w.r.t. updates
- synchronous or asynchronous w.r.t. updates in discrete neural nets

We're going to consider

- back-propagation nets
- Hopfield nets

## Type of neurons

- perceptron (with transfer function)
- threshold gate (with binary valued-input)
- majority gate (all  $w = 1$ ,  $b = \frac{|X|}{2}$ )

## Parameters for evaluating neural nets' complexity

- size
- depth
- weight

### Definition (Threshold function)

A boolean function  $t : \{0, 1\} \rightarrow 0, 1$  is a *threshold function* if it is computable by a linear threshold unit.

### Theorem

*Any threshold function on  $n$  variables can be computed by a threshold gate with integer weights  $w_i$  st :*

$$|w_i| \leq \frac{(n+1)^{(n+1)/2}}{2^n}, \forall i = 0, \dots, n.$$

This theorem defines a class of threshold gates of polynomially bounded weights.

### Theorem

*For infinitely many  $n$ , there are threshold functions on  $n$  variables whose computation by a single threshold gate requires weights as large as  $\frac{n^{n/2}}{2^n}$ .*

These threshold gates have not polynomially bounded gates which, on the other hand, are representable by  $O(s \log s)$  bits.

### Definition (Complexity classes)

We define five class of functions computable by specifically structured circuits.

- $NC^k$  : polynomial-size circuits of depth  $O(\log^k(n))$ , using bounded fan-in AND, OR, and NOT gates.
- $AC^k$  : polynomial-size circuits of depth  $O(\log^k(n))$ , using unbounded AND, OR, and NOT gates.
- $TC^k$  : threshold circuits of polynomial size and depth  $O(\log^k n)$ ,
- $TC_d^0$  threshold circuits of polynomial size and depth  $d$ ,
- $\widehat{TC}_d^0$  majority circuits of polynomial size and depth  $d$ .

### Theorem (Complexity classes' hierarchy)

$$AC^k \subseteq TC^k \subseteq NC^{k+1}.$$

$$\widehat{TC}_d^0 \subseteq TC_d^0 \subseteq \widehat{TC}_{d+1}^0.$$

We only know that  $AC^0 \subset TC^0$ , e.g. because of the majority function, and  $TC_1^0 \subset \widehat{TC}_2^0$  for the parity function.



## Definition (Threshold circuit loading)

Given  $\{(\vec{x}_1, b_1), \dots, (\vec{x}_m, b_m)\}$ , where each  $\vec{x}_i \in \{0, 1\}^n$  and each  $b_i \in \{0, 1\}$ , and a directed acyclic graph, can we find a weight assignment st given  $f$  the function computed by the resulting threshold circuit :  $f(\vec{x}_i) = b_i, \forall i = 1, \dots, m$ ?

## Definition (Threshold circuit minimization)

Given  $\{(\vec{x}_1, b_1), \dots, (\vec{x}_m, b_m)\}$ , where each  $\vec{x}_i \in \{0, 1\}^n$  and each  $b_i \in \{0, 1\}$ , and  $K \in \mathbb{N}$ , does exist a threshold circuit of at most  $K$  neurones st given  $f$  the function computed by the resulting threshold circuit :  $f(\vec{x}_i) = b_i, \forall i = 1, \dots, m$ ?

## Definition (Hopfield network)

Neural net in which

- all arcs are bidirectional
- each neuron is interconnected to all other neurons
- each neuron has a state
- each neuron's state is updated asynchronously
- each state has an "energy function" assigned :

$$E = -\frac{1}{2} \sum_{i,j} w_{ij} s_i s_j - \sum_i b y_i s_i$$

### Theorem

*In a symmetric Hopfield net with  $w_{ij}$  integer weights and  $n$  neurons, the convergence is obtain in*

$$3 \sum_{j < i} |w_{ij}| = O(n^2 \cdot \max_{i,j} |w_{ij}|)$$

*asynchronous changes of neuron states.*

### Definition (Attraction domain)

Let  $\vec{x}_i$  be a stable global state for a given Hopfield net. The *attraction domain* of  $\vec{x}_i$  is the set of all vectors which are guaranteed to converge to  $\vec{x}_i$ .

### Definition (Attraction radius)

The *attraction radius* of  $x_i$  is the largest Hamming distance from which all other vectors are guaranteed to converge to  $\vec{x}_i$ .

- Decide if a given symmetric, simple network has more than one stable state.
- Determine if a symmetric simple network converges to a given stable state from any other initial state.

- Decide if a given asymmetric net under synchronous updates will converge from any (or all) initial state.
- Computing the attraction radius of a given stable state in a symmetric simple net.

## Definition (Hopfield Net Loading Problem)

Given a set of vectors  $X = \{\vec{x}_1, \dots, \vec{x}_m\}$ , where each  $\vec{x}_i \in \{-1, 1\}^n$ , and a constant  $\rho$ , is there a symmetric net of  $n$  neurons that has  $X$  as stable states with attraction radii  $\geq \rho$ ?



The following questions are yet to be answered.

- Is the Hopfield Net Loading Problem in  $P$  or in  $NP$ ?
- Are the separations  $TC_d^0 \subseteq TC_{d+1}^0, \forall d \geq 2$  proper?

Thank you for your attention.