## Alejandro Santorum Varela – EDAT 3<sup>rd</sup> Assignment

## 21.

- a) Best key: a, c
- b) No. Becuase **d** and **b** depend on a part of the best key (primary key).
- c) R(b, o, a, q, c, d)  $\rightarrow$  R1( $\underline{c}$ ,  $\underline{a}$ , q, d) and R2( $\underline{c}$ , b, o) Now, we are descomposing R1 and R2 to satisfy the third normal form: R1  $\rightarrow$  R1.1( $\underline{a}$ , d) and R1.2( $\underline{c}$ , $\underline{a}$ , q) R2  $\rightarrow$  R2.1( $\underline{c}$ , b) and R2.2( $\underline{b}$ , o)
- d) No, because **o** depends on **b** that at the same time depends on **c**. For the same reason **q** depends on **a** and **c**, but **d** only depends on **a**.

## 22.

```
R = RESERVATION:
                                                 F = FLIGHTS
                                                 f = flightID
p = passengerName
SSN = SSN
                                                 t = time
f = flightID
                                                 d = departureAirport
                                                 a = arrivalAirport
d = departureAirport
a = arrivalAirport
                                                 dC = departureCity
t = time
                                                 aC = arrivalCity
date = date
c = cost
```

Examinating all the requirements we can assume those functional dependencies:

a) Candidates Key:

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R-Cand1 = {SSN, f, date}
F-Cand1 = {flightID}
    note: {time, d} also describes uniquely a row but it is not minimal.
R-PK = {SSN, f, date}
F-PK={flightID}
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b)

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\begin{array}{ll} \underline{R = RESERVATION} \colon & \underline{F = FLIGHTS} \\ SSN \to P & d \to dC \\ \{SSN, f, date\} \to c & a \to aC \\ \{f, date\} \to t, a, d & flightID \to \{t, d, a, dC, aC\} \\ \{t, date, d\} \to f, a & \{time, d\} \to \{flightID, a, dC, aC\} \end{array}
```

c) R and F satisfy first NF, but they don't satisfy  $2^{nd}$  and  $3^{rd}$ .

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d)
R = \{\underline{SSN}, p, \underline{f}, d, a, t, \underline{date}, c\}
R.1 = \{\underline{SSN}, p\}
R.2 = \{\underline{SSN}, \underline{f}, d, a, t, \underline{date}, c\}
R2.1 = \{\underline{SSN}, \underline{f}, \underline{date}, c\}
R2.2 = \{\underline{f}, \underline{date}, t, a, d\}
F = \{\underline{f}, t, d, a, dC, aC\}
F1 = \{\underline{f}, t, d, a\}
F2 = \{\underline{a}, aC\}
F3 = \{\underline{d}, dC\}
```

## 23.

Due to the fact all relations have to have a Primary Key, let's suppose three cases:

 $A \rightarrow B$ : A is the determinant and at the same time it's the Primary Key (which is one of the possible Candidates Keys). So it satisfies BCNF.

 $B \rightarrow A$ : B is the determinant and at the same time it's the Primary Key (which is one of the possible Candidates Keys). So it satisfies BCNF.

 $AB \rightarrow AB$ : Now AB is the determinant and at the same time it's the Primary Key (which is one of the possible Candidates Keys). So it satisfies BCNF.