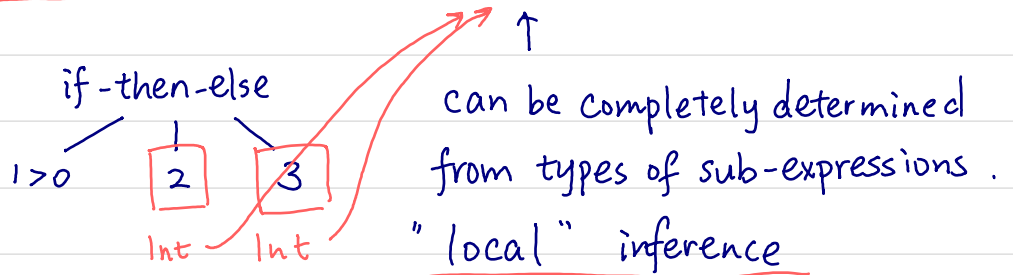


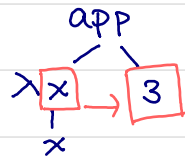
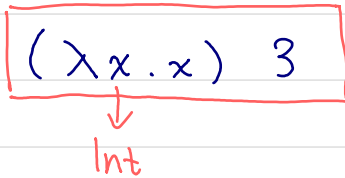
Type checking as type inference

$\Gamma \vdash \text{if } 1 > 0 \text{ then } 2 \text{ else } 3 : \text{Int}$

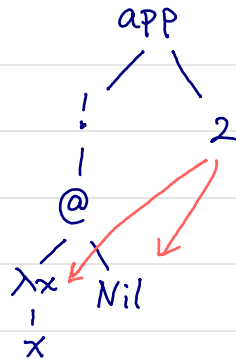


Question How to do inference if no annotation is provided?

Problem Local inference stops working.



$(!(\lambda x. x) @ \text{Nil}) \ 2$



Insights (1) Although we may not know exactly what an unknown type T is, we still have information that constrains what T can be.

(2) Constraints interact each other. If we consider all constraints at the same time, the constraint system may have a (unique) solution.

$$\left. \begin{array}{l} x + y = 0 \\ z - x = 2 \end{array} \right\} \Rightarrow \begin{array}{l} x = \dots \\ y = \dots \\ z = \dots \end{array}$$

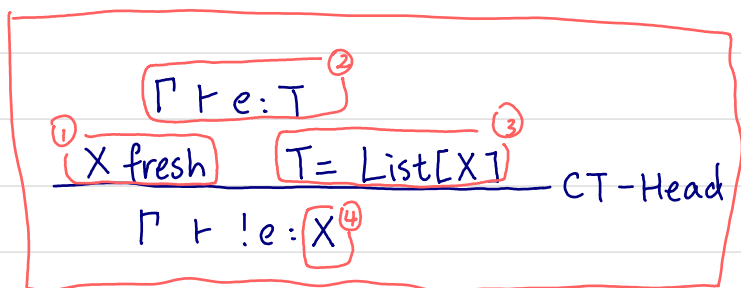
system of eqns solution

Type Inference

= constraint generation \Rightarrow constraint solving
(this section)

Constraint Generation:

Given an expr e , we can generate constraints mechanically, based on the shape of e .

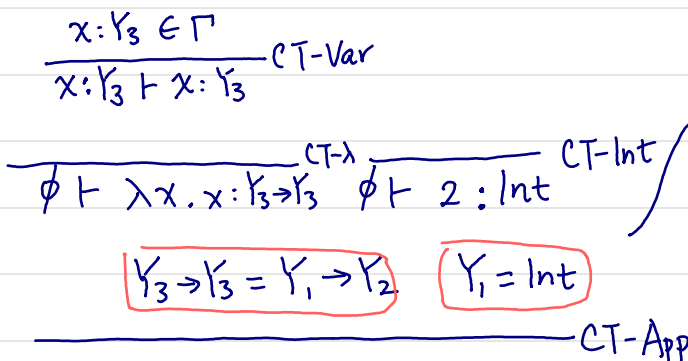
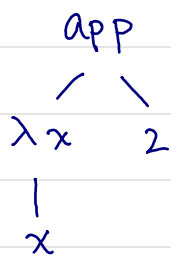


Components:

- (unknowns)
- ① A reservoir of type variables
- ② Recursion
- ③ Generated constraints
- ④ Type of the current expression

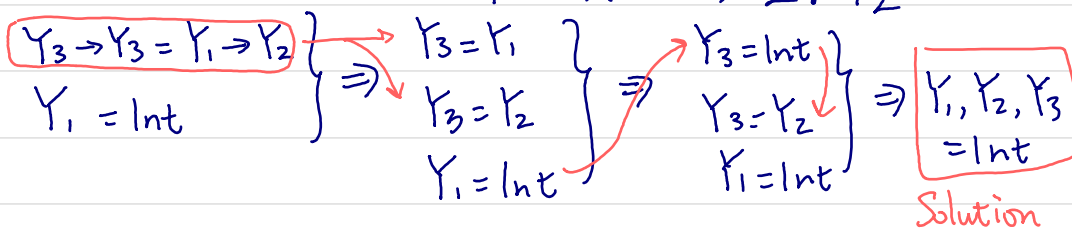
Example 1

" $(\lambda x. x) 2$ "



(Global)
Bag of Constraints

$Y_3 \rightarrow Y_3 = Y_1 \rightarrow Y_2$
$\text{Int} = Y_1$



Solution

Example 2

" $(!(\lambda x.x)@ Nil)$ 5 "

