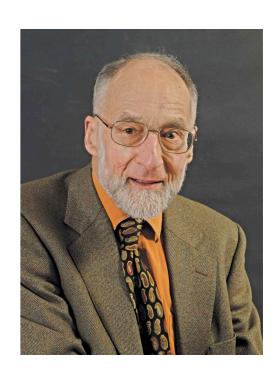
CS 162 Programming languages

Lecture 4: A crash course in OCaml

Yu Feng Winter 2020

History of ML

- ML = "Meta Language"
- Designed by Robin Milner @ Edinburgh
- Language to manipulate Theorems/Proofs
- Several dialects:
 - Standard" ML (of New Jersey)
 - French dialect with support for objects
 - State-of-the-art
 - Extensive library, tool, user support



Who is using OCaml

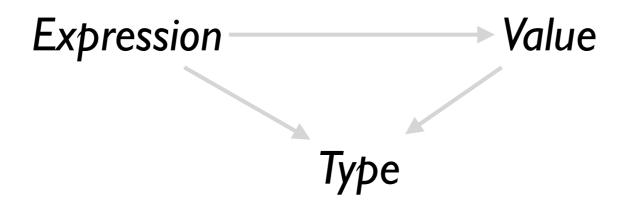




Bloomberg



ML's holy grail



- Everything is an expression
- Everything has a value
- Everything has a type

Interacting with ML

"Read-Eval-Print" Loop

Repeat:

- I. System reads expression e
- 2. System evaluates e to get value v
- 3. System prints value v and type t

What are these expressions, values and types?

Basic types

```
# 2;;
         # 2+3;;
                         Int
         # "hi";;
                                 "hi"
# "hi,"^"0Caml";;
                       String
                                "hi, OCaml"
            true;;
                                 true
         # 2>3;;
                                 false
                        Bool
```

Type errors

```
# "Hi," ^ 2;;
# (2+3) || 9;;
```

Untypable expression is rejected

- No casting or coercing
- Fancy algorithm to catch errors
- ML's single most powerful feature

Complex types: Lists

List operators:

- Cons (::): "cons" element to a list of same type
- append (@): only append two list of the same type
- Head (List.hd): return the head element of a nonempty list
- Tail (List.tl): return the tail of nonempty list

Syntax:

Lists = semicolon

Semantics:

Same type, unbounded number

Complex types: Tuples

Syntax:

• Lists = comma

Semantics:

Different type, fixed number

Variables and bindings

let
$$x = e$$

"Bind the value of expression e to the variable x"

```
# let x = 2+2;;
val x : int = 4
```

Variables and bindings

Later declared expressions can use x

Most recent "bound" value used for evaluation

```
# let x = 2+2;;
val x : int = 4
# let y = x * x * x;;
val y : int = 64
# let z = [x;y;x+y];;
val z : int list = [4;64;68]
```

Variables and bindings

Undeclared variables (i.e. without a value binding) are not accepted!

```
# let p = a + 1;
Characters 8-9:
    let p = a + 1 ;;
^ Unbound value a
```

Local bindings

for expressions using "temporary" variables

```
# let
    tempVar = x + 2 * y
    in
    tempVar * tempVar ;;
```

- tempVar is bound only inside expr body from in ...;
- Not visible ("in scope") outside

Complex types: functions

```
# let inc = fun x -> x+1;
val inc : int -> int = fn
# inc 0;
val it : int = 1
# inc 10;
val it : int = 11
```

How to evaluate a function app:

- Evaluate the argument
- Bind formal to arg value
- Evaluate the "body expr"

Complex types: functions

Wow! A function can return a function

```
# let lt = fun x -> fun y -> x < y;;
val lt : 'a -> 'a -> bool = fn
# let is5Lt = lt 5;
val is5lt : int -> bool = fn;;
# is5lt 10;;
val it : bool = true;
# is5lt 2;;
val it : bool = false;
```

Complex types: functions

A function can also take a function argument

```
# let neg = fun f -> fun x -> not (f x);
val lt : (a -> bool) -> a -> bool = fn
# let is5gte = neg is5lt;
val is5gte : int -> bool = fn
# is5gte 10;
val it : bool = false;
# is5gte 2;
val it : bool = true;
```

Put it together: a "filter" function

```
If arg matches this pattern
```

then use this body expr

Put it together: a "quicksort" function

```
# let partition f l = (filter f l, filter (neg f) l);;
val partition :('a->bool)->'a list->'a list * 'a list = fn
# let list1 = [1,31,12,4,7,2,10];
# partition is5lt list1;
val it : (int list * int list) = ([31,12,7,10],[1,2,10])
```

```
# let rec sort l =
    match l with
    [] -> []
    | (h::t) ->
    let (l,r) = partition ((<) h) t in
        (sort l)@(h::(sort r));;</pre>
```

TODOs by next lecture

- Start to work on hw1
- Hw2 will be out next week
- Come to the discussion session if you have questions
- 1st homework will be due in \sim 10 days